Isabelle/CTT — Constructive Type Theory with extensional equality and without universes

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Contents

theory *CTT*
imports *Pure*
begin

1 Constructive Type Theory: axiomatic basis

ML-file ‹~~/*src*/*Provers*/*typedsimp.ML*›
setup *Pure-Thy.old-appl-syntax-setup*

typedecl *i*
typedecl *t*
typedecl *o*

consts
— Types
F :: *t*
T :: *t* — *F* is empty, *T* contains one element
contr :: *i⇒i*
tt :: *i*
— Natural numbers
N :: *t*
succ :: *i⇒i*
rec :: [*i*, *i*, [*i*,*i*]⇒*i*] ⇒ *i*
— Unions
inl :: *i⇒i*
inr :: *i⇒i*
when :: [*i*, *i⇒i*, *i⇒i*]⇒*i*
— General Sum and Binary Product
Sum :: [*t*, *i⇒t*]⇒*t*
fst :: *i⇒i*
snd :: *i⇒i*
split :: [*i*, [*i*,*i*]⇒*i*] ⇒*i*
— General Product and Function Space
Prod :: [*t*, *i⇒t*]⇒*t*
— Types
Plus :: [*t*,*t*]⇒*t* (**infixr** + *40*)
— Equality type
Eq :: [*t*,*i*,*i*]⇒*t*
eq :: *i*
— Judgements
Type :: *t* ⇒ *prop* ((- *type*) [*10*] *5*)
Eqtype :: [*t*,*t*]⇒*prop* ((- =/ -) [*10*,*10*] *5*)
Elem :: [*i*, *t*]⇒*prop* ((- /: -) [*10*,*10*] *5*)
Eqelem :: [*i*,*i*,*t*]⇒*prop* ((- =/ - :/ -) [*10*,*10*,*10*] *5*)
Reduce :: [*i*,*i*]⇒*prop* (*Reduce*[-,-])

— Types

— Functions
lambda :: $(i \Rightarrow i) \Rightarrow i$ (**binder** $\boldsymbol{\lambda}$ *10*)
app :: $[i,i] \Rightarrow i$ (**infixl** ‘ *60*)
— Natural numbers
Zero :: i (*0*)
— Pairing
pair :: $[i,i] \Rightarrow i$ ((*1*<-,/->))

syntax
-PROD :: $[idt,t,t] \Rightarrow t$ ((*3*$\prod$ -:-./ -) *10*)
-SUM :: $[idt,t,t] \Rightarrow t$ ((*3*$\sum$ -:-./ -) *10*)
translations
$\prod x{:}A.\ B \rightleftharpoons CONST\ Prod(A, \lambda x.\ B)$
$\sum x{:}A.\ B \rightleftharpoons CONST\ Sum(A, \lambda x.\ B)$

abbreviation *Arrow* :: $[t,t] \Rightarrow t$ (**infixr** $\longrightarrow$ *30*)
where $A \longrightarrow B \equiv \prod$ -:$A.\ B$

abbreviation *Times* :: $[t,t] \Rightarrow t$ (**infixr** $\times$ *50*)
where $A \times B \equiv \sum$ -:$A.\ B$

Reduction: a weaker notion than equality; a hack for simplification. *Reduce*[*a*,*b*] means either that $a = b : A$ for some A or else that a and b are textually identical.

Does not verify $a{:}A$! Sound because only *trans-red* uses a *Reduce* premise. No new theorems can be proved about the standard judgements.

axiomatization
where
refl-red: $\bigwedge a.\ Reduce[a,a]$ **and**
red-if-equal: $\bigwedge a\ b\ A.\ a = b : A \Longrightarrow Reduce[a,b]$ **and**
trans-red: $\bigwedge a\ b\ c\ A.\ [\![a = b : A;\ Reduce[b,c]]\!] \Longrightarrow a = c : A$ **and**

— Reflexivity

refl-type: $\bigwedge A.\ A\ type \Longrightarrow A = A$ **and**
refl-elem: $\bigwedge a\ A.\ a : A \Longrightarrow a = a : A$ **and**

— Symmetry

sym-type: $\bigwedge A\ B.\ A = B \Longrightarrow B = A$ **and**
sym-elem: $\bigwedge a\ b\ A.\ a = b : A \Longrightarrow b = a : A$ **and**

— Transitivity

trans-type: $\bigwedge A\ B\ C.\ [\![A = B;\ B = C]\!] \Longrightarrow A = C$ **and**
trans-elem: $\bigwedge a\ b\ c\ A.\ [\![a = b : A;\ b = c : A]\!] \Longrightarrow a = c : A$ **and**

equal-types: $\bigwedge a\ A\ B.\ [\![a : A;\ A = B]\!] \Longrightarrow a : B$ **and**
equal-typesL: $\bigwedge a\ b\ A\ B.\ [\![a = b : A;\ A = B]\!] \Longrightarrow a = b : B$ **and**

— Substitution

subst-type: $\bigwedge a\ A\ B.\ [\![a : A;\ \bigwedge z.\ z{:}A \Longrightarrow B(z)\ type]\!] \Longrightarrow B(a)\ type$ **and**
subst-typeL: $\bigwedge a\ c\ A\ B\ D.\ [\![a = c : A;\ \bigwedge z.\ z{:}A \Longrightarrow B(z) = D(z)]\!] \Longrightarrow B(a) = D(c)$ **and**

subst-elem: $\bigwedge a\ b\ A\ B.\ [\![a : A;\ \bigwedge z.\ z{:}A \Longrightarrow b(z){:}B(z)]\!] \Longrightarrow b(a){:}B(a)$ **and**
subst-elemL:
$\bigwedge a\ b\ c\ d\ A\ B.\ [\![a = c : A;\ \bigwedge z.\ z{:}A \Longrightarrow b(z){=}d(z) : B(z)]\!] \Longrightarrow b(a){=}d(c) : B(a)$ **and**

— The type N – natural numbers

NF: $N\ type$ **and**
NI0: $0 : N$ **and**
NI-succ: $\bigwedge a.\ a : N \Longrightarrow succ(a) : N$ **and**
NI-succL: $\bigwedge a\ b.\ a = b : N \Longrightarrow succ(a) = succ(b) : N$ **and**

NE:
$\bigwedge p\ a\ b\ C.\ [\![p{:}\ N;\ a{:}\ C(0);\ \bigwedge u\ v.\ [\![u{:}\ N;\ v{:}\ C(u)]\!] \Longrightarrow b(u,v){:}\ C(succ(u))]\!]$
$\Longrightarrow rec(p,\ a,\ \lambda u\ v.\ b(u,v)) : C(p)$ **and**

NEL:
$\bigwedge p\ q\ a\ b\ c\ d\ C.\ [\![p = q : N;\ a = c : C(0);$
$\bigwedge u\ v.\ [\![u{:}\ N;\ v{:}\ C(u)]\!] \Longrightarrow b(u,v) = d(u,v){:}\ C(succ(u))]\!]$
$\Longrightarrow rec(p,\ a,\ \lambda u\ v.\ b(u,v)) = rec(q,c,d) : C(p)$ **and**

NC0:
$\bigwedge a\ b\ C.\ [\![a{:}\ C(0);\ \bigwedge u\ v.\ [\![u{:}\ N;\ v{:}\ C(u)]\!] \Longrightarrow b(u,v){:}\ C(succ(u))]\!]$
$\Longrightarrow rec(0,\ a,\ \lambda u\ v.\ b(u,v)) = a : C(0)$ **and**

NC-succ:
$\bigwedge p\ a\ b\ C.\ [\![p{:}\ N;\ \ a{:}\ C(0);\ \bigwedge u\ v.\ [\![u{:}\ N;\ v{:}\ C(u)]\!] \Longrightarrow b(u,v){:}\ C(succ(u))]\!] \Longrightarrow$
$rec(succ(p),\ a,\ \lambda u\ v.\ b(u,v)) = b(p,\ rec(p,\ a,\ \lambda u\ v.\ b(u,v))) : C(succ(p))$ **and**

— The fourth Peano axiom. See page 91 of Martin-Löf's book.
zero-ne-succ: $\bigwedge a.\ [\![a{:}\ N;\ 0 = succ(a) : N]\!] \Longrightarrow 0{:}\ F$ **and**

— The Product of a family of types

ProdF: $\bigwedge A\ B.\ [\![A\ type;\ \bigwedge x.\ x{:}A \Longrightarrow B(x)\ type]\!] \Longrightarrow \prod x{:}A.\ B(x)\ type$ **and**

ProdFL:
$\bigwedge A\ B\ C\ D.\ [\![A = C;\ \bigwedge x.\ x{:}A \Longrightarrow B(x) = D(x)]\!] \Longrightarrow \prod x{:}A.\ B(x) = \prod x{:}C.$

$D(x)$ **and**

ProdI:
$\bigwedge b\ A\ B.\ [\![A\ type;\ \bigwedge x.\ x{:}A \Longrightarrow b(x){:}B(x)]\!] \Longrightarrow \boldsymbol{\lambda} x.\ b(x) : \prod x{:}A.\ B(x)$ **and**

ProdIL: $\bigwedge b\ c\ A\ B.\ [\![A\ type;\ \bigwedge x.\ x{:}A \Longrightarrow b(x) = c(x) : B(x)]\!] \Longrightarrow$
$\boldsymbol{\lambda} x.\ b(x) = \boldsymbol{\lambda} x.\ c(x) : \prod x{:}A.\ B(x)$ **and**

ProdE: $\bigwedge p\ a\ A\ B.\ [\![p : \prod x{:}A.\ B(x);\ a : A]\!] \Longrightarrow p`a : B(a)$ **and**
ProdEL: $\bigwedge p\ q\ a\ b\ A\ B.\ [\![p = q{:}\ \prod x{:}A.\ B(x);\ a = b : A]\!] \Longrightarrow p`a = q`b : B(a)$ **and**

ProdC: $\bigwedge a\ b\ A\ B.\ [\![a : A;\ \bigwedge x.\ x{:}A \Longrightarrow b(x) : B(x)]\!] \Longrightarrow (\boldsymbol{\lambda} x.\ b(x))\ `\ a = b(a) : B(a)$ **and**

ProdC2: $\bigwedge p\ A\ B.\ p : \prod x{:}A.\ B(x) \Longrightarrow (\boldsymbol{\lambda} x.\ p`x) = p : \prod x{:}A.\ B(x)$ **and**

— The Sum of a family of types

SumF: $\bigwedge A\ B.\ [\![A\ type;\ \bigwedge x.\ x{:}A \Longrightarrow B(x)\ type]\!] \Longrightarrow \sum x{:}A.\ B(x)\ type$ **and**
SumFL: $\bigwedge A\ B\ C\ D.\ [\![A = C;\ \bigwedge x.\ x{:}A \Longrightarrow B(x) = D(x)]\!] \Longrightarrow \sum x{:}A.\ B(x) = \sum x{:}C.\ D(x)$ **and**

SumI: $\bigwedge a\ b\ A\ B.\ [\![a : A;\ b : B(a)]\!] \Longrightarrow <a,b> : \sum x{:}A.\ B(x)$ **and**
SumIL: $\bigwedge a\ b\ c\ d\ A\ B.\ [\![\ a = c : A;\ b = d : B(a)]\!] \Longrightarrow <a,b> = <c,d> : \sum x{:}A.\ B(x)$ **and**

SumE: $\bigwedge p\ c\ A\ B\ C.\ [\![p{:}\ \sum x{:}A.\ B(x);\ \bigwedge x\ y.\ [\![x{:}A;\ y{:}B(x)]\!] \Longrightarrow c(x,y){:}\ C(<x,y>)]\!]$
$\Longrightarrow split(p,\ \lambda x\ y.\ c(x,y)) : C(p)$ **and**

SumEL: $\bigwedge p\ q\ c\ d\ A\ B\ C.\ [\![p = q : \sum x{:}A.\ B(x);$
$\bigwedge x\ y.\ [\![x{:}A;\ y{:}B(x)]\!] \Longrightarrow c(x,y){=}d(x,y){:}\ C(<x,y>)]\!]$
$\Longrightarrow split(p,\ \lambda x\ y.\ c(x,y)) = split(q,\ \lambda x\ y.\ d(x,y)) : C(p)$ **and**

SumC: $\bigwedge a\ b\ c\ A\ B\ C.\ [\![a{:}\ A;\ \ b{:}\ B(a);\ \bigwedge x\ y.\ [\![x{:}A;\ y{:}B(x)]\!] \Longrightarrow c(x,y){:}\ C(<x,y>)]\!]$
$\Longrightarrow split(<a,b>,\ \lambda x\ y.\ c(x,y)) = c(a,b) : C(<a,b>)$ **and**

fst-def: $\bigwedge a.\ fst(a) \equiv split(a,\ \lambda x\ y.\ x)$ **and**
snd-def: $\bigwedge a.\ snd(a) \equiv split(a,\ \lambda x\ y.\ y)$ **and**

— The sum of two types

PlusF: $\bigwedge A\ B.\ [\![A\ type;\ B\ type]\!] \Longrightarrow A{+}B\ type$ **and**
PlusFL: $\bigwedge A\ B\ C\ D.\ [\![A = C;\ B = D]\!] \Longrightarrow A{+}B = C{+}D$ **and**

PlusI-inl: $\bigwedge a\ A\ B.\ [\![a : A;\ B\ type]\!] \Longrightarrow inl(a) : A{+}B$ **and**
PlusI-inlL: $\bigwedge a\ c\ A\ B.\ [\![a = c : A;\ B\ type]\!] \Longrightarrow inl(a) = inl(c) : A{+}B$ **and**

PlusI-inr: $\bigwedge b\ A\ B.\ [\![A\ type;\ b : B]\!] \Longrightarrow inr(b) : A+B$ **and**
PlusI-inrL: $\bigwedge b\ d\ A\ B.\ [\![A\ type;\ b = d : B]\!] \Longrightarrow inr(b) = inr(d) : A+B$ **and**

PlusE:
$\bigwedge p\ c\ d\ A\ B\ C.\ [\![p{:}\ A+B;$
$\bigwedge x.\ x{:}A \Longrightarrow c(x){:}\ C(inl(x));$
$\bigwedge y.\ y{:}B \Longrightarrow d(y){:}\ C(inr(y))\ ]\!] \Longrightarrow when(p,\ \lambda x.\ c(x),\ \lambda y.\ d(y)) : C(p)$ **and**

PlusEL:
$\bigwedge p\ q\ c\ d\ e\ f\ A\ B\ C.\ [\![p = q : A+B;$
$\bigwedge x.\ x{:}\ A \Longrightarrow c(x) = e(x) : C(inl(x));$
$\bigwedge y.\ y{:}\ B \Longrightarrow d(y) = f(y) : C(inr(y))]\!]$
$\Longrightarrow when(p,\ \lambda x.\ c(x),\ \lambda y.\ d(y)) = when(q,\ \lambda x.\ e(x),\ \lambda y.\ f(y)) : C(p)$ **and**

PlusC-inl:
$\bigwedge a\ c\ d\ A\ B\ C.\ [\![a{:}\ A;$
$\bigwedge x.\ x{:}A \Longrightarrow c(x){:}\ C(inl(x));$
$\bigwedge y.\ y{:}B \Longrightarrow d(y){:}\ C(inr(y))\ ]\!]$
$\Longrightarrow when(inl(a),\ \lambda x.\ c(x),\ \lambda y.\ d(y)) = c(a) : C(inl(a))$ **and**

PlusC-inr:
$\bigwedge b\ c\ d\ A\ B\ C.\ [\![b{:}\ B;$
$\bigwedge x.\ x{:}A \Longrightarrow c(x){:}\ C(inl(x));$
$\bigwedge y.\ y{:}B \Longrightarrow d(y){:}\ C(inr(y))]\!]$
$\Longrightarrow when(inr(b),\ \lambda x.\ c(x),\ \lambda y.\ d(y)) = d(b) : C(inr(b))$ **and**

— The type *Eq*

EqF: $\bigwedge a\ b\ A.\ [\![A\ type;\ a : A;\ b : A]\!] \Longrightarrow Eq(A,a,b)\ type$ **and**
EqFL: $\bigwedge a\ b\ c\ d\ A\ B.\ [\![A = B;\ a = c : A;\ b = d : A]\!] \Longrightarrow Eq(A,a,b) = Eq(B,c,d)$
and
EqI: $\bigwedge a\ b\ A.\ a = b : A \Longrightarrow eq : Eq(A,a,b)$ **and**
EqE: $\bigwedge p\ a\ b\ A.\ p : Eq(A,a,b) \Longrightarrow a = b : A$ **and**

— By equality of types, can prove $C(p)$ from $C(eq)$, an elimination rule
EqC: $\bigwedge p\ a\ b\ A.\ p : Eq(A,a,b) \Longrightarrow p = eq : Eq(A,a,b)$ **and**

— The type *F*

FF: *F type* **and**
FE: $\bigwedge p\ C.\ [\![p{:}\ F;\ C\ type]\!] \Longrightarrow contr(p) : C$ **and**
FEL: $\bigwedge p\ q\ C.\ [\![p = q : F;\ C\ type]\!] \Longrightarrow contr(p) = contr(q) : C$ **and**

— The type T
— Martin-Löf's book (page 68) discusses elimination and computation. Elim-

ination can be derived by computation and equality of types, but with an extra premise $C(x)$ type $x{:}T$. Also computation can be derived from elimination.

TF: T *type* **and**
TI: $tt : T$ **and**
TE: $\bigwedge p\ c\ C.\ [\![p : T;\ c : C(tt)]\!] \Longrightarrow c : C(p)$ **and**
TEL: $\bigwedge p\ q\ c\ d\ C.\ [\![p = q : T;\ c = d : C(tt)]\!] \Longrightarrow c = d : C(p)$ **and**
TC: $\bigwedge p.\ p : T \Longrightarrow p = tt : T$

1.1 Tactics and derived rules for Constructive Type Theory

Formation rules.

lemmas *form-rls = NF ProdF SumF PlusF EqF FF TF*
and *formL-rls = ProdFL SumFL PlusFL EqFL*

Introduction rules. OMITTED:

- *EqI*, because its premise is an *eqelem*, not an *elem*.

lemmas *intr-rls = NI0 NI-succ ProdI SumI PlusI-inl PlusI-inr TI*
and *intrL-rls = NI-succL ProdIL SumIL PlusI-inlL PlusI-inrL*

Elimination rules. OMITTED:

- *EqE*, because its conclusion is an *eqelem*, not an *elem*
- *TE*, because it does not involve a constructor.

lemmas *elim-rls = NE ProdE SumE PlusE FE*
and *elimL-rls = NEL ProdEL SumEL PlusEL FEL*

OMITTED: *eqC* are *TC* because they make rewriting loop: $p = un = un = \ldots$

lemmas *comp-rls = NC0 NC-succ ProdC SumC PlusC-inl PlusC-inr*

Rules with conclusion $a{:}A$, an elem judgement.

lemmas *element-rls = intr-rls elim-rls*

Definitions are (meta)equality axioms.

lemmas *basic-defs = fst-def snd-def*

Compare with standard version: B is applied to UNSIMPLIFIED expression!

lemma *SumIL2*: $[\![c = a : A;\ d = b : B(a)]\!] \Longrightarrow <c,d> = <a,b> : Sum(A,B)$
by (*rule sym-elem*) (*rule SumIL*; *rule sym-elem*)

lemmas *intrL2-rls = NI-succL ProdIL SumIL2 PlusI-inlL PlusI-inrL*

Exploit p:*Prod*(A,B) to create the assumption z:$B(a)$. A more natural form of product elimination.

lemma *subst-prodE*:
assumes p: *Prod*(A,B)
and a: A
and $\bigwedge z.\ z$: $B(a) \Longrightarrow c(z)$: $C(z)$
shows $c(p`a)$: $C(p`a)$
by (*rule assms ProdE*)+

1.2 Tactics for type checking

ML ‹
local

fun is-rigid-elem ***Const-*** ‹*Elem for a* -› = *not*(*is-Var* (*head-of a*))
| *is-rigid-elem* ***Const-*** ‹*Eqelem for a* - -› = *not*(*is-Var* (*head-of a*))
| *is-rigid-elem* ***Const-*** ‹*Type for a*› = *not*(*is-Var* (*head-of a*))
| *is-rigid-elem* - = *false*

in

(∗*Try solving a*:*A or a*=*b*:*A by assumption provided a is rigid*!∗)
fun test-assume-tac ctxt = *SUBGOAL* (*fn* (*prem*, *i*) =>
if is-rigid-elem (*Logic.strip-assums-concl prem*)
then assume-tac ctxt i else no-tac)

fun ASSUME ctxt tf i = *test-assume-tac ctxt i* *ORELSE* *tf i*

end
›

For simplification: type formation and checking, but no equalities between terms.

lemmas *routine-rls* = *form-rls formL-rls refl-type element-rls*

ML ‹
fun routine-tac rls ctxt prems =
ASSUME ctxt (*filt-resolve-from-net-tac ctxt 4* (*Tactic.build-net* (*prems* @ *rls*)));

(∗*Solve all subgoals A type using formation rules.* ∗)
val form-net = *Tactic.build-net* @{*thms form-rls*};
fun form-tac ctxt =
REPEAT-FIRST (*ASSUME ctxt* (*filt-resolve-from-net-tac ctxt 1 form-net*));

(∗*Type checking*: *solve a*:*A* (*a rigid*, *A flexible*) *by intro and elim rules.* ∗)
fun typechk-tac ctxt thms =
let val tac =
filt-resolve-from-net-tac ctxt 3
(*Tactic.build-net* (*thms* @ @{*thms form-rls*} @ @{*thms element-rls*}))

```
  in  REPEAT-FIRST (ASSUME ctxt tac)  end

(*Solve a:A (a flexible, A rigid) by introduction rules.
  Cannot use stringtrees (filt-resolve-tac) since
  goals like ?a:SUM(A,B) have a trivial head−string *)
fun intr-tac ctxt thms =
  let val tac =
    filt-resolve-from-net-tac ctxt 1
      (Tactic.build-net (thms @ @{thms form-rls} @ @{thms intr-rls}))
  in  REPEAT-FIRST (ASSUME ctxt tac)  end

(*Equality proving: solve a=b:A (where a is rigid) by long rules. *)
fun equal-tac ctxt thms =
  REPEAT-FIRST
    (ASSUME ctxt
      (filt-resolve-from-net-tac ctxt 3
        (Tactic.build-net (thms @ @{thms form-rls element-rls intrL-rls elimL-rls
refl-elem}))))
›

method-setup form = ‹Scan.succeed (fn ctxt => SIMPLE-METHOD (form-tac
ctxt))›
method-setup typechk = ‹Attrib.thms >> (fn ths => fn ctxt => SIMPLE-METHOD
(typechk-tac ctxt ths))›
method-setup intr = ‹Attrib.thms >> (fn ths => fn ctxt => SIMPLE-METHOD
(intr-tac ctxt ths))›
method-setup equal = ‹Attrib.thms >> (fn ths => fn ctxt => SIMPLE-METHOD
(equal-tac ctxt ths))›
```

1.3 Simplification

To simplify the type in a goal.

```
lemma replace-type: ⟦B = A; a : A⟧ ⟹ a : B
  apply (rule equal-types)
   apply (rule-tac [2] sym-type)
   apply assumption+
  done
```

Simplify the parameter of a unary type operator.

```
lemma subst-eqtyparg:
  assumes 1: a=c : A
    and 2: ⋀z. z:A ⟹ B(z) type
  shows B(a) = B(c)
  apply (rule subst-typeL)
   apply (rule-tac [2] refl-type)
   apply (rule 1)
  apply (erule 2)
  done
```

Simplification rules for Constructive Type Theory.

lemmas *reduction-rls = comp-rls [THEN trans-elem]*

ML ‹
```
(*Converts each goal e : Eq(A,a,b) into a=b:A for simplification.
  Uses other intro rules to avoid changing flexible goals.*)
val eqintr-net = Tactic.build-net @{thms EqI intr-rls}
fun eqintr-tac ctxt =
  REPEAT-FIRST (ASSUME ctxt (filt-resolve-from-net-tac ctxt 1 eqintr-net))

(** Tactics that instantiate CTT-rules.
    Vars in the given terms will be incremented!
    The (rtac EqE i) lets them apply to equality judgements. **)

fun NE-tac ctxt sp i =
  TRY (resolve-tac ctxt @{thms EqE} i) THEN
  Rule-Insts.res-inst-tac ctxt [(((p, 0), Position.none), sp)] [] @{thm NE} i

fun SumE-tac ctxt sp i =
  TRY (resolve-tac ctxt @{thms EqE} i) THEN
  Rule-Insts.res-inst-tac ctxt [(((p, 0), Position.none), sp)] [] @{thm SumE} i

fun PlusE-tac ctxt sp i =
  TRY (resolve-tac ctxt @{thms EqE} i) THEN
  Rule-Insts.res-inst-tac ctxt [(((p, 0), Position.none), sp)] [] @{thm PlusE} i

(** Predicate logic reasoning, WITH THINNING!!  Procedures adapted from NJ.
**)

(*Finds f:Prod(A,B) and a:A in the assumptions, concludes there is z:B(a) *)
fun add-mp-tac ctxt i =
  resolve-tac ctxt @{thms subst-prodE} i  THEN  assume-tac ctxt i  THEN  as-
sume-tac ctxt i

(*Finds P⟶Q and P in the assumptions, replaces implication by Q *)
fun mp-tac ctxt i = eresolve-tac ctxt @{thms subst-prodE} i  THEN  assume-tac
ctxt i

(*safe when regarded as predicate calculus rules*)
val safe-brls = sort (make-ord lessb)
    [ (true, @{thm FE}), (true,asm-rl),
      (false, @{thm ProdI}), (true, @{thm SumE}), (true, @{thm PlusE}) ]

val unsafe-brls =
    [ (false, @{thm PlusI-inl}), (false, @{thm PlusI-inr}), (false, @{thm SumI}),
      (true, @{thm subst-prodE}) ]

(*0 subgoals vs 1 or more*)
val (safe0-brls, safep-brls) =
```

List.partition (curry (op =) 0 o subgoals-of-brl) safe-brls

fun safestep-tac ctxt thms i =
form-tac ctxt ORELSE
resolve-tac ctxt thms i ORELSE
biresolve-tac ctxt safe0-brls i ORELSE mp-tac ctxt i ORELSE
DETERM (biresolve-tac ctxt safep-brls i)

fun safe-tac ctxt thms i = DEPTH-SOLVE-1 (safestep-tac ctxt thms i)

fun step-tac ctxt thms = safestep-tac ctxt thms ORELSE' biresolve-tac ctxt unsafe-brls

(∗*Fails unless it solves the goal*!∗)
fun pc-tac ctxt thms = DEPTH-SOLVE-1 o (step-tac ctxt thms)
›

method-setup *eqintr* = ‹*Scan.succeed (SIMPLE-METHOD o eqintr-tac)*›
method-setup *NE* = ‹
Scan.lift Parse.embedded-inner-syntax >> (fn s => fn ctxt => SIMPLE-METHOD' (NE-tac ctxt s))
›
method-setup *pc* = ‹*Attrib.thms >> (fn ths => fn ctxt => SIMPLE-METHOD' (pc-tac ctxt ths))*›
method-setup *add-mp* = ‹*Scan.succeed (SIMPLE-METHOD' o add-mp-tac)*›

ML-file ‹*rew.ML*›
method-setup *rew* = ‹*Attrib.thms >> (fn ths => fn ctxt => SIMPLE-METHOD (rew-tac ctxt ths))*›
method-setup *hyp-rew* = ‹*Attrib.thms >> (fn ths => fn ctxt => SIMPLE-METHOD (hyp-rew-tac ctxt ths))*›

1.4 The elimination rules for fst/snd

lemma *SumE-fst*: *p* : *Sum*(*A*,*B*) ⟹ *fst*(*p*) : *A*
apply (*unfold basic-defs*)
apply (*erule SumE*)
apply *assumption*
done

The first premise must be *p*:*Sum*(*A*,*B*)!!.

lemma *SumE-snd*:
assumes *major*: *p*: *Sum*(*A*,*B*)
and *A type*
and ⋀*x*. *x*:*A* ⟹ *B*(*x*) *type*
shows *snd*(*p*) : *B*(*fst*(*p*))
apply (*unfold basic-defs*)
apply (*rule major* [*THEN SumE*])
apply (*rule SumC* [*THEN subst-eqtyparg*, *THEN replace-type*])

apply (*typechk assms*)
done

2 The two-element type (booleans and conditionals)

definition *Bool* :: *t*
where *Bool* ≡ *T*+*T*

definition *true* :: *i*
where *true* ≡ *inl*(*tt*)

definition *false* :: *i*
where *false* ≡ *inr*(*tt*)

definition *cond* :: [*i*,*i*,*i*]⇒*i*
where *cond*(*a*,*b*,*c*) ≡ *when*(*a*, λ-. *b*, λ-. *c*)

lemmas *bool-defs* = *Bool-def true-def false-def cond-def*

2.1 Derivation of rules for the type *Bool*

Formation rule.

lemma *boolF*: *Bool type*
unfolding *bool-defs* **by** *typechk*

Introduction rules for *true*, *false*.

lemma *boolI-true*: *true* : *Bool*
unfolding *bool-defs* **by** *typechk*

lemma *boolI-false*: *false* : *Bool*
unfolding *bool-defs* **by** *typechk*

Elimination rule: typing of *cond*.

lemma *boolE*: ⟦*p*:*Bool*; *a* : *C*(*true*); *b* : *C*(*false*)⟧ ⟹ *cond*(*p*,*a*,*b*) : *C*(*p*)
unfolding *bool-defs*
apply (*typechk*; *erule TE*)
apply *typechk*
done

lemma *boolEL*: ⟦*p* = *q* : *Bool*; *a* = *c* : *C*(*true*); *b* = *d* : *C*(*false*)⟧
⟹ *cond*(*p*,*a*,*b*) = *cond*(*q*,*c*,*d*) : *C*(*p*)
unfolding *bool-defs*
apply (*rule PlusEL*)
apply (*erule asm-rl refl-elem* [*THEN TEL*])+
done

Computation rules for *true*, *false*.

lemma *boolC-true*: $[\![a : C(true);\ b : C(false)]\!] \Longrightarrow cond(true,a,b) = a : C(true)$
unfolding *bool-defs*
apply (*rule comp-rls*)
apply *typechk*
apply (*erule-tac* [!] *TE*)
apply *typechk*
done

lemma *boolC-false*: $[\![a : C(true);\ b : C(false)]\!] \Longrightarrow cond(false,a,b) = b : C(false)$
unfolding *bool-defs*
apply (*rule comp-rls*)
apply *typechk*
apply (*erule-tac* [!] *TE*)
apply *typechk*
done

3 Elementary arithmetic

3.1 Arithmetic operators and their definitions

definition *add* :: $[i,i]\Rightarrow i$ (**infixr** $\#+$ *65*)
where $a\#+b \equiv rec(a,\ b,\ \lambda u\ v.\ succ(v))$

definition *diff* :: $[i,i]\Rightarrow i$ (**infixr** $-$ *65*)
where $a-b \equiv rec(b,\ a,\ \lambda u\ v.\ rec(v,\ 0,\ \lambda x\ y.\ x))$

definition *absdiff* :: $[i,i]\Rightarrow i$ (**infixr** $|-|$ *65*)
where $a|-|b \equiv (a-b)\ \#+\ (b-a)$

definition *mult* :: $[i,i]\Rightarrow i$ (**infixr** $\#*$ *70*)
where $a\#*b \equiv rec(a,\ 0,\ \lambda u\ v.\ b\ \#+\ v)$

definition *mod* :: $[i,i]\Rightarrow i$ (**infixr** *mod* *70*)
where $a\ mod\ b \equiv rec(a,\ 0,\ \lambda u\ v.\ rec(succ(v)\ |-|\ b,\ 0,\ \lambda x\ y.\ succ(v)))$

definition *div* :: $[i,i]\Rightarrow i$ (**infixr** *div* *70*)
where $a\ div\ b \equiv rec(a,\ 0,\ \lambda u\ v.\ rec(succ(u)\ mod\ b,\ succ(v),\ \lambda x\ y.\ v))$

lemmas *arith-defs* = *add-def diff-def absdiff-def mult-def mod-def div-def*

3.2 Proofs about elementary arithmetic: addition, multiplication, etc.

3.2.1 Addition

Typing of *add*: short and long versions.

lemma *add-typing*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow a\ \#+\ b : N$
unfolding *arith-defs* **by** *typechk*

lemma *add-typingL*: $[\![a = c{:}N;\ b = d{:}N]\!] \Longrightarrow a \#+ b = c \#+ d : N$
unfolding *arith-defs* **by** *equal*

Computation for *add*: 0 and successor cases.

lemma *addC0*: $b{:}N \Longrightarrow 0 \#+ b = b : N$
unfolding *arith-defs* **by** *rew*

lemma *addC-succ*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow succ(a) \#+ b = succ(a \#+ b) : N$
unfolding *arith-defs* **by** *rew*

3.2.2 Multiplication

Typing of *mult*: short and long versions.

lemma *mult-typing*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow a \#* b : N$
unfolding *arith-defs* **by** (*typechk add-typing*)

lemma *mult-typingL*: $[\![a = c{:}N;\ b = d{:}N]\!] \Longrightarrow a \#* b = c \#* d : N$
unfolding *arith-defs* **by** (*equal add-typingL*)

Computation for *mult*: 0 and successor cases.

lemma *multC0*: $b{:}N \Longrightarrow 0 \#* b = 0 : N$
unfolding *arith-defs* **by** *rew*

lemma *multC-succ*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow succ(a) \#* b = b \#+ (a \#* b) : N$
unfolding *arith-defs* **by** *rew*

3.2.3 Difference

Typing of difference.

lemma *diff-typing*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow a - b : N$
unfolding *arith-defs* **by** *typechk*

lemma *diff-typingL*: $[\![a = c{:}N;\ b = d{:}N]\!] \Longrightarrow a - b = c - d : N$
unfolding *arith-defs* **by** *equal*

Computation for difference: 0 and successor cases.

lemma *diffC0*: $a{:}N \Longrightarrow a - 0 = a : N$
unfolding *arith-defs* **by** *rew*

Note: $rec(a,\ 0,\ \lambda z\ w.z)$ is $pred(a)$.

lemma *diff-0-eq-0*: $b{:}N \Longrightarrow 0 - b = 0 : N$
unfolding *arith-defs*
apply (*NE b*)
apply *hyp-rew*
done

Essential to simplify FIRST!! (Else we get a critical pair) $succ(a) - succ(b)$ rewrites to $pred(succ(a) - b)$.

lemma *diff-succ-succ*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow succ(a) - succ(b) = a - b : N$
unfolding *arith-defs*
apply *hyp-rew*
apply (*NE b*)
apply *hyp-rew*
done

3.3 Simplification

lemmas *arith-typing-rls* = *add-typing mult-typing diff-typing*
and *arith-congr-rls* = *add-typingL mult-typingL diff-typingL*

lemmas *congr-rls* = *arith-congr-rls intrL2-rls elimL-rls*

lemmas *arithC-rls* =
addC0 addC-succ
multC0 multC-succ
diffC0 diff-0-eq-0 diff-succ-succ

ML ‹
structure Arith-simp = *TSimpFun*(
val refl = @{*thm refl-elem*}
val sym = @{*thm sym-elem*}
val trans = @{*thm trans-elem*}
val refl-red = @{*thm refl-red*}
val trans-red = @{*thm trans-red*}
val red-if-equal = @{*thm red-if-equal*}
val default-rls = @{*thms arithC-rls comp-rls*}
val routine-tac = *routine-tac* @{*thms arith-typing-rls routine-rls*}
)

fun arith-rew-tac ctxt prems =
make-rew-tac ctxt (*Arith-simp.norm-tac ctxt* (@{*thms congr-rls*}, *prems*))

fun hyp-arith-rew-tac ctxt prems =
make-rew-tac ctxt
(*Arith-simp.cond-norm-tac ctxt* (*prove-cond-tac ctxt*, @{*thms congr-rls*}, *prems*))
›

method-setup *arith-rew* = ‹
Attrib.thms >> (*fn ths* => *fn ctxt* => *SIMPLE-METHOD* (*arith-rew-tac ctxt ths*))
›

method-setup *hyp-arith-rew* = ‹
Attrib.thms >> (*fn ths* => *fn ctxt* => *SIMPLE-METHOD* (*hyp-arith-rew-tac ctxt ths*))
›

3.4 Addition

Associative law for addition.

lemma *add-assoc*: $[\![a{:}N;\ b{:}N;\ c{:}N]\!] \Longrightarrow (a \#+ b) \#+ c = a \#+ (b \#+ c) : N$
 apply (*NE a*)
  apply *hyp-arith-rew*
 done

Commutative law for addition. Can be proved using three inductions. Must simplify after first induction! Orientation of rewrites is delicate.

lemma *add-commute*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow a \#+ b = b \#+ a : N$
 apply (*NE a*)
  apply *hyp-arith-rew*
 apply (*rule sym-elem*)
 prefer *2*
 apply (*NE b*)
  prefer *4*
  apply (*NE b*)
   apply *hyp-arith-rew*
 done

3.5 Multiplication

Right annihilation in product.

lemma *mult-0-right*: $a{:}N \Longrightarrow a \#* 0 = 0 : N$
 apply (*NE a*)
  apply *hyp-arith-rew*
 done

Right successor law for multiplication.

lemma *mult-succ-right*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow a \#* succ(b) = a \#+ (a \#* b) : N$
 apply (*NE a*)
  apply (*hyp-arith-rew add-assoc* [*THEN sym-elem*])
apply (*assumption* | *rule add-commute mult-typingL add-typingL intrL-rls refl-elem*)+
 done

Commutative law for multiplication.

lemma *mult-commute*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow a \#* b = b \#* a : N$
 apply (*NE a*)
  apply (*hyp-arith-rew mult-0-right mult-succ-right*)
 done

Addition distributes over multiplication.

lemma *add-mult-distrib*: $[\![a{:}N;\ b{:}N;\ c{:}N]\!] \Longrightarrow (a \#+ b) \#* c = (a \#* c) \#+ (b \#* c) : N$
 apply (*NE a*)
  apply (*hyp-arith-rew add-assoc* [*THEN sym-elem*])

done

Associative law for multiplication.

lemma *mult-assoc*: $[\![a{:}N;\ b{:}N;\ c{:}N]\!] \Longrightarrow (a \#* b) \#* c = a \#* (b \#* c) : N$
apply (*NE a*)
apply (*hyp-arith-rew add-mult-distrib*)
done

3.6 Difference

Difference on natural numbers, without negative numbers

- $a - b = 0$ iff $a \leq b$
- $a - b = succ(c)$ iff $a > b$

lemma *diff-self-eq-0*: $a{:}N \Longrightarrow a - a = 0 : N$
apply (*NE a*)
apply *hyp-arith-rew*
done

lemma *add-0-right*: $[\![c : N;\ 0 : N;\ c : N]\!] \Longrightarrow c \#+ 0 = c : N$
by (*rule addC0* [*THEN* [*3*] *add-commute* [*THEN trans-elem*]])

Addition is the inverse of subtraction: if $b \leq x$ then $b \#+ (x - b) = x$. An example of induction over a quantified formula (a product). Uses rewriting with a quantified, implicative inductive hypothesis.

schematic-goal *add-diff-inverse-lemma*:
$b{:}N \Longrightarrow ?a : \prod x{:}N.\ Eq(N,\ b-x,\ 0) \longrightarrow Eq(N,\ b \#+ (x-b),\ x)$
apply (*NE b*)
— strip one "universal quantifier" but not the "implication"
apply (*rule-tac* [*3*] *intr-rls*)
— case analysis on x in $succ(u) \leq x \longrightarrow succ(u) \#+ (x - succ(u)) = x$
prefer *4*
apply (*NE x*)
apply *assumption*
— Prepare for simplification of types – the antecedent $succ(u) \leq x$
apply (*rule-tac* [*2*] *replace-type*)
apply (*rule-tac* [*1*] *replace-type*)
apply *arith-rew*
— Solves first 0 goal, simplifies others. Two sugbgoals remain. Both follow by rewriting, (2) using quantified induction hyp.
apply *intr* — strips remaining $\prod$s
apply (*hyp-arith-rew add-0-right*)
apply *assumption*
done

Version of above with premise $b - a = 0$ i.e. $a \geq b$. Using *ProdE* does not work – for *?B*(*?a*) is ambiguous. Instead, *add-diff-inverse-lemma* states the desired induction scheme; the use of *THEN* below instantiates Vars in *ProdE* automatically.

lemma *add-diff-inverse*: $[\![a{:}N;\ b{:}N;\ b - a = 0 : N]\!] \Longrightarrow b \#+ (a-b) = a : N$
apply (*rule EqE*)
apply (*rule add-diff-inverse-lemma* [*THEN ProdE*, *THEN ProdE*])
apply (*assumption* | *rule EqI*)+
done

3.7 Absolute difference

Typing of absolute difference: short and long versions.

lemma *absdiff-typing*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow a \mid\!-\!\mid b : N$
unfolding *arith-defs* **by** *typechk*

lemma *absdiff-typingL*: $[\![a = c{:}N;\ b = d{:}N]\!] \Longrightarrow a \mid\!-\!\mid b = c \mid\!-\!\mid d : N$
unfolding *arith-defs* **by** *equal*

lemma *absdiff-self-eq-0*: $a{:}N \Longrightarrow a \mid\!-\!\mid a = 0 : N$
unfolding *absdiff-def* **by** (*arith-rew diff-self-eq-0*)

lemma *absdiffC0*: $a{:}N \Longrightarrow 0 \mid\!-\!\mid a = a : N$
unfolding *absdiff-def* **by** *hyp-arith-rew*

lemma *absdiff-succ-succ*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow succ(a) \mid\!-\!\mid succ(b) = a \mid\!-\!\mid b : N$
unfolding *absdiff-def* **by** *hyp-arith-rew*

Note how easy using commutative laws can be? ...not always...

lemma *absdiff-commute*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow a \mid\!-\!\mid b = b \mid\!-\!\mid a : N$
unfolding *absdiff-def*
apply (*rule add-commute*)
apply (*typechk diff-typing*)
done

If $a + b = 0$ then $a = 0$. Surprisingly tedious.

schematic-goal *add-eq0-lemma*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow ?c : Eq(N,a\#+b,0) \longrightarrow Eq(N,a,0)$
apply (*NE a*)
apply (*rule-tac* [*3*] *replace-type*)
apply *arith-rew*
apply *intr* — strips remaining $\prod$s
apply (*rule-tac* [*2*] *zero-ne-succ* [*THEN FE*])
apply (*erule-tac* [*3*] *EqE* [*THEN sym-elem*])
apply (*typechk add-typing*)
done

Version of above with the premise $a + b = 0$. Again, resolution instantiates variables in *ProdE*.

lemma *add-eq0*: $[\![a{:}N;\ b{:}N;\ a\ \#+\ b = 0 : N]\!] \Longrightarrow a = 0 : N$
  apply (*rule EqE*)
  apply (*rule add-eq0-lemma* [*THEN ProdE*])
    apply (*rule-tac* [*3*] *EqI*)
    apply *typechk*
  done

Here is a lemma to infer $a - b = 0$ and $b - a = 0$ from $a \mid\!-\!\mid b = 0$, below.

schematic-goal *absdiff-eq0-lem*:
  $[\![a{:}N;\ b{:}N;\ a \mid\!-\!\mid b = 0 : N]\!] \Longrightarrow ?a : Eq(N,\ a{-}b,\ 0) \times Eq(N,\ b{-}a,\ 0)$
  apply (*unfold absdiff-def*)
  apply *intr*
   apply *eqintr*
   apply (*rule-tac* [*2*] *add-eq0*)
     apply (*rule add-eq0*)
       apply (*rule-tac* [*6*] *add-commute* [*THEN trans-elem*])
         apply (*typechk diff-typing*)
  done

If $a \mid\!-\!\mid b = 0$ then $a = b$ proof: $a - b = 0$ and $b - a = 0$, so $b = a + (b - a) = a + 0 = a$.

lemma *absdiff-eq0*: $[\![a \mid\!-\!\mid b = 0 : N;\ a{:}N;\ b{:}N]\!] \Longrightarrow a = b : N$
  apply (*rule EqE*)
  apply (*rule absdiff-eq0-lem* [*THEN SumE*])
    apply *eqintr*
  apply (*rule add-diff-inverse* [*THEN sym-elem*, *THEN trans-elem*])
    apply (*erule-tac* [*3*] *EqE*)
   apply (*hyp-arith-rew add-0-right*)
  done

3.8 Remainder and Quotient

Typing of remainder: short and long versions.

lemma *mod-typing*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow a\ mod\ b : N$
  unfolding *mod-def* **by** (*typechk absdiff-typing*)

lemma *mod-typingL*: $[\![a = c{:}N;\ b = d{:}N]\!] \Longrightarrow a\ mod\ b = c\ mod\ d : N$
  unfolding *mod-def* **by** (*equal absdiff-typingL*)

Computation for *mod*: 0 and successor cases.

lemma *modC0*: $b{:}N \Longrightarrow 0\ mod\ b = 0 : N$
  unfolding *mod-def* **by** (*rew absdiff-typing*)

lemma *modC-succ*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow$
  $succ(a)\ mod\ b = rec(succ(a\ mod\ b) \mid\!-\!\mid b,\ 0,\ \lambda x\ y.\ succ(a\ mod\ b)) : N$
  unfolding *mod-def* **by** (*rew absdiff-typing*)

Typing of quotient: short and long versions.

lemma *div-typing*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow a\ div\ b : N$
unfolding *div-def* **by** (*typechk absdiff-typing mod-typing*)

lemma *div-typingL*: $[\![a = c{:}N;\ b = d{:}N]\!] \Longrightarrow a\ div\ b = c\ div\ d : N$
unfolding *div-def* **by** (*equal absdiff-typingL mod-typingL*)

lemmas *div-typing-rls* = *mod-typing div-typing absdiff-typing*

Computation for quotient: 0 and successor cases.

lemma *divC0*: $b{:}N \Longrightarrow 0\ div\ b = 0 : N$
unfolding *div-def* **by** (*rew mod-typing absdiff-typing*)

lemma *divC-succ*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow$
$succ(a)\ div\ b = rec(succ(a)\ mod\ b,\ succ(a\ div\ b),\ \lambda x\ y.\ a\ div\ b) : N$
unfolding *div-def* **by** (*rew mod-typing*)

Version of above with same condition as the *mod* one.

lemma *divC-succ2*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow$
$succ(a)\ div\ b = rec(succ(a\ mod\ b)\ |-|\ b,\ succ(a\ div\ b),\ \lambda x\ y.\ a\ div\ b) : N$
apply (*rule divC-succ* [*THEN trans-elem*])
apply (*rew div-typing-rls modC-succ*)
apply (*NE succ* (*a mod b*) |−|*b*)
apply (*rew mod-typing div-typing absdiff-typing*)
done

For case analysis on whether a number is 0 or a successor.

lemma *iszero-decidable*: $a{:}N \Longrightarrow rec(a,\ inl(eq),\ \lambda ka\ kb.\ inr(<ka,\ eq>)) :$
$Eq(N,a,0) + (\sum x{:}N.\ Eq(N,a,\ succ(x)))$
apply (*NE a*)
apply (*rule-tac* [*3*] *PlusI-inr*)
apply (*rule-tac* [*2*] *PlusI-inl*)
apply *eqintr*
apply *equal*
done

Main Result. Holds when b is 0 since $a\ mod\ 0 = a$ and $a\ div\ 0 = 0$.

lemma *mod-div-equality*: $[\![a{:}N;\ b{:}N]\!] \Longrightarrow a\ mod\ b\ \#+\ (a\ div\ b)\ \#*\ b = a : N$
apply (*NE a*)
apply (*arith-rew div-typing-rls modC0 modC-succ divC0 divC-succ2*)
apply (*rule EqE*)
— case analysis on $succ(u\ mod\ b)\ |-|\ b$
apply (*rule-tac a1* = *succ* (*u mod b*) |−| *b* **in** *iszero-decidable* [*THEN PlusE*])
apply (*erule-tac* [*3*] *SumE*)
apply (*hyp-arith-rew div-typing-rls modC0 modC-succ divC0 divC-succ2*)
— Replace one occurrence of b by $succ(u\ mod\ b)$. Clumsy!
apply (*rule add-typingL* [*THEN trans-elem*])
apply (*erule EqE* [*THEN absdiff-eq0*, *THEN sym-elem*])
apply (*rule-tac* [*3*] *refl-elem*)

apply (*hyp-arith-rew div-typing-rls*)
done

end

4 Easy examples: type checking and type deduction

theory *Typechecking*
imports *../CTT*
begin

4.1 Single-step proofs: verifying that a type is well-formed

schematic-goal *?A type*
apply (*rule form-rls*)
done

schematic-goal *?A type*
apply (*rule form-rls*)
back
apply (*rule form-rls*)
apply (*rule form-rls*)
done

schematic-goal $\prod z{:}?A\ .\ N + ?B(z)$ *type*
apply (*rule form-rls*)
apply (*rule form-rls*)
apply (*rule form-rls*)
apply (*rule form-rls*)
apply (*rule form-rls*)
done

4.2 Multi-step proofs: Type inference

lemma $\prod w{:}N.\ N + N$ *type*
apply *form*
done

schematic-goal $<0,\ succ(0)> : ?A$
apply *intr*
done

schematic-goal $\prod w{:}N\ .\ Eq(?A,w,w)$ *type*
apply *typechk*
done

schematic-goal $\prod x{:}N\ .\ \prod y{:}N\ .\ Eq(?A,x,y)$ *type*
apply *typechk*

done

typechecking an application of fst

schematic-goal (**λ***u*. *split*(*u*, *λv w*. *v*)) ‘ <*0*, *succ*(*0*)> : *?A*
apply *typechk*
done

typechecking the predecessor function

schematic-goal **λ***n*. *rec*(*n*, *0*, *λx y*. *x*) : *?A*
apply *typechk*
done

typechecking the addition function

schematic-goal **λ***n*. **λ***m*. *rec*(*n*, *m*, *λx y*. *succ*(*y*)) : *?A*
apply *typechk*
done

method-setup *N* =
‹*Scan.succeed* (*fn ctxt* => *SIMPLE-METHOD* (*TRYALL* (*resolve-tac ctxt* @{*thms NF*}))))›

schematic-goal **λ***w*. <*w*,*w*> : *?A*
apply *typechk*
apply *N*
done

schematic-goal **λ***x*. **λ***y*. *x* : *?A*
apply *typechk*
apply *N*
done

typechecking fst (as a function object)

schematic-goal **λ***i*. *split*(*i*, *λj k*. *j*) : *?A*
apply *typechk*
apply *N*
done

end

5 Examples with elimination rules

theory *Elimination*
imports *../CTT*
begin

This finds the functions fst and snd!

schematic-goal [*folded basic-defs*]: *A type* $\Longrightarrow$ *?a* : $(A \times A) \longrightarrow A$

apply *pc*
done

schematic-goal [*folded basic-defs*]: A *type* $\Longrightarrow$ *?a* : $(A \times A) \longrightarrow A$
apply *pc*
back
done

Double negation of the Excluded Middle

schematic-goal A *type* $\Longrightarrow$ *?a* : $((A + (A \longrightarrow F)) \longrightarrow F) \longrightarrow F$
apply *intr*
apply (*rule ProdE*)
apply *assumption*
apply *pc*
done

schematic-goal $[\![A$ *type*; B *type*$]\!] \Longrightarrow$ *?a* : $(A \times B) \longrightarrow (B \times A)$
apply *pc*
done

Binary sums and products

schematic-goal $[\![A$ *type*; B *type*; C *type*$]\!] \Longrightarrow$ *?a* : $(A + B \longrightarrow C) \longrightarrow (A \longrightarrow C) \times (B \longrightarrow C)$
apply *pc*
done

schematic-goal $[\![A$ *type*; B *type*; C *type*$]\!] \Longrightarrow$ *?a* : $A \times (B + C) \longrightarrow (A \times B + A \times C)$
apply *pc*
done

schematic-goal
assumes A *type*
and $\bigwedge x.\ x{:}A \Longrightarrow B(x)$ *type*
and $\bigwedge x.\ x{:}A \Longrightarrow C(x)$ *type*
shows *?a* : $(\sum x{:}A.\ B(x) + C(x)) \longrightarrow (\sum x{:}A.\ B(x)) + (\sum x{:}A.\ C(x))$
apply (*pc assms*)
done

Construction of the currying functional

schematic-goal $[\![A$ *type*; B *type*; C *type*$]\!] \Longrightarrow$ *?a* : $(A \times B \longrightarrow C) \longrightarrow (A \longrightarrow (B \longrightarrow C))$
apply *pc*
done

schematic-goal

assumes A *type*
and $\bigwedge x.\ x{:}A \Longrightarrow B(x)$ *type*
and $\bigwedge z.\ z{:}\ (\sum x{:}A.\ B(x)) \Longrightarrow C(z)$ *type*
shows *?a* : $\prod f{:}\ (\prod z : (\sum x{:}A\ .\ B(x))\ .\ C(z)).$
$(\prod x{:}A\ .\ \prod y{:}B(x)\ .\ C(<x,y>))$
apply (*pc assms*)
done

Martin-Löf (1984), page 48: axiom of sum-elimination (uncurry)

schematic-goal $[\![A$ *type*; B *type*; C *type*$]\!] \Longrightarrow$ *?a* : $(A \longrightarrow (B \longrightarrow C)) \longrightarrow (A \times B \longrightarrow C)$
apply *pc*
done

schematic-goal
assumes A *type*
and $\bigwedge x.\ x{:}A \Longrightarrow B(x)$ *type*
and $\bigwedge z.\ z{:}\ (\sum x{:}A\ .\ B(x)) \Longrightarrow C(z)$ *type*
shows *?a* : $(\prod x{:}A\ .\ \prod y{:}B(x)\ .\ C(<x,y>))$
$\longrightarrow (\prod z : (\sum x{:}A\ .\ B(x))\ .\ C(z))$
apply (*pc assms*)
done

Function application

schematic-goal $[\![A$ *type*; B *type*$]\!] \Longrightarrow$ *?a* : $((A \longrightarrow B) \times A) \longrightarrow B$
apply *pc*
done

Basic test of quantifier reasoning

schematic-goal
assumes A *type*
and B *type*
and $\bigwedge x\ y.\ [\![x{:}A;\ y{:}B]\!] \Longrightarrow C(x,y)$ *type*
shows
?a : $(\sum y{:}B\ .\ \prod x{:}A\ .\ C(x,y))$
$\longrightarrow (\prod x{:}A\ .\ \sum y{:}B\ .\ C(x,y))$
apply (*pc assms*)
done

Martin-Löf (1984) pages 36-7: the combinator S

schematic-goal
assumes A *type*
and $\bigwedge x.\ x{:}A \Longrightarrow B(x)$ *type*
and $\bigwedge x\ y.\ [\![x{:}A;\ y{:}B(x)]\!] \Longrightarrow C(x,y)$ *type*
shows *?a* : $(\prod x{:}A.\ \prod y{:}B(x).\ C(x,y))$
$\longrightarrow (\prod f{:}\ (\prod x{:}A.\ B(x)).\ \prod x{:}A.\ C(x,\ f`x))$
apply (*pc assms*)

done

Martin-Löf (1984) page 58: the axiom of disjunction elimination

schematic-goal
  assumes *A type*
    and *B type*
    and $\bigwedge z.\ z{:}\ A+B \Longrightarrow C(z)\ type$
  shows *?a* : $(\prod x{:}A.\ C(inl(x))) \longrightarrow (\prod y{:}B.\ C(inr(y)))$
        $\longrightarrow (\prod z{:}\ A+B.\ C(z))$
apply (*pc assms*)
done

schematic-goal [*folded basic-defs*]:
  $[\![A\ type;\ B\ type;\ C\ type]\!] \Longrightarrow ?a : (A \longrightarrow B \times C) \longrightarrow (A \longrightarrow B) \times (A \longrightarrow C)$
apply *pc*
done

AXIOM OF CHOICE! Delicate use of elimination rules

schematic-goal
  assumes *A type*
    and $\bigwedge x.\ x{:}A \Longrightarrow B(x)\ type$
    and $\bigwedge x\ y.\ [\![x{:}A;\ y{:}B(x)]\!] \Longrightarrow C(x,y)\ type$
  shows *?a* : $(\prod x{:}A.\ \sum y{:}B(x).\ C(x,y)) \longrightarrow (\sum f{:}\ (\prod x{:}A.\ B(x)).\ \prod x{:}A.\ C(x,$
f`x))
apply (*intr assms*)
prefer *2* **apply** *add-mp*
prefer *2* **apply** *add-mp*
apply (*erule SumE-fst*)
apply (*rule replace-type*)
apply (*rule subst-eqtyparg*)
apply (*rule comp-rls*)
apply (*rule-tac* [*4*] *SumE-snd*)
apply (*typechk SumE-fst assms*)
done

Axiom of choice. Proof without fst, snd. Harder still!

schematic-goal [*folded basic-defs*]:
  assumes *A type*
    and $\bigwedge x.\ x{:}A \Longrightarrow B(x)\ type$
    and $\bigwedge x\ y.\ [\![x{:}A;\ y{:}B(x)]\!] \Longrightarrow C(x,y)\ type$
  shows *?a* : $(\prod x{:}A.\ \sum y{:}B(x).\ C(x,y)) \longrightarrow (\sum f{:}\ (\prod x{:}A.\ B(x)).\ \prod x{:}A.\ C(x,$
f`x))
apply (*intr assms*)

apply (*rule ProdE* [*THEN SumE*])
apply *assumption*
apply *assumption*
apply *assumption*

apply (*rule replace-type*)
apply (*rule subst-eqtyparg*)
apply (*rule comp-rls*)
apply (*erule-tac* [*4*] *ProdE* [*THEN SumE*])
apply (*typechk assms*)
apply (*rule replace-type*)
apply (*rule subst-eqtyparg*)
apply (*rule comp-rls*)
apply (*typechk assms*)
apply *assumption*
done

Example of sequent-style deduction

schematic-goal
  assumes *A type*
    and *B type*
    and $\bigwedge z.\ z{:}A \times B \Longrightarrow C(z)$ *type*
  shows *?a* : $(\sum z{:}A \times B.\ C(z)) \longrightarrow (\sum u{:}A.\ \sum v{:}B.\ C(<u,v>))$
apply (*rule intr-rls*)
apply (*tactic* ‹*biresolve-tac* ***context*** *safe-brls 2*›)

apply (*rule-tac* [*2*] *a* = *y* **in** *ProdE*)
apply (*typechk assms*)
apply (*rule SumE*, *assumption*)
apply *intr*
defer *1*
apply *assumption+*
apply (*typechk assms*)
done

end

6 Equality reasoning by rewriting

theory *Equality*
imports *../CTT*
begin

lemma *split-eq*: *p* : *Sum*(*A*,*B*) $\Longrightarrow$ *split*(*p*,*pair*) = *p* : *Sum*(*A*,*B*)
apply (*rule EqE*)
apply (*rule elim-rls*, *assumption*)
apply *rew*
done

lemma *when-eq*: $[\![A\ type;\ B\ type;\ p : A{+}B]\!] \Longrightarrow$ *when*(*p*,*inl*,*inr*) = *p* : *A* + *B*
apply (*rule EqE*)
apply (*rule elim-rls*, *assumption*)
apply *rew*
done

lemma $p{:}N \Longrightarrow rec(p,0,\ \lambda y\ z.\ succ(y)) = p : N$
apply (*rule EqE*)
apply (*rule elim-rls*, *assumption*)
apply *rew*
done

lemma $p{:}N \Longrightarrow rec(p,0,\ \lambda y\ z.\ succ(z)) = p : N$
apply (*rule EqE*)
apply (*rule elim-rls*, *assumption*)
apply *hyp-rew*
done

lemma $[\![a{:}N;\ b{:}N;\ c{:}N]\!]$
$\Longrightarrow rec(rec(a,\ b,\ \lambda x\ y.\ succ(y)),\ c,\ \lambda x\ y.\ succ(y)) =$
$rec(a,\ rec(b,\ c,\ \lambda x\ y.\ succ(y)),\ \lambda x\ y.\ succ(y)) : N$
apply (*NE a*)
apply *hyp-rew*
done

lemma $p : Sum(A,B) \Longrightarrow <split(p,\lambda x\ y.\ x),\ split(p,\lambda x\ y.\ y)> = p : Sum(A,B)$
apply (*rule EqE*)
apply (*rule elim-rls*, *assumption*)
apply (*tactic* ‹*DEPTH-SOLVE-1* (*rew-tac* ***context*** [])›)
done

lemma $[\![a : A;\ b : B]\!] \Longrightarrow (\boldsymbol{\lambda} u.\ split(u,\ \lambda v\ w.<w,v>))\ `<a,b> = <b,a> : \sum x{:}B.$
A
apply *rew*
done

lemma $(\boldsymbol{\lambda} f.\ \boldsymbol{\lambda} x.\ f`(f`x))\ `(\boldsymbol{\lambda} u.\ split(u,\ \lambda v\ w.<w,v>)) =$
$\boldsymbol{\lambda} x.\ x\ :\ \prod x{:}(\sum y{:}N.\ N).\ (\sum y{:}N.\ N)$
apply (*rule reduction-rls*)
apply (*rule-tac* [*3*] *intrL-rls*)
apply (*rule-tac* [*4*] *EqE*)
apply (*erule-tac* [*4*] *SumE*)

apply *rew*
done

end

7 Synthesis examples, using a crude form of narrowing

theory *Synthesis*
imports *../CTT*
begin

discovery of predecessor function

schematic-goal $?a : \sum pred{:}?A\ .\ Eq(N,\ pred`0,\ 0) \times (\prod n{:}N.\ Eq(N,\ pred\ `\ succ(n),\ n))$
apply *intr*
apply *eqintr*
apply (*rule-tac* [*3*] *reduction-rls*)
apply (*rule-tac* [*5*] *comp-rls*)
apply *rew*
done

the function fst as an element of a function type

schematic-goal [*folded basic-defs*]:
$A\ type \Longrightarrow ?a{:}\ \sum f{:}?B\ .\ \prod i{:}A.\ \prod j{:}A.\ Eq(A,\ f\ `\ <i,j>,\ i)$
apply *intr*
apply *eqintr*
apply (*rule-tac* [*2*] *reduction-rls*)
apply (*rule-tac* [*4*] *comp-rls*)
apply *typechk*

now put in A everywhere

apply *assumption+*
done

An interesting use of the eliminator, when

schematic-goal $?a : \prod i{:}N.\ Eq(?A,\ ?b(inl(i)),\ <0\ \ ,\ \ i>) \times Eq(?A,\ ?b(inr(i)),\ <succ(0),\ i>)$
apply *intr*
apply *eqintr*
apply (*rule comp-rls*)
apply *rew*
done

schematic-goal $?a : \prod i{:}N.\ Eq(?A(i),\ ?b(inl(i)),\ <0\ \ ,\ \ i>) \times Eq(?A(i),\ ?b(inr(i)),\ <succ(0),i>)$
oops

A tricky combination of when and split

schematic-goal [*folded basic-defs*]:
$?a : \prod i{:}N.\ \prod j{:}N.\ Eq(?A,\ ?b(inl(<i,j>)),\ i) \times Eq(?A,\ ?b(inr(<i,j>)),\ j)$

apply *intr*
apply *eqintr*
apply (*rule PlusC-inl* [*THEN trans-elem*])
apply (*rule-tac* [*4*] *comp-rls*)
apply (*rule-tac* [*7*] *reduction-rls*)
apply (*rule-tac* [*10*] *comp-rls*)
apply *typechk*
done

schematic-goal $?a : \prod i{:}N.\ \prod j{:}N.\ Eq(?A(i,j),\ ?b(inl(<i,j>)),\ i)$
$\times\quad Eq(?A(i,j),\ ?b(inr(<i,j>)),\ j)$
oops

schematic-goal $?a : \prod i{:}N.\ \prod j{:}N.\ Eq(N,\ ?b(inl(<i,j>)),\ i)$
$\times\quad Eq(N,\ ?b(inr(<i,j>)),\ j)$
oops

Deriving the addition operator

schematic-goal [*folded arith-defs*]:
$?c : \prod n{:}N.\ Eq(N,\ ?f(0,n),\ n)$
$\times\ (\prod m{:}N.\ Eq(N,\ ?f(succ(m),\ n),\ succ(?f(m,n))))$
apply *intr*
apply *eqintr*
apply (*rule comp-rls*)
apply *rew*
done

The addition function – using explicit lambdas

schematic-goal [*folded arith-defs*]:
$?c : \sum plus : ?A\ .$
$\prod x{:}N.\ Eq(N,\ plus`0`x,\ x)$
$\times\ (\prod y{:}N.\ Eq(N,\ plus`succ(y)`x,\ succ(plus`y`x)))$
apply *intr*
apply *eqintr*
apply (*tactic resolve-tac* ***context*** [*TSimp.split-eqn*] *3*)
apply (*tactic SELECT-GOAL* (*rew-tac* ***context*** []) *4*)
apply (*tactic resolve-tac* ***context*** [*TSimp.split-eqn*] *3*)
apply (*tactic SELECT-GOAL* (*rew-tac* ***context*** []) *4*)
apply (*rule-tac* [*3*] $p = y$ **in** *NC-succ*)

apply *rew*
done

end