

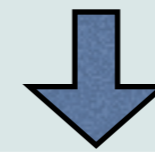
Isabelle and Proof General: Preview

Lawrence C. Paulson
University of Cambridge



Getting Started

- Install Isabelle, following instructions on the [download page](#).
- Install [Proof General](#).
- Proof General requires the editor [XEmacs](#) to be installed.
- If you have not used XEmacs before, practice on plain text files before attempting proofs!
- Launch Isabelle from the command line.
- Here, Isabelle has been installed at `/usr/local` and is used to open one of the standard theories.



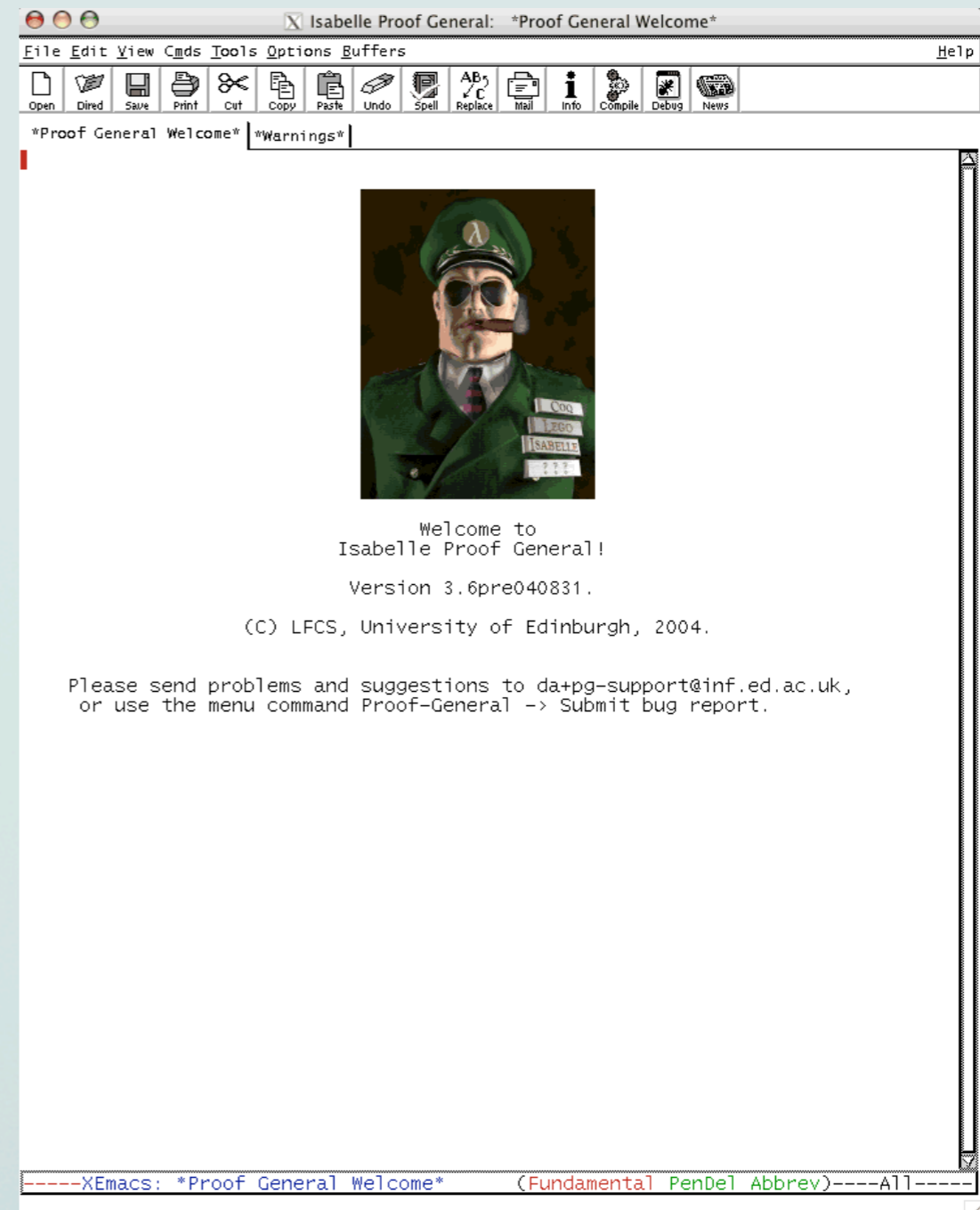
Launching Isabelle at the
UNIX command line

```
xterm
/usr/local/Isabelle: ./bin/Isabelle src/HOL/NumberTheory/Fib.thy&
[6] 5837
/usr/local/Isabelle: █
```



Proof General

- Proof General launches within XEmacs.
- If you don't see this splash screen, Proof General is not correctly installed.





The Theory File

- The theory opens in Proof General.
- Theory files visited from XEmacs also open in Proof General.
- Syntax colouring distinguishes constants, types, keywords, etc.
- The toolbar gives quick access to basic proof operations.
- This theory defines the Fibonacci function and proves theorems about it.

```
Fib.thy
(* ID:          $Id: Fib.thy,v 1.11 2005/01/14 11:00:27 nipkow Exp $
   Author:      Lawrence C Paulson, Cambridge University Computer Laboratory
   Copyright    1997 University of Cambridge
*)

header {* The Fibonacci function *}

theory Fib = Primes:

text {*
  Fibonacci numbers: proofs of laws taken from:
  R. L. Graham, D. E. Knuth, O. Patashnik. Concrete Mathematics.
  (Addison-Wesley, 1989)
*}

\bigskip
*}

consts fib :: "nat => nat"
recdef fib "measure (\lambda. x)"
  zero:   "fib 0 = 0"
  one:    "fib (Suc 0) = Suc 0"
  Suc_Suc: "fib (Suc (Suc x)) = fib x + fib (Suc x)"
[]
text {*
  \medskip The difficulty in these proofs is to ensure that the
  induction hypotheses are applied before the definition of @{term
  fib}. Towards this end, the @{term fib} equations are not declared
  to the Simplifier and are applied very selectively at first.
*}

text{*We disable @{text fib.Suc_Suc} for simplification ...*}
declare fib.Suc_Suc [simp del]

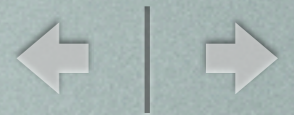
text{*...then prove a version that has a more restrictive pattern.*}
lemma fib_Suc3: "fib (Suc (Suc (Suc n))) = fib (Suc n) + fib (Suc (Suc n))"
  by (rule fib.Suc_Suc)

text {* \medskip Concrete Mathematics, page 280 *}

lemma fib_add: "fib (Suc (n + k)) = fib (Suc k) * fib (Suc n) + fib k * fib n"
  apply (induct n rule: fib.induct)
  prefer 3
  txt {* simplify the LHS just enough to apply the induction hypotheses *}
  apply (simp add: fib_Suc3)
  apply (simp_all add: fib.Suc_Suc add_mult_distrib2)
  done

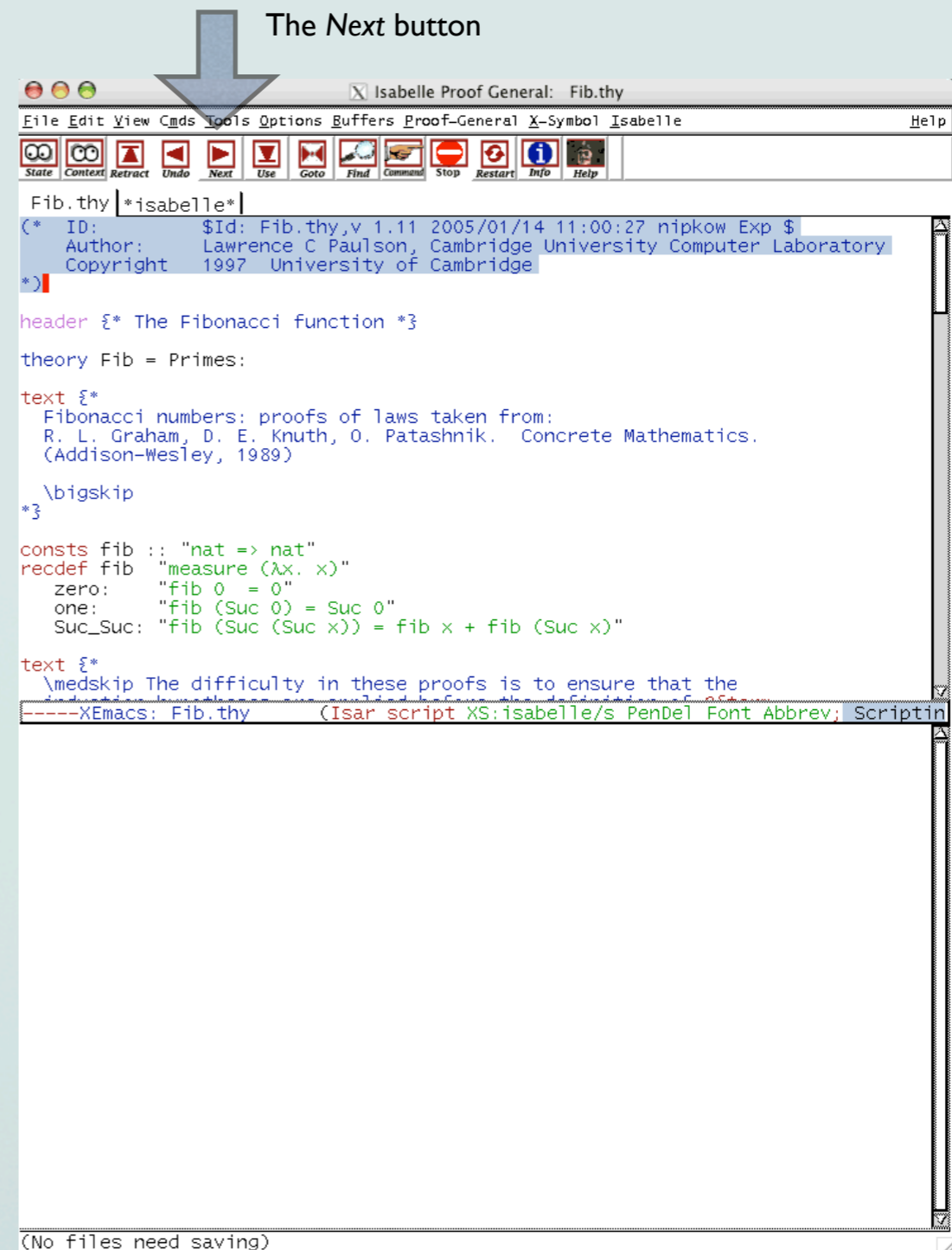
lemma fib_Suc_neq_0: "fib (Suc n) \neq 0"
  apply (induct n rule: fib.induct)
  apply (simp_all add: fib.Suc_Suc)
  done

-----XEmacs: Fib.thy (Isar script XS:isabelle/s PenDel Font Abbrev;)-----Top-
```



Basic Navigation

- A theory file contains definitions, proofs, LaTeX markup, and general commands.
- Clicking on *Next* starts Isabelle and processes the first item: a comment.
- Repeated clicks on *Next* step through the definitions.
- Proof General highlights material that has been processed in blue.





Jumping Ahead

- You can click anywhere in the theory and then click on *Goto*.
- *Goto* can even go backward, undoing declarations and commands. (To undo a single command, use the *Undo* button.)
- The header command is processed quickly, but the theory command refers to another theory.
- While Isabelle is working, Proof General highlights the corresponding text in pink.

The *Undo* button

The *Goto* button

```
Fib.thy |*isabelle*|
(* ID:          $Id: Fib.thy,v 1.11 2005/01/14 11:00:27 nipkow Exp $
   Author:      Lawrence C Paulson, Cambridge University Computer Laboratory
   Copyright    1997 University of Cambridge
*)

header {* The Fibonacci function *}

theory Fib = Primes:

text {*
  Fibonacci numbers: proofs of laws taken from:
  R. L. Graham, D. E. Knuth, O. Patashnik. Concrete Mathematics.
  (Addison-Wesley, 1989)
  \bigskip
*}

consts fib :: "nat => nat"
recdef fib "measure (\lambda. x)"
  zero:  "fib 0 = 0"
  one:   "fib (Suc 0) = Suc 0"
  Suc_Suc: "fib (Suc (Suc x)) = fib x + fib (Suc x)"

text {*
  \medskip The difficulty in these proofs is to ensure that the
  \medskip
  -----XEmacs: Fib.thy (Isar script XS:isabelle/s PenDeI Font Abbrev; Scriptin

Simple arithmetic decision procedure failed.
Now trying full Presburger arithmetic...

[Isabelle] ### Search depth = 1
```



Running a Proof

- We are about to replay a small proof relating the Fibonacci function, addition and multiplication.
- Processing the `lemma` command displays one subgoal in the proof window.
- The commands `lemma`, `theorem` and `corollary` are essentially equivalent.

```
Fib.thy |*isabelle*|
text{*We disable @{text fib.Suc_Suc} for simplification ...*}
declare fib.Suc_Suc [simp del]

text{*... then prove a version that has a more restrictive pattern.*}
lemma fib_Suc3: "fib (Suc (Suc (Suc n))) = fib (Suc n) + fib (Suc (Suc n))"
  by (rule fib.Suc_Suc)

text {* \medskip Concrete Mathematics, page 280 *}

lemma fib_add: "fib (Suc (n + k)) = fib (Suc k) * fib (Suc n) + fib k * fib n"
  apply (induct n rule: fib.induct)
  prefer 3
  txt {* simplify the LHS just enough to apply the induction hypotheses *}
  apply (simp add: fib_Suc3)
  apply (simp_all add: fib.Suc_Suc add_mult_distrib2)
  done

lemma fib_Suc_neq_0: "fib (Suc n)  $\neq$  0"
  apply (induct n rule: fib.induct)
  apply (simp_all add: fib.Suc_Suc)
  done

lemma fib_Suc_gr_0: "0 < fib (Suc n)"
  by (insert fib_Suc_neq_0 [of n], simp)

-----XEmacs: Fib.thy (Isar script XS:isabelle/s PenDel Font Abbrev; Scriptin
proof (prove): step 0

fixed variables: n, k

goal (lemma (fib_add), 1 subgoal):
1. fib (Suc (n + k)) = fib (Suc k) * fib (Suc n) + fib k * fib n
```



Performing Induction

- The first command performs induction on n using the rule `fib.induct`.
- Isabelle produced this rule while processing the recursive definition of the Fibonacci function.
- The proof window now displays three subgoals.

```
Isabelle Proof General: Fib.thy
File Edit View Cnds Tools Options Buffers Proof-General X-Symbol Isabelle Help
State Context Retract Undo Next Use Goto Find Command Stop Restart Info Help
Fib.thy |*isabelle*|
text{*We disable @{text fib.Suc_Suc} for simplification ...*}
declare fib.Suc_Suc [simp del]
text{*... then prove a version that has a more restrictive pattern.*}
lemma fib_Suc3: "fib (Suc (Suc (Suc n))) = fib (Suc n) + fib (Suc (Suc n))"
  by (rule fib.Suc_Suc)
text {* \medskip Concrete Mathematics, page 280 *}
lemma fib_add: "fib (Suc (n + k)) = fib (Suc k) * fib (Suc n) + fib k * fib n"
  apply (induct n rule: fib.induct)
  prefer 3
  txt {* simplify the LHS just enough to apply the induction hypotheses *}
  apply (simp add: fib_Suc3)
  apply (simp_all add: fib.Suc_Suc add_mult_distrib2)
  done
lemma fib_Suc_neq_0: "fib (Suc n)  $\neq$  0"
  apply (induct n rule: fib.induct)
  apply (simp_all add: fib.Suc_Suc)
  done
lemma fib_Suc_gr_0: "0 < fib (Suc n)"
  by (insert fib_Suc_neq_0 [of n], simp)
-----XEmacs: Fib.thy (Isar script XS:isabelle/s PenDel Font Abbrev; Scriptin
proof (prove): step 1
fixed variables: n, k
goal (lemma (fib_add), 3 subgoals):
1. fib (Suc (0 + k)) = fib (Suc k) * fib (Suc 0) + fib k * fib 0
2. fib (Suc (Suc 0 + k)) =
   fib (Suc k) * fib (Suc (Suc 0)) + fib k * fib (Suc 0)
3.  $\forall x$ . [fib (Suc (Suc x + k)) =
   fib (Suc k) * fib (Suc (Suc x)) + fib k * fib (Suc x);
   fib (Suc (x + k)) = fib (Suc k) * fib (Suc x) + fib k * fib x]
 $\Rightarrow$  fib (Suc (Suc (Suc x) + k)) =
   fib (Suc k) * fib (Suc (Suc (Suc x))) + fib k * fib (Suc (Suc x))
```




A Rewriting Step

- The third subgoal is selected: prefer 3.
- Then, it is simplified with the help of a rewrite rule called fib_Suc3.
- This subgoal is still rather complicated!

```
Isabelle Proof General: Fib.thy
File Edit View Cnds Tools Options Buffers Proof-General X-Symbol Isabelle Help
State Context Retract Undo Next Use Goto Find Command Stop Restart Info Help
Fib.thy |*isabelle*|
text{*We disable @{text fib.Suc_Suc} for simplification ...*}
declare fib.Suc_Suc [simp del]

text{*... then prove a version that has a more restrictive pattern.*}
lemma fib_Suc3: "fib (Suc (Suc (Suc n))) = fib (Suc n) + fib (Suc (Suc n))"
  by (rule fib.Suc_Suc)

text {* \medskip Concrete Mathematics, page 280 *}

lemma fib_add: "fib (Suc (n + k)) = fib (Suc k) * fib (Suc n) + fib k * fib n"
  apply (induct n rule: fib.induct)
  prefer 3
  txt {* simplify the LHS just enough to apply the induction hypotheses *}
  apply (simp add: fib_Suc3)
  apply (simp_all add: fib.Suc_Suc add_mult_distrib2)
  done

lemma fib_Suc_neq_0: "fib (Suc n) ≠ 0"
  apply (induct n rule: fib.induct)
  apply (simp_all add: fib.Suc_Suc)
  done

lemma fib_Suc_gr_0: "0 < fib (Suc n)"
  by (insert fib_Suc_neq_0 [of n], simp)

-----XEmacs: Fib.thy (Isar script XS:isabelle/s PenDel Font Abbrev; Scriptin
proof (prove): step 4
fixed variables: n, k
goal (lemma (fib_add), 3 subgoals):
1.  $\forall x. [\text{fib} (\text{Suc} (\text{Suc} (x + k))) = \text{fib} (\text{Suc} k) * \text{fib} (\text{Suc} (\text{Suc} x)) + \text{fib} k * \text{fib} (\text{Suc} x); \text{fib} (\text{Suc} (x + k)) = \text{fib} (\text{Suc} k) * \text{fib} (\text{Suc} x) + \text{fib} k * \text{fib} x] \Rightarrow \text{fib} (\text{Suc} k) * \text{fib} (\text{Suc} x) + \text{fib} k * \text{fib} x + (\text{fib} (\text{Suc} k) * \text{fib} (\text{Suc} (\text{Suc} x)) + \text{fib} k * \text{fib} (\text{Suc} x)) = \text{fib} (\text{Suc} k) * (\text{fib} (\text{Suc} x) + \text{fib} (\text{Suc} (\text{Suc} x))) + \text{fib} k * \text{fib} (\text{Suc} (\text{Suc} x))$ 
2.  $\text{fib} (\text{Suc} (0 + k)) = \text{fib} (\text{Suc} k) * \text{fib} (\text{Suc} 0) + \text{fib} k * \text{fib} 0$ 
3.  $\text{fib} (\text{Suc} (\text{Suc} 0 + k)) = \text{fib} (\text{Suc} k) * \text{fib} (\text{Suc} (\text{Suc} 0)) + \text{fib} k * \text{fib} (\text{Suc} 0)$ 
```



Finishing the Proof

- Next, all three subgoals are simplified, with the help of the rewrite rules shown.
- The simplifier automatically includes hundreds of other rewrite rules, as well as various decision procedures.
- This time, no subgoals remain.

```
Isabelle Proof General: Fib.thy
File Edit View Cnds Tools Options Buffers Proof-General X-Symbol Isabelle Help
State Context Retract Undo Next Use Goto Find Command Stop Restart Info Help
Fib.thy |*isabelle*|
text{*We disable @{text fib.Suc_Suc} for simplification ...*}
declare fib.Suc_Suc [simp del]

text{*... then prove a version that has a more restrictive pattern.*}
lemma fib_Suc3: "fib (Suc (Suc (Suc n))) = fib (Suc n) + fib (Suc (Suc n))"
  by (rule fib.Suc_Suc)

text {* \medskip Concrete Mathematics, page 280 *}

lemma fib_add: "fib (Suc (n + k)) = fib (Suc k) * fib (Suc n) + fib k * fib n"
  apply (induct n rule: fib.induct)
  prefer 3
  txt {* simplify the LHS just enough to apply the induction hypotheses *}
  apply (simp add: fib_Suc3)
  apply (simp_all add: fib.Suc_Suc add_mult_distrib2)
  done

lemma fib_Suc_neq_0: "fib (Suc n)  $\neq$  0"
  apply (induct n rule: fib.induct)
  apply (simp_all add: fib.Suc_Suc)
  done

lemma fib_Suc_gr_0: "0 < fib (Suc n)"
  by (insert fib_Suc_neq_0 [of n], simp)

-----XEmacs: Fib.thy (Isar script XS:isabelle/s PenDel Font Abbrev; Scriptin
proof (prove): step 5
fixed variables: n, k
goal (lemma (fib_add)):
No subgoals!
```



Storing the Theorem

- The done command causes Isabelle to accept the proof, storing the theorem.
- If you were proving this theorem for the first time, you would try various commands right in the editor buffer. You would use *Undo* frequently!
- Once you have succeeded, the file will hold the final version of your proof.
- Using *Undo* on a done command moves the cursor above its proof. Isabelle “forgets” the theorem.

```
Isabelle Proof General: Fib.thy
File Edit View Cnds Tools Options Buffers Proof-General X-Symbol Isabelle Help
State Context Retract Undo Next Use Goto Find Command Stop Restart Info Help
Fib.thy |*isabelle*|
text{*We disable @{text fib.Suc_Suc} for simplification ...*}
declare fib.Suc_Suc [simp del]

text{*... then prove a version that has a more restrictive pattern.*}
lemma fib_Suc3: "fib (Suc (Suc (Suc n))) = fib (Suc n) + fib (Suc (Suc n))"
  by (rule fib.Suc_Suc)

text {* \medskip Concrete Mathematics, page 280 *}

lemma fib_add: "fib (Suc (n + k)) = fib (Suc k) * fib (Suc n) + fib k * fib n"
  apply (induct n rule: fib.induct)
  prefer 3
  txt {* simplify the LHS just enough to apply the induction hypotheses *}
  apply (simp add: fib_Suc3)
  apply (simp_all add: fib.Suc_Suc add_mult_distrib2)
  done

lemma fib_Suc_neq_0: "fib (Suc n) ≠ 0"
  apply (induct n rule: fib.induct)
  apply (simp_all add: fib.Suc_Suc)
  done

lemma fib_Suc_gr_0: "0 < fib (Suc n)"
  by (insert fib_Suc_neq_0 [of n], simp)

-----XEmacs: Fib.thy (Isar script XS:isabelle/s PenDel Font Abbrev; Scriptin
lemma
  fib_add:
    fib (Suc (?n + ?k)) = fib (Suc ?k) * fib (Suc ?n) + fib ?k * fib ?n
```



Processing a Theory

- To run a theory right to the end, click on the *Use* button.
- Now the rest of the theory appears in pink until Isabelle can process it.

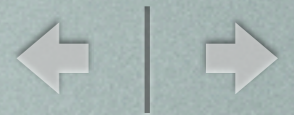
The *Use* button

```
Fib.thy |*isabelle*|
lemma fib_gcd: "fib (gcd (m, n)) = gcd (fib m, fib n)" -- {* Law 6.111 *}
  apply (induct m n rule: gcd_induct)
  apply (simp_all add: gcd_non_0 gcd_commute gcd_fib_mod)
  done

theorem fib_mult_eq_setsum:
  "fib (Suc n) * fib n = (∑k ∈ {..n}. fib k * fib k)"
  apply (induct n rule: fib.induct)
  apply (auto simp add: atMost_Suc fib.Suc_Suc)
  apply (simp add: add_mult_distrib add_mult_distrib2)
  done

end

--**-XEmacs: Fib.thy      (Isar script XS:isabelle/s PenDel Font Abbrev; Scriptin
lemma
  fib_add:
    fib (Suc (?n + ?k)) = fib (Suc ?k) * fib (Suc ?n) + fib ?k * fib ?n
```



Stop!

- Proof taking too long? Simplifier's looping? Clicked the wrong button? Just click on *Stop*.
- If things behave weirdly after this, perhaps Proof General has got out of sync with Isabelle.
- To get back into sync, use *Goto* to go back to the start of the current proof.
- You can use *Revert* to go back to the top of the theory file.

The Revert button

The Stop button

```
Fib.thy |*isabelle*|
lemma fib_gcd: "fib (gcd (m, n)) = gcd (fib m, fib n)" -- {* Law 6.111 *}
  apply (induct m n rule: gcd_induct)
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theorem fib_mult_eq_setsum:
  "fib (Suc n) * fib n = (∑k ∈ {..n}. fib k * fib k)"
  apply (induct n rule: fib.induct)
  apply (auto simp add: atMost_Suc fib.Suc_Suc)
  apply (simp add: add_mult_distrib add_mult_distrib2)
  done

end

--**-XEmacs: Fib.thy      (Isar script XS:isabelle/s PenDel Font Abbrev; Scriptin
*** Interrupt.
*** At command "apply".

Interrupt: script management may be in an inconsistent state
         (but it's probably okay)
```

Use C-c C-. to jump to end of processed region



Where Am I?

- If a proof fails—or is interrupted—in a long theory file, how do we locate the critical spot?
- You could simply scroll through the file until you find the end of the blue region.
- To jump right there, use the menu item Proof General > Goto Locked End. The key combination CTRL/C- . does the same thing.
- The proof was interrupted during a call to presburger, an arithmetic decision procedure.

```
Fib.thy |*isabelle*|
\medskip Concrete Mathematics, page 278: Cassini's identity. The proof is
much easier using integers, not natural numbers!
*3
lemma fib_Cassini_int:
  "int (fib (Suc (Suc n)) * fib n) =
  (if n mod 2 = 0 then int (fib (Suc n) * fib (Suc n)) - 1
  else int (fib (Suc n) * fib (Suc n)) + 1)"
  apply (induct n rule: fib.induct)
  apply (simp add: fib.Suc_Suc)
  apply (simp add: fib.Suc_Suc mod_Suc)
  apply (simp add: fib.Suc_Suc add_mult_distrib add_mult_distrib2
  mod_Suc zmult_int [symmetric])
  apply presburger
  done

text{*We now obtain a version for the natural numbers via the coercion
function @{term int}.*3}
theorem fib_Cassini:
  "fib (Suc (Suc n)) * fib n =
  (if n mod 2 = 0 then fib (Suc n) * fib (Suc n) - 1
  else fib (Suc n) * fib (Suc n) + 1)"
  apply (rule int_int_eq [THEN iffD1])
  apply (simp add: fib_Cassini_int)
  apply (subst zdiff_int [symmetric])
  apply (simp add: fib_Suc_eq_0 [of n] simp_all)
--**XEmacs: Fib.thy (Isar script XS:isabelle/s PenDel Font Abbrev; Scriptin
*** Interrupt.
*** At command "apply".

Interrupt: script management may be in an inconsistent state
        (but it's probably okay)

Mark set
```



The Proof State

- Clicking on the *State* button reveals the proof state at the given point.
- Here, there was one subgoal left when the proof was interrupted.

The *State* button

```

Isabelle Proof General: Fib.thy
File Edit View Ccmds Tools Options Buffers Proof-General X-Symbol Isabelle Help
State Context Retract Undo Next Use Goto Find Command Stop Restart Info Help

Fib.thy |*isabelle*|
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  apply (induct n rule: fib.induct)
  apply (simp add: fib.Suc_Suc)
  apply (simp add: fib.Suc_Suc mod_Suc)
  apply (simp add: fib.Suc_Suc add_mult_distrib add_mult_distrib2
    mod_Suc zmult_int [symmetric])
  apply presburger
  done

text{*We now obtain a version for the natural numbers via the coercion
function @{term int}.*3}
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  "fib (Suc (Suc n)) * fib n =
   (if n mod 2 = 0 then fib (Suc n) * fib (Suc n) - 1
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  apply (rule int_int_eq [THEN iffD1])
  apply (simp add: fib_Cassini_int)
  apply (subst zdiff_int [symmetric])
  apply (simp add: fib_Suc_eq_0 [of n] simp_all)
--**XEmacs: Fib.thy (Isar script XS:isabelle/s PenDel Font Abbrev; Scriptin
proof (prove): step 4
fixed variables: n
goal (lemma (fib_Cassini_int), 1 subgoal):
1.  $\forall x. [2 * (\text{int } (\text{fib } (\text{Suc } x)) * \text{int } (\text{fib } (\text{Suc } x))) +
\text{int } (\text{fib } x) * \text{int } (\text{fib } (\text{Suc } x)) =
(\text{if } (\text{if } \text{Suc } 0 = x \bmod 2 \text{ then } 0 \text{ else } \text{Suc } (x \bmod 2)) = 0
\text{ then } \text{int } (\text{fib } (\text{Suc } (\text{Suc } x)) * \text{fib } (\text{Suc } (\text{Suc } x))) - 1
\text{ else } \text{int } (\text{fib } (\text{Suc } (\text{Suc } x)) * \text{fib } (\text{Suc } (\text{Suc } x))) + 1];
\text{int } (\text{fib } x) * \text{int } (\text{fib } x) + \text{int } (\text{fib } (\text{Suc } x)) * \text{int } (\text{fib } x) =
(\text{if } x \bmod 2 = 0 \text{ then } \text{int } (\text{fib } (\text{Suc } x) * \text{fib } (\text{Suc } x)) - 1
\text{ else } \text{int } (\text{fib } (\text{Suc } x) * \text{fib } (\text{Suc } x)) + 1)]
\Rightarrow \text{Suc } 0 \neq x \bmod 2 \rightarrow x \bmod 2 = 0$ 
```

Use C-c C-o to rotate output buffers; C-c C-w to clear response & trace.



Finding Theorems

- Isabelle provides thousands of lemmas. How do you find the ones you need? One way is to click the *Find* button.
- Then, type some constants—or entire terms—into the XEmacs minibuffer.
- Type the term " $(_+_)*_ = _$ ", including the quotation marks, to see all theorems containing an instance of this pattern.
- The pattern " $_+_$ " matches all terms containing the infix operator $+$ because $_$ matches any term.

The *Find* button

The screenshot shows the Isabelle XEmacs editor interface. The title bar reads "Isabelle Proof General: Fib.thy". The menu bar includes "File Edit View Ccmds Tools Options Buffers Proof-General X-Symbol Isabelle Help". The toolbar contains various icons, with the "Find" icon (magnifying glass) highlighted. The main text area shows the content of "Fib.thy", including a lemma definition for "fib_Cassini_int" and a theorem "fib_Cassini". The minibuffer at the bottom displays the search results for the term "fib Suc 0".

```

Fib.thy |*isabelle*|
\medskip Concrete Mathematics, page 278: Cassini's identity. The proof is
much easier using integers, not natural numbers!
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lemma fib_Cassini_int:
  "int (fib (Suc (Suc n)) * fib n) =
   (if n mod 2 = 0 then int (fib (Suc n) * fib (Suc n)) - 1
    else int (fib (Suc n) * fib (Suc n)) + 1)"
  apply (induct n rule: fib.induct)
  apply (simp add: fib.Suc_Suc)
  apply (simp add: fib.Suc_Suc mod_Suc)
  apply (simp add: fib.Suc_Suc add_mult_distrib add_mult_distrib2
    mod_Suc zmult_int [symmetric])
  apply presburger
  done

text{*We now obtain a version for the natural numbers via the coercion
function @{term int}.*3
theorem fib_Cassini:
  "fib (Suc (Suc n)) * fib n =
   (if n mod 2 = 0 then fib (Suc n) * fib (Suc n) - 1
    else fib (Suc n) * fib (Suc n) + 1)"
  apply (rule int_int_eq [THEN iffD1])
  apply (simp add: fib_Cassini_int)
  apply (subst zdiff_int [symmetric])
  done
--**XEmacs: Fib.thy (Isar script XS:isabelle/s PenDel Font Abbrev; Scriptin
proof (prove): step 4
fixed variables: n
goal (lemma (fib_Cassini_int), 1 subgoal):
1.  $\forall x. [2 * (int (fib (Suc x)) * int (fib (Suc x))) +
  int (fib x) * int (fib (Suc x)) =
  (if (if Suc 0 = x mod 2 then 0 else Suc (x mod 2)) = 0
  then int (fib (Suc (Suc x)) * fib (Suc (Suc x))) - 1
  else int (fib (Suc (Suc x)) * fib (Suc (Suc x))) + 1);
  int (fib x) * int (fib x) + int (fib (Suc x)) * int (fib x) =
  (if x mod 2 = 0 then int (fib (Suc x) * fib (Suc x)) - 1
  else int (fib (Suc x) * fib (Suc x)) + 1)]
 $\Rightarrow$  Suc 0  $\neq$  x mod 2  $\rightarrow$  x mod 2 = 0
Find theorems containing: fib Suc 0$ 
```

Type some constants



Theorems Found

- The response buffer lists the theorems containing *all* of the listed constants.
- If you are lucky, there will be just a few rather than hundreds!
- The more patterns you type, the fewer theorems will be displayed.
- During the search, variables mentioned in the current goal are viewed as constants.

```
Fib.thy |*isabelle*|
\medskip Concrete Mathematics, page 278: Cassini's identity. The proof is
much easier using integers, not natural numbers!
*3
lemma fib_Cassini_int:
  "int (fib (Suc (Suc n)) * fib n) =
  (if n mod 2 = 0 then int (fib (Suc n) * fib (Suc n)) - 1
  else int (fib (Suc n) * fib (Suc n)) + 1)"
  apply (induct n rule: fib.induct)
  apply (simp add: fib.Suc_Suc)
  apply (simp add: fib.Suc_Suc mod_Suc)
  apply (simp add: fib.Suc_Suc add_mult_distrib add_mult_distrib2
  mod_Suc zmult_int [symmetric])
  apply presburger
  done

text{*We now obtain a version for the natural numbers via the coercion
function @{term int}.*3}
theorem fib_Cassini:
  "fib (Suc (Suc n)) * fib n =
  (if n mod 2 = 0 then fib (Suc n) * fib (Suc n) - 1
  else fib (Suc n) * fib (Suc n) + 1)"
  apply (rule int_int_eq [THEN iffD1])
  apply (simp add: fib_Cassini_int)
  apply (subst zdiff_int [symmetric])
  done

--**XEmacs: Fib.thy (Isar script XS:isabelle/s PenDel Font Abbrev; Scriptin
Facts containing constants "0" "Suc" "fib":

Fib.fib.one: fib (Suc 0) = Suc 0
Fib.fib.simps:
  fib 0 = 0
  fib (Suc 0) = Suc 0
  fib (Suc (Suc ?x)) = fib ?x + fib (Suc ?x)
Fib.fib_Suc_gr_0: 0 < fib (Suc ?n)
Fib.fib_Suc_neq_0: fib (Suc ?n) ≠ 0
Fib.fib_def:
  fib =
  wfrec (measure (λx. x))
  (λfib. nat_case 0 (nat_case (Suc 0) (λv. fib v + fib (Suc v))))
```



The Isabelle Menu

- The Isabelle menu gives access to Isabelle commands and information.
- Isabelle > Show me... provides other ways of finding theorems: matching rules and matching rewrites.
- In the example, the current subgoal has the form $x \leq y$, and matching rules displays all known theorems that can prove a conclusion of that form.

The screenshot shows the Emacs editor with the Isabelle menu open. The menu items include: Logics, Show me..., Display draft, Favourites, Settings, Start Isabelle, Exit Isabelle, Help, and a search bar. The search results show a list of theorems that can prove the goal $x \leq y$.

```

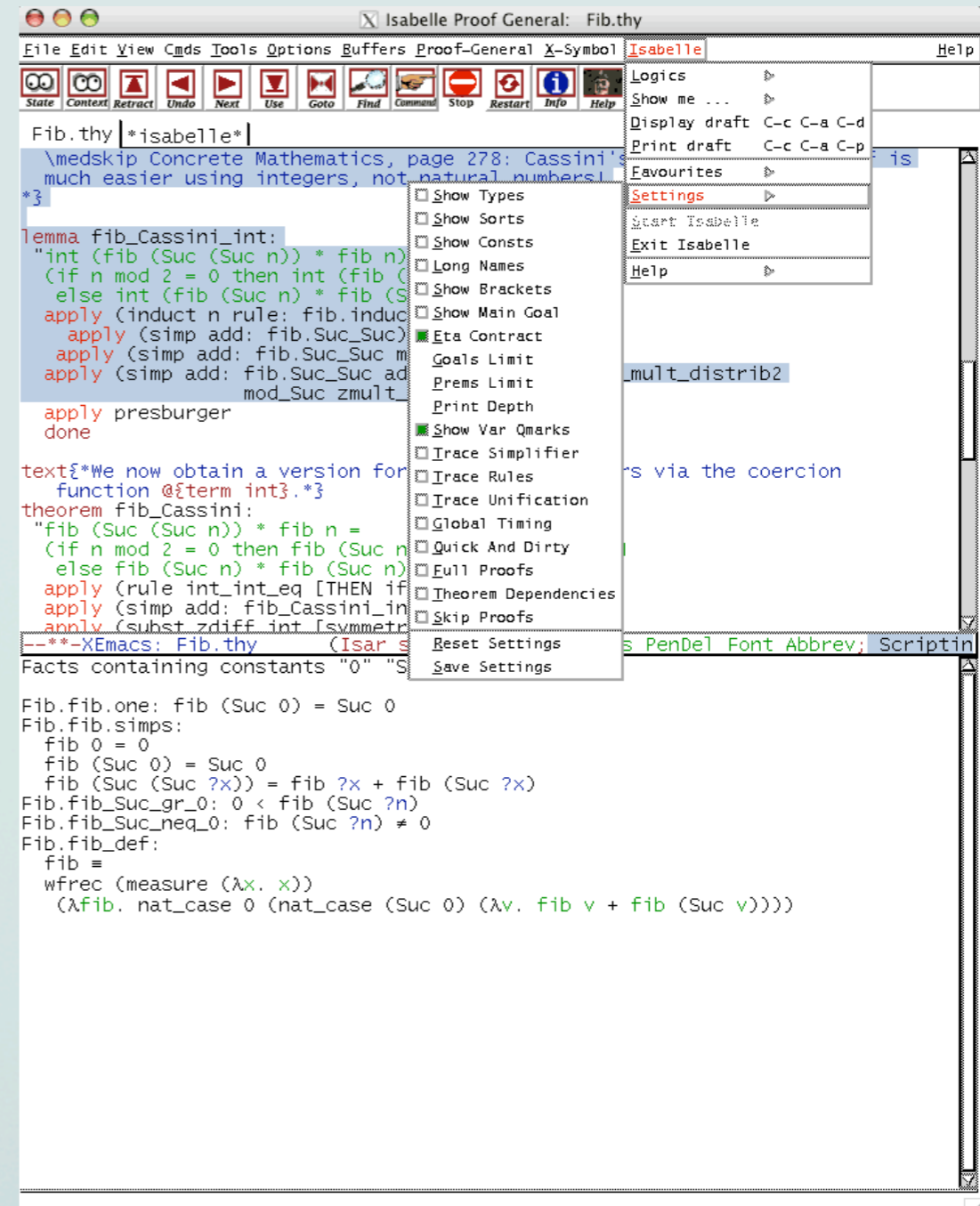
----XEmacs: Fib.thy (Isar script XS:isabelle/s PenDel Font Abbrev; Scriptin
[[?a = ?f ?b; ?b ≤ ?c;  $\forall x y. x \leq y \Rightarrow ?f x \leq ?f y$ ]  $\Rightarrow ?a \leq ?f ?c$ 
Set.ord_le_eq_subst:
[[?a ≤ ?b; ?f ?b = ?c;  $\forall x y. x \leq y \Rightarrow ?f x \leq ?f y$ ]  $\Rightarrow ?f ?a \leq ?c$ 
Set.order_subst2:
[[?a ≤ ?b; ?f ?b ≤ ?c;  $\forall x y. x \leq y \Rightarrow ?f x \leq ?f y$ ]  $\Rightarrow ?f ?a \leq ?c$ 
Set.order_subst1:
[[?a ≤ ?f ?b; ?b ≤ ?c;  $\forall x y. x \leq y \Rightarrow ?f x \leq ?f y$ ]  $\Rightarrow ?a \leq ?f ?c$ 
Set.ord_eq_le_trans: [[?a = ?b; ?b ≤ ?c]  $\Rightarrow ?a \leq ?c$ 
Set.ord_le_eq_trans: [[?a ≤ ?b; ?b = ?c]  $\Rightarrow ?a \leq ?c$ 
Orderings.order_less_imp_le:  $?x < ?y \Rightarrow ?x \leq ?y$ 
Orderings.order_eq_refl:  $?x = ?y \Rightarrow ?x \leq ?y$ 
Orderings.order_order_trans: [[?x ≤ ?y; ?y ≤ ?z]  $\Rightarrow ?x \leq ?z$ 
Orderings.order_axioms_2: [[?x ≤ ?y; ?y ≤ ?z]  $\Rightarrow ?x \leq ?z$ 
OrderedGroup.add_le_imp_le_right:  $?a + ?c \leq ?b + ?c \Rightarrow ?a \leq ?b$ 
OrderedGroup.pordered_ab_semigroup_add_imp_le.add_le_imp_le_left:
 $?c + ?a \leq ?c + ?b \Rightarrow ?a \leq ?b$ 
OrderedGroup.pordered_ab_semigroup_add_imp_le.axioms:
 $?c + ?a \leq ?c + ?b \Rightarrow ?a \leq ?b$ 
Divides.dvd_imp_le: [[?k dvd ?n; 0 < ?n]  $\Rightarrow ?k \leq ?n$ 
Divides.unique_quotient_lemma:
[[?b * ?q' + ?r' ≤ ?b * ?q + ?r; 0 < ?b; ?r < ?b]  $\Rightarrow ?q' \leq ?q$ 
Power.power_dvd_imp_le: [[?i ^ ?m dvd ?i ^ ?n; 1 < ?i]  $\Rightarrow ?m \leq ?n$ 
Power.power_le_imp_le_base:
[[?a ^ Suc ?n ≤ ?b ^ Suc ?n; (0:?'a) ≤ ?a; (0:?'a) ≤ ?b]  $\Rightarrow ?a \leq ?b$ 
Power.power_le_imp_le_exp: [[(1:?'a) < ?a; ?a ^ ?m ≤ ?a ^ ?n]  $\Rightarrow ?m \leq ?n$ 
Nat.Suc_leI:  $?m < ?n \Rightarrow \text{Suc } ?m \leq ?n$ 

```



Settings

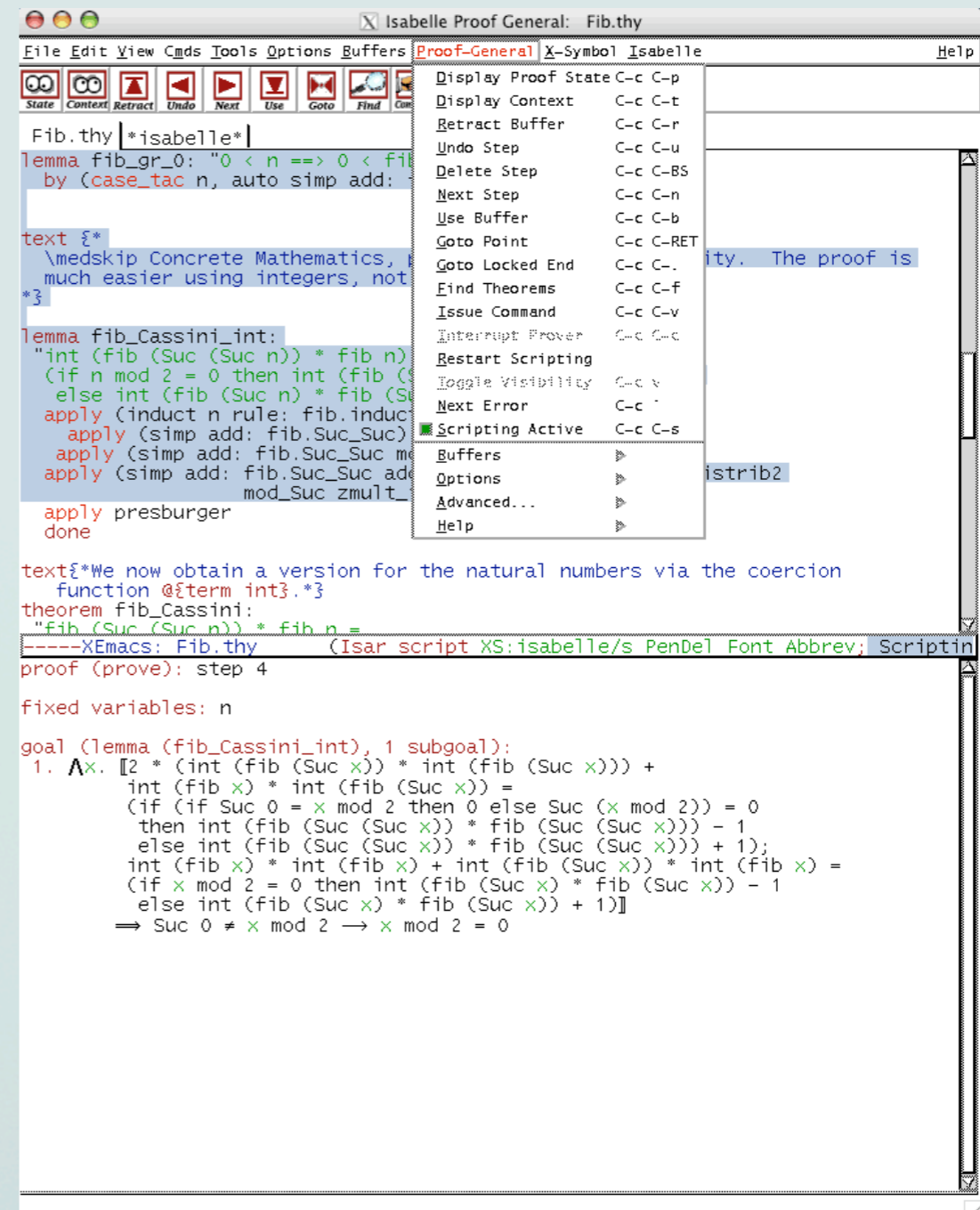
- The menu Isabelle > Settings can request the display of types, execution times, and various traces.
- There are printing options to suit special situations, such as enormous subgoals.
- Use Show Types and Show Sorts to cause more type information to be displayed.
- The various Show options make the output more verbose, but more explicit, and are helpful for diagnosing problems.

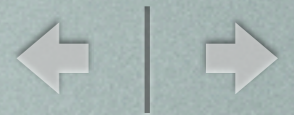




The PG Menu

- The Proof General menu gives access to many commands.
- The main commands are available from the toolbar. A notable exception is Goto Locked End.
- Choose Proof General > Buffers > Trace to see tracing output.





Mathematical Symbols

- Proof General uses the X-Symbol package to display mathematical symbols such as $\lambda \leq \neq \in \notin \cup$ and \cap
- The package is included with Proof General, but may need to be switched on.
- If X-Symbol mode is off, Proof General will display ASCII escape sequences, as shown on the right.

```

emacs: Fib.thy
File Edit View Cnds Tools Options Buffers Proof-General Isabelle Help
State Context Retract Undo Next Use Goto Find Command Stop Restart Info Help
Fib.thy
(* ID: $Id: Fib.thy,v 1.12 2005/03/25 15:20:57 paulson Exp $
  Author: Lawrence C Paulson, Cambridge University Computer Laboratory
  Copyright 1997 University of Cambridge
*)

header {* The Fibonacci function *}

theory Fib = Primes:

text {*
  Fibonacci numbers: proofs of laws taken from:
  R. L. Graham, D. E. Knuth, O. Patashnik. Concrete Mathematics.
  (Addison-Wesley, 1989)
  \bigskip
  *}

consts fib :: "nat => nat"
reodef fib "measure (\<lambda>x. x)"
zero: "fib 0 = 0"
one: "fib (Suc 0) = Suc 0"
Suc_Suc: "fib (Suc (Suc x)) = fib x + fib (Suc x)"

text {*
  \medskip The difficulty in these proofs is to ensure that the
  induction hypotheses are applied before the definition of @term
  fib}. Towards this end, the @term fib} equations are not declared
  to the Simplifier and are applied very selectively at first.
  *}

text{*We disable @text fib.Suc_Suc} for simplification ...*}
declare fib.Suc_Suc [simp del]

text{*...then prove a version that has a more restrictive pattern.*}
lemma fib_Suc3: "fib (Suc (Suc (Suc n))) = fib (Suc n) + fib (Suc (Suc n))"
  by (rule fib.Suc_Suc)

text {* \medskip Concrete Mathematics, page 280 *}

lemma fib_add: "fib (Suc (n + k)) = fib (Suc k) * fib (Suc n) + fib k * fib n"
  apply (induct n rule: fib.induct)
  prefer 3
  txt {* simplify the LHS just enough to apply the induction hypotheses *}
  apply (simp add: fib_Suc3)
  apply (simp_all add: fib.Suc_Suc add_mult_distrib2)
  done

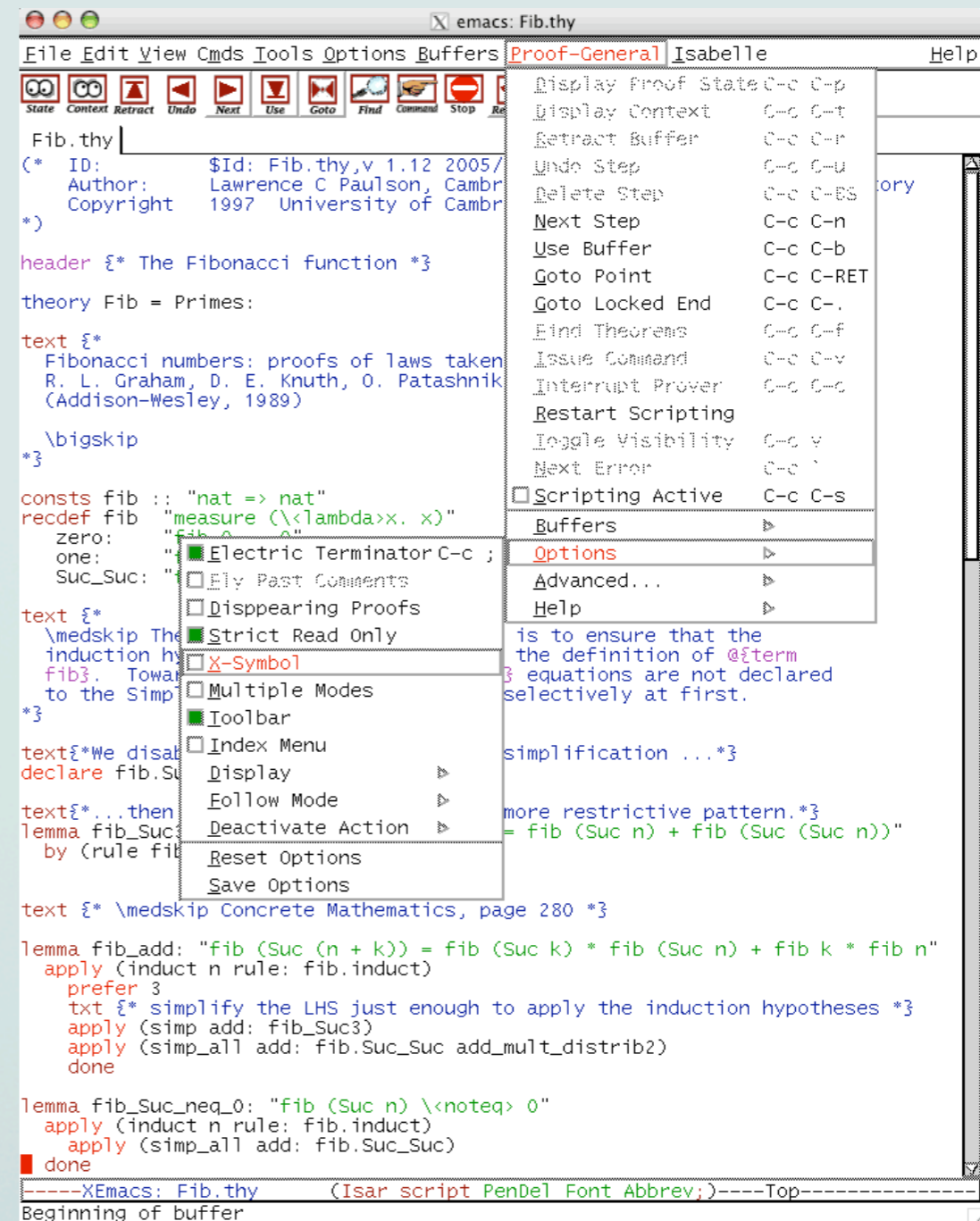
lemma fib_Suc_neq_0: "fib (Suc n) \<noteq> 0"
  apply (induct n rule: fib.induct)
  apply (simp_all add: fib.Suc_Suc)
  done
-----XEmacs: Fib.thy (Isar script PenDel Font Abbrev;)-----Top-----
Beginning of buffer

```



Enabling Symbols

- To enable X-Symbol mode, select the menu item Proof General > Options > X-Symbol.
- Then, make this setting permanent using Proof General > Options > Save Options.
- Take the time to explore the many other options and settings on offer.





Go Forth and Prove!

- Try out this theory yourself: you will find it in `src/HOL/NumberTheory/Fib.thy`.
- For more information on Isabelle, read the [documentation](#).
- For more information on Proof General, see its [user manual](#).
- Have fun!

```
Fib.thy |*isabelle*|
apply (induct m n rule: gcd_induct)
apply (simp_all add: gcd_non_0 gcd_commute gcd_fib_mod)
done

theorem fib_mult_eq_setsum:
  "fib (Suc n) * fib n = (∑k ∈ {..n}. fib k * fib k)"
  apply (induct n rule: fib.induct)
  apply (auto simp add: atMost_Suc fib.Suc_Suc)
  apply (simp add: add_mult_distrib add_mult_distrib2)
done

end
```

-----XEmacs: Fib.thy (Isar script XS:isabelle/s PenDel Font Abbrev; Scriptin