

**theory** *EX-Mult* = *SALTimeFWInst*:

## 1 Time Bounded Multiplication

We show that a machine program that multiplies two numbers A and B is type safe, contains no overflows and does not require more than MAXTIME computation steps provided the input for A and B is less than MXA or MXB respectively.

### 1.1 Variables

We use A and B to store the input values

**constdefs**

*A* :: *nat* — Faktor1

*A* ≡ 0

*B* :: *nat* — Faktor2

*B* ≡ 1

We use C as counter and R as result variable

**constdefs**

*C*::*nat* — Counter

*C* ≡ 2

*R* :: *nat* — Result

*R* ≡ 3

Both inputs should not be greater than 3.

**constdefs**

*MXA*::*nat*

*MXA* ≡ 3

*MXB*::*nat*

*MXB* ≡ 3

We start the program with these values. They could be changed to other values within [0,MXA] or [0,MXB] without affecting the safety proof

**constdefs**

*a0*::*nat*

*a0* ≡ 3

*b0*::*nat*

*b0* ≡ 3

**constdefs**

*prog* :: *SALprogram*

*prog* ≡

[  
  (0,[(*SET R 0, None*),  
     (*SET C 0, None*)],

```

      (SET A a0,None),
      (SET B b0,None),
      (JMPEQ C A 4, Some (λ (p,m,e). ∃ a b r c. (m A)=NAT a ∧ (m B) =NAT
b ∧ (m C)=NAT c ∧ (m R)=NAT r
      ∧ a <= MXA ∧ b <= MXB ∧ c <= a ∧ (r
= c * b) ∧ (tim e) = 4 + c*4)),
      (INC C, None),
      (ADD R B,None),
      (JMPB 3,None),
      (HALT, None)
    ]
  )
]

```

**lemma** [code]:

```

prog =
[
  (0,[
    (SET R 0, None),
    (SET C 0, None),
    (SET A a0, None),
    (SET B b0, None),
    (JMPEQ C A 4, Some (term (λ (pc,m,e). ∃ a b r c. (m A)=NAT a ∧ (m B)
=NAT b ∧ (m C)=NAT c ∧ (m R)=NAT r
    ∧ a <= MXA ∧ b <= MXB ∧ c <= a ∧ (r
= c * b) ∧ (tim e) = 4 + c*4))),
    (INC C, None),
    (ADD R B,None),
    (JMPB 3,None),
    (HALT, None)
  ]
)
]

```

**apply** (simp only: term-def prog-def)  
**done**

## 1.2 Verification Condition

**generate-code** (vcgSALT.ML) [term-of]

vcg = vcgSALT

prg = prog

This is the verification condition the ML program we generated for the VCG above yields.

**constdefs**

vc::SALform

vc ≡ (%s. (((%s. (((% (pc, m, e). ((pc = ((0::nat), (0::nat))) & ((cs e) =  
[[(0::nat), m]]) & ((h e) = [] & (ALL X. ((m X) = ILLEGAL)))))) s) -->  
((%s. (((%s. (((% (pc, m, e). ((0::nat) <= MAX)) s) & ((%s. (((% (pc, m, e).  
((tim e) < MAXTIME)) s) & ((%s. True) s))) s)) s) & ((%s. (((%s. (((% (pc,





The actual proof for the  $vc$  is almost entirely by automatic simplification. They only point where human insight is needed, is where we use lemma about monotonicity of multiplication

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lemma mult-mono:
   $\llbracket a \leq (b::nat); c \leq (d::nat) \rrbracket \implies a * c \leq b * d$ 
apply (simp add: mult-le-mono)
done

```

```

theorem vc-prog-holds:
  provable prog vc
apply (simp add: provable-def valid-def)
apply (rule allI)+
apply (rename-tac pn i m e)
apply (simp only: vc-def)
apply (rule impI)
apply (subgoal-tac sysinv (((pn,i),m,e),prog))
  prefer 2
  apply (cut-tac wf-prog)
  apply (cut-tac sysinv-pres)
  apply (erule-tac x=prog in allE)
  apply (erule-tac x=((pn,i),m,e) in allE)
  apply (drule isafeP-mono)
  apply (drule mp, assumption)+
  apply assumption
apply (subgoal-tac sysinv2 (((pn,i),m,e),prog))
  prefer 2
  apply (cut-tac wf-prog)
  apply (cut-tac sysinv2-pres)
  apply (erule-tac x=prog in allE)
  apply (erule-tac x=((pn,i),m,e) in allE)
  apply (drule isafeP-mono)
  apply (drule mp, assumption)+
  apply assumption
apply (subgoal-tac  $\exists c \text{css. } cs \ e = c\#\text{css}$ )
  prefer 2
  apply (rule classical)
  apply (subgoal-tac cs e = [])
  prefer 2
  apply (rule classical)
  apply (simp only: neq-Nil-conv)
  apply (simp add: sysinv.simps)
apply (erule exE)+

```

— main proof

```

apply (case-tac prog,((pn,i),m,e)  $\models$  (initF prog))
apply (simp add: prog-simps)

```

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— case not (initF prg)
apply (drule isafeP-mono)
apply (erule isafeP-elim)
apply (simp add: prog-simps)

apply (simp add: prog-simps)
apply (rule conjI)
apply (rule impI)+
apply (erule conjE | erule exE)+
apply arith

apply (rule impI)
apply (erule conjE | erule exE)+
apply (simp add: lift-def nat-number)
apply (subgoal-tac ca * b ≤ (Suc (Suc 0)) * (Suc (Suc (Suc 0))))
prefer 2
apply (rule mult-mono)
apply simp
apply simp
apply simp
done

end

```