
An introduction to recursion and induction

A recursive datatype: toy lists

```
datatype 'a list = Nil | Cons 'a ('a list)
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Nil: empty list

Cons x xs: head $x :: 'a$, tail $xs :: 'a\ list$

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Cons x xs: head $x :: 'a$, tail $xs :: 'a \text{ list}$

A toy list: *Cons False (Cons True Nil)*

A recursive datatype: toy lists

datatype 'a list = Nil | Cons 'a ('a list)

Nil: empty list

Cons x xs: head $x :: 'a$, tail $xs :: 'a \text{ list}$

A toy list: *Cons False (Cons True Nil)*

Predefined lists: *[False, True]*

Structural induction on lists

P xs holds for all lists xs if

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$P\ xs$ holds for all lists xs if

- $P\ Nil$

Structural induction on lists

$P\ xs$ holds for all lists xs if

- $P\ Nil$
- and for arbitrary x and xs , $P\ xs$ implies $P\ (Cons\ x\ xs)$

A recursive function: append

Definition by *primitive recursion*:

primrec $app :: 'a\ list \Rightarrow 'a\ list \Rightarrow 'a\ list$ **where**
 $app\ Nil\ ys = ?$ |
 $app\ (Cons\ x\ xs)\ ys = ??$

A recursive function: append

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1 rule per constructor

Recursive calls must drop the constructor \implies Termination

Concrete syntax

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Except for single identifiers, e.g. `'a`

... normally not shown on slides

Demo: append and reverse

Proofs

General schema:

```
lemma name : " . . . "  
apply ( . . . )  
apply ( . . . )  
⋮  
done
```

If the lemma is suitable as a simplification rule:

```
lemma name[simp] : " . . . "
```

Proof methods

- **Structural induction**
 - Format: *(induct x)*
x must be a free variable in the first subgoal.
The type of x must be a datatype.
 - Effect: generates 1 new subgoal per constructor
- **Simplification and a bit of logic**
 - Format: *auto*
 - Effect: tries to solve as many subgoals as possible using simplification and basic logical reasoning.

Top down proofs

sorry

“completes” any proof.

Suitable for top down developments:

Assume lemmas first, prove them later.

Disproving

quickcheck

tries to find counterexample by random testing