

theory *EX-SmartCardPurse-fa = SALOverflowFWInst:*

1 Smart Card Purse - fully annotated

This program adds a credit C to a balance B if the new balance $B + C$ does not exceed an upper bound MAX (the highest number the safety policy accepts). To check this condition a procedure is called which sets a flag F to $NAT\ 0$ if this condition is violated

1.1 Program Variables

constdefs

$B :: nat$ — balance
 $B \equiv 0$
 $C :: nat$ — credit
 $C \equiv 1$
 $M :: nat$ — maximum
 $M \equiv 2$
 $F :: nat$ — flag
 $F \equiv 4$
 $P :: nat$ — return address storage
 $P \equiv 5$
 $A :: nat$ — auxiliary variable
 $A \equiv 6$

— initial values

constdefs

$B0 :: nat$
 $B0 \equiv 4$ — maximal initial balance
 $C0 :: nat$
 $C0 \equiv 2$ — initial credit

Program Code

constdefs

$prog :: SALprogram$
 $prog \equiv [$
 $(0, [(SET\ B\ B0, None),$
 $(SET\ C\ C0, None),$
 $(SET\ A\ 0, None),$
 $(CALL\ P\ 1, Some\ (\lambda\ (pc,m,e).(\exists\ b. m\ B = NAT\ b \wedge b \leq B0)$
 $\wedge m\ B = NAT\ B0$
 $\wedge m\ C = NAT\ C0$
 $\wedge m\ A = NAT\ 0)),$
 $(JMPL\ F\ A\ 2, Some\ (\lambda\ (pc,m,e).(\exists\ b. m\ B = NAT\ b \wedge b \leq B0)$
 $\wedge m\ B = NAT\ B0$
 $\wedge m\ C = NAT\ C0$
 $\wedge m\ A = NAT\ 0$
 $])$

$$\begin{aligned}
& \wedge (\exists f. m F = NAT f \\
& \quad \wedge (f=0 \longrightarrow (\exists b. m B = NAT b \wedge ((MAX \\
- C0) < b))) \\
& \quad) \\
& \quad), \\
& \quad (ADD B C, None), \\
& \quad (JMPB 0, Some TrueF)], \\
(1, [(SET M MAX, Some (\lambda (pc,m,e). m P = RA (incA (\hat{p}c e)) \\
& \quad \wedge (\exists b. m B = NAT b) \\
& \quad \wedge (\exists c. m C = NAT c) \\
& \quad \wedge (\forall x. x \neq P \longrightarrow m x = (\hat{m} e) x))), \\
(SUB M C, Some (\lambda (pc,m,e). m P = RA (incA (\hat{p}c e)) \\
& \quad \wedge (\exists b. m B = NAT b) \\
& \quad \wedge (\exists c. m C = NAT c) \\
& \quad \wedge (m M = NAT MAX) \\
& \quad \wedge (\forall x. x \neq M \wedge x \neq P \longrightarrow m x = (\hat{m} e) x) \\
& \quad)]), \\
(SET F 0, Some (\lambda (pc,m,e). m P = RA (incA (\hat{p}c e)) \\
& \quad \wedge (\exists b. m B = NAT b) \\
& \quad \wedge (\exists c. m C = NAT c \wedge m M = NAT (MAX - c)) \\
& \quad \wedge (\forall x. x \neq M \wedge x \neq P \longrightarrow m x = (\hat{m} e) x) \\
& \quad)]), \\
(JMPL M B 2, Some (\lambda (pc,m,e). m P = RA (incA (\hat{p}c e)) \\
& \quad \wedge (\exists b. m B = NAT b) \\
& \quad \wedge (\exists c. m C = NAT c \wedge m M = NAT (MAX - c)) \\
& \quad \wedge (\forall x. x \neq F \wedge x \neq M \wedge x \neq P \longrightarrow m x = (\hat{m} \\
e) x) \\
& \quad \wedge (m F = NAT 0) \\
& \quad)]), \\
(SET F 1, Some (\lambda (pc,m,e). m P = RA (incA (\hat{p}c e)) \\
& \quad \wedge (\exists b. m B = NAT b) \\
& \quad \wedge (\exists c. m C = NAT c) \\
& \quad \wedge (\forall x. x \neq F \wedge x \neq M \wedge x \neq P \longrightarrow m x = (\hat{m} \\
e) x) \\
& \quad)]), \\
(RET P, Some (\lambda (pc,m,e). (\forall x. x \neq F \wedge x \neq M \wedge x \neq P \longrightarrow m x = \\
(\hat{m} e) x) \\
& \quad \wedge (m F = NAT 0 \vee m F = NAT 1) \\
& \quad \wedge (\exists c b. m C = NAT c \wedge m B = NAT b \\
& \quad \wedge (m F = NAT 0 \longrightarrow (MAX - c) < b)) \\
& \quad)])])
\end{aligned}$$

1.2 The Verification Condition

Generate an ML file `vcgSAL.ML` that contains an executable VCG and the program source.

generate-code (`vcgSAL.ML`) [*term-of*]
`vcg = vcgSAL`

prg = prog

This is the output our ML program for the VCG yield for this example (copy and paste).

constdefs

vc::SALform

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vc ≡ (%s. (((%s. (((% (pc, m, e). ((pc = ((0::nat), (0::nat))) & ((cs e) = [((0::nat),
m)])) & ((h e) = []) & (ALL X. ((m X) = ILLEGAL)))))) s) --> (%s. (((% (pc,
m, e). ((Suc (Suc (Suc (Suc (0::nat)))))) <= MAX)) s) & (%s. (((%s. (((% (pc, m,
e). (pc = ((0::nat), (0::nat)))))) s) --> (% (pc, m, e). ((%s. (((% (pc, m, e). ((Suc
(Suc (0::nat))) <= MAX)) s) & (%s. (((%s. (((% (pc, m, e). (pc = ((0::nat), (Suc
(0::nat)))))) s) --> (% (pc, m, e). ((%s. (((% (pc, m, e). ((0::nat) <= MAX)) s)
& (%s. (((%s. (((% (pc, m, e). (pc = ((0::nat), (Suc (Suc (0::nat)))))) s) -->
((% (pc, m, e). ((%s. (((%s. True) s) & (%s. (((% (pc, m, e). ((EX b. ((m B)
= (NAT b)) & (b <= B0)) & ((m B) = (NAT B0)) & ((m C) = (NAT C0))
& ((m A) = (NAT (0::nat)))))) s) & (%s. True) s))) s))) ((0::nat), (Suc (Suc
(Suc (Suc (0::nat)))))) (update m (Suc (Suc (Suc (Suc (Suc (Suc (0::nat)))))) (NAT
(0::nat))), (h-update ((h e) @ [((0::nat), (Suc (Suc (0::nat))))] e)))) s)) s) &
((%s. True) s))) s))) ((0::nat), (Suc (Suc (0::nat))), (update m (Suc (0::nat))
(NAT (Suc (Suc (0::nat))))), (h-update ((h e) @ [((0::nat), (Suc (0::nat))] e))))
s)) s) & ((%s. True) s))) s))) ((0::nat), (Suc (0::nat)), (update m (0::nat)
(NAT (Suc (Suc (Suc (Suc (0::nat)))))) (h-update ((h e) @ [((0::nat), (0::nat))]
e)))) s)) s) & (%s. True) s))) s))) s) & (%s. (((%s. (((%s. (((%s. (((%s.
((%s. True) s) & (%s. (((% (pc, m, e). ((EX b. ((m B) = (NAT b)) & (b <=
B0)) & ((m B) = (NAT B0)) & ((m C) = (NAT C0)) & ((m A) = (NAT
(0::nat)))))) s) & (%s. True) s))) s)) s) & (%s. (((% (pc, m, e). (pc = ((0::nat),
(Suc (Suc (Suc (0::nat)))))) s) &
((%s. True) s))) s)) s) --> (% (pc, m, e). ((%s. (((% (pc, m, e). ((Suc (Suc (Suc
(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (0::nat))))))))) s) <= MAX)) s) & (%s. (((% (pc, m, e). ((m P) = (RA (incA (callpc e)))) & ((EX
b. ((m B) = (NAT b)) & ((EX c. ((m C) = (NAT c)) & (ALL x. ((x ~ = P)
--> ((m x) = (callmem e x)))))) s) & (%s. True) s))) s))) ((Suc (0::nat)),
(0::nat), (update m (Suc (Suc (Suc (Suc (Suc (0::nat)))))) (RA ((0::nat), (Suc
(Suc (Suc (Suc (0::nat)))))) (cs-update (((length (h e)), m) # (cs e)) (h-update
((h e) @ [((0::nat), (Suc (Suc (Suc (0::nat))))] e)))) s)) s) & (%s. True) s)))
s) & (%s. (((%s. (((%s. (((%s. (((%s. (((% (pc, m, e). ((EX n. ((m (Suc (Suc
(Suc (Suc (0::nat)))))) = (NAT n)) & (EX n. ((m (Suc (Suc (Suc (Suc (Suc
(Suc (0::nat)))))) = (NAT n)))) s) & (%s. (((% (pc, m, e). ((EX b. ((m B)
= (NAT b)) & (b <= B0)) & ((m B) = (NAT B0)) & ((m C) = (NAT C0))
& ((m A) = (NAT (0::nat))) & (EX f. ((m F) = (NAT f)) & (f = (0::nat))
--> (EX b. ((m B) = (NAT b)) &
((MAX - C0) < b)))))) s) & (%s. True) s))) s)) s) & (%s. (((% (pc, m,
e). ((EX n n'. ((m (Suc (Suc (Suc (Suc (0::nat)))))) = (NAT n)) & ((m (Suc
(Suc (Suc (Suc (Suc (Suc (0::nat)))))) = (NAT n')) & (n < n')))) & (pc =
((0::nat), (Suc (Suc (Suc (Suc (0::nat)))))) s) & (%s. True) s))) s)) s) -->
((% (pc, m, e). ((%s. (((%s. True) s) & (%s. (((%s. True) s) & (%s. True) s)))
s))) ((0::nat), (Suc (Suc (Suc (Suc (Suc (Suc (0::nat)))))) m, (h-update ((h
e) @ [((0::nat), (Suc (Suc (Suc (Suc (0::nat))))] e)))) s)) s) & (%s. (((%s.
((%s. (((%s. (((% (pc, m, e). ((EX n. ((m (Suc (Suc (Suc (Suc (0::nat)))))) =

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$(NAT\ n))) \ \& \ (EX\ n.\ ((m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))))) = (NAT\ n)))) \ s) \ \& \ ((\%s.\ ((\%(pc,\ m,\ e).\ ((EX\ b.\ ((m\ B) = (NAT\ b)) \ \& \ (b \leq B0))) \ \& \ ((m\ B) = (NAT\ B0)) \ \& \ ((m\ C) = (NAT\ C0)) \ \& \ ((m\ A) = (NAT\ (0::nat))) \ \&$

$(EX\ f.\ (((m\ F) = (NAT\ f)) \ \& \ (f = (0::nat)) \ \longrightarrow \ (EX\ b.\ (((m\ B) = (NAT\ b)) \ \& \ ((MAX - C0) < b)))))) \ s) \ \& \ ((\%s.\ True) \ s)) \ s)) \ s) \ \& \ ((\%s.\ ((\%(pc,\ m,\ e).\ ((EX\ n\ n'. (((m\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))) = (NAT\ n)) \ \& \ ((m\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))) = (NAT\ n')) \ \& \ (\sim\ (n < n')))) \ \& \ (pc = (0::nat), (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))) \ s) \ \& \ ((\%s.\ True) \ s)) \ s)) \ s) \ \longrightarrow \ ((\%(pc,\ m,\ e).\ ((\%s.\ ((\%(pc,\ m,\ e).\ ((EX\ n.\ ((m\ (0::nat)) = (NAT\ n)) \ \& \ (EX\ n.\ ((m\ (Suc\ (0::nat)) = (NAT\ n)) \ \& \ (EX\ n.\ (((lift\ op + (m\ (0::nat))\ (m\ (Suc\ (0::nat)))) = (NAT\ n)) \ \& \ (n \leq MAX)))))) \ s) \ \& \ ((\%s.\ ((\%s.\ ((\%(pc,\ m,\ e).\ (pc = (0::nat), (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))) \ s) \ \longrightarrow \ ((\%(pc,\ m,\ e).\ ((\%s.\ ((\%s.\ True) \ s) \ \& \ ((\%s.\ ((\%s.\ True) \ s) \ \& \ ((\%s.\ True) \ s)) \ s)) \ (((0::nat), (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))))), (update\ m\ (0::nat)\ (lift\ op + (m\ (0::nat))\ (m\ (Suc\ (0::nat))))), (h-update\ ((h\ e) @ [((0::nat), (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))) \ e])) \ s)) \ s) \ \& \ ((\%s.\ True) \ s)) \ s)) \ (((0::nat), (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))) \ e)) \ s)) \ s) \ \& \ ((\%s.\ True) \ s)) \ s)) \ s) \ \& \ ((\%s.\ ((\%s.\ ((\%s.\ ((\%s.\ ((\%s.\ ((\%s.\ True) \ s) \ \& \ ((\%s.\ ((\%s.\ True) \ s) \ \& \ ((\%s.\ True) \ s)) \ s)) \ s) \ \& \ ((\%s.\ ((\%(pc,\ m,\ e).\ (pc = (0::nat), (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))) \ s) \ \& \ ((\%s.\ True) \ s)) \ s)) \ s) \ \longrightarrow \ ((\%(pc,\ m,\ e).\ ((\%s.\ ((\%s.\ True) \ s) \ \& \ ((\%s.\ ((\%s.\ True) \ s) \ \& \ ((\%s.\ True) \ s)) \ s)) \ (((0::nat), (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))) \ e)) \ s)) \ s) \ \& \ ((\%s.\ True) \ s)) \ s)) \ s) \ \& \ ((\%s.\ ((\%s.\ ((\%s.\ ((\%s.\ ((\%s.\ ((\%s.\ ((\%s.\ True) \ s) \ \& \ ((\%s.\ ((\%s.\ True) \ s) \ \& \ ((\%s.\ True) \ s)) \ s)) \ s) \ \& \ ((\%s.\ ((\%(pc,\ m,\ e).\ ((Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))) \ s) \ \leq \ MAX)) \ s) \ \& \ ((\%s.\ ((\%(pc,\ m,\ e).\ ((m\ P) = (RA\ (incA\ (callpc\ e)))) \ \& \ ((EX\ b.\ ((m\ B) = (NAT\ b)) \ \& \ (EX\ c.\ ((m\ C) = (NAT\ c)) \ \& \ (ALL\ x.\ ((x \sim = P) \ \longrightarrow \ ((m\ x) = (callmem\ e\ x)))))) \ s) \ \& \ ((\%s.\ True) \ s)) \ s)) \ s) \ \& \ ((\%s.\ ((\%(pc,\ m,\ e).\ (pc = (Suc\ (0::nat), (0::nat))) \ s) \ \& \ ((\%s.\ True) \ s)) \ s)) \ s) \ \longrightarrow \ ((\%(pc,\ m,\ e).\ ((\%s.\ ((\%(pc,\ m,\ e).\ ((EX\ n.\ ((m\ (Suc\ (Suc\ (0::nat))) = (NAT\ n)) \ \& \ (EX\ n.\ ((m\ (Suc\ (0::nat))) = (NAT\ n)) \ s) \ \& \ ((\%s.\ ((\%(pc,\ m,\ e).\ ((m\ P) = (RA\ (incA\ (callpc\ e)))) \ \& \ ((EX\ b.\ ((m\ B) = (NAT\ b)) \ \& \ (EX\ c.\ ((m\ C) = (NAT\ c)) \ \& \ ((m\ M) = (NAT\ MAX)) \ \& \ (ALL\ x.\ ((x \sim = M) \ \& \ (x \sim = P)) \ \longrightarrow \ ((m\ x) = (callmem\ e\ x)))))) \ s) \ \& \ ((\%s.\ True) \ s)) \ s)) \ (((Suc\ (0::nat)), (Suc\ (0::nat)), (update\ m\ (Suc\ (Suc\ (0::nat)))\ (NAT\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (0::nat)))))) \ s) \ \& \ ((\%s.\ True) \ s)) \ s) \ \& \ ((\%s.\ ((\%s.\ ((\%s.\ ((\%s.\ ((\%s.\ ((\%s.\ ((\%(pc,\ m,\ e).\ ((EX\ n.\ ((m\ (Suc\ (Suc\ (0::nat))) = (NAT\ n)) \ \& \ (EX\ n.\ ((m\ (Suc\ (0::nat))) = (NAT\ n)) \ s) \ \& \ ((\%s.\ ((\%(pc,\ m,\ e).\ ((m\ P) = (RA\ (incA\ (callpc\ e)))) \ \& \ ((EX\ b.\ ((m\ B) = (NAT\ b)) \ \& \ (EX\ c.\ ((m\ C) = (NAT\ c)) \ \& \ ((m\ M) = (NAT\ MAX)) \ \& \ (ALL\ x.\ ((x \sim = M) \ \& \ (x \sim = P)) \ \longrightarrow \ ((m\ x) = (callmem\ e\ x)))))) \ s) \ \& \ ((\%s.\ True) \ s)) \ s)) \ s) \ \& \ ((\%s.\ ((\%(pc,\ m,\ e).\ (pc = (Suc\ (0::nat), (Suc\ (0::nat))) \ s) \ \& \ ((\%s.\ True) \ s)) \ s)) \ s) \ \longrightarrow \ ((\%(pc,\ m,\ e).\ ((\%s.\ ((\%(pc,\ m,\ e).\ ((0::nat) \leq \ MAX)) \ s) \ \& \ ((\%s.\ ((\%(pc,\ m,\ e).\ ((m\ P) = (RA\ (incA\ (callpc\ e)))) \ \& \ ((EX\ b.\ ((m\ B) = (NAT\ b)) \ \& \ ((EX\ c.\ ((m\ C) = (NAT\ c)) \ \& \ ((m\ M) = (NAT\ (MAX - c)))))) \ \& \ (ALL\ x.\ ((x \sim = M) \ \& \ (x \sim = P)) \ \longrightarrow \ ((m$

Then, we prove the verification condition

Manually Guided Proof

lemma *vc-prog-holds*:

provable prog vc

— start up

apply (*simp add: provable-def valid-def | rule allI | rule impI*)+

apply (*rename-tac pn i m e*)

apply (*cut-tac wf-prog*)

apply (*drule vc-proof-startup*)

apply *assumption*

apply (*erule conjE | erule exE*)+

apply (*simp only: vc-def*)

— main proof

apply (*case-tac prog,((pn,i),m,e) \models (initF prog)*)

apply (*rule context-conjI*)

— (initF prg) implies (isafeF prg (ipc prg))

apply (*simp add: initF-def valid-def B-def MAX-def B0-def A-def C-def C0-def update-def fun-upd-apply*)

apply (*simp only: initF-def valid-def split-def fst-conv snd-conv*)

apply (*erule conjE*)+

apply *simp*

— case not (initF prg)

apply (*erule isafeP-elim*s)

apply *simp*

apply (*rule conjI*)

apply (*rule impI*)

apply (*simp add: split-def fst-conv snd-conv MAX-def C0-def B0-def B-def A-def C-def update-def fun-upd-apply*)

apply (*rule conjI*)

apply (*simp add: split-def fst-conv snd-conv MAX-def C0-def B0-def B-def A-def C-def P-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append*)

apply (*rule conjI*)

apply (*simp add: split-def fst-conv snd-conv MAX-def C0-def B0-def B-def A-def C-def P-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def*)

apply (*rule conjI*)

apply (*case-tac css*)

apply (*simp add: split-def fst-conv snd-conv MAX-def C0-def B0-def B-def A-def C-def P-def M-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def*)

— case "css = a list"

apply (*simp add: split-def fst-conv snd-conv MAX-def C0-def B0-def B-def A-def C-def P-def M-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def*)

apply (*rule conjI*)
apply (*simp add: split-def fst-conv snd-conv*)
apply (*rule impI*)
apply (*erule conjE | erule exE*)
apply (*drule-tac t=s'' in sym*)
apply (*case-tac css*)

apply (*simp add: split-def fst-conv snd-conv MAX-def C0-def B0-def B-def A-def C-def P-def M-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def sysinv.simps*)

apply (*simp add: split-def fst-conv snd-conv MAX-def C0-def B0-def B-def A-def C-def P-def M-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def sysinv.simps sysinv2.simps*)

apply (*rule conjI*)
apply (*simp add: split-def fst-conv snd-conv*)
apply (*rule impI*)
apply (*erule conjE | erule exE*)
apply (*drule-tac t=s'' in sym*)
apply (*case-tac css*)
apply (*simp add: split-def fst-conv snd-conv MAX-def C0-def B0-def B-def A-def C-def P-def M-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def*)
apply (*simp add: nat-number split-def fst-conv snd-conv MAX-def C0-def B0-def B-def A-def C-def P-def F-def M-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def*)

apply (*rule conjI*)
apply (*simp add: split-def fst-conv snd-conv*)
apply (*rule impI*)
apply (*erule conjE | erule exE*)
apply (*drule-tac t=s'' in sym*)
apply (*case-tac css*)
apply (*simp add: split-def fst-conv snd-conv MAX-def C0-def B0-def B-def A-def C-def P-def M-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def*)

apply (*simp add: nat-number split-def fst-conv snd-conv MAX-def C0-def B0-def B-def A-def C-def P-def F-def M-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def nth-append lift-def Let-def*)

apply (*rule conjI*)
apply (*simp add: split-def fst-conv snd-conv*)

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apply (drule-tac t=s'' in sym)
apply (case-tac css)
apply (simp add: split-def fst-conv snd-conv MAX-def C0-def B0-def B-def A-def
C-def P-def M-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def
nth-append lift-def Let-def)
apply (simp add: nat-number split-def fst-conv snd-conv MAX-def C0-def B0-def
B-def A-def C-def P-def F-def M-def update-def fun-upd-apply incA-def callpc-def
callmem-def callenv-def nth-append lift-def Let-def)

apply (rule conjI)
apply (rule impI)
apply (erule conjE | erule exE)+
apply fastsimp

apply fastsimp

apply (rule conjI)
apply (simp add: split-def fst-conv snd-conv)
apply (drule-tac t=s'' in sym)
apply (case-tac css)
apply (simp add: split-def fst-conv snd-conv MAX-def C0-def B0-def B-def A-def
C-def P-def M-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def
nth-append lift-def Let-def)
apply (simp add: nat-number split-def fst-conv snd-conv MAX-def C0-def B0-def
B-def A-def C-def P-def F-def M-def update-def fun-upd-apply incA-def callpc-def
callmem-def callenv-def nth-append lift-def Let-def)
apply fastsimp

apply simp
apply (rule impI)
apply (erule conjE | erule exE)+
apply (drule-tac t=s'' in sym)
apply (case-tac css)
apply (simp add: split-def fst-conv snd-conv MAX-def C0-def B0-def B-def A-def
C-def P-def M-def update-def fun-upd-apply incA-def callpc-def callmem-def callenv-def
nth-append lift-def Let-def)
apply (simp add: nat-number split-def fst-conv snd-conv MAX-def C0-def B0-def
B-def A-def C-def P-def F-def M-def update-def fun-upd-apply incA-def callpc-def
callmem-def callenv-def nth-append lift-def Let-def)
apply fastsimp
done

end

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