

theory *DeepLogic* = *FiniteMap* + *SALSyntax*:

1 Syntax and Manipulation of formulae

datatype

tval = *ILLEGAL* | *NAT nat* | *POS pos*

constdefs

MAX :: *nat*

MAX \equiv 15

datatype

vtype = *Nat* | *Pos*

types *var* = *nat*

datatype *expr* = *V var* | *Lv var* |

C tval |

Pc |

Rp |

Tm |

Add expr expr (**infix** \oplus 65) |

Minus expr expr (**infix** \ominus 65) |

Mult expr expr (**infix** \odot 65) |

Deref expr |

Ifeq expr expr expr expr (*IF* - \doteq - *THEN* - *ELSE* - [61,61,60,60] 60) |

Old expr

datatype *form* = *T* |

F |

And form list (\wedge - 70) |

Imp form form (**infix** \supset 65) |

Neg form |

Eq expr expr (**infix** \doteq 65) |

Leq expr expr (**infix** \preceq 65) |

Less expr expr (**infix** \prec 65) |

Ty expr vtype |

Forall var form ((\exists [*-./* -] [0, 10] 10) |

Yields expr tval list

datatype *mode* = *NoPt* | *Mv var var*

consts

substE::mode \Rightarrow (*expr* $\sim\sim$ \triangleright *expr*) \Rightarrow *expr* \Rightarrow *expr*

consts

$changedvars :: (expr \rightsquigarrow expr) \Rightarrow (var \rightsquigarrow expr)$

primrec

$changedvars [] = []$
 $changedvars (m\#ms) = (case (fst m)$
 $of V v \Rightarrow [(v, snd m)]$
 $| Lv v \Rightarrow []$
 $| C tv \Rightarrow []$
 $| Pc \Rightarrow []$
 $| Rp \Rightarrow []$
 $| Tm \Rightarrow []$
 $| Add e1 e2 \Rightarrow []$
 $| Minus e1 e2 \Rightarrow []$
 $| Mult e1 e2 \Rightarrow []$
 $| Deref e \Rightarrow []$
 $| Ifeq e1 e2 e3 e4 \Rightarrow []$
 $| Old e \Rightarrow [])\@(changedvars ms)$

primrec substE:

$substE m em (V v) = (case m$
 $of NoPt \Rightarrow em ? = (V v)$
 $| Mv s t \Rightarrow (Ifeq (V t) (C (NAT v)) (Deref (V s)) (em ? = (V$
 $v))))$
 $substE m em (Lv v) = (Lv v)$
 $substE m em (C tv) = (C tv)$
 $substE m em Pc = em ? = Pc$
 $substE m em Rp = em ? = Rp$
 $substE m em Tm = em ? = Tm$
 $substE m em (Add e1 e2) = Add (substE m em e1) (substE m em e2)$
 $substE m em (Minus e1 e2) = Minus (substE m em e1) (substE m em e2)$
 $substE m em (Mult e1 e2) = Mult (substE m em e1) (substE m em e2)$
 $substE m em (Deref e1) = (let e1' = (substE m em e1);$
 $res = (foldl (\lambda e (a, b). (Ifeq e1' (C (NAT a)) b e)) (Deref$
 $e1') (changedvars em))$
 $of NoPt \Rightarrow res$
 $| Mv s t \Rightarrow Ifeq (V t) e1' (Deref (V s)) res))$
 $substE m em (Ifeq e1 e2 e3 e4) = Ifeq (substE m em e1) (substE m em e2) (substE$
 $m em e3) (substE m em e4)$
 $substE m em (Old e1) = Old e1$

consts

$intVarsE :: expr \Rightarrow var list$

primrec

$intVarsE (V n) = []$
 $intVarsE (Lv v) = [v]$

$intVarsE (C tv) = []$
 $intVarsE Pc = []$
 $intVarsE Rp = []$
 $intVarsE Tm = []$
 $intVarsE (Add e1 e2) = (intVarsE e1) @ (intVarsE e2)$
 $intVarsE (Minus e1 e2) = (intVarsE e1) @ (intVarsE e2)$
 $intVarsE (Mult e1 e2) = (intVarsE e1) @ (intVarsE e2)$
 $intVarsE (Deref e1) = intVarsE e1$
 $intVarsE (Ifeq e1 e2 e3 e4) = (intVarsE e1) @ (intVarsE e2) @ (intVarsE e3) @ (intVarsE e4)$
 $intVarsE (Old e) = intVarsE e$

consts

$freeIntVars :: form \Rightarrow var\ list$
 $freeIntVarsL :: form\ list \Rightarrow var\ list$

primrec

$freeIntVarsL [] = []$
 $freeIntVarsL (f \# fs) = (freeIntVars f) @ (freeIntVarsL fs)$
 $freeIntVars T = []$
 $freeIntVars F = []$
 $freeIntVars (And fs) = freeIntVarsL fs$
 $freeIntVars (Imp f1 f2) = (freeIntVars f1) @ (freeIntVars f2)$
 $freeIntVars (Neg f) = freeIntVars f$
 $freeIntVars (Eq e1 e2) = (intVarsE e1) @ (intVarsE e2)$
 $freeIntVars (Leq e1 e2) = (intVarsE e1) @ (intVarsE e2)$
 $freeIntVars (Less e1 e2) = (intVarsE e1) @ (intVarsE e2)$
 $freeIntVars (Ty e tv) = (intVarsE e)$
 $freeIntVars (Forall n f) = filter (\lambda v. v \neq n) (freeIntVars f)$
 $freeIntVars (Yields e vs) = (intVarsE e)$

constdefs $del-id-rns :: (var \sim\sim > expr) \Rightarrow (var \sim\sim > expr)$
 $del-id-rns rm == filter (\lambda (v,e). e \neq V v) rm$

consts

$modevars :: mode \Rightarrow var\ list$

primrec

$modevars NoPt = []$
 $modevars (Mv s t) = [s,t]$

consts

$substF :: mode \Rightarrow (expr \sim\sim > expr) \Rightarrow form \Rightarrow form$
 $substFL :: mode \Rightarrow (expr \sim\sim > expr) \Rightarrow form\ list \Rightarrow form\ list$

primrec

$substFL m em [] = []$
 $substFL m em (f \# fs) = (substF m em f) \# (substFL m em fs)$

$substF\ m\ em\ T = T$
 $substF\ m\ em\ F = F$
 $substF\ m\ em\ (And\ fs) = And\ (substFL\ m\ em\ fs)$
 $substF\ m\ em\ (Imp\ f1\ f2) = Imp\ (substF\ m\ em\ f1)\ (substF\ m\ em\ f2)$
 $substF\ m\ em\ (Neg\ f) = Neg\ (substF\ m\ em\ f)$
 $substF\ m\ em\ (Eq\ e1\ e2) = (Eq\ (substE\ m\ em\ e1)\ (substE\ m\ em\ e2))$
 $substF\ m\ em\ (Leq\ e1\ e2) = (Leq\ (substE\ m\ em\ e1)\ (substE\ m\ em\ e2))$
 $substF\ m\ em\ (Less\ e1\ e2) = (Less\ (substE\ m\ em\ e1)\ (substE\ m\ em\ e2))$
 $substF\ m\ em\ (Ty\ ex\ vt) = (Ty\ (substE\ m\ em\ ex)\ vt)$
 $substF\ m\ em\ (Forall\ v\ f) = (Forall\ v\ (substF\ m\ em\ f))$
 $substF\ m\ em\ (Yields\ ex\ vs) = (Yields\ (substE\ m\ em\ ex)\ vs)$

constdefs

$nv::tval \Rightarrow nat$

$nv\ v \equiv (case\ v\ of\ ILLEGAL \Rightarrow (MAX+1) \mid NAT\ n \Rightarrow n \mid POS\ r \Rightarrow (MAX+1))$

constdefs

$rv::tval \Rightarrow pos$

$rv\ v \equiv (case\ v\ of\ ILLEGAL \Rightarrow (0,0) \mid NAT\ n \Rightarrow (0,0) \mid POS\ r \Rightarrow r)$

end