

theory *EX-ListRev* = *SALMemFWInst*:

1 SAL Example: List Reversal

We analyse a program that reverts a list. We show this this program does not exceed its granted memory.

List are represented as chains of natural numbers finished by *NAT 0*, which acts as null pointer. Each number is the address of the next element. The address of the first element is expected the base location *B*. For example the list $[a,b,c]$ is represented as $m\ bse = NAT\ a$, $m\ a = NAT\ b$, $m\ b = NAT\ c$, $m\ c = NAT\ 0$.

constdefs

null::nat — stores NAT 0
null $\equiv 0$
bs :: nat — base pointer
bs $\equiv 1$
nxt:: nat
nxt $\equiv 2$ — stores pointer to the next list element
buf::nat — saves list elements temporarily
buf $\equiv 3$
rt:: nat — stores return locations
rt $\equiv 4$

Since MOV operates indirectly we also have to store the locations of our variables. We do this with so called location variables. The main procedure initialises them to point to the corresponding variable.

constdefs

lnull::nat — initialised with NAT null
lnull $\equiv 5$
lbs::nat — initialised with NAT bs
lbs $\equiv 6$
lnxt::nat — initialised with NAT nxt
lnxt $\equiv 7$
lbuf::nat — initialised with NAT buf
lbuf $\equiv 8$

Arguments

constdefs

e1::nat
e1 $\equiv 10$
e2::nat
e2 $\equiv 11$

1.1 Program with Annotations

consts *List::(loc \Rightarrow tval) \Rightarrow loc \Rightarrow loc list \Rightarrow bool*

primrec

$List\ m\ l\ [] = (m\ l = NAT\ 0)$
 $List\ m\ l\ (l'\#ls) = (0 < l \wedge (m\ l = NAT\ l') \wedge (List\ m\ l'\ ls))$

constdefs

$prog :: SALprogram$
 $prog \equiv [$
 $(0, [(SET\ lnull\ null, None),$
 $(SET\ lbs\ bs, None),$
 $(SET\ lnext\ next, None),$
 $(SET\ lbuf\ buf, None),$
 $(SET\ null\ 0, None),$
 $(SET\ bs\ e1, None),$
 $(SET\ e1\ e2, None),$
 $(SET\ e2\ 0, None),$
 $(CALL\ rt\ 1, Some\ (\lambda(p,m,e). m\ lnull = NAT\ 0 \wedge m\ bs = NAT\ e1 \wedge$
 $m\ lnull = NAT\ null \wedge m\ lbs = NAT\ bs \wedge$
 $m\ lnext = NAT\ next \wedge m\ lbuf = NAT\ buf \wedge$
 $(List\ m\ bs\ [e1,e2])))$),
 $(HALT, Some\ TrueF)]$),
 $(1, [$
 $(MOV\ bs\ lnext, Some\ (\lambda(p,m,e). m\ lnull = NAT\ null \wedge m\ lbs = NAT\ bs \wedge$
 $m\ lnext = NAT\ next \wedge m\ lbuf = NAT\ buf \wedge m\ lnull = NAT\ 0$
 \wedge
 $m\ rt = RA\ (incA\ (\overline{pc}\ e)) \wedge$
 $(\exists n. m\ bs = NAT\ n \wedge 10 \leq n \wedge n < MAXMEM) \wedge$
 $(\exists lst. List\ m\ bs\ lst \wedge List\ (\overline{m}\ e)\ bs\ lst \wedge ((\forall l \in (set\ lst). 10 \leq$
 $l \wedge l < MAXMEM))))$),
 $(MOV\ lnull\ bs, None),$
 $(JMPEQ\ next\ null\ 6, Some\ (\lambda(p,m,e). m\ lnull = NAT\ null \wedge m\ lbs = NAT\ bs$
 \wedge
 $m\ lnext = NAT\ next \wedge m\ lbuf = NAT\ buf \wedge m\ lnull = NAT\ 0 \wedge$
 $m\ rt = RA\ (incA\ (\overline{pc}\ e)) \wedge$
 $(\exists n. m\ bs = NAT\ n \wedge 10 \leq n \wedge n < MAXMEM) \wedge$
 $(\exists l1. List\ m\ bs\ l1 \wedge (\exists l2. List\ m\ next\ l2 \wedge$
 $(\exists l3. List\ (\overline{m}\ e)\ bs\ l3 \wedge l3 = (rev\ l1)\ @\ l2 \wedge (\forall l \in (set\ l3).$
 $10 \leq l \wedge l < MAXMEM))))$),
 $(MOV\ next\ lbuf, None),$
 $(MOV\ lbs\ next, None),$
 $(MOV\ lnext\ lbs, None),$
 $(MOV\ lbuf\ lnext, None),$
 $(JMPB\ 5, None),$
 $(RET\ rt, Some\ (\lambda(p,m,e). m\ lnull = NAT\ null \wedge m\ lbs = NAT\ bs \wedge$
 $m\ lnext = NAT\ next \wedge m\ lbuf = NAT\ buf \wedge m\ lnull = NAT\ 0 \wedge$
 $m\ rt = RA\ (incA\ (\overline{pc}\ e)) \wedge$
 $(\exists lst. List\ m\ bs\ lst \wedge List\ (\overline{m}\ e)\ bs\ (rev\ lst))))$)]

constdefs

& (10::nat) <= n & n < MAXMEM) & (EX l1. List m bs l1 & (EX l2. List
m nxt l2 & (EX l3. List (callmem e) bs l3 & l3 = rev l1 @ l2 & (ALL l:set l3.
(10::nat) <= l & l < MAXMEM)))))) s & (%s. True) s) s) ((Suc (0::nat), Suc
(Suc (0::nat))), update m ta (m sa), env.h-update (env.h e @ [(Suc (0::nat), Suc
(0::nat))]) e) (%ra. False) (m (Suc (0::nat)))) (%ra. False) (m (Suc (Suc (Suc
(Suc (Suc (0::nat)))))))) s) s & (%s. True) s) s) ((Suc (0::nat), Suc (0::nat)),
update m ta (m sa), env.h-update (env.h e @ [(Suc (0::nat), 0::nat)]) e) (%ra.
False) (m (Suc (Suc (Suc (Suc (Suc (Suc (Suc (0::nat)))))))))) (%ra. False) (m
(Suc (0::nat))) s) s & (%s. True) s) s & (%s. (%s. (%s. (%s. (%s. (%s. (%pc,
m, e). (EX n. m (Suc (Suc (0::nat))) = NAT n) & (EX n. m (0::nat) = NAT n))
s & (%s. (%p, m, e). Suc (Suc (0::nat)) < MAXMEM & (0::nat) < MAXMEM)
s & (%s. True) s) s) s & (%s. (%p, m, e). m lnull = NAT null & m lbs = NAT
bs & m lnext = NAT nxt & m lbuf = NAT buf & m null = NAT (0::nat) &
m rt = RA (incA (callpc e)) & (EX n. m bs = NAT n & (10::nat) <= n & n
< MAXMEM) & (EX l1. List m bs l1 & (EX l2. List m nxt l2 & (EX l3. List
(callmem e) bs l3 & l3 = rev l1 @ l2 & (ALL l:set l3. (10::nat) <= l & l <
MAXMEM)))))) s & (%s. True) s) s) s & (%s. (%pc, m, e). (EX n n'. m (Suc
(Suc (0::nat))) = NAT n & m (0::nat) = NAT n' & n = n') & pc = (Suc (0::nat),
Suc (Suc (0::nat)))) s & (%s. True) s) s) s -->
(%pc, m, e). (%s. (%s. (%pc, m, e). EX pn' i'. m (Suc (Suc (Suc (Suc (0::nat))))))
= RA (pn', Suc i') & (EX k m' cl css. env.cs e = (k, m') # cl # css & (pn', i') =
env.h e ! k)) s & (%s. (%p, m, e). Suc (Suc (Suc (Suc (0::nat)))) < MAXMEM)
s & (%s. True) s) s) s & (%s. (%p, m, e). m lnull = NAT null & m lbs = NAT
bs & m lnext = NAT nxt & m lbuf = NAT buf & m null = NAT (0::nat) & m rt
= RA (incA (callpc e)) & (EX lst. List m bs lst & List (callmem e) bs (rev lst)))
s & (%s. True) s) s) ((Suc (0::nat), Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(0::nat)))))))))) m, env.h-update (env.h e @ [(Suc (0::nat), Suc (Suc (0::nat)))]
e) s) s & (%s. (%s. (%s. (%s. (%s. (%pc, m, e). (EX n. m (Suc (Suc (0::nat)))
= NAT n) & (EX n. m (0::nat) = NAT n)) s & (%s. (%p, m, e). Suc (Suc
(0::nat)) < MAXMEM & (0::nat) < MAXMEM) s & (%s. True) s) s) s & (%s.
(%p, m, e). m lnull = NAT null & m lbs = NAT bs & m lnext = NAT nxt & m
lbuf = NAT buf & m null = NAT (0::nat) & m rt = RA (incA (callpc e)) & (EX
n. m bs = NAT n & (10::nat) <= n & n < MAXMEM) & (EX l1. List m bs l1
& (EX l2. List m nxt l2 & (EX l3. List (callmem e) bs l3 & l3 = rev l1 @ l2 &
(ALL l:set l3. (10::nat) <= l & l < MAXMEM)))))) s & (%s. True) s) s) s &
(%s. (%pc, m, e). (EX n n'. m (Suc (Suc (0::nat))) = NAT n & m (0::nat) =
NAT n' & n ~ = n') & pc = (Suc (0::nat), Suc (Suc (0::nat)))) s & (%s. True) s)
s) s --> (%pc, m, e). (%s. (%s. (%pc, m, e). (EX sa. m (Suc (Suc (0::nat)))
= NAT sa) & (EX ta. m (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (0::nat))))))))
= NAT ta)) s & (%s. (%p, m, e). (ALL ns. m (Suc (Suc (0::nat))) = NAT ns
--> ns < MAXMEM) & (ALL nt. m (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(0::nat)))))))))) = NAT nt --> nt < MAXMEM)) s & (%s. True) s) s) s & (%s.
(%s. (%pc, m, e). pc = (Suc (0::nat), Suc (Suc (Suc (0::nat)))))) s -->
(%pc, m, e). tval-case False (%sa. tval-case False (%ta. (%s. (%s. (%pc, m, e).
(EX sa. m (Suc (Suc (Suc (Suc (Suc (Suc (0::nat)))))))) = NAT sa) & (EX ta. m
(Suc (Suc (0::nat))) = NAT ta)) s & (%s. (%p, m, e). (ALL ns. m (Suc (Suc
(Suc (Suc (Suc (Suc (0::nat)))))))) = NAT ns --> ns < MAXMEM) & (ALL nt.
m (Suc (Suc (0::nat))) = NAT nt --> nt < MAXMEM)) s & (%s. True) s) s) s
s & (%s. (%s. (%pc, m, e). pc = (Suc (0::nat), Suc (Suc (Suc (Suc (0::nat))))))

$s \dashv\vdash (\% (pc, m, e). \text{ tval-case False } (\% sa. \text{ tval-case False } (\% ta. (\% s. (\% s. (\% (pc, m, e). (EX\ sa.\ m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))))))) = NAT\ sa) \& (EX\ ta.\ m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))))) = NAT\ ta))\ s \& (\% s. (\% (p, m, e). (ALL\ ns.\ m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))))) = NAT\ ns \dashv\vdash ns < MAXMEM) \& (ALL\ nt.\ m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))))) = NAT\ nt \dashv\vdash nt < MAXMEM))\ s \& (\% s. True)\ s)\ s \& (\% s. (\% s. (\% (pc, m, e). pc = (Suc\ (0::nat), Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat))))))))\ s \dashv\vdash (\% (pc, m, e). \text{ tval-case False } (\% sa. \text{ tval-case False } (\% ta. (\% s. (\% s. (\% (pc, m, e). (EX\ sa.\ m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))))))) = NAT\ sa) \& (EX\ ta.\ m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))))) = NAT\ ta))\ s \& (\% s. (\% (p, m, e). (ALL\ ns.\ m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))))) = NAT\ ns \dashv\vdash ns < MAXMEM) \& (ALL\ nt.\ m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))))) = NAT\ nt \dashv\vdash nt < MAXMEM))\ s \& (\% s. True)\ s)\ s \& (\% s. (\% s. (\% (pc, m, e). pc = (Suc\ (0::nat), Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat))))))))\ s \dashv\vdash (\% (pc, m, e). \text{ tval-case False } (\% sa. \text{ tval-case False } (\% ta. (\% s. (\% s. (\% s. True)\ s \& (\% s. (\% (p, m, e). True)\ s \& (\% s. True)\ s)\ s \& (\% s. (\% s. (\% (pc, m, e). pc = (Suc\ (0::nat), Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat))))))))\ s \dashv\vdash (\% (pc, m, e). (\% s. (\% s. (\% (pc, m, e). (EX\ n.\ m\ (Suc\ (Suc\ (0::nat))) = NAT\ n) \& (EX\ n.\ m\ (0::nat) = NAT\ n))\ s \& (\% s. (\% (p, m, e). Suc\ (Suc\ (0::nat)) < MAXMEM \& (0::nat) < MAXMEM)\ s \& (\% s. True)\ s)\ s \& (\% s. (\% (p, m, e). m\ lnull = NAT\ lnull \& m\ lbs = NAT\ bs \& m\ lnext = NAT\ next \& m\ lbuf = NAT\ buf \& m\ null = NAT\ (0::nat) \& m\ rt = RA\ (incA\ (callpc\ e)) \& (EX\ n.\ m\ bs = NAT\ n \& (10::nat) <= n \& n < MAXMEM) \& (EX\ l1.\ List\ m\ bs\ l1 \& (EX\ l2.\ List\ m\ next\ l2 \& (EX\ l3.\ List\ (callmem\ e)\ bs\ l3 \& l3 = rev\ l1\ @\ l2 \& (ALL\ l:set\ l3. (10::nat) <= l \& l < MAXMEM))))\ s \& (\% s. True)\ s)\ s)\ ((Suc\ (0::nat), Suc\ (Suc\ (0::nat))), m, env.h-update\ (env.h\ e\ @\ [(Suc\ (0::nat), Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat))))))))])\ e))\ s \& (\% s. True)\ s)\ s)\ ((Suc\ (0::nat), Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat))))))))\ e))\ update\ m\ ta\ (m\ sa), env.h-update\ (env.h\ e\ @\ [(Suc\ (0::nat), Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat))))))))])\ e))\ (%ra. False)\ (m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat))))))))))\ (%ra. False)\ (m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat))))))))))\ s)\ s \& (\% s. True)\ s)\ s)\ ((Suc\ (0::nat), Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat))))))))\ e))\ update\ m\ ta\ (m\ sa), env.h-update\ (env.h\ e\ @\ [(Suc\ (0::nat), Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat))))))))])\ e))\ (%ra. False)\ (m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat))))))))))\ (%ra. False)\ (m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat))))))))))\ s)\ s \& (\% s. True)\ s)\ s)\ ((Suc\ (0::nat), Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat))))))))\ e))\ update\ m\ ta\ (m\ sa), env.h-update\ (env.h\ e\ @\ [(Suc\ (0::nat), Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat))))))))])\ e))\ (%ra. False)\ (m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat))))))))))\ (%ra. False)\ (m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat))))))))))\ s)\ s \& (\% s. True)\ s)\ s)\ ((Suc\ (0::nat), Suc\ (Suc\ (Suc\ (0::nat))), m, env.h-update\ (env.h\ e\ @\ [(Suc\ (0::nat), Suc\ (Suc\ (0::nat))])\ e))\ s \& (\% s. True)\ s)\ s \& (\% s. (\% s. (\% s. (\% (pc, m, e). EX\ pn'\ i'. m\ (Suc\ (Suc\ (Suc\ (0::nat)))))) = RA\ (pn', Suc\ i') \& (EX\ k\ m'\ cl\ css.\ env.cs\ e = (k, m') \# cl \# css \& (pn', i') = env.h\ e\ ! k))\ s \& (\% s. (\% (p, m, e). Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))) < MAXMEM)\ s \& (\% s. True)\ s)\ s \& (\% s. (\% (p, m, e).$

$m \text{ lnull} = \text{NAT null} \ \& \ m \text{ lbs} = \text{NAT bs} \ \& \ m \text{ lnext} = \text{NAT next} \ \& \ m \text{ lbuf} = \text{NAT buf}$
 $\ \& \ m \text{ null} = \text{NAT } (0::\text{nat}) \ \& \ m \text{ rt} = \text{RA } (\text{incA } (\text{callpc } e)) \ \& \ (\text{EX lst. List}$
 $m \text{ bs lst} \ \& \ \text{List } (\text{callmem } e) \text{ bs } (\text{rev lst})) \ s \ \& \ (\%s. \ \text{True}) \ s) \ s) \ \& \ (\%s. \ (\%s.$
 $(\%(\text{pc}, m, e). m \ (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (0::\text{nat})))))) = \text{RA } (0::\text{nat}, \text{Suc } (\text{Suc } (\text{Suc}$
 $(\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (0::\text{nat})))))))))) \ \& \ \text{pc} = (\text{Suc } (0::\text{nat}), \text{Suc } (\text{Suc}$
 $(\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (0::\text{nat})))))))))) \ s \ \& \ (\%s. \ (\%(\text{pc}, m, e). \ (\%(\text{p}, m,$
 $e). m \ \text{null} = \text{NAT } (0::\text{nat}) \ \& \ m \ \text{bs} = \text{NAT } e1 \ \& \ m \ \text{lnull} = \text{NAT null} \ \& \ m \ \text{lbs} =$
 $\ \text{NAT bs} \ \& \ m \ \text{lnext} = \text{NAT next} \ \& \ m \ \text{lbuf} = \text{NAT buf} \ \& \ \text{List } m \ \text{bs} \ [e1, e2]) \ (\text{let } (k,$
 $m' = \text{hd } (\text{env.cs } e); \ \text{cs}' = \text{tl } (\text{env.cs } e); \ h' = \text{take } k \ (\text{env.h } e) \ \text{in } (\text{env.h } e \ ! \ k,$
 $m', \ \text{env.h-update } h' \ (\text{env.cs-update } \text{cs}' \ e))) \ s \ \& \ (\%s. \ \text{True}) \ s) \ s) \ \& \ (\%s. \ \text{True})$
 $s) \ s) \ s \ \text{---} \> \ (\%(\text{pc}, m, e). \ (\%s. \ (\%s. \ (\%s. \ \text{True}) \ s \ \& \ (\%s. \ (\%(\text{p}, m, e). \ \text{True}) \ s$
 $\ \& \ (\%s. \ \text{True}) \ s) \ s) \ \& \ (\%s. \ (\%s. \ \text{True}) \ s \ \& \ (\%s. \ \text{True}) \ s) \ s) \ ((0::\text{nat}, \text{Suc } (\text{Suc}$
 $(\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (0::\text{nat})))))))))) \ , \ \text{env.cs-update } (\text{tl } (\text{env.cs}$
 $e)) \ (\text{env.h-update } (\text{env.h } e \ @ \ [(\text{Suc } (0::\text{nat}), \text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc}$
 $(\text{Suc } (0::\text{nat})))))))))) \] \ e))) \ s) \ s \ \& \ (\%s. \ \text{True}) \ s) \ s \ \& \ (\%s. \ \text{True}) \ s) \ s) \ s) \ s) \ s$

1.2 Verifying the program

lemma *forall-switch4*:

$\forall w \ x \ y \ z. P \ w \ x \ y \ z \implies \forall z. \forall x. \forall y. \forall w. P \ w \ x \ y \ z$

apply *simp*

done

lemma *List-unique*:

$\bigwedge ls \ ls' \ x. 0 \notin \text{set } (ls \ @ \ ls') \implies \text{List } m \ x \ ls \implies \text{List } m \ x \ ls' \implies ls = ls'$

apply (*subgoal-tac* $\exists k. k = \text{size } (ls \ @ \ ls')$)

prefer 2

apply *simp*

apply (*erule exE*)

apply (*erule rev-mp*)⁺

apply (*simp only: atomize-all*)

apply (*rule forall-switch4*)

apply (*rule allI*)

apply (*rule-tac n=k in nat-less-induct*)

apply (*rule allI*)⁺

apply (*rule impI*)⁺

apply (*case-tac ls*)

apply (*case-tac ls'*)

apply *simp*

apply *simp*

apply (*case-tac ls'*)

apply *simp*

apply (*erule-tac x=length (list @ lista) in allE*)

apply *simp*

apply (*erule-tac x=lista in allE*)

apply (*erule-tac x=a in allE*)

apply (*erule-tac x=list in allE*)

apply *simp*

done

lemma *List-split*:
 $\bigwedge x. \text{List } m \ x \ (ls@a\#ls') \implies \text{List } m \ a \ ls'$
apply (*induct* *ls*)
apply *simp*
apply *atomize*
apply (*erule-tac* *x=aa in allE*)
apply *simp*
done

First we ensure that the program is wellformed.

lemma *wf-prog*:
wf prog
apply (*simp add: wf-def domC-prog checkPos.simps prog-def Let-def split-def*
fst-conv snd-conv cmd.simps ret-succs.simps callpoints-def isCall-def anF.simps)
done

Then, we prove the verification condition

lemma *List-m-upd*: $\bigwedge z. x \neq z \implies x \notin (\text{set } lst) \implies (\text{List } (m(x := y)) \ z \ lst) =$
 $(\text{List } m \ z \ lst)$
apply (*induct* *lst*)
apply *simp*
apply *simp*
done

lemmas *prog-ids = bs-def lbs-def null-def lnull-def rt-def buf-def lbuf-def nrt-def*
lnrt-def e1-def e2-def MAX-def MAXMEM-def

lemmas *vc-simps = prog-ids split-def fst-conv snd-conv update-def callstate-def*
incA-def callpc-def callmem-def nth-append nat-number

lemma *vc-prog-holds*:
provable prog vc
— *start up*
apply (*simp add: provable-def valid-def | rule allI | rule impI*)
apply (*rename-tac pn i m e*)
apply (*cut-tac wf-prog*)
apply (*drule isafeP-mono*)
apply (*drule vc-proof-startup*)
apply *assumption*
apply (*erule conjE | erule exE*)
apply (*simp only: vc-def*)

— *main proof*
apply (*case-tac prog,((pn,i),m,e) |=(initF prog)*)


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— (initF prg) implies (isafeF prg (ipc prg))
apply (rule conjI, rule impI)
apply (simp add: initF-def valid-def vc-simps)

apply (simp add: initF-def valid-def vc-simps)

— case not (initF prg)
apply (erule isafeP-elims)
apply simp

apply (rule conjI)
— initF - -  $\zeta$  isafe (0,0)
apply (rule impI)
apply (drule-tac t=s'' in sym)
apply (simp add: vc-simps)

— (0,5) - -  $\zeta$  (1,0)
apply (drule-tac t=s'' in sym)
apply (rule conjI)
apply (simp add: vc-simps Let-def)
apply (rule impI)
apply (erule conjE | erule exE)+
apply (simp add: vc-simps)
apply (rule-tac x=[e1,e2] in exI)
apply (simp add: vc-simps)

apply (simp only: simp-thms split-def fst-conv snd-conv)
— (1,0) nach (1,2)
apply (rule conjI)
apply (rule impI)
apply (erule conjE | erule exE)+
apply (simp add: vc-simps)
apply (rule conjI)
apply (rule impI)
apply (case-tac lst)
apply (simp add: List.simps)
apply (simp add: List.simps)
apply (case-tac list)
apply (simp add: List.simps)
apply (rule-tac x=[a] in exI)
apply (rule context-conjI)
apply (simp add: List-m-upd)
apply arith
apply (rule-tac x=[] in exI)
apply simp

apply (simp add: List.simps)
apply (rule-tac x=[a] in exI)
apply (simp add: List-m-upd)

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apply (rule conjI)
apply arith
apply (rule impI)
apply (rule-tac x=aa#lista in exI)
apply simp
apply (rule conjI)
apply (rule impI)
apply arith

apply (erule conjE | erule exE)+
apply (subgoal-tac 0  $\notin$  set (a#aa#lista))
prefer 2
apply simp
apply (rule classical)
apply (erule-tac x=0 in ballE)
apply simp

apply simp
apply (subgoal-tac distinct (a#aa#lista))
prefer 2
apply (rule-tac m=m and x=a in List-distinct)
apply simp
apply (rule classical)
apply (erule-tac x=0 in ballE)
apply simp

apply simp
apply (rule impI)
apply (subgoal-tac List (m(Suc (Suc 0) := NAT aa, a := NAT 0)) aa lista = List
(m(Suc (Suc 0) := NAT aa)) aa lista)
prefer 2
apply (rule List-m-upd)
apply simp

apply simp
apply (subgoal-tac List (m(Suc (Suc 0) := NAT aa)) aa lista = List m aa lista)
prefer 2
apply (rule List-m-upd)
apply simp

apply (rule classical)
apply (erule-tac x=Suc (Suc 0) in ballE)
apply simp

apply simp
apply simp

apply (rule impI)
apply (case-tac css)

```

apply *simp*

apply (*simp add: Let-def*)

— (1,2) nach (1,8)

apply (*rule conjI*)

apply (*rule impI*)

apply (*erule conjE | erule exE*)+

apply (*simp add: vc-simps*)

apply (*case-tac css*)

apply (*simp add: Let-def*)

apply (*simp add: Let-def*)

apply (*rule conjI*)

apply (*rule-tac x=fst c in exI*)

apply *simp*

apply (*rule-tac x=snd c in exI*)

apply *simp*

apply (*rule-tac x=l1 in exI*)

apply *simp*

apply (*subgoal-tac l2 = []*)

prefer 2

apply (*rule classical*)

apply *simp*

apply (*simp add: neq-Nil-conv*)

apply (*erule exE*)+

apply *simp*

apply *simp*

— (1,2) nach (1,2)

apply (*rule impI*)

apply (*erule conjE | erule exE*)+

apply (*case-tac css*)

apply *simp*

apply (*rule conjI*)

apply (*simp add: vc-simps*)

apply (*case-tac l2*)

apply *simp*

apply *simp*

apply (*simp add: Let-def vc-simps cases*)

apply (*rule conjI*)

apply (*rule impI*)

apply (*rule conjI*)

apply (*case-tac l2*)

apply simp

apply simp

apply (rule-tac x=l1 in exI)

apply (rule conjI)

**apply (subgoal-tac List (m(Suc na := m na, Suc 0 := NAT nc, na := m na))
(Suc 0) l1 = List (m(Suc na := m na, Suc 0 := NAT nc)) (Suc 0) l1))**

prefer 2

apply (rule List-m-upd)

apply simp

apply (rule classical)

apply (erule-tac x=na in ballE)

apply simp

apply simp

apply (subgoal-tac m(Suc na := m na, Suc 0 := NAT nc) = m(Suc na := m na))

prefer 2

apply (rule ext)

apply simp

apply simp

apply (subgoal-tac List (m(Suc na := NAT na)) (Suc 0) l1 = List m (Suc 0) l1)

prefer 2

apply (rule List-m-upd)

apply simp

apply (rule classical)

apply (erule-tac x=Suc na in ballE)

apply simp

apply simp

apply simp

apply (rule-tac x=l2 in exI)

**apply (subgoal-tac List (m(Suc na := m na, Suc 0 := NAT nc, na := m na)) na
l2 = List m (Suc (Suc 0)) l2)**

prefer 2

**apply (subgoal-tac m(Suc na := m na, Suc 0 := NAT nc, na := m na) = m(Suc
na := m na, Suc 0 := NAT nc))**

prefer 2

apply (rule ext)

apply simp

apply simp

**apply (subgoal-tac List (m(Suc na := NAT na, Suc 0 := NAT nc)) na l2 =
List (m(Suc na := NAT na)) na l2)**

prefer 2

apply (rule List-m-upd)

apply simp

apply (*rule classical*)
apply (*erule-tac x=Suc 0 in ballE*)
apply *simp*

apply *simp*
apply *simp*
apply (*subgoal-tac List (m(Suc na := NAT na)) na l2 = List m na l2*)
prefer 2
apply (*rule List-m-upd*)
apply *simp*

apply (*rule classical*)
apply (*erule-tac x=Suc na in ballE*)
apply *simp*

apply *simp*
apply *simp*
apply *simp*

apply (*rule impI*)
apply (*rule conjI*)
apply (*case-tac l2*)
apply *simp*

apply *simp*
apply *arith*

apply (*rule impI*)
apply (*rule conjI*)
apply (*case-tac l2*)
apply *simp*

apply *simp*
apply *arith*

apply (*rule impI*)
apply (*rule conjI*)
apply (*case-tac l2*)
apply *simp*

apply *simp*
apply *arith*

apply (*rule impI*)
apply (*rule conjI*)
apply (*case-tac l2*)
apply *simp*

apply *simp*

apply *arith*

apply (*rule impI*)
apply (*rule conjI*)
apply (*rule impI*)
apply (*case-tac l2*)
apply *simp*

apply *simp*

apply (*rule impI*)
apply (*rule conjI*)
apply (*rule impI*)
apply (*case-tac l2*)
apply *simp*

apply *simp*

apply (*rule impI*)
apply (*case-tac l2*)
apply *simp*

apply *simp*
apply (*rule conjI*)
apply (*case-tac lista*)
apply *simp*

apply *simp*

apply (*rule-tac x=aa#l1 in exI*)
apply (*rule conjI*)
apply *simp*
apply (*case-tac l1*)
apply *simp*

apply *simp*
apply (*rule conjI*)
apply (*rule impI*)
apply (*simp only:*)
apply (*erule conjE*)
apply (*erule-tac P=?P ≤ Suc 0 in rev-mp*)
apply (*simp (no-asm)*)

apply (*rule impI*)
apply (*subgoal-tac distinct (Suc 0 #(rev listb @ ab # aa # lista))*)
prefer 2
apply (*rule-tac m=snd c and x=Suc 0 in List-distinct*)
apply *simp*
apply (*rule conjI*)

apply (*erule conjE*)
apply (*erule-tac P=?P ≤ ab in rev-mp*)
apply (*simp (no-asm)*)

apply (*rule conjI*)
apply (*rule classical*)
apply (*erule conjE*)
apply (*erule-tac x=0 in ballE*)
apply (*erule conjE*)
apply (*erule-tac P=?P ≤ 0 in rev-mp*)
apply (*simp (no-asm)*)

apply (*erule-tac P=0 ∉ ?P in rev-mp*)
apply (*simp (no-asm)*)

apply (*erule conjE*)
apply (*rule classical*)
apply (*erule-tac x=0 in ballE*)
apply (*erule conjE*)
apply (*erule-tac P=?P ≤ 0 in rev-mp*)
apply (*simp (no-asm)*)

apply (*erule-tac P=0 ∉ ?P in rev-mp*)
apply (*simp (no-asm)*)

apply *assumption*

apply (*erule conjE*)
apply (*subgoal-tac List (m(Suc (Suc (Suc 0)) := m aa, aa := NAT ab, Suc 0 := NAT aa, Suc (Suc 0) := m aa)) ab listb = List (m(Suc (Suc (Suc 0)) := m aa, aa := NAT ab, Suc 0 := NAT aa)) ab listb*)
prefer 2
apply (*rule List-m-upd*)
apply (*erule-tac P=?P ≤ ab in rev-mp*)
apply (*simp (no-asm)*)

apply (*rule classical*)
apply (*erule-tac x=Suc (Suc 0) in ballE*)
apply (*simp only:*)
apply (*erule conjE*)
apply (*erule-tac P=?P ≤ Suc (Suc 0) in rev-mp*)
apply (*simp (no-asm)*)
apply (*erule-tac P=Suc (Suc 0) ∉ ?P in rev-mp*)
apply (*simp (no-asm)*)

apply (*subgoal-tac List (m(Suc (Suc (Suc 0)) := m aa, aa := NAT ab, Suc 0 := NAT aa)) ab listb = List (m(Suc (Suc (Suc 0)) := m aa, aa := NAT ab)) ab listb*)
prefer 2

apply (*rule List-m-upd*)
apply (*erule-tac P=?P ≤ ab in rev-mp*)
apply (*simp (no-asm)*)

apply (*rule classical*)
apply (*erule-tac x=Suc 0 in ballE*)
apply (*simp only:*)
apply (*erule conjE*)
apply (*erule-tac P=?P ≤ Suc 0 in rev-mp*)
apply (*simp (no-asm)*)
apply (*erule-tac P=Suc 0 ∉ ?P in rev-mp*)
apply (*simp (no-asm)*)

apply (*subgoal-tac List (m(Suc (Suc (Suc 0)) := m aa, aa := NAT ab)) ab listb*)
 =

List (m(Suc (Suc (Suc 0)) := m aa)) ab listb
prefer 2
apply (*rule List-m-upd*)
apply (*erule-tac P=distinct ?P in rev-mp*)
apply (*simp (no-asm)*)
apply (*rule impI*)
apply (*erule conjE*)+
apply (*rule not-sym*)
apply *assumption*

apply (*erule-tac P=distinct ?P in rev-mp*)
apply (*simp (no-asm)*)

apply (*subgoal-tac List (m(Suc (Suc (Suc 0)) := m aa)) ab listb = List m ab listb*)
prefer 2

apply (*rule List-m-upd*)
apply (*erule-tac P=?P ≤ ab in rev-mp*)
apply (*simp (no-asm)*)

apply (*rule classical*)
apply (*erule-tac x=Suc (Suc (Suc 0)) in ballE*)
apply (*simp only:*)
apply (*erule conjE*)
apply (*erule-tac P=?P ≤ Suc (Suc (Suc 0)) in rev-mp*)
apply (*simp (no-asm)*)
apply (*erule-tac P=Suc (Suc (Suc 0)) ∉ ?P in rev-mp*)
apply (*simp (no-asm)*)
apply (*simp only:*)

apply (*rule-tac x=lista in exI*)
apply (*rule conjI*)
apply (*case-tac lista*)
apply *simp*


```

apply simp
apply (subgoal-tac distinct (Suc 0 #(rev l1 @ aa # ab # listb)))
prefer 2
apply (rule-tac m=snd c and x=Suc 0 in List-distinct)
apply simp
apply (rule conjI)
apply (erule conjE)+
apply (erule-tac P=?P ≤ ab in rev-mp)
apply (simp (no-asm))

apply (rule conjI)
apply (rule classical)
apply (erule conjE)+
apply (erule-tac x=0 in ballE)
apply (erule conjE)
apply (erule-tac P=?P ≤ 0 in rev-mp)
apply (simp (no-asm))

apply (erule-tac P=0 ∉ ?P in rev-mp)
apply (simp (no-asm))

apply (erule conjE)+
apply (rule classical)
apply (erule-tac x=0 in ballE)
apply (erule conjE)
apply (erule-tac P=?P ≤ 0 in rev-mp)
apply (simp (no-asm))

apply (erule-tac P=0 ∉ ?P in rev-mp)
apply (simp (no-asm))

apply assumption

apply (erule conjE)+
apply (subgoal-tac List (m(Suc (Suc (Suc 0)) := NAT ab, aa := NAT nc, Suc 0
:= NAT aa, Suc (Suc 0) := NAT ab)) ab listb = List (m(Suc (Suc (Suc 0)) :=
NAT ab, aa := NAT nc, Suc 0 := NAT aa)) ab
listb)
prefer 2
apply (rule List-m-upd)
apply (erule-tac P=?P ≤ ab in rev-mp)
apply (simp (no-asm))

apply (rule classical)
apply (erule-tac x=Suc (Suc 0) in ballE)
apply (simp only.)
apply (erule conjE)
apply (erule-tac P=?P ≤ Suc (Suc 0) in rev-mp)
apply (simp (no-asm))

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apply (erule-tac P=Suc (Suc 0)  $\notin$  ?P in rev-mp)
apply (simp (no-asm))

apply (subgoal-tac List (m(Suc (Suc (Suc 0)) := NAT ab, aa := NAT nc, Suc 0
:= NAT aa)) ab listb =
List (m(Suc (Suc (Suc 0)) := NAT ab, aa := NAT nc)) ab listb)
prefer 2
apply (rule List-m-upd)
apply (erule-tac P=?P  $\leq$  ab in rev-mp)
apply (simp (no-asm))

apply (rule classical)
apply (erule-tac x=Suc 0 in ballE)
apply (simp only:)
apply (erule conjE)
apply (erule-tac P=?P  $\leq$  Suc 0 in rev-mp)
apply (simp (no-asm))
apply (erule-tac P=Suc 0  $\notin$  ?P in rev-mp)
apply (simp (no-asm))

apply (subgoal-tac List (m(Suc (Suc (Suc 0)) := NAT ab, aa := NAT nc)) ab
listb =
List (m(Suc (Suc (Suc 0)) := NAT ab)) ab listb)
prefer 2
apply (rule List-m-upd)
apply (erule-tac P=distinct ?P in rev-mp)
apply (simp (no-asm))

apply (erule-tac P=distinct ?P in rev-mp)
apply (simp (no-asm))

apply (subgoal-tac List (m(Suc (Suc (Suc 0)) := NAT ab)) ab listb = List m ab
listb)
prefer 2
apply (rule List-m-upd)
apply (erule-tac P=?P  $\leq$  ab in rev-mp)
apply (simp (no-asm))

apply (rule classical)
apply (erule-tac x=Suc (Suc (Suc 0)) in ballE)
apply (simp only:)
apply (erule conjE)
apply (erule-tac P=?P  $\leq$  Suc (Suc (Suc 0)) in rev-mp)
apply (simp (no-asm))
apply (erule-tac P=Suc (Suc (Suc 0))  $\notin$  ?P in rev-mp)
apply (simp (no-asm))
apply (simp only:)

apply (rule conjI)

```

apply *simp*

apply *simp*

apply (*erule conjE*)+

apply *assumption*

done

end