

theory *EX-RecMultMem* = *SALMemFWInst*:

1 SAL Example: Multiplication

This example is about a program that multiplies variable C1 with C2 and writes the result to C3.

The variables C1, C2, C3 are just abbreviations for the addresses 1, 2 and 3.

constdefs

C1 :: *nat* — cell1
C1 ≡ 1
C2 :: *nat* — cell2
C2 ≡ 2
C3 :: *nat* — cell3
C3 ≡ 3

At some point in the program we need to check whether a variable contains number one or zero. Since conditional jumps in SAL expect their arguments in form of variables, we introduce variables Zero and One. The program uses them to the constants NAT 0 and NAT 1.

constdefs

Zero::*nat* — stores NAT 0
Zero ≡ 4
One::*nat* — stores NAT 1
One ≡ 5

Our program will do multiplication by recursively calling a procedure that decrements C1 and adds C2 to C3. Whenever this procedure is called, the return address, which is the current address incremented by one, is dumped in a variable r. The address of r, which is 6, is kept in the variable Ra.

constdefs

R :: *nat* — return address buffer
R ≡ 6
Ra :: *nat* — stores the address of R
Ra ≡ 7

The program maintains a stack of return addresses. This stack starts at address Soff and grows towards higher addresses. The variable S contains the stack pointer. This is the address, where the next element of the stack is going to be stored.

constdefs

S::*nat* — stack pointer
S ≡ 8
Soff::*nat* — stack base
Soff ≡ 9

Since SAL programs are executed on a default initial state, we have to simulate the input phase of the program by SET operations. In this example we analyse the program for inputs arg1 and arg2.

```

constdefs
  arg1::nat
  arg1≡2
  arg2::nat
  arg2≡3

```

1.1 Program Code

Next, we present the program without annotations.

```

prog = [ (0, [ Procedure 0 (Startup)
  SET Zero 0, Initialise Zero with NAT 0.
  SET One 1, Initialise One with NAT 1.
  SET RA r, Initialise RA with NAT r (r = 6).
  SET S Soff, Initialise S with NAT Soff (Soff = 9).
  SET C1 arg1, Initialise Argument C1 with NAT arg1.
  SET C2 arg2, Initialise Argument C2 with NAT arg2.
  SET C3 0, Initialise Result Variable C3 with NAT 0.
  CALL r 1, Start multiplication procedure.
  HALT ]), Stop execution

(1, [ Procedure 1 (Multiplication)
  MOV Ra S, Push the return address onto the Stack.
  JMPEQ C1 Zero, If C1 is NAT 0 we are done, otherwise . . .
  SUB C1 One,
  . . . we decrement C1
  ADD C3 C2, and add C2 to the result variable C3.
  INC S, We increment the stack pointer, before .
  CALL r 1, we do the recursive call.
  SUB S One, After return we restore the old stack pointer
  MOV S Ra, and copy the return address from the stack to r.
  RET r ] ) ] We finish this call by returning to the caller.

```

1.2 Program with Annotations

```

constdefs
  prog :: SALprogram
  prog ≡ [
(0, [ (SET Zero 0, None),
      (SET One 1, None),
      (SET Ra R, None),
      (SET S Soff, None),

```

(SET C1 arg1, None),
 (SET C2 arg2, None),
 (SET C3 0, None),
 (CALL R 1, Some ($\lambda(pc, m, e). m C1 = NAT\ arg1 \wedge$
 $m C2 = NAT\ arg2 \wedge$
 $m C3 = NAT\ 0 \wedge$
 $m Zero = NAT\ 0 \wedge$
 $m One = NAT\ 1 \wedge$
 $m Ra = NAT\ R \wedge$
 $m S = NAT\ Soff \wedge$
 $(\forall Z. Z \notin \{C1, C2, C3, Zero, One, Ra, S\} \longrightarrow m Z = (\overline{m}\ e)$
 $Z))),$
 (JMPB 0, Some ($\lambda(pc, m, e). m C3 = NAT\ (arg1 * arg2))))),$
 (1, [(MOV Ra S, Some ($\lambda(pc, m, e). (\exists n1. m C1 = NAT\ n1 \wedge$
 $(\exists n2. m C2 = NAT\ n2 \wedge$
 $(\exists n3. m C3 = NAT\ n3 \wedge$
 $n3 + (n1 * n2) \leq MAX)))) \wedge$
 $m Zero = NAT\ 0 \wedge$
 $m One = NAT\ 1 \wedge$
 $m Ra = NAT\ R \wedge$
 $(\exists s. m S = NAT\ s \wedge Soff \leq s \wedge$
 $(\exists n1. m C1 = NAT\ n1 \wedge s + n1 \leq MAX)) \wedge$
 $m R = RA\ (incA\ (\overline{pc}\ e)) \wedge$
 $(\forall Z. Z \neq R \longrightarrow m Z = (\overline{m}\ e)\ Z))),$
 (JMPEQ C1 Zero 7, None),
 (SUB C1 One, None),
 (ADD C3 C2, None),
 (INC S, None),
 (CALL R 1, Some ($\lambda(pc, m, e). (\exists n1. m C1 = NAT\ n1 \wedge$
 $(\exists n2. m C2 = NAT\ n2 \wedge$
 $(\exists n3. m C3 = NAT\ n3 \wedge$
 $(\exists n1'. (\overline{m}\ e)\ C1 = NAT\ n1' \wedge$
 $(\exists n2'. (\overline{m}\ e)\ C2 = NAT\ n2' \wedge$
 $(\exists n3'. (\overline{m}\ e)\ C3 = NAT\ n3' \wedge$
 $0 < n1' \wedge n1 = (n1' - (1 :: nat)) \wedge$
 $n2 = n2' \wedge n3 = n3' + n2' \wedge$
 $n3 + (n1 * n2) \leq MAX)))))) \wedge$
 $m Zero = NAT\ 0 \wedge$
 $m One = NAT\ 1 \wedge$
 $m Ra = NAT\ R \wedge$
 $(\exists s. m S = NAT\ s \wedge Soff \leq s \wedge$
 $(\exists s'. (\overline{m}\ e)\ S = NAT\ s' \wedge$
 $s = (s' + (1 :: nat)) \wedge$
 $m s' = RA\ (incA\ (\overline{pc}\ e)) \wedge$
 $(\exists n1. m C1 = NAT\ n1 \wedge s + n1 \leq MAX)) \wedge$
 $(\forall Z. Z \neq R \wedge Z \neq S \wedge Z \neq C1 \wedge Z \neq C3 \wedge$
 $(\exists s. m S = NAT\ s \wedge (Suc\ Z) < s) \longrightarrow m Z = (\overline{m}\ e)$
 $Z))),$
 (SUB S One, Some ($\lambda(pc, m, e). \exists n3. m C3 = NAT\ n3 \wedge$

$$\begin{aligned}
& (\exists n3'. (\overline{m} \ e) \ C3 = NAT \ n3' \wedge \\
& (\exists n1'. (\overline{m} \ e) \ C1 = NAT \ n1' \wedge \\
& (\exists n2'. (\overline{m} \ e) \ C2 = NAT \ n2' \wedge \\
& \quad n3 = n3' + (n1' * n2'))) \wedge \\
& m \ Zero = NAT \ 0 \wedge \\
& m \ One = NAT \ 1 \wedge \\
& m \ Ra = NAT \ R \wedge \\
& (\exists s. m \ S = NAT \ s \wedge \text{Soff} \leq s \wedge \\
& \quad (\exists s'. (\overline{m} \ e) \ S = NAT \ s' \wedge \\
& \quad \quad s = (s' + (1 :: nat)) \wedge \\
& \quad \quad m \ s' = RA \ (\text{incA} \ (\overline{pc} \ e))) \wedge \\
& \quad (\exists n1. m \ C1 = NAT \ n1 \wedge s + n1 \leq MAX)) \wedge \\
& (\forall Z. Z \neq R \wedge Z \neq S \wedge Z \neq C1 \wedge Z \neq C3 \wedge (\exists s. m \ S = NAT \\
& s \wedge (Suc \ Z) < s) \\
& \quad \longrightarrow m \ Z = (\overline{m} \ e) \ Z)), \\
& (MOV \ S \ Ra, None), \\
& (RET \ R, Some \ (\lambda(pc, m, e). \exists n3. m \ C3 = NAT \ n3 \wedge \\
& \quad (\exists n3'. (\overline{m} \ e) \ C3 = NAT \ n3' \wedge \\
& \quad (\exists n1'. (\overline{m} \ e) \ C1 = NAT \ n1' \wedge \\
& \quad (\exists n2'. (\overline{m} \ e) \ C2 = NAT \ n2' \wedge \\
& \quad \quad n3 = n3' + (n1' * n2'))) \wedge \\
& \quad m \ Zero = NAT \ 0 \wedge \\
& \quad m \ One = NAT \ 1 \wedge \\
& \quad m \ Ra = NAT \ R \wedge \\
& \quad m \ R = RA \ (\text{incA} \ (\overline{pc} \ e)) \wedge \\
& \quad (\exists s. m \ S = NAT \ s \wedge \text{Soff} \leq s \wedge \\
& \quad \quad (\exists n1. m \ C1 = NAT \ n1 \wedge s + n1 \leq MAX)) \wedge \\
& \quad (\forall Z. Z \neq R \wedge Z \neq C1 \wedge Z \neq C3 \wedge (\exists s. m \ S = NAT \ s \wedge Z \\
& < s) \\
& \quad \longrightarrow m \ Z = (\overline{m} \ e) \ Z)) \\
& \quad] \\
&]
\end{aligned}$$

constdefs

vc :: SALform

$$\begin{aligned}
vc \equiv & (\%s. (((\%s. (((\%pc, m, e). ((pc = ((0::nat), (0::nat))) \& (((env.cs \ e) = \\
& [((0::nat), m)]) \& (((env.h \ e) = [])) \& (ALL \ X. ((m \ X) = ILLEGAL)))))) s) \longrightarrow \\
& ((\%s. (((\%s. (((\%pc, m, e). ((0::nat) <= MAX)) s) \& ((\%s. (((\%p, m, e). ((Suc \\
& (Suc (Suc (Suc (0::nat)))))) < MAXMEM)) s) \& ((\%s. True) s))) s))) s) \& ((\%s. \\
& (((\%s. (((\%pc, m, e). (pc = ((0::nat), (0::nat)))) s) \longrightarrow ((\%pc, m, e). ((\%s. \\
& (((\%s. (((\%pc, m, e). ((Suc (0::nat)) <= MAX)) s) \& ((\%s. (((\%p, m, e). ((Suc \\
& (Suc (Suc (Suc (Suc (0::nat)))))) < MAXMEM)) s) \& ((\%s. True) s))) s))) s) \\
& \& ((\%s. (((\%s. (((\%pc, m, e). (pc = ((0::nat), (Suc (0::nat)))))) s) \longrightarrow \\
& ((\%pc, m, e). ((\%s. (((\%s. (((\%pc, m, e). ((Suc (Suc (Suc (Suc (Suc (Suc (Suc \\
& (0::nat)))))) <= MAX)) s) \& ((\%s. (((\%p, m, e). ((Suc (Suc (Suc (Suc (Suc \\
& (Suc (Suc (0::nat)))))) < MAXMEM)) s) \& ((\%s. True) s))) s))) s) \& ((\%s.
\end{aligned}$$

$\& (pc = ((Suc\ 0::nat), (Suc\ 0::nat))))\ s) \dashrightarrow ((\% (pc, m, e). ((\%s. (((\%s. ((\% (pc, m, e). (EX\ pn'\ i'. (((m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ 0::nat)))))))) = (RA\ (pn', (Suc\ i')))) \& (EX\ k\ m'\ cl\ css. (((env.cs\ e) = ((k, m') \# (cl \# css))) \&$
 $((pn', i') = ((env.h\ e) ! k))))))\ s) \& ((\%s. (((\% (p, m, e). ((Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ 0::nat)))))) < MAXMEM))\ s) \& ((\%s. True)\ s)))\ s)))\ s) \& ((\%s. (((\% (pc, m, e). (EX\ n3. (((m\ C3) = (NAT\ n3)) \& ((EX\ n3'. (((callmem\ e\ C3) = (NAT\ n3')) \& (EX\ n1'. (((callmem\ e\ C1) = (NAT\ n1')) \& (EX\ n2'. (((callmem\ e\ C2) = (NAT\ n2')) \& (n3 = (n3' + n1' * n2'))))))))\ s) \& ((m\ Zero) = (NAT\ 0::nat)) \& ((m\ One) = (NAT\ 1::nat)) \& ((m\ Ra) = (NAT\ R)) \& ((m\ R) = (RA\ (incA\ (callpc\ e)))) \& ((EX\ s. (((m\ S) = (NAT\ s)) \& ((Soff\ <= s) \& (EX\ n1. (((m\ C1) = (NAT\ n1)) \& ((s + n1) <= MAX))))))\ s) \& (ALL\ Z. (((Z \sim = R) \& ((Z \sim = C1) \& ((Z \sim = C3) \&$
 $(EX\ s. (((m\ S) = (NAT\ s)) \& (Z < s))))))\ s) \dashrightarrow ((m\ Z) = (callmem\ e\ Z))))))\ s) \& ((\%s. True)\ s)))\ s)))\ s) \& ((\%s. (((\% (pc, m, e). ((EX\ n\ n'. (((m\ (Suc\ 0::nat)) = (NAT\ n)) \& ((m\ (Suc\ (Suc\ (Suc\ (Suc\ 0::nat)))) = (NAT\ n')) \& (n \sim = n')))) \& (pc = ((Suc\ 0::nat), (Suc\ 0::nat))))\ s) \dashrightarrow ((\% (pc, m, e). ((\%s. (((\% (pc, m, e). ((EX\ n. ((m\ (Suc\ 0::nat)) = (NAT\ n)) \& (EX\ n. ((m\ (Suc\ (Suc\ (Suc\ (Suc\ 0::nat)))) = (NAT\ n))))\ s) \&$
 $((\%s. (((\% (p, m, e). (((Suc\ 0::nat) < MAXMEM) \& ((Suc\ (Suc\ (Suc\ (Suc\ (Suc\ 0::nat)))) < MAXMEM))\ s) \& ((\%s. True)\ s)))\ s)))\ s) \& ((\%s. (((\% (pc, m, e). (pc = ((Suc\ 0::nat), (Suc\ (Suc\ 0::nat))))\ s) \dashrightarrow ((\% (pc, m, e). ((\%s. (((\% (pc, m, e). ((EX\ n. ((m\ (Suc\ (Suc\ (Suc\ 0::nat)))) = (NAT\ n)) \& ((EX\ n. ((m\ (Suc\ (Suc\ 0::nat))) = (NAT\ n)) \& (EX\ n. (((lift\ op + (m\ (Suc\ (Suc\ (Suc\ 0::nat)))) (m\ (Suc\ (Suc\ 0::nat)))) = (NAT\ n)) \& (n <= MAX))))))\ s) \& ((\%s. (((\% (p, m, e). (((Suc\ (Suc\ (Suc\ 0::nat))) < MAXMEM) \& ((Suc\ (Suc\ 0::nat)) < MAXMEM))\ s) \& ((\%s. True)\ s)))\ s)))\ s) \& ((\%s. (((\% (pc, m, e). (pc = ((Suc\ 0::nat), (Suc\ (Suc\ (Suc\ 0::nat))))\ s) \dashrightarrow$
 $((\% (pc, m, e). ((\%s. (((\%s. (((\% (pc, m, e). (EX\ n. (((m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ 0::nat)))))) = (NAT\ n)) \& (n < MAX))))\ s) \& ((\%s. (((\% (p, m, e). ((Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ 0::nat)))))) < MAXMEM))\ s) \& ((\%s. True)\ s)))\ s)))\ s) \& ((\%s. (((\% (pc, m, e). (pc = ((Suc\ 0::nat), (Suc\ (Suc\ (Suc\ (Suc\ 0::nat))))\ s) \dashrightarrow ((\% (pc, m, e). ((\%s. (((\%s. (((\%s. True)\ s) \& ((\%s. (((\% (p, m, e). ((Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ 0::nat)))))) < MAXMEM))\ s) \& ((\%s. True)\ s)))\ s)))\ s) \& ((\%s. (((\% (pc, m, e). ((EX\ n1. (((m\ C1) = (NAT\ n1)) \& (EX\ n2. (((m\ C2) = (NAT\ n2)) \& (EX\ n3. (((m\ C3) = (NAT\ n3)) \& (EX\ n1'. (((callmem\ e\ C1) = (NAT\ n1')) \& (EX\ n2'. (((callmem\ e\ C2) = (NAT\ n2')) \& (EX\ n3'. (((callmem\ e\ C3) = (NAT\ n3')) \& (((0::nat) < n1') \&$
 $((n1 = (n1' - (1::nat)) \& ((n2 = n2') \& ((n3 = (n3' + n2')) \& ((n3 + (n1 * n2) <= MAX))))))\ s) \& ((m\ Zero) = (NAT\ 0::nat)) \& ((m\ One) = (NAT\ 1::nat)) \& ((m\ Ra) = (NAT\ R)) \& ((EX\ s. (((m\ S) = (NAT\ s)) \& ((Soff\ <= s) \& ((EX\ s'. (((callmem\ e\ S) = (NAT\ s')) \& ((s = (s' + (1::nat))) \& ((m\ s') = (RA\ (incA\ (callpc\ e))))))\ s) \& (EX\ n1. (((m\ C1) = (NAT\ n1)) \& ((s + n1) <= MAX))))))\ s) \& (ALL\ Z. (((Z \sim = R) \& ((Z \sim = S) \& ((Z \sim = C1) \& ((Z \sim = C3) \& (EX\ s. (((m\ S) = (NAT\ s)) \& ((Suc\ Z) < s))))))\ s) \dashrightarrow ((m\ Z)$

= (callmem e Z)))))) s) &
 ((%s. True) s)) s)) (((Suc (0::nat)), (Suc (Suc (Suc (Suc (Suc (0::nat))))))),
 (update m (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (0::nat)))))))))) (lift op +
 (m (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (0::nat)))))))))) (NAT (1::nat))),
 (env.h-update ((env.h e) @ (((Suc (0::nat)), (Suc (Suc (Suc (Suc (0::nat)))))))
 e)))) s)) s) & ((%s. True) s)) s)) (((Suc (0::nat)), (Suc (Suc (Suc (Suc (0::nat))))))),
 (update m (Suc (Suc (Suc (0::nat)))) (lift op + (m (Suc (Suc (Suc (0::nat))))))
 (m (Suc (Suc (0::nat))))), (env.h-update ((env.h e) @ (((Suc (0::nat)), (Suc (Suc
 (Suc (0::nat)))))) e)))) s)) s) &
 ((%s. True) s)) s)) (((Suc (0::nat)), (Suc (Suc (Suc (0::nat))))), (update m (Suc
 (0::nat)) (lift op - (m (Suc (0::nat))) (m (Suc (Suc (Suc (Suc (Suc (0::nat))))))))),
 (env.h-update ((env.h e) @ (((Suc (0::nat)), (Suc (Suc (0::nat)))))) e)))) s)) s)
 & ((%s. True) s)) s)) (((Suc (0::nat)), (Suc (Suc (0::nat))), m, (env.h-update
 ((env.h e) @ (((Suc (0::nat)), (Suc (0::nat)))) e)))) s)) s) & ((%s. True) s))
 s)) s)) (((Suc (0::nat)), (Suc (0::nat))), (update m ta (m sa)), (env.h-update
 ((env.h e) @ (((Suc (0::nat)), (0::nat))) e))) (%ra. False) (m (Suc (Suc (Suc
 (Suc (Suc (Suc (Suc (Suc (0::nat)))))))))) (%ra. False) (m (Suc (Suc (Suc (Suc
 (Suc (Suc (Suc (0::nat)))))))))) s)) s) &
 ((%s. True) s)) s) & ((%s. (((%s. (((%s. (((%s. (((%s. (((%s. (((%s. True) s) &
 (((%s. (((%p, m, e). ((Suc (Suc (Suc (Suc (Suc (Suc (0::nat)))))))) < MAXMEM))
 s) & ((%s. True) s)) s)) s) & ((%s. (((%pc, m, e). ((EX n1. (((m C1) = (NAT
 n1)) & (EX n2. (((m C2) = (NAT n2)) & (EX n3. (((m C3) = (NAT n3)) &
 (EX n1'. (((callmem e C1) = (NAT n1')) & (EX n2'. (((callmem e C2) = (NAT
 n2')) & (EX n3'. (((callmem e C3) = (NAT n3')) & ((0::nat) < n1') & ((n1 =
 (n1' - (1::nat)) & ((n2 = n2') & ((n3 = (n3' + n2')) & ((n3 + (n1 * n2))
 <= MAX)))))))))))))) & (((m Zero) = (NAT (0::nat))) & (((m One) = (NAT
 (1::nat))) & (((m Ra) = (NAT R)) & ((EX s. (((m S) = (NAT s)) & ((Soff <=
 s) & ((EX s'. (((callmem e S) = (NAT s')) & ((s = (s' + (1::nat))) & ((m s') =
 (RA (incA (callpc e)))))) &
 (EX n1. (((m C1) = (NAT n1)) & ((s + n1) <= MAX)))))) & (ALL Z. (((Z
 ~ = R) & ((Z ~ = S) & ((Z ~ = C1) & ((Z ~ = C3) & (EX s. (((m S) = (NAT
 s) & ((Suc Z) < s)))))) --> ((m Z) = (callmem e Z)))))) s) & ((%s.
 True) s)) s)) s) & ((%s. (((%pc, m, e). (pc = ((Suc (0::nat)), (Suc (Suc
 (Suc (Suc (Suc (0::nat)))))))) s) & ((%s. True) s)) s)) s) --> ((%pc, m,
 e). ((%s. (((%s. (((%pc, m, e). ((EX sa. ((m (Suc (Suc (Suc (Suc (Suc (Suc (Suc
 (Suc (0::nat)))))))) = (NAT sa))) & (EX ta. ((m (Suc (Suc (Suc (Suc (Suc (Suc
 (Suc (Suc (0::nat)))))))) = (NAT ta)))) s) & ((%s. (((%p, m, e). (ALL ns.
 (((m (Suc (Suc (Suc (Suc (Suc (Suc (Suc (0::nat)))))))) = (NAT ns)) --> (ns
 < MAXMEM))) &
 (ALL nt. (((m (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (0::nat)))))))) = (NAT
 nt)) --> (nt < MAXMEM)))) s) & ((%s. True) s)) s)) s) & ((%s. (((%pc,
 m, e). ((EX n1. (((m C1) = (NAT n1)) & (EX n2. (((m C2) = (NAT n2)) &
 (EX n3. (((m C3) = (NAT n3)) & ((n3 + (n1 * n2)) <= MAX)))))) & (((m
 Zero) = (NAT (0::nat))) & (((m One) = (NAT (1::nat))) & (((m Ra) = (NAT
 R)) & ((EX s. (((m S) = (NAT s)) & ((Soff <= s) & (EX n1. (((m C1) = (NAT
 n1)) & ((s + n1) <= MAX)))))) & (((m R) = (RA (incA (callpc e)))) & (ALL
 Z. ((Z ~ = R) -->
 ((m Z) = (callmem e Z)))))) s) & ((%s. True) s)) s)) (((Suc (0::nat)),
 (0::nat)), (update m (Suc (Suc (Suc (Suc (Suc (Suc (0::nat))))))) (RA ((Suc

$(0::\text{nat}), (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (0::\text{nat}))))))))))), (\text{env.cs-update} (((\text{length}$
 $(\text{env.h e}), m) \# (\text{env.cs e})) (\text{env.h-update} ((\text{env.h e}) @ [((\text{Suc} (0::\text{nat}), (\text{Suc} (\text{Suc}$
 $(\text{Suc} (\text{Suc} (\text{Suc} (0::\text{nat})))))))] e)))) s))) s) \& ((\%s. \text{True}) s))) s) \& ((\%s. (((\%s.$
 $((\%s. (((\%s. (((\%s. (((\%s. (((\%pc, m, e). ((EX n. ((m (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc}$
 $(\text{Suc} (\text{Suc} (\text{Suc} (0::\text{nat}))))))))) = (\text{NAT } n)) \& (EX n. ((m (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc}$
 $(\text{Suc} (0::\text{nat})))))) = (\text{NAT } n)))))) s) \&$
 $((\%s. (((\%p, m, e). (((\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (0::\text{nat}))))))))))$
 $< \text{MAXMEM}) \& ((\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (0::\text{nat})))))) < \text{MAXMEM})) s) \&$
 $((\%s. \text{True}) s))) s))) s) \& ((\%s. (((\%pc, m, e). (EX n3. (((m C3) = (\text{NAT } n3))$
 $\& ((EX n3'. (((\text{callmem } e C3) = (\text{NAT } n3')) \& (EX n1'. (((\text{callmem } e C1) =$
 $(\text{NAT } n1')) \& (EX n2'. (((\text{callmem } e C2) = (\text{NAT } n2')) \& (n3 = (n3' + (n1' *$
 $n2')))))))) \& (((m \text{Zero}) = (\text{NAT } (0::\text{nat}))) \& (((m \text{One}) = (\text{NAT } (1::\text{nat}))) \&$
 $((m \text{Ra}) = (\text{NAT } R)) \& ((EX s. (((m S) = (\text{NAT } s)) \& ((\text{Soff } \leq s) \& ((EX$
 $s'. (((\text{callmem } e S) = (\text{NAT } s')) \& ((s = (s' + (1::\text{nat}))) \& ((m s') = (\text{RA } (\text{incA}$
 $(\text{callpc } e)))))) \& (EX n1. (((m C1) = (\text{NAT } n1)) \& ((s + n1) \leq \text{MAX}))))))$
 $\& (\text{ALL } Z. (((Z \sim = R) \&$
 $((Z \sim = S) \& ((Z \sim = C1) \& ((Z \sim = C3) \& (EX s. (((m S) = (\text{NAT } s)) \&$
 $((\text{Suc } Z) < s)))))) --> ((m Z) = (\text{callmem } e Z)))))) s) \& ((\%s. \text{True})$
 $s))) s) \& ((\%s. (((\%pc, m, e). (pc = ((\text{Suc} (0::\text{nat}), (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc}$
 $(\text{Suc} (\text{Suc} (0::\text{nat}))))))))) s) \& ((\%s. \text{True}) s))) s) --> ((\%pc, m, e).$
 $((\%s. (((\%s. (((\%pc, m, e). ((EX sa. ((m (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc}$
 $(\text{Suc} (0::\text{nat}))))))))) = (\text{NAT } sa)) \& (EX ta. ((m (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc}$
 $(\text{Suc} (\text{Suc} (0::\text{nat}))))))))) = (\text{NAT } ta)))) s) \& ((\%s. (((\%p, m, e). ((\text{ALL } ns.$
 $((m (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (0::\text{nat}))))))))) = (\text{NAT } ns))$
 $--> (ns < \text{MAXMEM})) \& (\text{ALL } nt. (((m (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc}$
 $(0::\text{nat}))))))))) = (\text{NAT } nt)) -->$
 $(nt < \text{MAXMEM})))) s) \& ((\%s. \text{True}) s))) s) \& ((\%s. (((\%s. (((\%pc, m,$
 $e). (pc = ((\text{Suc} (0::\text{nat}), (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (0::\text{nat})))))))))$
 $s) --> ((\%pc, m, e). (\text{tval-case False } \%sa. (\text{tval-case False } \%ta. (((\%s. (((\%s.$
 $((\%pc, m, e). (EX pn' i'. (((m (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (0::\text{nat})))))) =$
 $(\text{RA } (pn', (\text{Suc } i')))) \& (EX k m' cl css. (((\text{env.cs e}) = ((k, m') \# (cl \# css)) \&$
 $((pn', i') = ((\text{env.h e}) ! k)))))) s) \& ((\%s. (((\%p, m, e). ((\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc}$
 $(\text{Suc} (\text{Suc} (0::\text{nat})))))) < \text{MAXMEM})) s) \&$
 $((\%s. \text{True}) s))) s))) s) \& ((\%s. (((\%pc, m, e). (EX n3. (((m C3) = (\text{NAT } n3))$
 $\& ((EX n3'. (((\text{callmem } e C3) = (\text{NAT } n3')) \& (EX n1'. (((\text{callmem } e C1) =$
 $(\text{NAT } n1')) \& (EX n2'. (((\text{callmem } e C2) = (\text{NAT } n2')) \& (n3 = (n3' + (n1' *$
 $n2')))))))) \& (((m \text{Zero}) = (\text{NAT } (0::\text{nat}))) \& (((m \text{One}) = (\text{NAT } (1::\text{nat}))) \&$
 $((m \text{Ra}) = (\text{NAT } R)) \& (((m R) = (\text{RA } (\text{incA } (\text{callpc } e)))) \& ((EX s. (((m S)$
 $= (\text{NAT } s)) \& ((\text{Soff } \leq s) \& (EX n1. (((m C1) = (\text{NAT } n1)) \& ((s + n1) \leq$
 $\text{MAX})))))) \& (\text{ALL } Z. (((Z \sim = R) \& ((Z \sim = C1) \& ((Z \sim = C3) \& (EX s. (((m$
 $S) = (\text{NAT } s)) \& (Z < s)))))) -->$
 $((m Z) = (\text{callmem } e Z)))))) s) \& ((\%s. \text{True}) s))) s) (((\text{Suc} (0::\text{nat}),$
 $(\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (0::\text{nat}))))))))))), (\text{update } m \text{ ta } (m$
 $sa)), (\text{env.h-update} ((\text{env.h e}) @ [((\text{Suc} (0::\text{nat}), (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc}$
 $(\text{Suc} (0::\text{nat})))))))] e)))] (%ra. \text{False}) (m (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc}$
 $(0::\text{nat}))))))))) (%ra. \text{False}) (m (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (0::\text{nat})))))))))$
 $s))) s) \&$
 $((\%s. \text{True}) s))) s) (((\text{Suc} (0::\text{nat}), (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (0::\text{nat}))))))$
 $))) (\text{update } m (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (0::\text{nat})))))))))) (\text{lift op } - (m$


```

apply (simp add: provable-def valid-def | rule allI | rule impI)+
apply (rename-tac pn i m e)
apply (cut-tac wf-prog)
apply (drule isafeP-mono)
apply (drule vc-proof-startup)
apply assumption
apply (erule conjE | erule exE)+
apply (simp only: vc-def)

```

```

apply (simp add: numeral-0-eq-0 [THEN sym] del: numeral-0-eq-0 numeral-1-eq-1)
apply (simp only: numeral-0-eq-0 numeral-1-eq-1)
apply (simp only: suc-eq-plus1 env-simp)

```

— main proof

```

apply (case-tac prog,((pn,i),m,e) |= (initF prog))
apply (rule conjI')
  — (initF prg) implies (isafeF prg (ipc prg))
apply (rule impI)
apply (erule conjE)+
apply (simp (no-asm-simp) add: initF-def valid-def C1-def C2-def C3-def Zero-def
One-def arg1-def arg2-def S-def Soff-def Ra-def R-def MAX-def MAXMEM-def
update-def fun-upd-apply Let-def arg1-def arg2-def)
apply (rule allI)
apply (rule impI)
apply (simp add: update-def nth-append callmem-def callstate-def)

```

```

apply (simp only: initF-def valid-def split-def fst-conv snd-conv)
apply (erule conjE)+
apply (simp add: initF-def valid-def C1-def C2-def C3-def Zero-def One-def arg1-def
arg2-def S-def Soff-def Ra-def R-def MAX-def MAXMEM-def update-def fun-upd-apply
del:Let-def)

```

— case not (initF prg)

```

thm isafeP-elim
apply (erule isafeP-elim)
apply simp

```

```

apply (rule conjI)
  — initF - - ; isafe (0,0)
apply (rule impI)
apply (erule conjE)+
apply (simp add: split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def
C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def update-def
fun-upd-apply)
apply (rule allI)
apply (rule impI)
apply (erule conjE)+

```

apply (*simp add: callmem-def*)

apply (*rule conjI*)

— (0,6) - - ι (1,0)

apply (*rule impI*)

apply (*erule conjE* | *erule exE*)+

apply (*simp only: MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def
arg1-def arg2-def S-def Soff-def Ra-def R-def*)

apply (*rule conjI*)

apply (*rule-tac x=6 in exI*)

apply (*simp add: update-def*)

apply (*rule conjI*)

apply (*rule-tac x=9 in exI*)

apply (*simp add: update-def*)

apply (*rule conjI*)

apply (*rule allI*)

apply (*simp add: update-def*)

apply (*rule conjI*)

apply (*rule allI*)

apply (*simp add: update-def*)

apply (*rule conjI*)

apply (*rule-tac x=2 in exI*)

apply (*simp add: update-def*)

apply (*simp add: update-def*)

apply (*rule conjI*)

apply (*simp add: callpc-def callstate-def nth-append incA-def*)

apply (*rule allI*)

apply (*rule impI*)

apply (*simp add: callmem-def*)

apply (*rule conjI*)

— (1,0) - - ι (1,5)

apply (*rule impI*)

apply (*erule conjE* | *erule exE*)+

apply (*simp add: numeral-0-eq-0 [THEN sym] del: numeral-0-eq-0 numeral-1-eq-1*)

apply (*simp only: numeral-0-eq-0 numeral-1-eq-1*)

apply (*rule conjI*)

apply (*rule-tac x=n1 in exI*)

apply (*simp add: MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def
arg1-def arg2-def S-def Soff-def Ra-def R-def update-def fun-upd-apply*)

apply (*rule conjI*)

apply (*rule-tac* $x=0$ **in** exI)
apply (*simp add:* *MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def*
arg1-def arg2-def S-def Soff-def Ra-def R-def update-def fun-upd-apply)

apply (*rule conjI*)
apply (*simp add:* *MAXMEM-def*)

apply (*rule conjI*)
apply (*simp add:* *MAXMEM-def*)

apply (*rule conjI*)
apply (*rule impI*)
apply (*rule conjI'*)
apply (*case-tac css*)
apply (*simp add:* *Let-def split-def fst-conv snd-conv*)

— $css =$ a list

apply (*simp add:* *Let-def split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def*
C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def callmem-def
callpc-def callenv-def incA-def)
apply (*rule-tac* $x=fst$ ($h e ! fst c$) **in** exI)
apply (*rule-tac* $x=snd$ ($h e ! fst c$) **in** exI)
apply (*rule conjI*)
apply (*simp add:* *update-def fun-upd-apply*)

apply (*rule-tac* $x=fst c$ **in** exI)
apply (*rule conjI'*)
apply (*rule-tac* $x=snd c$ **in** exI)
apply *simp*

apply (*simp add:* *nth-append*)

apply (*rule conjI*)
apply (*simp add:* *MAXMEM-def*)

apply (*rule-tac* $x=n3$ **in** exI)
apply (*rule conjI'*)
apply (*simp add:* *MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def*
arg1-def arg2-def S-def Soff-def Ra-def R-def update-def fun-upd-apply)

apply (*case-tac css*)
apply (*simp add:* *Let-def split-def fst-conv snd-conv*)

apply (*simp add:* *Let-def split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def*
C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def callmem-def
callpc-def callenv-def incA-def update-def)
apply (*erule conjE* | *erule exE*)
apply (*simp add:* *nth-append*)

```

apply (case-tac css)
apply (simp add: Let-def split-def fst-conv snd-conv)

apply (rule impI)
apply (simp add: Let-def split-def fst-conv snd-conv)
apply (erule exE | erule conjE)+

apply (rule conjI')
apply (rule-tac x=n in exI)
apply simp

apply (rule conjI)
apply (rule-tac x=Suc 0 in exI)
apply (simp add: One-def S-def Soff-def update-def)

apply (rule conjI)
apply (simp add: MAXMEM-def)

apply (rule conjI)
apply (simp add: MAXMEM-def)

apply (rule conjI)
apply (rule-tac x=n3 in exI)
apply (simp add: split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def
C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def update-def
fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def
nth-append update-def)

apply (rule conjI)
apply (rule-tac x=n2 in exI)
apply (simp add: split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def
C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def update-def
fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def
nth-append update-def)

apply (rule conjI)
apply (rule-tac x=n3 + n2 in exI)
apply (simp add: split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def
C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def update-def
fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def
nth-append update-def)
apply (subgoal-tac n2 ≤ n * n2)
prefer 2
apply (rule nat-mult-mono)
apply assumption
apply arith

apply (rule conjI)

```

apply (*simp add: MAXMEM-def*)

apply (*rule conjI*)
apply (*simp add: MAXMEM-def*)

apply (*rule conjI'*)
apply (*rule-tac x=s in exI*)
apply (*simp add: split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def nth-append update-def*)

apply (*rule conjI*)
apply (*simp add: MAXMEM-def*)

apply (*rule conjI*)
apply (*simp add: MAXMEM-def*)

— Problem: Stackueberlauf

apply (*rule conjI*)
apply (*rule-tac x=n1-(1::nat) in exI*)
apply (*rule conjI*)
apply (*simp add: split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def nth-append update-def*)

apply (*rule-tac x=n2 in exI*)
apply (*rule conjI*)
apply (*simp add: split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def nth-append update-def*)

apply (*rule-tac x=n3+n2 in exI*)
apply (*rule conjI*)
apply (*simp add: split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def nth-append update-def*)

apply (*rule-tac x=n1 in exI*)
apply (*rule conjI*)
apply (*simp add: split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def nth-append update-def*)

apply (*rule-tac x=n2 in exI*)

apply (*rule conjI*)
apply (*simp add: split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def nth-append update-def*)

apply (*rule-tac x=n3 in exI*)
apply (*rule conjI*)
apply (*simp add: split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def nth-append update-def*)

apply (*rule conjI'*)
apply (*subgoal-tac n1≠0*)
prefer 2
apply (*subgoal-tac n=n1*)
prefer 2
apply (*simp add: update-def MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def*)
apply (*subgoal-tac n'=0*)
prefer 2
apply (*simp add: update-def MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def*)
apply simp
apply simp

apply simp
apply (*simp only: diff-mult-distrib*)
apply simp

apply (*rule conjI*)
apply (*simp add: split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def nth-append update-def*)

apply (*rule conjI*)
apply (*simp add: split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def nth-append update-def*)

apply (*rule conjI*)
apply (*simp add: split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def nth-append update-def*)

apply (*rule conjI'*)
apply (*rule-tac x=s+(1::nat) in exI*)
apply (*simp add: split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def nth-append update-def*)
apply *arith*

apply (*rule allI*)
apply (*rule impI*)
apply (*erule conjE | erule exE*)
apply (*simp add: split-def fst-conv snd-conv MAX-def MAXMEM-def C1-def C2-def C3-def Zero-def One-def arg1-def arg2-def S-def Soff-def Ra-def R-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def nth-append update-def*)

apply (*rule conjI*)
— (1,5) nach (1,0)
apply (*rule impI*)
apply (*erule conjE | erule exE*)
apply (*rule conjI*)
apply (*rule-tac x=R in exI*)
apply (*simp add: R-def Ra-def update-def fun-upd-apply*)

apply (*rule conjI*)
apply (*rule-tac x=s in exI*)
apply (*simp (no-asm) add: S-def update-def fun-upd-apply*)
apply (*erule-tac P=m S = NAT s in rev-mp*)
apply (*simp (no-asm) add: S-def*)

apply (*rule conjI*)
apply (*rule allI*)
apply (*simp add: update-def MAXMEM-def R-def Ra-def*)

apply (*rule conjI*)
apply (*rule allI*)
apply (*erule-tac P=s + n1a ≤ MAX in rev-mp*)
apply (*erule-tac P=m S = NAT s in rev-mp*)
apply (*simp (no-asm) add: MAXMEM-def MAX-def S-def update-def*)

apply (*rule conjI*)
apply (*rule-tac x=n1 in exI*)
apply (*rule conjI*)
apply (*erule-tac P=m C1 = NAT n1 in rev-mp*)
apply (*simp (no-asm) add: C1-def update-def*)

apply (*rule-tac x=n2 in exI*)
apply (*rule conjI*)
apply (*erule-tac P=m C2 = NAT n2 in rev-mp*)

```

apply (simp (no-asm) add: C2-def update-def)

apply (rule-tac x=n3 in exI)
apply (erule-tac P=m C3 = NAT n3 in rev-mp)
apply (simp add: C3-def update-def)

apply (rule conjI)
apply (erule-tac P=m Zero = NAT 0 in rev-mp)
apply (simp (no-asm) add: Zero-def update-def)

apply (rule conjI)
apply (erule-tac P=m One = NAT 1 in rev-mp)
apply (simp (no-asm) add: One-def update-def)

apply (rule conjI)
apply (erule-tac P=m Ra = NAT R in rev-mp)
apply (simp (no-asm) add: Ra-def update-def)

apply (rule conjI)
apply (rule-tac x=s in exI)
apply (rule conjI)
apply (erule-tac P=m S = NAT s in rev-mp)
apply (simp (no-asm) add: S-def update-def)

apply (rule conjI)
apply assumption

apply (rule-tac x=n1a in exI)
apply (rule conjI)
apply (erule-tac P=m C1 = NAT n1a in rev-mp)
apply (simp (no-asm) add: C1-def update-def)
apply assumption

apply (rule conjI)
apply (simp (no-asm) add: R-def update-def incA-def split-def fst-conv snd-conv  

nth-append callpc-def)

apply (rule allI)
apply (simp (no-asm) add: R-def update-def callmem-def)

apply (rule conjI)
— (1,6) nach (1,8)
apply (rule impI)
apply (erule conjE | erule exE)+

apply (rule conjI)
apply (rule-tac x=s'a in exI)
apply (erule-tac P=m S = NAT s in rev-mp)
apply (erule-tac P=m One = NAT 1 in rev-mp)

```

```

apply (erule-tac  $P=s = s'a + 1$  in rev-mp)
apply (simp (no-asm) add: S-def One-def lift-def update-def)

apply (rule conjI)
apply (rule-tac  $x=R$  in exI)
apply (erule-tac  $P=m Ra = NAT R$  in rev-mp)
apply (simp (no-asm) add: Ra-def update-def)

apply (subgoal-tac ( $m[8 \mapsto lift\ op - (m\ 8)\ (m\ 5)]$ )  $8 = NAT\ s'a$ )
prefer 2
apply (erule-tac  $P=m S = NAT\ s$  in rev-mp)
apply (erule-tac  $P=s = s'a + 1$  in rev-mp)
apply (erule-tac  $P=m One = NAT\ 1$  in rev-mp)
apply (simp (no-asm) add: One-def S-def lift-def update-def)
apply (erule-tac  $P=(m[8 \mapsto lift\ op - (m\ 8)\ (m\ 5)]$ )  $8 = NAT\ s'a$  in rev-mp)
apply (simp (no-asm) add: cases)
apply (rule impI)
apply (subgoal-tac ( $m[8 \mapsto lift\ op - (m\ 8)\ (m\ 5)]$ )
  ( $Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ 0)))))) = NAT\ 6$ ))
prefer 2
apply (erule-tac  $P=m Ra = NAT R$  in rev-mp)
apply (simp (no-asm) add: Ra-def R-def update-def nat-number)
apply (erule-tac  $P=(m[8 \mapsto lift\ op - (m\ 8)\ (m\ 5)]$ )( $Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ 0)))))) = NAT\ 6$  in rev-mp)
apply (simp (no-asm) add: cases)
apply (rule impI)
apply (rule conjI)
apply (erule-tac  $P=s + n1 \leq MAX$  in rev-mp)
apply (erule-tac  $P=s = s'a + 1$  in rev-mp)
apply (simp (no-asm) add: MAX-def MAXMEM-def)

apply (rule conjI)
apply (rule allI)
apply (erule-tac  $P=m Ra = NAT R$  in rev-mp)
apply (simp (no-asm) add: Ra-def R-def update-def MAXMEM-def)

apply (rule conjI)
apply (rule-tac  $x=fst\ (\tilde{p}c\ e)$  in exI)
apply (rule-tac  $x=snd\ (\tilde{p}c\ e)$  in exI)
apply (rule conjI)
apply (erule-tac  $P=m s'a = RA\ (incA\ (\tilde{p}c\ e))$  in rev-mp)
apply (simp (no-asm) add: nat-number update-def incA-def)
apply (rule conjI)
apply (rule impI)
apply (erule-tac  $P=m S = NAT\ s$  in rev-mp)
apply (simp (no-asm) add: S-def nat-number)
apply (rule impI)+
apply (simp (no-asm) add: split-def)

```

```

apply (case-tac css)
apply (simp add: Let-def split-def fst-conv snd-conv)

apply (simp add: Let-def split-def fst-conv snd-conv)
apply (erule conjE)+
apply (rule-tac x=fst c in exI)
apply (rule conjI)
apply (rule-tac x=snd c in exI)
apply (simp (no-asm))

apply (erule-tac P=fst c < length (h e) in rev-mp)
apply (erule-tac P=cs e = c # a # list in rev-mp)
apply (simp (no-asm) add: nth-append callpc-def)

apply (rule conjI)
apply (simp (no-asm) add: MAXMEM-def)

apply (rule-tac x=n3 in exI)
apply (rule conjI)
apply (simp add: C3-def update-def)

apply (rule conjI)
apply (rule-tac x=n3' in exI)
apply (simp add: C3-def update-def callmem-def)

apply (rule conjI)
apply (simp add: Zero-def update-def)

apply (rule conjI)
apply (simp add: One-def update-def)

apply (rule conjI)
apply (simp add: Ra-def R-def update-def)

apply (rule conjI)
apply (simp add: R-def Ra-def incA-def callpc-def update-def)
apply (rule conjI)
apply (rule classical)
apply (simp add: S-def)

apply (rule impI)
apply (case-tac css)
apply (simp add: Let-def split-def fst-conv snd-conv)

apply (simp add: Let-def split-def nth-append)

apply (rule conjI)
apply (rule-tac x=s'a in exI)
apply (rule conjI)

```

```

apply (erule-tac  $P=m S = NAT s$  in rev-mp)
apply (erule-tac  $P=s = s'a + 1$  in rev-mp)
apply (erule-tac  $P=m One = NAT 1$  in rev-mp)
apply (simp (no-asm) add: S-def update-def One-def lift-def)
apply (rule conjI)
— Idee: Wenn  $S=Soff=9$ , dann  $s'a=8$ , dann  $m 8 = NAT$  und  $m 8 = RA$  .. Typ-
konflikt
apply (case-tac  $s = 9$ )
apply (subgoal-tac  $s'a = 8$ )
  prefer 2
  apply arith
apply (erule-tac  $P=m S = NAT s$  in rev-mp)
apply (erule-tac  $P=m s'a = RA (incA \hat{p}c e)$  in rev-mp)
apply (erule-tac  $P=s'a = 8$  in rev-mp)
apply (simp (no-asm) add: S-def)

apply (erule-tac  $P=Soff \leq s$  in rev-mp)
apply (erule-tac  $P=s \neq 9$  in rev-mp)
apply (erule-tac  $P=s=s'a + 1$  in rev-mp)
apply (simp (no-asm) add: S-def Soff-def)

apply (rule-tac  $x=n1$  in exI)
apply (simp add: C1-def update-def)

apply (rule allI)
apply (rule impI)+

apply (erule conjE | erule exE)+
apply (subgoal-tac  $sa=s'a$ )
  prefer 2
  apply (erule-tac  $P=(m[8 \mapsto lift\ op - (m\ 8)\ (m\ 5)])\ 8 = NAT\ s'a$  in rev-mp)
  apply (erule-tac  $P=(m[8 \mapsto lift\ op - (m\ 8)\ (m\ 5)])[6 \mapsto$ 
     $(m[8 \mapsto lift\ op - (m\ 8)\ (m\ 5)])\ s'a]$ 
     $S =$ 
     $NAT\ sa$  in rev-mp)
  apply (simp (no-asm) add: update-def S-def)

apply (erule-tac  $P=Z \neq R$  in rev-mp)
apply (case-tac  $Z = S$ )
apply (erule-tac  $P=Z = S$  in rev-mp)
apply (erule-tac  $P=(m[8 \mapsto lift\ op - (m\ 8)\ (m\ 5)])\ 8 = NAT\ s'a$  in rev-mp)
apply (erule-tac  $P=\tilde{m}\ e\ S = NAT\ s'a$  in rev-mp)
apply (simp (no-asm) add: update-def S-def callmem-def)

apply (erule-tac  $P=Z \neq S$  in rev-mp)
apply (simp (no-asm) add: S-def R-def update-def)
apply (rule impI)+
apply (erule-tac  $x=Z$  in allE)
apply (subgoal-tac  $(\exists s. m\ S = NAT\ s \wedge Z + 1 < s)$ )

```

```

prefer 2
apply (rule-tac x=s in exI)
apply simp
apply (simp add: S-def R-def callmem-def)

apply (rule conjI)
— (1,8) nach (0,8)
apply (rule impI)
apply (erule conjE | erule exE)+
apply (simp add: Let-def split-def fst-conv snd-conv)
apply (erule conjE)+
apply (simp add: callmem-def)

apply (rule impI)
apply (erule conjE | erule exE)+
apply (simp add: Let-def split-def fst-conv snd-conv)
apply (erule conjE | erule exE)+
apply (rule conjI)
apply (rule-tac x=s in exI)
apply (erule-tac P=m S = NAT s in rev-mp)
apply (simp (no-asm) add: S-def)

apply (rule conjI)
apply (rule-tac x=Suc 0 in exI)
apply (erule-tac P=m One = NAT (Suc 0) in rev-mp)
apply (simp (no-asm) add: One-def)

apply (rule conjI)
apply (simp (no-asm) add: MAXMEM-def)

apply (rule conjI)
apply (rule-tac x=n3'a in exI)
apply (rule conjI)

apply (erule-tac P= $\overline{m}$  e(|cs := (a, b) # cssa, h := take k (h e)|) C3 = NAT n3'a
in rev-mp)
apply (simp (no-asm) add: callmem-def)

apply (rule-tac x=n1'a in exI)
apply (rule conjI)
apply (erule-tac P= $\overline{m}$  e(|cs := (a, b) # cssa, h := take k (h e)|) C1 = NAT n1'a
in rev-mp)
apply (simp (no-asm) add: callmem-def)

apply (rule-tac x=n2'a in exI)
apply (simp add: callmem-def)
apply (erule-tac t=n3' in sym)
apply (erule-tac t=n1' in sym)
apply (simp add: diff-mult-distrib)

```

```

apply(rule conjI)
apply (rule-tac x=s'a in exI)
apply (simp add: callmem-def callpc-def Let-def split-def fst-conv snd-conv)
apply (case-tac cssa)
apply (drule-tac s=(Suc 0,5) in sym)
apply (simp add: Let-def split-def fst-conv snd-conv)

apply (simp add: Let-def split-def fst-conv snd-conv nth-append)
apply (rule conjI')
apply (subgoal-tac m S = m' S)
  prefer 2
  apply (erule-tac x=S in allE)
  apply (simp add: S-def R-def C1-def C3-def Soff-def)
apply simp
apply (subgoal-tac m s'a = m' s'a)
  prefer 2
  apply (erule-tac x=s'a in allE)
  apply (simp add: Soff-def R-def C1-def C3-def)
apply simp

apply (rule allI)
apply (rule impI)
apply (erule conjE)+
apply (subgoal-tac s = Suc s'a)
  prefer 2
  apply (erule-tac x=S in allE)
  apply (subgoal-tac m S = m' S)
  prefer 2
  apply (subgoal-tac S ≠ R ∧ S ≠ C1 ∧ S ≠ C3 ∧ S < s)
  prefer 2
  apply (erule-tac P=Soff ≤ s in rev-mp)
  apply (simp (no-asm) add: Soff-def S-def R-def C1-def C3-def Soff-def)
  apply (drule mp, assumption)
  apply (simp add: callmem-def)
apply simp
apply (erule-tac x=Z in allE)
apply (erule-tac x=Z in allE)
apply (subgoal-tac m Z =  $\tilde{m}$  e Z)
  prefer 2
  apply simp
apply (subgoal-tac  $\tilde{m}$  e Z = b Z)
  prefer 2
  apply (simp add: callmem-def)
apply (simp add: callmem-def)
done

```


end