

Isabelle/HOL-NSA — Non-Standard Analysis

October 8, 2017

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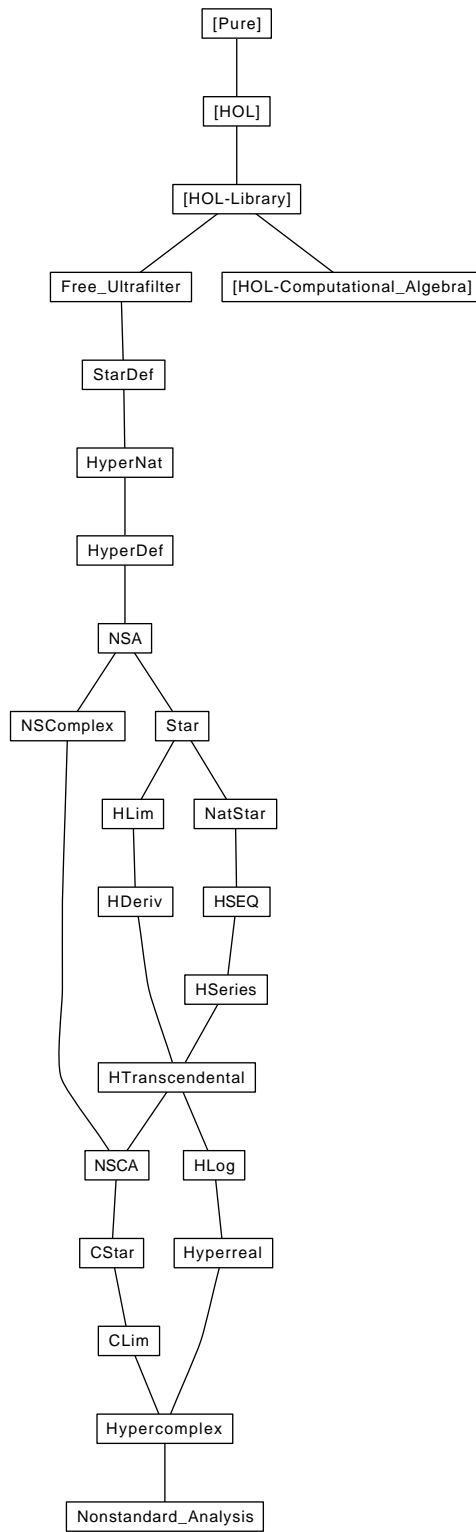
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1 Filters and Ultrafilters

```
theory Free-Ultrafilter
  imports HOL-Library.Infinite-Set
begin
```

1.1 Definitions and basic properties

1.1.1 Ultrafilters

```
locale ultrafilter =
  fixes F :: 'a filter
  assumes proper: F ≠ bot
  assumes ultra: eventually P F ∨ eventually (λx. ¬ P x) F
begin
```

```
lemma eventually-imp-frequently: frequently P F ⇒ eventually P F
  using ultra[of P] by (simp add: frequently-def)
```

```
lemma frequently-eq-eventually: frequently P F = eventually P F
  using eventually-imp-frequently eventually-frequently[OF proper] ..
```

```
lemma eventually-disj-iff: eventually (λx. P x ∨ Q x) F ↔ eventually P F ∨
  eventually Q F
  unfolding frequently-eq-eventually[symmetric] frequently-disj-iff ..
```

```
lemma eventually-all-iff: eventually (λx. ∀ y. P x y) F = (∀ Y. eventually (λx. P
  x (Y x)) F)
  using frequently-all[of P F] by (simp add: frequently-eq-eventually)
```

```
lemma eventually-imp-iff: eventually (λx. P x → Q x) F ↔ (eventually P F
  → eventually Q F)
  using frequently-imp-iff[of P Q F] by (simp add: frequently-eq-eventually)
```

```
lemma eventually-iff-iff: eventually (λx. P x ↔ Q x) F ↔ (eventually P F
  ↔ eventually Q F)
  unfolding iff-conv-conj-imp eventually-conj-iff eventually-imp-iff by simp
```

```
lemma eventually-not-iff: eventually (λx. ¬ P x) F ↔ ¬ eventually P F
  unfolding not-eventually frequently-eq-eventually ..
```

```
end
```

1.2 Maximal filter = Ultrafilter

A filter F is an ultrafilter iff it is a maximal filter, i.e. whenever G is a filter and $F \subseteq G$ then $F = G$

Lemma that shows existence of an extension to what was assumed to be a maximal filter. Will be used to derive contradiction in proof of property of

ultrafilter.

lemma *extend-filter*: $\text{frequently } P F \implies \text{inf } F (\text{principal } \{x. P x\}) \neq \text{bot}$
by (*simp add: trivial-limit-def eventually-inf-principal not-eventually*)

lemma *max-filter-ultrafilter*:

assumes $F \neq \text{bot}$

assumes *max*: $\bigwedge G. G \neq \text{bot} \implies G \leq F \implies F = G$

shows *ultrafilter* F

proof

show *eventually* $P F \vee (\forall_F x \text{ in } F. \neg P x)$ **for** P

proof (*rule disjCI*)

assume $\neg (\forall_F x \text{ in } F. \neg P x)$

then have $\text{inf } F (\text{principal } \{x. P x\}) \neq \text{bot}$

by (*simp add: not-eventually extend-filter*)

then have $F: F = \text{inf } F (\text{principal } \{x. P x\})$

by (*rule max*) *simp*

show *eventually* $P F$

by (*subst F*) (*simp add: eventually-inf-principal*)

qed

qed fact

lemma *le-filter-frequently*: $F \leq G \iff (\forall P. \text{frequently } P F \longrightarrow \text{frequently } P G)$

unfolding *frequently-def le-filter-def*

apply *auto*

apply (*erule-tac x= $\lambda x. \neg P x$ in allE*)

apply *auto*

done

lemma (*in ultrafilter*) *max-filter*:

assumes $G: G \neq \text{bot}$

and *sub*: $G \leq F$

shows $F = G$

proof (*rule antisym*)

show $F \leq G$

using *sub*

by (*auto simp: le-filter-frequently[of F] frequently-eq-eventually le-filter-def[of G]*)

intro!: eventually-frequently G proper)

qed fact

1.3 Ultrafilter Theorem

lemma *ex-max-ultrafilter*:

fixes $F :: 'a \text{ filter}$

assumes $F: F \neq \text{bot}$

shows $\exists U \leq F. \text{ultrafilter } U$

proof –

let $?X = \{G. G \neq \text{bot} \wedge G \leq F\}$

let $?R = \{(b, a). a \neq \text{bot} \wedge a \leq b \wedge b \leq F\}$

have *bot-notin-R*: $c \in \text{Chains } ?R \implies \text{bot} \notin c$ **for** c
by (*auto simp: Chains-def*)

have [*simp*]: $\text{Field } ?R = ?X$
by (*auto simp: Field-def bot-unique*)

have $\exists m \in \text{Field } ?R. \forall a \in \text{Field } ?R. (m, a) \in ?R \longrightarrow a = m$ (**is** $\exists m \in ?A. ?B m$)
proof (*rule Zorns-po-lemma*)
show *Partial-order* $?R$
by (*auto simp: partial-order-on-def preorder-on-def antisym-def refl-on-def trans-def Field-def bot-unique*)
show $\forall C \in \text{Chains } ?R. \exists u \in \text{Field } ?R. \forall a \in C. (a, u) \in ?R$
proof (*safe intro!: botI del: notI*)
fix c
assume $c: c \in \text{Chains } ?R$

have *Inf-c*: $\text{Inf } c \neq \text{bot} \implies \text{Inf } c \leq F$ **if** $c \neq \{\}$
proof –
from c **that** **have** $\text{Inf } c = \text{bot} \iff (\exists x \in c. x = \text{bot})$
unfolding *trivial-limit-def* **by** (*intro eventually-Inf-base*) (*auto simp: Chains-def*)
with c **show** $\text{Inf } c \neq \text{bot}$
by (*simp add: bot-notin-R*)
from c **obtain** x **where** $x \in c$ **by** *auto*
with c **show** $\text{Inf } c \leq F$
by (*auto intro!: Inf-lower2[of x] simp: Chains-def*)
qed

then **have** [*simp*]: $\text{inf } F (\text{Inf } c) = (\text{if } c = \{\} \text{ then } F \text{ else } \text{Inf } c)$
using c **by** (*auto simp add: inf-absorb2*)

from c **show** $\text{inf } F (\text{Inf } c) \neq \text{bot}$
by (*simp add: F Inf-c*)
from c **show** $\text{inf } F (\text{Inf } c) \in \text{Field } ?R$
by (*simp add: Chains-def Inf-c F*)

assume $x \in c$
with c **show** $\text{inf } F (\text{Inf } c) \leq x \leq F$
by (*auto intro: Inf-lower simp: Chains-def*)
qed

then **obtain** U **where** $U: U \in ?A ?B U ..$
show *?thesis*
proof
from U **show** $U \leq F \wedge \text{ultrafilter } U$
by (*auto intro!: max-filter-ultrafilter*)
qed

qed

1.3.1 Free Ultrafilters

There exists a free ultrafilter on any infinite set.

locale *freeultrafilter* = *ultrafilter* +

assumes *infinite*: *eventually P F* \implies *infinite* {*x*. *P x*}

begin

lemma *finite*: *finite* {*x*. *P x*} \implies \neg *eventually P F*

by (*erule contrapos-pn*) (*erule infinite*)

lemma *finite'*: *finite* {*x*. \neg *P x*} \implies *eventually P F*

by (*drule finite*) (*simp add: not-eventually frequently-eq-eventually*)

lemma *le-cofinite*: *F* \leq *cofinite*

by (*intro filter-leI*)

(*auto simp add: eventually-cofinite not-eventually frequently-eq-eventually dest!*:
finite)

lemma *singleton*: \neg *eventually* ($\lambda x. x = a$) *F*

by (*rule finite*) *simp*

lemma *singleton'*: \neg *eventually* (*op = a*) *F*

by (*rule finite*) *simp*

lemma *ultrafilter*: *ultrafilter F* ..

end

lemma *freeultrafilter-Ex*:

assumes [*simp*]: *infinite* (*UNIV* :: 'a set)

shows $\exists U :: 'a$ filter. *freeultrafilter U*

proof –

from *ex-max-ultrafilter*[*of cofinite* :: 'a filter]

obtain *U* :: 'a filter **where** *U* \leq *cofinite ultrafilter U*

by *auto*

interpret *ultrafilter U* **by** *fact*

have *freeultrafilter U*

proof

fix *P*

assume *eventually P U*

with *proper* **have** *frequently P U*

by (*rule eventually-frequently*)

then **have** *frequently P cofinite*

using $\langle U \leq \text{cofinite} \rangle$ **by** (*simp add: le-filter-frequently*)

then **show** *infinite* {*x*. *P x*}

by (*simp add: frequently-cofinite*)

qed

then **show** *?thesis* ..

qed

end

2 Construction of Star Types Using Ultrafilters

```
theory StarDef
  imports Free-Ultrafilter
begin
```

2.1 A Free Ultrafilter over the Naturals

```
definition FreeUltrafilterNat :: nat filter (U)
  where U = (SOME U. freeultrafilter U)
```

```
lemma freeultrafilter-FreeUltrafilterNat: freeultrafilter U
  apply (unfold FreeUltrafilterNat-def)
  apply (rule someI-ex)
  apply (rule freeultrafilter-Ex)
  apply (rule infinite-UNIV-nat)
done
```

```
interpretation FreeUltrafilterNat: freeultrafilter U
  by (rule freeultrafilter-FreeUltrafilterNat)
```

2.2 Definition of *star* type constructor

```
definition starrel :: ((nat ⇒ 'a) × (nat ⇒ 'a)) set
  where starrel = {(X, Y). eventually (λn. X n = Y n) U}
```

```
definition star = (UNIV :: (nat ⇒ 'a) set) // starrel
```

```
typedef 'a star = star :: (nat ⇒ 'a) set set
  by (auto simp: star-def intro: quotientI)
```

```
definition star-n :: (nat ⇒ 'a) ⇒ 'a star
  where star-n X = Abs-star (starrel “ {X})
```

```
theorem star-cases [case-names star-n, cases type: star]:
  obtains X where x = star-n X
  by (cases x) (auto simp: star-n-def star-def elim: quotientE)
```

```
lemma all-star-eq: (∀ x. P x) ⟷ (∀ X. P (star-n X))
  apply auto
  apply (rule-tac x = x in star-cases)
  apply simp
done
```

```
lemma ex-star-eq: (∃ x. P x) ⟷ (∃ X. P (star-n X))
  apply auto
```

```

apply (rule-tac x=x in star-cases)
apply auto
done

```

Proving that *starrel* is an equivalence relation.

```

lemma starrel-iff [iff]: (X, Y) ∈ starrel ↔ eventually (λn. X n = Y n) U
by (simp add: starrel-def)

```

```

lemma equiv-starrel: equiv UNIV starrel

```

```

proof (rule equivI)

```

```

  show refl starrel by (simp add: refl-on-def)

```

```

  show sym starrel by (simp add: sym-def eq-commute)

```

```

  show trans starrel by (intro transI) (auto elim: eventually-elim2)

```

```

qed

```

```

lemmas equiv-starrel-iff = eq-equiv-class-iff [OF equiv-starrel UNIV-I UNIV-I]

```

```

lemma starrel-in-star: starrel“{x} ∈ star
by (simp add: star-def quotientI)

```

```

lemma star-n-eq-iff: star-n X = star-n Y ↔ eventually (λn. X n = Y n) U
by (simp add: star-n-def Abs-star-inject starrel-in-star equiv-starrel-iff)

```

2.3 Transfer principle

This introduction rule starts each transfer proof.

```

lemma transfer-start: P ≡ eventually (λn. Q) U ⇒ Trueprop P ≡ Trueprop Q
by (simp add: FreeUltrafilterNat.proper)

```

Standard principles that play a central role in the transfer tactic.

```

definition Ifun :: ('a ⇒ 'b) star ⇒ 'a star ⇒ 'b star ((- ★/ -) [300, 301] 300)
where Ifun f ≡
  λx. Abs-star (⋃ F ∈ Rep-star f. ⋃ X ∈ Rep-star x. starrel“{λn. F n (X n)})

```

```

lemma Ifun-congruent2: congruent2 starrel starrel (λF X. starrel“{λn. F n (X n)})

```

```

by (auto simp add: congruent2-def equiv-starrel-iff elim!: eventually-rev-mp)

```

```

lemma Ifun-star-n: star-n F ★ star-n X = star-n (λn. F n (X n))

```

```

by (simp add: Ifun-def star-n-def Abs-star-inverse starrel-in-star
  UN-equiv-class2 [OF equiv-starrel equiv-starrel Ifun-congruent2])

```

```

lemma transfer-Ifun: f ≡ star-n F ⇒ x ≡ star-n X ⇒ f ★ x ≡ star-n (λn. F n (X n))

```

```

by (simp only: Ifun-star-n)

```

```

definition star-of :: 'a ⇒ 'a star

```

```

where star-of x ≡ star-n (λn. x)

```

Initialize transfer tactic.

ML-file *transfer.ML*

method-setup *transfer* =
 ⟨*Attrib.thms* >> (fn *ths* => fn *ctxt* => *SIMPLE-METHOD'* (*Transfer-Principle.transfer-tac*
ctxt ths))
transfer principle

Transfer introduction rules.

lemma *transfer-ex* [*transfer-intro*]:
 $(\bigwedge X. p \text{ (star-} n \text{ } X) \equiv \text{eventually } (\lambda n. P \ n \ (X \ n)) \ \mathcal{U}) \implies$
 $\exists x::'a \text{ star. } p \ x \equiv \text{eventually } (\lambda n. \exists x. P \ n \ x) \ \mathcal{U}$
by (*simp only: ex-star-eq eventually-ex*)

lemma *transfer-all* [*transfer-intro*]:
 $(\bigwedge X. p \text{ (star-} n \text{ } X) \equiv \text{eventually } (\lambda n. P \ n \ (X \ n)) \ \mathcal{U}) \implies$
 $\forall x::'a \text{ star. } p \ x \equiv \text{eventually } (\lambda n. \forall x. P \ n \ x) \ \mathcal{U}$
by (*simp only: all-star-eq FreeUltrafilterNat.eventually-all-iff*)

lemma *transfer-not* [*transfer-intro*]: $p \equiv \text{eventually } P \ \mathcal{U} \implies \neg p \equiv \text{eventually}$
 $(\lambda n. \neg P \ n) \ \mathcal{U}$
by (*simp only: FreeUltrafilterNat.eventually-not-iff*)

lemma *transfer-conj* [*transfer-intro*]:
 $p \equiv \text{eventually } P \ \mathcal{U} \implies q \equiv \text{eventually } Q \ \mathcal{U} \implies p \wedge q \equiv \text{eventually } (\lambda n. P \ n \wedge$
 $Q \ n) \ \mathcal{U}$
by (*simp only: eventually-conj-iff*)

lemma *transfer-disj* [*transfer-intro*]:
 $p \equiv \text{eventually } P \ \mathcal{U} \implies q \equiv \text{eventually } Q \ \mathcal{U} \implies p \vee q \equiv \text{eventually } (\lambda n. P \ n \vee$
 $Q \ n) \ \mathcal{U}$
by (*simp only: FreeUltrafilterNat.eventually-disj-iff*)

lemma *transfer-imp* [*transfer-intro*]:
 $p \equiv \text{eventually } P \ \mathcal{U} \implies q \equiv \text{eventually } Q \ \mathcal{U} \implies p \longrightarrow q \equiv \text{eventually } (\lambda n. P \ n$
 $\longrightarrow Q \ n) \ \mathcal{U}$
by (*simp only: FreeUltrafilterNat.eventually-imp-iff*)

lemma *transfer-iff* [*transfer-intro*]:
 $p \equiv \text{eventually } P \ \mathcal{U} \implies q \equiv \text{eventually } Q \ \mathcal{U} \implies p = q \equiv \text{eventually } (\lambda n. P \ n$
 $= Q \ n) \ \mathcal{U}$
by (*simp only: FreeUltrafilterNat.eventually-iff-iff*)

lemma *transfer-if-bool* [*transfer-intro*]:
 $p \equiv \text{eventually } P \ \mathcal{U} \implies x \equiv \text{eventually } X \ \mathcal{U} \implies y \equiv \text{eventually } Y \ \mathcal{U} \implies$
 $(\text{if } p \text{ then } x \text{ else } y) \equiv \text{eventually } (\lambda n. \text{if } P \ n \text{ then } X \ n \text{ else } Y \ n) \ \mathcal{U}$
by (*simp only: if-bool-eq-conj transfer-conj transfer-imp transfer-not*)

lemma *transfer-eq* [*transfer-intro*]:

$x \equiv \text{star-n } X \implies y \equiv \text{star-n } Y \implies x = y \equiv \text{eventually } (\lambda n. X n = Y n) \mathcal{U}$
by (*simp only: star-n-eq-iff*)

lemma *transfer-if* [*transfer-intro*]:

$p \equiv \text{eventually } (\lambda n. P n) \mathcal{U} \implies x \equiv \text{star-n } X \implies y \equiv \text{star-n } Y \implies$
 $(\text{if } p \text{ then } x \text{ else } y) \equiv \text{star-n } (\lambda n. \text{if } P n \text{ then } X n \text{ else } Y n)$
by (*rule eq-reflection*) (*auto simp: star-n-eq-iff transfer-not elim!: eventually-mono*)

lemma *transfer-fun-eq* [*transfer-intro*]:

$(\bigwedge X. f (\text{star-n } X) = g (\text{star-n } X) \equiv \text{eventually } (\lambda n. F n (X n) = G n (X n))$
 $\mathcal{U}) \implies$
 $f = g \equiv \text{eventually } (\lambda n. F n = G n) \mathcal{U}$
by (*simp only: fun-eq-iff transfer-all*)

lemma *transfer-star-n* [*transfer-intro*]: $\text{star-n } X \equiv \text{star-n } (\lambda n. X n)$

by (*rule reflexive*)

lemma *transfer-bool* [*transfer-intro*]: $p \equiv \text{eventually } (\lambda n. p) \mathcal{U}$

by (*simp add: FreeUltrafilterNat.proper*)

2.4 Standard elements

definition *Standard* :: 'a star set

where *Standard* = *range star-of*

Transfer tactic should remove occurrences of *star-of*.

setup $\langle \text{Transfer-Principle.add-const } @\{\text{const-name star-of}\} \rangle$

lemma *star-of-inject*: $\text{star-of } x = \text{star-of } y \longleftrightarrow x = y$

by *transfer* (*rule refl*)

lemma *Standard-star-of* [*simp*]: $\text{star-of } x \in \text{Standard}$

by (*simp add: Standard-def*)

2.5 Internal functions

Transfer tactic should remove occurrences of *Ifun*.

setup $\langle \text{Transfer-Principle.add-const } @\{\text{const-name Ifun}\} \rangle$

lemma *Ifun-star-of* [*simp*]: $\text{star-of } f \star \text{star-of } x = \text{star-of } (f x)$

by *transfer* (*rule refl*)

lemma *Standard-Ifun* [*simp*]: $f \in \text{Standard} \implies x \in \text{Standard} \implies f \star x \in \text{Standard}$

by (*auto simp add: Standard-def*)

Nonstandard extensions of functions.

definition *starfun* :: ('a \Rightarrow 'b) \Rightarrow 'a star \Rightarrow 'b star (**f* - [80] 80*)

where *starfun* $f \equiv \lambda x. \text{star-of } f \star x$

definition *starfun2* :: (*'a* \Rightarrow *'b* \Rightarrow *'c*) \Rightarrow *'a* *star* \Rightarrow *'b* *star* \Rightarrow *'c* *star* (**f2** - [80] 80)

where *starfun2* *f* \equiv $\lambda x y. \text{star-of } f \star x \star y$

declare *starfun-def* [*transfer-unfold*]

declare *starfun2-def* [*transfer-unfold*]

lemma *starfun-star-n*: (**f** *f*) (*star-n* *X*) = *star-n* ($\lambda n. f$ (*X* *n*))

by (*simp only: starfun-def star-of-def Ifun-star-n*)

lemma *starfun2-star-n*: (**f2** *f*) (*star-n* *X*) (*star-n* *Y*) = *star-n* ($\lambda n. f$ (*X* *n*) (*Y* *n*))

by (*simp only: starfun2-def star-of-def Ifun-star-n*)

lemma *starfun-star-of* [*simp*]: (**f** *f*) (*star-of* *x*) = *star-of* (*f* *x*)

by *transfer (rule refl)*

lemma *starfun2-star-of* [*simp*]: (**f2** *f*) (*star-of* *x*) = **f** *f* *x*

by *transfer (rule refl)*

lemma *Standard-starfun* [*simp*]: $x \in \text{Standard} \Longrightarrow \text{starfun } f \ x \in \text{Standard}$

by (*simp add: starfun-def*)

lemma *Standard-starfun2* [*simp*]: $x \in \text{Standard} \Longrightarrow y \in \text{Standard} \Longrightarrow \text{starfun2 } f \ x \ y \in \text{Standard}$

by (*simp add: starfun2-def*)

lemma *Standard-starfun-iff*:

assumes *inj*: $\bigwedge x y. f \ x = f \ y \Longrightarrow x = y$

shows $\text{starfun } f \ x \in \text{Standard} \longleftrightarrow x \in \text{Standard}$

proof

assume $x \in \text{Standard}$

then show $\text{starfun } f \ x \in \text{Standard}$ by *simp*

next

from *inj* have *inj'*: $\bigwedge x y. \text{starfun } f \ x = \text{starfun } f \ y \Longrightarrow x = y$

by *transfer*

assume $\text{starfun } f \ x \in \text{Standard}$

then obtain *b* where *b*: $\text{starfun } f \ x = \text{star-of } b$

unfolding *Standard-def* ..

then have $\exists x. \text{starfun } f \ x = \text{star-of } b$..

then have $\exists a. f \ a = b$ by *transfer*

then obtain *a* where $f \ a = b$..

then have $\text{starfun } f \ (\text{star-of } a) = \text{star-of } b$ by *transfer*

with *b* have $\text{starfun } f \ x = \text{starfun } f \ (\text{star-of } a)$ by *simp*

then have $x = \text{star-of } a$ by (*rule inj'*)

then show $x \in \text{Standard}$ by (*simp add: Standard-def*)

qed

lemma *Standard-starfun2-iff*:

assumes *inj*: $\bigwedge a b a' b'. f a b = f a' b' \implies a = a' \wedge b = b'$

shows $starfun2 f x y \in Standard \iff x \in Standard \wedge y \in Standard$

proof

assume $x \in Standard \wedge y \in Standard$

then show $starfun2 f x y \in Standard$ **by** *simp*

next

have *inj'*: $\bigwedge x y z w. starfun2 f x y = starfun2 f z w \implies x = z \wedge y = w$

using *inj* **by** *transfer*

assume $starfun2 f x y \in Standard$

then obtain *c* **where** $c: starfun2 f x y = star-of c$

unfolding *Standard-def* **..**

then have $\exists x y. starfun2 f x y = star-of c$ **by** *auto*

then have $\exists a b. f a b = c$ **by** *transfer*

then obtain *a b* **where** $f a b = c$ **by** *auto*

then have $starfun2 f (star-of a) (star-of b) = star-of c$ **by** *transfer*

with *c* **have** $starfun2 f x y = starfun2 f (star-of a) (star-of b)$ **by** *simp*

then have $x = star-of a \wedge y = star-of b$ **by** (rule *inj'*)

then show $x \in Standard \wedge y \in Standard$ **by** (*simp add: Standard-def*)

qed

2.6 Internal predicates

definition *unstar* :: $bool \Rightarrow bool$

where $unstar b \iff b = star-of True$

lemma *unstar-star-n*: $unstar (star-n P) \iff eventually P \mathcal{U}$

by (*simp add: unstar-def star-of-def star-n-eq-iff*)

lemma *unstar-star-of [simp]*: $unstar (star-of p) = p$

by (*simp add: unstar-def star-of-inject*)

Transfer tactic should remove occurrences of *unstar*.

setup $\langle Transfer-Principle.add-const @\{const-name\} unstar \rangle$

lemma *transfer-unstar [transfer-intro]*: $p \equiv star-n P \implies unstar p \equiv eventually P \mathcal{U}$

by (*simp only: unstar-star-n*)

definition *starP* :: $('a \Rightarrow bool) \Rightarrow 'a \Rightarrow bool$ (**p** - [80] 80)

where $*p* P = (\lambda x. unstar (star-of P \star x))$

definition *starP2* :: $('a \Rightarrow 'b \Rightarrow bool) \Rightarrow 'a \Rightarrow 'b \Rightarrow bool$ (**p2** - [80] 80)

where $*p2* P = (\lambda x y. unstar (star-of P \star x \star y))$

declare *starP-def* [*transfer-unfold*]

declare *starP2-def* [*transfer-unfold*]

lemma *starP-star-n*: $(*p* P) (\text{star-n } X) = \text{eventually } (\lambda n. P (X n)) \mathcal{U}$
by (*simp only*: *starP-def star-of-def Ifun-star-n unstar-star-n*)

lemma *starP2-star-n*: $(*p2* P) (\text{star-n } X) (\text{star-n } Y) = (\text{eventually } (\lambda n. P (X n) (Y n))) \mathcal{U}$
by (*simp only*: *starP2-def star-of-def Ifun-star-n unstar-star-n*)

lemma *starP-star-of [simp]*: $(*p* P) (\text{star-of } x) = P x$
by *transfer (rule refl)*

lemma *starP2-star-of [simp]*: $(*p2* P) (\text{star-of } x) = *p* P x$
by *transfer (rule refl)*

2.7 Internal sets

definition *Iset* :: 'a set star \Rightarrow 'a star set
where *Iset* $A = \{ x. (*p2* op \in) x A \}$

lemma *Iset-star-n*: $(\text{star-n } X \in \text{Iset } (\text{star-n } A)) = (\text{eventually } (\lambda n. X n \in A n)) \mathcal{U}$
by (*simp add*: *Iset-def starP2-star-n*)

Transfer tactic should remove occurrences of *Iset*.

setup $\langle \text{Transfer-Principle.add-const } @\{ \text{const-name } \text{Iset} \} \rangle$

lemma *transfer-mem [transfer-intro]*:
 $x \equiv \text{star-n } X \Longrightarrow a \equiv \text{Iset } (\text{star-n } A) \Longrightarrow x \in a \equiv \text{eventually } (\lambda n. X n \in A n) \mathcal{U}$
by (*simp only*: *Iset-star-n*)

lemma *transfer-Collect [transfer-intro]*:
 $(\bigwedge X. p (\text{star-n } X) \equiv \text{eventually } (\lambda n. P n (X n)) \mathcal{U}) \Longrightarrow$
 $\text{Collect } p \equiv \text{Iset } (\text{star-n } (\lambda n. \text{Collect } (P n)))$
by (*simp add*: *atomize-eq set-eq-iff all-star-eq Iset-star-n*)

lemma *transfer-set-eq [transfer-intro]*:
 $a \equiv \text{Iset } (\text{star-n } A) \Longrightarrow b \equiv \text{Iset } (\text{star-n } B) \Longrightarrow a = b \equiv \text{eventually } (\lambda n. A n = B n) \mathcal{U}$
by (*simp only*: *set-eq-iff transfer-all transfer-iff transfer-mem*)

lemma *transfer-ball [transfer-intro]*:
 $a \equiv \text{Iset } (\text{star-n } A) \Longrightarrow (\bigwedge X. p (\text{star-n } X) \equiv \text{eventually } (\lambda n. P n (X n)) \mathcal{U}) \Longrightarrow$
 $\forall x \in a. p x \equiv \text{eventually } (\lambda n. \forall x \in A n. P n x) \mathcal{U}$
by (*simp only*: *Ball-def transfer-all transfer-imp transfer-mem*)

lemma *transfer-bex [transfer-intro]*:
 $a \equiv \text{Iset } (\text{star-n } A) \Longrightarrow (\bigwedge X. p (\text{star-n } X) \equiv \text{eventually } (\lambda n. P n (X n)) \mathcal{U}) \Longrightarrow$
 $\exists x \in a. p x \equiv \text{eventually } (\lambda n. \exists x \in A n. P n x) \mathcal{U}$
by (*simp only*: *Bex-def transfer-ex transfer-conj transfer-mem*)

lemma *transfer-Iset* [*transfer-intro*]: $a \equiv \text{star-n } A \implies \text{Iset } a \equiv \text{Iset } (\text{star-n } (\lambda n. A \ n))$
by *simp*

Nonstandard extensions of sets.

definition *starset* :: 'a set \Rightarrow 'a star set (**s** - [80] 80)
where *starset* A = *Iset* (*star-of* A)

declare *starset-def* [*transfer-unfold*]

lemma *starset-mem*: *star-of* $x \in \text{*s* } A \longleftrightarrow x \in A$
by *transfer* (*rule refl*)

lemma *starset-UNIV*: $\text{*s* } (\text{UNIV}::'a \text{ set}) = (\text{UNIV}::'a \text{ star set})$
by (*transfer UNIV-def*) (*rule refl*)

lemma *starset-empty*: $\text{*s* } \{\} = \{\}$
by (*transfer empty-def*) (*rule refl*)

lemma *starset-insert*: $\text{*s* } (\text{insert } x \ A) = \text{insert } (\text{star-of } x) \ (\text{*s* } A)$
by (*transfer insert-def Un-def*) (*rule refl*)

lemma *starset-Un*: $\text{*s* } (A \cup B) = \text{*s* } A \cup \text{*s* } B$
by (*transfer Un-def*) (*rule refl*)

lemma *starset-Int*: $\text{*s* } (A \cap B) = \text{*s* } A \cap \text{*s* } B$
by (*transfer Int-def*) (*rule refl*)

lemma *starset-Compl*: $\text{*s* } \neg A = \neg (\text{*s* } A)$
by (*transfer Compl-eq*) (*rule refl*)

lemma *starset-diff*: $\text{*s* } (A - B) = \text{*s* } A - \text{*s* } B$
by (*transfer set-diff-eq*) (*rule refl*)

lemma *starset-image*: $\text{*s* } (f \ ' A) = (\text{*f* } f) \ ' (\text{*s* } A)$
by (*transfer image-def*) (*rule refl*)

lemma *starset-vimage*: $\text{*s* } (f \ -' A) = (\text{*f* } f) \ -' (\text{*s* } A)$
by (*transfer vimage-def*) (*rule refl*)

lemma *starset-subset*: $(\text{*s* } A \subseteq \text{*s* } B) \longleftrightarrow A \subseteq B$
by (*transfer subset-eq*) (*rule refl*)

lemma *starset-eq*: $(\text{*s* } A = \text{*s* } B) \longleftrightarrow A = B$
by *transfer* (*rule refl*)

lemmas *starset-simps* [*simp*] =
starset-mem *starset-UNIV*

```

starset-empty  starset-insert
starset-Un     starset-Int
starset-Compl  starset-diff
starset-image  starset-vimage
starset-subset starset-eq

```

2.8 Syntactic classes

```

instantiation star :: (zero) zero
begin
  definition star-zero-def: 0  $\equiv$  star-of 0
  instance ..
end

```

```

instantiation star :: (one) one
begin
  definition star-one-def: 1  $\equiv$  star-of 1
  instance ..
end

```

```

instantiation star :: (plus) plus
begin
  definition star-add-def: (op +)  $\equiv$  *f2* (op +)
  instance ..
end

```

```

instantiation star :: (times) times
begin
  definition star-mult-def: (op *)  $\equiv$  *f2* (op *)
  instance ..
end

```

```

instantiation star :: (uminus) uminus
begin
  definition star-minus-def: uminus  $\equiv$  *f* uminus
  instance ..
end

```

```

instantiation star :: (minus) minus
begin
  definition star-diff-def: (op -)  $\equiv$  *f2* (op -)
  instance ..
end

```

```

instantiation star :: (abs) abs
begin
  definition star-abs-def: abs  $\equiv$  *f* abs
  instance ..
end

```

```

instantiation star :: (sgn) sgn
begin
  definition star-sgn-def: sgn  $\equiv$  *f* sgn
  instance ..
end

instantiation star :: (divide) divide
begin
  definition star-divide-def: divide  $\equiv$  *f2* divide
  instance ..
end

instantiation star :: (inverse) inverse
begin
  definition star-inverse-def: inverse  $\equiv$  *f* inverse
  instance ..
end

instance star :: (Rings.dvd) Rings.dvd ..

instantiation star :: (modulo) modulo
begin
  definition star-mod-def: (op mod)  $\equiv$  *f2* (op mod)
  instance ..
end

instantiation star :: (ord) ord
begin
  definition star-le-def: (op  $\leq$ )  $\equiv$  *p2* (op  $\leq$ )
  definition star-less-def: (op  $<$ )  $\equiv$  *p2* (op  $<$ )
  instance ..
end

lemmas star-class-defs [transfer-unfold] =
  star-zero-def  star-one-def
  star-add-def   star-diff-def  star-minus-def
  star-mult-def  star-divide-def star-inverse-def
  star-le-def    star-less-def  star-abs-def    star-sgn-def
  star-mod-def

Class operations preserve standard elements.

lemma Standard-zero: 0  $\in$  Standard
  by (simp add: star-zero-def)

lemma Standard-one: 1  $\in$  Standard
  by (simp add: star-one-def)

lemma Standard-add:  $x \in$  Standard  $\implies y \in$  Standard  $\implies x + y \in$  Standard

```

by (simp add: star-add-def)

lemma *Standard-diff*: $x \in \text{Standard} \implies y \in \text{Standard} \implies x - y \in \text{Standard}$
by (simp add: star-diff-def)

lemma *Standard-minus*: $x \in \text{Standard} \implies -x \in \text{Standard}$
by (simp add: star-minus-def)

lemma *Standard-mult*: $x \in \text{Standard} \implies y \in \text{Standard} \implies x * y \in \text{Standard}$
by (simp add: star-mult-def)

lemma *Standard-divide*: $x \in \text{Standard} \implies y \in \text{Standard} \implies x / y \in \text{Standard}$
by (simp add: star-divide-def)

lemma *Standard-inverse*: $x \in \text{Standard} \implies \text{inverse } x \in \text{Standard}$
by (simp add: star-inverse-def)

lemma *Standard-abs*: $x \in \text{Standard} \implies |x| \in \text{Standard}$
by (simp add: star-abs-def)

lemma *Standard-mod*: $x \in \text{Standard} \implies y \in \text{Standard} \implies x \bmod y \in \text{Standard}$
by (simp add: star-mod-def)

lemmas *Standard-simps* [simp] =
 Standard-zero *Standard-one*
 Standard-add *Standard-diff* *Standard-minus*
 Standard-mult *Standard-divide* *Standard-inverse*
 Standard-abs *Standard-mod*

star-of preserves class operations.

lemma *star-of-add*: $\text{star-of } (x + y) = \text{star-of } x + \text{star-of } y$
by transfer (rule refl)

lemma *star-of-diff*: $\text{star-of } (x - y) = \text{star-of } x - \text{star-of } y$
by transfer (rule refl)

lemma *star-of-minus*: $\text{star-of } (-x) = - \text{star-of } x$
by transfer (rule refl)

lemma *star-of-mult*: $\text{star-of } (x * y) = \text{star-of } x * \text{star-of } y$
by transfer (rule refl)

lemma *star-of-divide*: $\text{star-of } (x / y) = \text{star-of } x / \text{star-of } y$
by transfer (rule refl)

lemma *star-of-inverse*: $\text{star-of } (\text{inverse } x) = \text{inverse } (\text{star-of } x)$
by transfer (rule refl)

lemma *star-of-mod*: $\text{star-of } (x \bmod y) = \text{star-of } x \bmod \text{star-of } y$

by *transfer* (rule *refl*)

lemma *star-of-abs*: $star-of\ |x| = |star-of\ x|$
by *transfer* (rule *refl*)

star-of preserves numerals.

lemma *star-of-zero*: $star-of\ 0 = 0$
by *transfer* (rule *refl*)

lemma *star-of-one*: $star-of\ 1 = 1$
by *transfer* (rule *refl*)

star-of preserves orderings.

lemma *star-of-less*: $(star-of\ x < star-of\ y) = (x < y)$
by *transfer* (rule *refl*)

lemma *star-of-le*: $(star-of\ x \leq star-of\ y) = (x \leq y)$
by *transfer* (rule *refl*)

lemma *star-of-eq*: $(star-of\ x = star-of\ y) = (x = y)$
by *transfer* (rule *refl*)

As above, for 0.

lemmas *star-of-0-less* = *star-of-less* [*of 0, simplified star-of-zero*]
lemmas *star-of-0-le* = *star-of-le* [*of 0, simplified star-of-zero*]
lemmas *star-of-0-eq* = *star-of-eq* [*of 0, simplified star-of-zero*]

lemmas *star-of-less-0* = *star-of-less* [*of - 0, simplified star-of-zero*]
lemmas *star-of-le-0* = *star-of-le* [*of - 0, simplified star-of-zero*]
lemmas *star-of-eq-0* = *star-of-eq* [*of - 0, simplified star-of-zero*]

As above, for 1.

lemmas *star-of-1-less* = *star-of-less* [*of 1, simplified star-of-one*]
lemmas *star-of-1-le* = *star-of-le* [*of 1, simplified star-of-one*]
lemmas *star-of-1-eq* = *star-of-eq* [*of 1, simplified star-of-one*]

lemmas *star-of-less-1* = *star-of-less* [*of - 1, simplified star-of-one*]
lemmas *star-of-le-1* = *star-of-le* [*of - 1, simplified star-of-one*]
lemmas *star-of-eq-1* = *star-of-eq* [*of - 1, simplified star-of-one*]

lemmas *star-of-simps* [*simp*] =
star-of-add *star-of-diff* *star-of-minus*
star-of-mult *star-of-divide* *star-of-inverse*
star-of-mod *star-of-abs*
star-of-zero *star-of-one*
star-of-less *star-of-le* *star-of-eq*
star-of-0-less *star-of-0-le* *star-of-0-eq*
star-of-less-0 *star-of-le-0* *star-of-eq-0*

star-of-1-less star-of-1-le star-of-1-eq
star-of-less-1 star-of-le-1 star-of-eq-1

2.9 Ordering and lattice classes

```
instance star :: (order) order
apply intro-classes
  apply (transfer, rule less-le-not-le)
  apply (transfer, rule order-refl)
  apply (transfer, erule (1) order-trans)
  apply (transfer, erule (1) order-antisym)
done
```

```
instantiation star :: (semilattice-inf) semilattice-inf
begin
  definition star-inf-def [transfer-unfold]: inf  $\equiv$  *f2* inf
  instance by (standard; transfer) auto
end
```

```
instantiation star :: (semilattice-sup) semilattice-sup
begin
  definition star-sup-def [transfer-unfold]: sup  $\equiv$  *f2* sup
  instance by (standard; transfer) auto
end
```

```
instance star :: (lattice) lattice ..
```

```
instance star :: (distrib-lattice) distrib-lattice
  by (standard; transfer) (auto simp add: sup-inf-distrib1)
```

```
lemma Standard-inf [simp]:  $x \in \text{Standard} \implies y \in \text{Standard} \implies \text{inf } x \ y \in \text{Standard}$ 
  by (simp add: star-inf-def)
```

```
lemma Standard-sup [simp]:  $x \in \text{Standard} \implies y \in \text{Standard} \implies \text{sup } x \ y \in \text{Standard}$ 
  by (simp add: star-sup-def)
```

```
lemma star-of-inf [simp]:  $\text{star-of } (\text{inf } x \ y) = \text{inf } (\text{star-of } x) (\text{star-of } y)$ 
  by transfer (rule refl)
```

```
lemma star-of-sup [simp]:  $\text{star-of } (\text{sup } x \ y) = \text{sup } (\text{star-of } x) (\text{star-of } y)$ 
  by transfer (rule refl)
```

```
instance star :: (linorder) linorder
  by (intro-classes, transfer, rule linorder-linear)
```

```
lemma star-max-def [transfer-unfold]:  $\text{max} = \text{*f2* max}$ 
  apply (rule ext, rule ext)
```

```

apply (unfold max-def, transfer, fold max-def)
apply (rule refl)
done

```

```

lemma star-min-def [transfer-unfold]: min = *f2* min
apply (rule ext, rule ext)
apply (unfold min-def, transfer, fold min-def)
apply (rule refl)
done

```

```

lemma Standard-max [simp]:  $x \in \text{Standard} \implies y \in \text{Standard} \implies \max x y \in \text{Standard}$ 
by (simp add: star-max-def)

```

```

lemma Standard-min [simp]:  $x \in \text{Standard} \implies y \in \text{Standard} \implies \min x y \in \text{Standard}$ 
by (simp add: star-min-def)

```

```

lemma star-of-max [simp]:  $\text{star-of} (\max x y) = \max (\text{star-of } x) (\text{star-of } y)$ 
by transfer (rule refl)

```

```

lemma star-of-min [simp]:  $\text{star-of} (\min x y) = \min (\text{star-of } x) (\text{star-of } y)$ 
by transfer (rule refl)

```

2.10 Ordered group classes

```

instance star :: (semigroup-add) semigroup-add
by (intro-classes, transfer, rule add.assoc)

```

```

instance star :: (ab-semigroup-add) ab-semigroup-add
by (intro-classes, transfer, rule add commute)

```

```

instance star :: (semigroup-mult) semigroup-mult
by (intro-classes, transfer, rule mult.assoc)

```

```

instance star :: (ab-semigroup-mult) ab-semigroup-mult
by (intro-classes, transfer, rule mult commute)

```

```

instance star :: (comm-monoid-add) comm-monoid-add
by (intro-classes, transfer, rule comm-monoid-add-class.add-0)

```

```

instance star :: (monoid-mult) monoid-mult
apply intro-classes
apply (transfer, rule mult-1-left)
apply (transfer, rule mult-1-right)
done

```

```

instance star :: (power) power ..

```



```

instance star :: (comm-monoid-mult) comm-monoid-mult
  by (intro-classes, transfer, rule mult-1)

instance star :: (cancel-semigroup-add) cancel-semigroup-add
  apply intro-classes
  apply (transfer, erule add-left-imp-eq)
  apply (transfer, erule add-right-imp-eq)
  done

instance star :: (cancel-ab-semigroup-add) cancel-ab-semigroup-add
  by intro-classes (transfer, simp add: diff-diff-eq)+

instance star :: (cancel-comm-monoid-add) cancel-comm-monoid-add ..

instance star :: (ab-group-add) ab-group-add
  apply intro-classes
  apply (transfer, rule left-minus)
  apply (transfer, rule diff-conv-add-uminus)
  done

instance star :: (ordered-ab-semigroup-add) ordered-ab-semigroup-add
  by (intro-classes, transfer, rule add-left-mono)

instance star :: (ordered-cancel-ab-semigroup-add) ordered-cancel-ab-semigroup-add
  ..

instance star :: (ordered-ab-semigroup-add-imp-le) ordered-ab-semigroup-add-imp-le
  by (intro-classes, transfer, rule add-le-imp-le-left)

instance star :: (ordered-comm-monoid-add) ordered-comm-monoid-add ..
instance star :: (ordered-ab-semigroup-monoid-add-imp-le) ordered-ab-semigroup-monoid-add-imp-le
  ..
instance star :: (ordered-cancel-comm-monoid-add) ordered-cancel-comm-monoid-add
  ..
instance star :: (ordered-ab-group-add) ordered-ab-group-add ..

instance star :: (ordered-ab-group-add-abs) ordered-ab-group-add-abs
  by intro-classes (transfer, simp add: abs-ge-self abs-leI abs-triangle-ineq)+

instance star :: (linordered-cancel-ab-semigroup-add) linordered-cancel-ab-semigroup-add
  ..

```

2.11 Ring and field classes

```

instance star :: (semiring) semiring
  by (intro-classes; transfer) (fact distrib-right distrib-left)+

instance star :: (semiring-0) semiring-0
  by (intro-classes; transfer) simp-all

```

```

instance star :: (semiring-0-cancel) semiring-0-cancel ..

instance star :: (comm-semiring) comm-semiring
  by (intro-classes; transfer) (fact distrib-right)

instance star :: (comm-semiring-0) comm-semiring-0 ..
instance star :: (comm-semiring-0-cancel) comm-semiring-0-cancel ..

instance star :: (zero-neq-one) zero-neq-one
  by (intro-classes; transfer) (fact zero-neq-one)

instance star :: (semiring-1) semiring-1 ..
instance star :: (comm-semiring-1) comm-semiring-1 ..

declare dvd-def [transfer-refold]

instance star :: (comm-semiring-1-cancel) comm-semiring-1-cancel
  by (intro-classes; transfer) (fact right-diff-distrib^)

instance star :: (semiring-no-zero-divisors) semiring-no-zero-divisors
  by (intro-classes; transfer) (fact no-zero-divisors)

instance star :: (semiring-1-no-zero-divisors) semiring-1-no-zero-divisors ..

instance star :: (semiring-no-zero-divisors-cancel) semiring-no-zero-divisors-cancel
  by (intro-classes; transfer) simp-all

instance star :: (semiring-1-cancel) semiring-1-cancel ..
instance star :: (ring) ring ..
instance star :: (comm-ring) comm-ring ..
instance star :: (ring-1) ring-1 ..
instance star :: (comm-ring-1) comm-ring-1 ..
instance star :: (semidom) semidom ..

instance star :: (semidom-divide) semidom-divide
  by (intro-classes; transfer) simp-all

instance star :: (semiring-div) semiring-div
  by (intro-classes; transfer) (simp-all add: div-mult-mod-eq)

instance star :: (ring-no-zero-divisors) ring-no-zero-divisors ..
instance star :: (ring-1-no-zero-divisors) ring-1-no-zero-divisors ..
instance star :: (idom) idom ..
instance star :: (idom-divide) idom-divide ..

instance star :: (division-ring) division-ring
  by (intro-classes; transfer) (simp-all add: divide-inverse)

```

```

instance star :: (field) field
  by (intro-classes; transfer) (simp-all add: divide-inverse)

instance star :: (ordered-semiring) ordered-semiring
  by (intro-classes; transfer) (fact mult-left-mono mult-right-mono)+

instance star :: (ordered-cancel-semiring) ordered-cancel-semiring ..

instance star :: (linordered-semiring-strict) linordered-semiring-strict
  by (intro-classes; transfer) (fact mult-strict-left-mono mult-strict-right-mono)+

instance star :: (ordered-comm-semiring) ordered-comm-semiring
  by (intro-classes; transfer) (fact mult-left-mono)

instance star :: (ordered-cancel-comm-semiring) ordered-cancel-comm-semiring ..

instance star :: (linordered-comm-semiring-strict) linordered-comm-semiring-strict
  by (intro-classes; transfer) (fact mult-strict-left-mono)

instance star :: (ordered-ring) ordered-ring ..

instance star :: (ordered-ring-abs) ordered-ring-abs
  by (intro-classes; transfer) (fact abs-eq-mult)

instance star :: (abs-if) abs-if
  by (intro-classes; transfer) (fact abs-if)

instance star :: (linordered-ring-strict) linordered-ring-strict ..
instance star :: (ordered-comm-ring) ordered-comm-ring ..

instance star :: (linordered-semidom) linordered-semidom
  by (intro-classes; transfer) (fact zero-less-one le-add-diff-inverse2)+

instance star :: (linordered-idom) linordered-idom
  by (intro-classes; transfer) (fact sgn-if)

instance star :: (linordered-field) linordered-field ..

```

2.12 Power

lemma star-power-def [transfer-unfold]: $(op \wedge) \equiv \lambda x n. (*f* (\lambda x. x \wedge n)) x$

proof (rule eq-reflection, rule ext, rule ext)

show $x \wedge n = (*f* (\lambda x. x \wedge n)) x$ **for** $n :: nat$ **and** $x :: 'a$ star

proof (induct n arbitrary: x)

case 0

have $\bigwedge x :: 'a$ star. $(*f* (\lambda x. 1)) x = 1$

by transfer simp

then show ?case by simp

next

```

case (Suc n)
have  $\bigwedge x::'a \text{ star. } x * (*f* (\lambda x::'a. x \wedge n)) x = (*f* (\lambda x::'a. x * x \wedge n)) x$ 
  by transfer simp
with Suc show ?case by simp
qed
qed

```

lemma *Standard-power* [simp]: $x \in \text{Standard} \implies x \wedge n \in \text{Standard}$
by (simp add: star-power-def)

lemma *star-of-power* [simp]: $\text{star-of } (x \wedge n) = \text{star-of } x \wedge n$
by transfer (rule refl)

2.13 Number classes

instance *star* :: (numeral) numeral ..

lemma *star-numeral-def* [transfer-unfold]: $\text{numeral } k = \text{star-of } (\text{numeral } k)$
by (induct k) (simp-all only: numeral.simps star-of-one star-of-add)

lemma *Standard-numeral* [simp]: $\text{numeral } k \in \text{Standard}$
by (simp add: star-numeral-def)

lemma *star-of-numeral* [simp]: $\text{star-of } (\text{numeral } k) = \text{numeral } k$
by transfer (rule refl)

lemma *star-of-nat-def* [transfer-unfold]: $\text{of-nat } n = \text{star-of } (\text{of-nat } n)$
by (induct n) simp-all

lemmas *star-of-compare-numeral* [simp] =
star-of-less [of numeral k, simplified star-of-numeral]
star-of-le [of numeral k, simplified star-of-numeral]
star-of-eq [of numeral k, simplified star-of-numeral]
star-of-less [of - numeral k, simplified star-of-numeral]
star-of-le [of - numeral k, simplified star-of-numeral]
star-of-eq [of - numeral k, simplified star-of-numeral]
star-of-less [of - numeral k, simplified star-of-numeral]
star-of-le [of - numeral k, simplified star-of-numeral]
star-of-eq [of - numeral k, simplified star-of-numeral]
star-of-less [of - numeral k, simplified star-of-numeral]
star-of-le [of - numeral k, simplified star-of-numeral]
star-of-eq [of - numeral k, simplified star-of-numeral] **for** k

lemma *Standard-of-nat* [simp]: $\text{of-nat } n \in \text{Standard}$
by (simp add: star-of-nat-def)

lemma *star-of-of-nat* [simp]: $\text{star-of } (\text{of-nat } n) = \text{of-nat } n$
by transfer (rule refl)

lemma *star-of-int-def* [*transfer-unfold*]: $of\text{-}int\ z = star\text{-}of\ (of\text{-}int\ z)$
by (*rule int-diff-cases* [*of z*]) *simp*

lemma *Standard-of-int* [*simp*]: $of\text{-}int\ z \in Standard$
by (*simp add: star-of-int-def*)

lemma *star-of-of-int* [*simp*]: $star\text{-}of\ (of\text{-}int\ z) = of\text{-}int\ z$
by *transfer* (*rule refl*)

instance *star* :: (*semiring-char-0*) *semiring-char-0*

proof

have *inj* (*star-of* :: $'a \Rightarrow 'a\ star$)
by (*rule injI*) *simp*
then have *inj* (*star-of* \circ *of-nat* :: $nat \Rightarrow 'a\ star$)
using *inj-of-nat* **by** (*rule inj-comp*)
then show *inj* (*of-nat* :: $nat \Rightarrow 'a\ star$)
by (*simp add: comp-def*)

qed

instance *star* :: (*ring-char-0*) *ring-char-0* ..

instance *star* :: (*semiring-parity*) *semiring-parity*

apply *intro-classes*
apply (*transfer*, *rule odd-one*)
apply (*transfer*, *erule* (1) *odd-even-add*)
apply (*transfer*, *erule even-multD*)
apply (*transfer*, *erule odd-ex-decrement*)
done

instance *star* :: (*semiring-div-parity*) *semiring-div-parity*

apply *intro-classes*
apply (*transfer*, *rule parity*)
apply (*transfer*, *rule one-mod-two-eq-one*)
apply (*transfer*, *rule zero-not-eq-two*)
done

instantiation *star* :: (*semiring-numeral-div*) *semiring-numeral-div*

begin

definition *divmod-star* :: $num \Rightarrow num \Rightarrow 'a\ star \times 'a\ star$

where *divmod-star-def*: $divmod\text{-}star\ m\ n = (numeral\ m\ div\ numeral\ n, numeral\ m\ mod\ numeral\ n)$

definition *divmod-step-star* :: $num \Rightarrow 'a\ star \times 'a\ star \Rightarrow 'a\ star \times 'a\ star$

where *divmod-step-star* *l qr* =
 $(let\ (q, r) = qr$
in if $r \geq numeral\ l$ **then** $(2 * q + 1, r - numeral\ l)$ **else** $(2 * q, r)$)

instance

```

proof
  show  $\text{divmod } m \ n = (\text{numeral } m \ \text{div} \ \text{numeral } n :: 'a \ \text{star}, \text{numeral } m \ \text{mod} \ \text{numeral } n)$  for  $m \ n$ 
    by (fact divmod-star-def)
  show  $\text{divmod-step } l \ qr =$ 
    (let  $(q, r) = qr$ 
     in if  $r \geq \text{numeral } l$  then  $(2 * q + 1, r - \text{numeral } l)$  else  $(2 * q, r)$ )
    for  $l$  and  $qr :: 'a \ \text{star} \times 'a \ \text{star}$ 
    by (fact divmod-step-star-def)
qed (transfer,
  fact
  semiring-numeral-div-class.div-less
  semiring-numeral-div-class.mod-less
  semiring-numeral-div-class.div-positive
  semiring-numeral-div-class.mod-less-eq-dividend
  semiring-numeral-div-class.pos-mod-bound
  semiring-numeral-div-class.pos-mod-sign
  semiring-numeral-div-class.mod-mult2-eq
  semiring-numeral-div-class.div-mult2-eq
  semiring-numeral-div-class.discrete)+

end

declare divmod-algorithm-code [where  $?'a = 'a :: \text{semiring-numeral-div star}$ , code]

```

2.14 Finite class

```

lemma starset-finite:  $\text{finite } A \implies *s* \ A = \text{star-of } 'A$ 
  by (erule finite-induct) simp-all

instance star :: (finite) finite
  apply intro-classes
  apply (subst starset-UNIV [symmetric])
  apply (subst starset-finite [OF finite])
  apply (rule finite-imageI [OF finite])
  done

```

end

3 Hypernatural numbers

```

theory HyperNat
  imports StarDef
begin

  type-synonym hypnat = nat star

  abbreviation hypnat-of-nat ::  $\text{nat} \Rightarrow \text{nat star}$ 
    where hypnat-of-nat  $\equiv \text{star-of}$ 

```

definition $hSuc :: hypnat \Rightarrow hypnat$
where $hSuc\text{-def}$ [*transfer-unfold*]: $hSuc = *f* Suc$

3.1 Properties Transferred from Naturals

lemma $hSuc\text{-not-zero}$ [*iff*]: $\bigwedge m. hSuc\ m \neq 0$
by *transfer* (*rule Suc-not-Zero*)

lemma $zero\text{-not-}hSuc$ [*iff*]: $\bigwedge m. 0 \neq hSuc\ m$
by *transfer* (*rule Zero-not-Suc*)

lemma $hSuc\text{-}hSuc\text{-eq}$ [*iff*]: $\bigwedge m\ n. hSuc\ m = hSuc\ n \longleftrightarrow m = n$
by *transfer* (*rule nat.inject*)

lemma $zero\text{-less-}hSuc$ [*iff*]: $\bigwedge n. 0 < hSuc\ n$
by *transfer* (*rule zero-less-Suc*)

lemma $hypnat\text{-minus-zero}$ [*simp*]: $\bigwedge z::hypnat. z - z = 0$
by *transfer* (*rule diff-self-eq-0*)

lemma $hypnat\text{-diff-0-eq-0}$ [*simp*]: $\bigwedge n::hypnat. 0 - n = 0$
by *transfer* (*rule diff-0-eq-0*)

lemma $hypnat\text{-add-is-0}$ [*iff*]: $\bigwedge m\ n::hypnat. m + n = 0 \longleftrightarrow m = 0 \wedge n = 0$
by *transfer* (*rule add-is-0*)

lemma $hypnat\text{-diff-diff-left}$: $\bigwedge i\ j\ k::hypnat. i - j - k = i - (j + k)$
by *transfer* (*rule diff-diff-left*)

lemma $hypnat\text{-diff-commute}$: $\bigwedge i\ j\ k::hypnat. i - j - k = i - k - j$
by *transfer* (*rule diff-commute*)

lemma $hypnat\text{-diff-add-inverse}$ [*simp*]: $\bigwedge m\ n::hypnat. n + m - n = m$
by *transfer* (*rule diff-add-inverse*)

lemma $hypnat\text{-diff-add-inverse2}$ [*simp*]: $\bigwedge m\ n::hypnat. m + n - n = m$
by *transfer* (*rule diff-add-inverse2*)

lemma $hypnat\text{-diff-cancel}$ [*simp*]: $\bigwedge k\ m\ n::hypnat. (k + m) - (k + n) = m - n$
by *transfer* (*rule diff-cancel*)

lemma $hypnat\text{-diff-cancel2}$ [*simp*]: $\bigwedge k\ m\ n::hypnat. (m + k) - (n + k) = m - n$
by *transfer* (*rule diff-cancel2*)

lemma $hypnat\text{-diff-add-0}$ [*simp*]: $\bigwedge m\ n::hypnat. n - (n + m) = 0$
by *transfer* (*rule diff-add-0*)

lemma $hypnat\text{-diff-mult-distrib}$: $\bigwedge k\ m\ n::hypnat. (m - n) * k = (m * k) - (n * k)$

k)

by *transfer* (rule *diff-mult-distrib*)

lemma *hypnat-diff-mult-distrib2*: $\bigwedge k m n :: \text{hypnat}. k * (m - n) = (k * m) - (k * n)$

by *transfer* (rule *diff-mult-distrib2*)

lemma *hypnat-le-zero-cancel* [*iff*]: $\bigwedge n :: \text{hypnat}. n \leq 0 \longleftrightarrow n = 0$

by *transfer* (rule *le-0-eq*)

lemma *hypnat-mult-is-0* [*simp*]: $\bigwedge m n :: \text{hypnat}. m * n = 0 \longleftrightarrow m = 0 \vee n = 0$

by *transfer* (rule *mult-is-0*)

lemma *hypnat-diff-is-0-eq* [*simp*]: $\bigwedge m n :: \text{hypnat}. m - n = 0 \longleftrightarrow m \leq n$

by *transfer* (rule *diff-is-0-eq*)

lemma *hypnat-not-less0* [*iff*]: $\bigwedge n :: \text{hypnat}. \neg n < 0$

by *transfer* (rule *not-less0*)

lemma *hypnat-less-one* [*iff*]: $\bigwedge n :: \text{hypnat}. n < 1 \longleftrightarrow n = 0$

by *transfer* (rule *less-one*)

lemma *hypnat-add-diff-inverse*: $\bigwedge m n :: \text{hypnat}. \neg m < n \implies n + (m - n) = m$

by *transfer* (rule *add-diff-inverse*)

lemma *hypnat-le-add-diff-inverse* [*simp*]: $\bigwedge m n :: \text{hypnat}. n \leq m \implies n + (m - n) = m$

by *transfer* (rule *le-add-diff-inverse*)

lemma *hypnat-le-add-diff-inverse2* [*simp*]: $\bigwedge m n :: \text{hypnat}. n \leq m \implies (m - n) + n = m$

by *transfer* (rule *le-add-diff-inverse2*)

declare *hypnat-le-add-diff-inverse2* [*OF order-less-imp-le*]

lemma *hypnat-le0* [*iff*]: $\bigwedge n :: \text{hypnat}. 0 \leq n$

by *transfer* (rule *le0*)

lemma *hypnat-le-add1* [*simp*]: $\bigwedge x n :: \text{hypnat}. x \leq x + n$

by *transfer* (rule *le-add1*)

lemma *hypnat-add-self-le* [*simp*]: $\bigwedge x n :: \text{hypnat}. x \leq n + x$

by *transfer* (rule *le-add2*)

lemma *hypnat-add-one-self-less* [*simp*]: $x < x + 1$ **for** $x :: \text{hypnat}$

by (*fact less-add-one*)

lemma *hypnat-neq0-conv* [*iff*]: $\bigwedge n :: \text{hypnat}. n \neq 0 \longleftrightarrow 0 < n$

by *transfer* (rule *neq0-conv*)

lemma *hypnat-gt-zero-iff*: $0 < n \longleftrightarrow 1 \leq n$ **for** $n :: \text{hypnat}$
by (*auto simp add: linorder-not-less [symmetric]*)

lemma *hypnat-gt-zero-iff2*: $0 < n \longleftrightarrow (\exists m. n = m + 1)$ **for** $n :: \text{hypnat}$
by (*auto intro!: add-nonneg-pos exI[of - n - 1] simp: hypnat-gt-zero-iff*)

lemma *hypnat-add-self-not-less*: $\neg x + y < x$ **for** $x y :: \text{hypnat}$
by (*simp add: linorder-not-le [symmetric] add.commute [of x]*)

lemma *hypnat-diff-split*: $P (a - b) \longleftrightarrow (a < b \longrightarrow P 0) \wedge (\forall d. a = b + d \longrightarrow P d)$
for $a b :: \text{hypnat}$
— elimination of $-$ on *hypnat*
proof (*cases a < b rule: case-split*)
case *True*
then show *?thesis*
by (*auto simp add: hypnat-add-self-not-less order-less-imp-le hypnat-diff-is-0-eq [THEN iffD2]*)
next
case *False*
then show *?thesis*
by (*auto simp add: linorder-not-less dest: order-le-less-trans*)
qed

3.2 Properties of the set of embedded natural numbers

lemma *of-nat-eq-star-of* [*simp*]: *of-nat* = *star-of*
proof
show *of-nat n = star-of n* **for** n
by *transfer simp*
qed

lemma *Nats-eq-Standard*: (*Nats* :: *nat star set*) = *Standard*
by (*auto simp: Nats-def Standard-def*)

lemma *hypnat-of-nat-mem-Nats* [*simp*]: *hypnat-of-nat* $n \in \text{Nats}$
by (*simp add: Nats-eq-Standard*)

lemma *hypnat-of-nat-one* [*simp*]: *hypnat-of-nat* (*Suc* 0) = 1
by *transfer simp*

lemma *hypnat-of-nat-Suc* [*simp*]: *hypnat-of-nat* (*Suc* n) = *hypnat-of-nat* $n + 1$
by *transfer simp*

lemma *of-nat-eq-add* [*rule-format*]: $\forall d :: \text{hypnat}. \text{of-nat } m = \text{of-nat } n + d \longrightarrow d \in \text{range of-nat}$
apply (*induct n*)
apply (*auto simp add: add.assoc*)

```

apply (case-tac x)
apply (auto simp add: add.commute [of 1])
done

```

lemma *Nats-diff* [simp]: $a \in \text{Nats} \implies b \in \text{Nats} \implies a - b \in \text{Nats}$ **for** $a\ b :: \text{hypnat}$
by (simp add: Nats-eq-Standard)

3.3 Infinite Hypernatural Numbers – *HNatInfinite*

The set of infinite hypernatural numbers.

definition *HNatInfinite* :: *hypnat set*
where *HNatInfinite* = $\{n. n \notin \text{Nats}\}$

lemma *Nats-not-HNatInfinite-iff*: $x \in \text{Nats} \longleftrightarrow x \notin \text{HNatInfinite}$
by (simp add: *HNatInfinite-def*)

lemma *HNatInfinite-not-Nats-iff*: $x \in \text{HNatInfinite} \longleftrightarrow x \notin \text{Nats}$
by (simp add: *HNatInfinite-def*)

lemma *star-of-neq-HNatInfinite*: $N \in \text{HNatInfinite} \implies \text{star-of } n \neq N$
by (auto simp add: *HNatInfinite-def* *Nats-eq-Standard*)

lemma *star-of-Suc-lessI*: $\bigwedge N. \text{star-of } n < N \implies \text{star-of } (\text{Suc } n) \neq N \implies \text{star-of } (\text{Suc } n) < N$
by transfer (rule *Suc-lessI*)

lemma *star-of-less-HNatInfinite*:
assumes $N: N \in \text{HNatInfinite}$
shows $\text{star-of } n < N$

proof (induct n)

case 0

from N **have** $\text{star-of } 0 \neq N$

by (rule *star-of-neq-HNatInfinite*)

then show ?case **by** simp

next

case ($\text{Suc } n$)

from N **have** $\text{star-of } (\text{Suc } n) \neq N$

by (rule *star-of-neq-HNatInfinite*)

with Suc **show** ?case

by (rule *star-of-Suc-lessI*)

qed

lemma *star-of-le-HNatInfinite*: $N \in \text{HNatInfinite} \implies \text{star-of } n \leq N$
by (rule *star-of-less-HNatInfinite* [THEN *order-less-imp-le*])

3.3.1 Closure Rules

lemma *Nats-less-HNatInfinite*: $x \in \text{Nats} \implies y \in \text{HNatInfinite} \implies x < y$

by (auto simp add: Nats-def star-of-less-HNatInfinite)

lemma *Nats-le-HNatInfinite*: $x \in \text{Nats} \implies y \in \text{HNatInfinite} \implies x \leq y$
 by (rule *Nats-less-HNatInfinite* [THEN *order-less-imp-le*])

lemma *zero-less-HNatInfinite*: $x \in \text{HNatInfinite} \implies 0 < x$
 by (simp add: *Nats-less-HNatInfinite*)

lemma *one-less-HNatInfinite*: $x \in \text{HNatInfinite} \implies 1 < x$
 by (simp add: *Nats-less-HNatInfinite*)

lemma *one-le-HNatInfinite*: $x \in \text{HNatInfinite} \implies 1 \leq x$
 by (simp add: *Nats-le-HNatInfinite*)

lemma *zero-not-mem-HNatInfinite* [simp]: $0 \notin \text{HNatInfinite}$
 by (simp add: *HNatInfinite-def*)

lemma *Nats-downward-closed*: $x \in \text{Nats} \implies y \leq x \implies y \in \text{Nats}$ **for** $x\ y :: \text{hypnat}$
 apply (simp only: *linorder-not-less* [symmetric])
 apply (erule *contrapos-np*)
 apply (drule *HNatInfinite-not-Nats-iff* [THEN *iffD2*])
 apply (erule (1) *Nats-less-HNatInfinite*)
 done

lemma *HNatInfinite-upward-closed*: $x \in \text{HNatInfinite} \implies x \leq y \implies y \in \text{HNatInfinite}$
 apply (simp only: *HNatInfinite-not-Nats-iff*)
 apply (erule *contrapos-nn*)
 apply (erule (1) *Nats-downward-closed*)
 done

lemma *HNatInfinite-add*: $x \in \text{HNatInfinite} \implies x + y \in \text{HNatInfinite}$
 apply (erule *HNatInfinite-upward-closed*)
 apply (rule *hypnat-le-add1*)
 done

lemma *HNatInfinite-add-one*: $x \in \text{HNatInfinite} \implies x + 1 \in \text{HNatInfinite}$
 by (rule *HNatInfinite-add*)

lemma *HNatInfinite-diff*: $x \in \text{HNatInfinite} \implies y \in \text{Nats} \implies x - y \in \text{HNatInfinite}$
 apply (frule (1) *Nats-le-HNatInfinite*)
 apply (simp only: *HNatInfinite-not-Nats-iff*)
 apply (erule *contrapos-nn*)
 apply (drule (1) *Nats-add*, simp)
 done

lemma *HNatInfinite-is-Suc*: $x \in \text{HNatInfinite} \implies \exists y. x = y + 1$ **for** $x :: \text{hypnat}$
 apply (rule-tac $x = x - (1::\text{hypnat})$ **in** *exI*)
 apply (simp add: *Nats-le-HNatInfinite*)
 done

3.4 Existence of an infinite hypernatural number

ω is in fact an infinite hypernatural number = [$\langle 1, 2, 3, \dots \rangle$]

definition $whn :: hypnat$

where $hypnat\text{-}\omega\text{-def}: whn = star\text{-}n (\lambda n::nat. n)$

lemma $hypnat\text{-of-nat-neq-whn}: hypnat\text{-of-nat } n \neq whn$

by ($simp$ add: $FreeUltrafilterNat.singleton' hypnat\text{-}\omega\text{-def } star\text{-of-def } star\text{-}n\text{-eq-iff}$)

lemma $whn\text{-neq-hypnat-of-nat}: whn \neq hypnat\text{-of-nat } n$

by ($simp$ add: $FreeUltrafilterNat.singleton hypnat\text{-}\omega\text{-def } star\text{-of-def } star\text{-}n\text{-eq-iff}$)

lemma $whn\text{-not-Nats}$ [$simp$]: $whn \notin Nats$

by ($simp$ add: $Nats\text{-def } image\text{-def } whn\text{-neq-hypnat-of-nat}$)

lemma $HNatInfinite-whn$ [$simp$]: $whn \in HNatInfinite$

by ($simp$ add: $HNatInfinite\text{-def}$)

lemma $lemma\text{-unbounded-set}$ [$simp$]: $eventually (\lambda n::nat. m < n) \mathcal{U}$

by ($rule$ $filter\text{-leD}[OF FreeUltrafilterNat.le\text{-cofinite}]$)

($auto$ $simp$ add: $cofinite\text{-eq-sequentially } eventually\text{-at-top-dense}$)

lemma $hypnat\text{-of-nat-eq}: hypnat\text{-of-nat } m = star\text{-}n (\lambda n::nat. m)$

by ($simp$ add: $star\text{-of-def}$)

lemma $SHNat\text{-eq}: Nats = \{n. \exists N. n = hypnat\text{-of-nat } N\}$

by ($simp$ add: $Nats\text{-def } image\text{-def}$)

lemma $Nats\text{-less-whn}: n \in Nats \implies n < whn$

by ($simp$ add: $Nats\text{-less-HNatInfinite}$)

lemma $Nats\text{-le-whn}: n \in Nats \implies n \leq whn$

by ($simp$ add: $Nats\text{-le-HNatInfinite}$)

lemma $hypnat\text{-of-nat-less-whn}$ [$simp$]: $hypnat\text{-of-nat } n < whn$

by ($simp$ add: $Nats\text{-less-whn}$)

lemma $hypnat\text{-of-nat-le-whn}$ [$simp$]: $hypnat\text{-of-nat } n \leq whn$

by ($simp$ add: $Nats\text{-le-whn}$)

lemma $hypnat\text{-zero-less-hypnat-}\omega$ [$simp$]: $0 < whn$

by ($simp$ add: $Nats\text{-less-whn}$)

lemma $hypnat\text{-one-less-hypnat-}\omega$ [$simp$]: $1 < whn$

by ($simp$ add: $Nats\text{-less-whn}$)

3.4.1 Alternative characterization of the set of infinite hypernaturals

$HNatInfinite = \{N. \forall n \in \mathbb{N}. n < N\}$

lemma *HNatInfinite-FreeUltrafilterNat-lemma*:
assumes $\forall N::nat. eventually (\lambda n. f n \neq N) \mathcal{U}$
shows $eventually (\lambda n. N < f n) \mathcal{U}$
apply (*induct N*)
using *assms*
apply (*drule-tac x = 0 in spec, simp*)
using *assms*
apply (*drule-tac x = Suc N in spec*)
apply (*auto elim: eventually-elim2*)
done

lemma *HNatInfinite-iff*: $HNatInfinite = \{N. \forall n \in Nats. n < N\}$
apply (*safe intro!: Nats-less-HNatInfinite*)
apply (*auto simp add: HNatInfinite-def*)
done

3.4.2 Alternative Characterization of $HNatInfinite$ using Free Ultrafilter

lemma *HNatInfinite-FreeUltrafilterNat*:
 $star-n X \in HNatInfinite \implies \forall u. eventually (\lambda n. u < X n) \mathcal{U}$
apply (*auto simp add: HNatInfinite-iff SHNat-eq*)
apply (*drule-tac x=star-of u in spec, simp*)
apply (*simp add: star-of-def star-less-def starP2-star-n*)
done

lemma *FreeUltrafilterNat-HNatInfinite*:
 $\forall u. eventually (\lambda n. u < X n) \mathcal{U} \implies star-n X \in HNatInfinite$
by (*auto simp add: star-less-def starP2-star-n HNatInfinite-iff SHNat-eq hypnat-of-nat-eq*)

lemma *HNatInfinite-FreeUltrafilterNat-iff*:
 $(star-n X \in HNatInfinite) = (\forall u. eventually (\lambda n. u < X n) \mathcal{U})$
by (*rule iffI [OF HNatInfinite-FreeUltrafilterNat FreeUltrafilterNat-HNatInfinite]*)

3.5 Embedding of the Hypernaturals into other types

definition *of-hypnat* :: $hypnat \Rightarrow 'a::semiring-1-cancel star$
where *of-hypnat-def* [*transfer-unfold*]: $of-hypnat = *f*$ *of-nat*

lemma *of-hypnat-0* [*simp*]: $of-hypnat 0 = 0$
by *transfer (rule of-nat-0)*

lemma *of-hypnat-1* [*simp*]: $of-hypnat 1 = 1$
by *transfer (rule of-nat-1)*

lemma *of-hypnat-hSuc*: $\bigwedge m. \text{of-hypnat } (hSuc\ m) = 1 + \text{of-hypnat } m$
by *transfer* (*rule of-nat-Suc*)

lemma *of-hypnat-add* [*simp*]: $\bigwedge m\ n. \text{of-hypnat } (m + n) = \text{of-hypnat } m + \text{of-hypnat } n$
by *transfer* (*rule of-nat-add*)

lemma *of-hypnat-mult* [*simp*]: $\bigwedge m\ n. \text{of-hypnat } (m * n) = \text{of-hypnat } m * \text{of-hypnat } n$
by *transfer* (*rule of-nat-mult*)

lemma *of-hypnat-less-iff* [*simp*]:
 $\bigwedge m\ n. \text{of-hypnat } m < (\text{of-hypnat } n :: 'a :: \text{linordered-semidom star}) \iff m < n$
by *transfer* (*rule of-nat-less-iff*)

lemma *of-hypnat-0-less-iff* [*simp*]:
 $\bigwedge n. 0 < (\text{of-hypnat } n :: 'a :: \text{linordered-semidom star}) \iff 0 < n$
by *transfer* (*rule of-nat-0-less-iff*)

lemma *of-hypnat-less-0-iff* [*simp*]: $\bigwedge m. \neg (\text{of-hypnat } m :: 'a :: \text{linordered-semidom star}) < 0$
by *transfer* (*rule of-nat-less-0-iff*)

lemma *of-hypnat-le-iff* [*simp*]:
 $\bigwedge m\ n. \text{of-hypnat } m \leq (\text{of-hypnat } n :: 'a :: \text{linordered-semidom star}) \iff m \leq n$
by *transfer* (*rule of-nat-le-iff*)

lemma *of-hypnat-0-le-iff* [*simp*]: $\bigwedge n. 0 \leq (\text{of-hypnat } n :: 'a :: \text{linordered-semidom star})$
by *transfer* (*rule of-nat-0-le-iff*)

lemma *of-hypnat-le-0-iff* [*simp*]: $\bigwedge m. (\text{of-hypnat } m :: 'a :: \text{linordered-semidom star}) \leq 0 \iff m = 0$
by *transfer* (*rule of-nat-le-0-iff*)

lemma *of-hypnat-eq-iff* [*simp*]:
 $\bigwedge m\ n. \text{of-hypnat } m = (\text{of-hypnat } n :: 'a :: \text{linordered-semidom star}) \iff m = n$
by *transfer* (*rule of-nat-eq-iff*)

lemma *of-hypnat-eq-0-iff* [*simp*]: $\bigwedge m. (\text{of-hypnat } m :: 'a :: \text{linordered-semidom star}) = 0 \iff m = 0$
by *transfer* (*rule of-nat-eq-0-iff*)

lemma *HNatInfinite-of-hypnat-gt-zero*:
 $N \in \text{HNatInfinite} \implies (0 :: 'a :: \text{linordered-semidom star}) < \text{of-hypnat } N$
by (*rule ccontr*) (*simp add: linorder-not-less*)

end

4 Construction of Hyperreals Using Ultrafilters

```

theory HyperDef
  imports Complex-Main HyperNat
begin

type-synonym hypreal = real star

abbreviation hypreal-of-real :: real  $\Rightarrow$  real star
  where hypreal-of-real  $\equiv$  star-of

abbreviation hypreal-of-hypnat :: hypnat  $\Rightarrow$  hypreal
  where hypreal-of-hypnat  $\equiv$  of-hypnat

definition omega :: hypreal ( $\omega$ )
  where  $\omega = \text{star-n } (\lambda n. \text{real } (\text{Suc } n))$ 
  — an infinite number = [ $\langle 1, 2, 3, \dots \rangle$ ]

definition epsilon :: hypreal ( $\varepsilon$ )
  where  $\varepsilon = \text{star-n } (\lambda n. \text{inverse } (\text{real } (\text{Suc } n)))$ 
  — an infinitesimal number = [ $\langle 1, 1/2, 1/3, \dots \rangle$ ]

4.1 Real vector class instances

instantiation star :: (scaleR) scaleR
begin
  definition star-scaleR-def [transfer-unfold]: scaleR r  $\equiv$  *f* (scaleR r)
  instance ..
end

lemma Standard-scaleR [simp]:  $x \in \text{Standard} \implies \text{scaleR } r \ x \in \text{Standard}$ 
  by (simp add: star-scaleR-def)

lemma star-of-scaleR [simp]:  $\text{star-of } (\text{scaleR } r \ x) = \text{scaleR } r \ (\text{star-of } x)$ 
  by transfer (rule refl)

instance star :: (real-vector) real-vector
proof
  fix a b :: real
  show  $\bigwedge x y :: 'a \text{ star}. \text{scaleR } a \ (x + y) = \text{scaleR } a \ x + \text{scaleR } a \ y$ 
    by transfer (rule scaleR-right-distrib)
  show  $\bigwedge x :: 'a \text{ star}. \text{scaleR } (a + b) \ x = \text{scaleR } a \ x + \text{scaleR } b \ x$ 
    by transfer (rule scaleR-left-distrib)
  show  $\bigwedge x :: 'a \text{ star}. \text{scaleR } a \ (\text{scaleR } b \ x) = \text{scaleR } (a * b) \ x$ 
    by transfer (rule scaleR-scaleR)
  show  $\bigwedge x :: 'a \text{ star}. \text{scaleR } 1 \ x = x$ 
    by transfer (rule scaleR-one)
qed

instance star :: (real-algebra) real-algebra

```

proof

fix $a :: \text{real}$

show $\bigwedge x y :: 'a \text{ star}. \text{scaleR } a \ x * y = \text{scaleR } a \ (x * y)$
by *transfer (rule mult-scaleR-left)*

show $\bigwedge x y :: 'a \text{ star}. x * \text{scaleR } a \ y = \text{scaleR } a \ (x * y)$
by *transfer (rule mult-scaleR-right)*

qed

instance *star* :: (*real-algebra-1*) *real-algebra-1* ..

instance *star* :: (*real-div-algebra*) *real-div-algebra* ..

instance *star* :: (*field-char-0*) *field-char-0* ..

instance *star* :: (*real-field*) *real-field* ..

lemma *star-of-real-def* [*transfer-unfold*]: *of-real* $r = \text{star-of}$ (*of-real* r)
by (*unfold of-real-def, transfer, rule refl*)

lemma *Standard-of-real* [*simp*]: *of-real* $r \in \text{Standard}$
by (*simp add: star-of-real-def*)

lemma *star-of-of-real* [*simp*]: *star-of* (*of-real* r) = *of-real* r
by *transfer (rule refl)*

lemma *of-real-eq-star-of* [*simp*]: *of-real* = *star-of*
proof

show *of-real* $r = \text{star-of}$ r **for** $r :: \text{real}$
by *transfer simp*

qed

lemma *Reals-eq-Standard*: ($\mathbb{R} :: \text{hypreal set}$) = *Standard*
by (*simp add: Reals-def Standard-def*)

4.2 Injection from *hypreal*

definition *of-hypreal* :: *hypreal* \Rightarrow $'a :: \text{real-algebra-1}$ *star*
where [*transfer-unfold*]: *of-hypreal* = $*f*$ *of-real*

lemma *Standard-of-hypreal* [*simp*]: $r \in \text{Standard} \implies \text{of-hypreal } r \in \text{Standard}$
by (*simp add: of-hypreal-def*)

lemma *of-hypreal-0* [*simp*]: *of-hypreal* $0 = 0$
by *transfer (rule of-real-0)*

lemma *of-hypreal-1* [*simp*]: *of-hypreal* $1 = 1$
by *transfer (rule of-real-1)*

lemma *of-hypreal-add* [*simp*]: $\bigwedge x y. \text{of-hypreal } (x + y) = \text{of-hypreal } x + \text{of-hypreal}$

y
by *transfer* (*rule of-real-add*)

lemma *of-hypreal-minus* [*simp*]: $\bigwedge x. \text{of-hypreal } (- x) = - \text{of-hypreal } x$
by *transfer* (*rule of-real-minus*)

lemma *of-hypreal-diff* [*simp*]: $\bigwedge x y. \text{of-hypreal } (x - y) = \text{of-hypreal } x - \text{of-hypreal } y$
by *transfer* (*rule of-real-diff*)

lemma *of-hypreal-mult* [*simp*]: $\bigwedge x y. \text{of-hypreal } (x * y) = \text{of-hypreal } x * \text{of-hypreal } y$
by *transfer* (*rule of-real-mult*)

lemma *of-hypreal-inverse* [*simp*]:
 $\bigwedge x. \text{of-hypreal } (\text{inverse } x) =$
 $\text{inverse } (\text{of-hypreal } x :: 'a::\{\text{real-div-algebra, division-ring}\} \text{star})$
by *transfer* (*rule of-real-inverse*)

lemma *of-hypreal-divide* [*simp*]:
 $\bigwedge x y. \text{of-hypreal } (x / y) =$
 $(\text{of-hypreal } x / \text{of-hypreal } y :: 'a::\{\text{real-field, field}\} \text{star})$
by *transfer* (*rule of-real-divide*)

lemma *of-hypreal-eq-iff* [*simp*]: $\bigwedge x y. (\text{of-hypreal } x = \text{of-hypreal } y) = (x = y)$
by *transfer* (*rule of-real-eq-iff*)

lemma *of-hypreal-eq-0-iff* [*simp*]: $\bigwedge x. (\text{of-hypreal } x = 0) = (x = 0)$
by *transfer* (*rule of-real-eq-0-iff*)

4.3 Properties of *starrel*

lemma *lemma-starrel-refl* [*simp*]: $x \in \text{starrel } \{x\}$
by (*simp add: starrel-def*)

lemma *starrel-in-hypreal* [*simp*]: $\text{starrel } \{x\} : \text{star}$
by (*simp add: star-def starrel-def quotient-def, blast*)

declare *Abs-star-inject* [*simp*] *Abs-star-inverse* [*simp*]
declare *equiv-starrel* [*THEN eq-equiv-class-iff, simp*]

4.4 *hypreal-of-real*: the Injection from *real* to *hypreal*

lemma *inj-star-of*: *inj star-of*
by (*rule inj-onI simp*)

lemma *mem-Rep-star-iff*: $X \in \text{Rep-star } x \longleftrightarrow x = \text{star-n } X$
by (*cases x (simp add: star-n-def)*)

lemma *Rep-star-star-n-iff* [simp]: $X \in \text{Rep-star} (\text{star-n } Y) \longleftrightarrow \text{eventually } (\lambda n. Y\ n = X\ n) \mathcal{U}$

by (*simp add: star-n-def*)

lemma *Rep-star-star-n*: $X \in \text{Rep-star} (\text{star-n } X)$

by *simp*

4.5 Properties of *star-n*

lemma *star-n-add*: $\text{star-n } X + \text{star-n } Y = \text{star-n } (\lambda n. X\ n + Y\ n)$

by (*simp only: star-add-def starfun2-star-n*)

lemma *star-n-minus*: $-\text{star-n } X = \text{star-n } (\lambda n. -(X\ n))$

by (*simp only: star-minus-def starfun-star-n*)

lemma *star-n-diff*: $\text{star-n } X - \text{star-n } Y = \text{star-n } (\lambda n. X\ n - Y\ n)$

by (*simp only: star-diff-def starfun2-star-n*)

lemma *star-n-mult*: $\text{star-n } X * \text{star-n } Y = \text{star-n } (\lambda n. X\ n * Y\ n)$

by (*simp only: star-mult-def starfun2-star-n*)

lemma *star-n-inverse*: $\text{inverse } (\text{star-n } X) = \text{star-n } (\lambda n. \text{inverse } (X\ n))$

by (*simp only: star-inverse-def starfun-star-n*)

lemma *star-n-le*: $\text{star-n } X \leq \text{star-n } Y = \text{eventually } (\lambda n. X\ n \leq Y\ n) \mathcal{U}$

by (*simp only: star-le-def starP2-star-n*)

lemma *star-n-less*: $\text{star-n } X < \text{star-n } Y = \text{eventually } (\lambda n. X\ n < Y\ n) \mathcal{U}$

by (*simp only: star-less-def starP2-star-n*)

lemma *star-n-zero-num*: $0 = \text{star-n } (\lambda n. 0)$

by (*simp only: star-zero-def star-of-def*)

lemma *star-n-one-num*: $1 = \text{star-n } (\lambda n. 1)$

by (*simp only: star-one-def star-of-def*)

lemma *star-n-abs*: $|\text{star-n } X| = \text{star-n } (\lambda n. |X\ n|)$

by (*simp only: star-abs-def starfun-star-n*)

lemma *hypreal-omega-gt-zero* [simp]: $0 < \omega$

by (*simp add: omega-def star-n-zero-num star-n-less*)

4.6 Existence of Infinite Hyperreal Number

Existence of infinite number not corresponding to any real number. Use assumption that member \mathcal{U} is not finite.

A few lemmas first.

lemma *lemma-omega-empty-singleton-disj*:

$\{n::\text{nat}. x = \text{real } n\} = \{\} \vee (\exists y. \{n::\text{nat}. x = \text{real } n\} = \{y\})$
by force

lemma *lemma-finite-omega-set*: *finite* $\{n::\text{nat}. x = \text{real } n\}$
using *lemma-omega-empty-singleton-disj* [of *x*] **by auto**

lemma *not-ex-hypreal-of-real-eq-omega*: $\nexists x. \text{hypreal-of-real } x = \omega$
apply (*simp add: omega-def star-of-def star-n-eq-iff*)
apply *clarify*
apply (*rule-tac x2=x-1 in lemma-finite-omega-set [THEN FreeUltrafilterNat.finite, THEN notE]*)
apply (*erule eventually-mono*)
apply *auto*
done

lemma *hypreal-of-real-not-eq-omega*: *hypreal-of-real* $x \neq \omega$
using *not-ex-hypreal-of-real-eq-omega* **by auto**

Existence of infinitesimal number also not corresponding to any real number.

lemma *lemma-epsilon-empty-singleton-disj*:
 $\{n::\text{nat}. x = \text{inverse}(\text{real}(\text{Suc } n))\} = \{\} \vee (\exists y. \{n::\text{nat}. x = \text{inverse}(\text{real}(\text{Suc } n))\} = \{y\})$
by auto

lemma *lemma-finite-epsilon-set*: *finite* $\{n. x = \text{inverse}(\text{real}(\text{Suc } n))\}$
using *lemma-epsilon-empty-singleton-disj* [of *x*] **by auto**

lemma *not-ex-hypreal-of-real-eq-epsilon*: $\nexists x. \text{hypreal-of-real } x = \varepsilon$
by (*auto simp: epsilon-def star-of-def star-n-eq-iff*
lemma-finite-epsilon-set [THEN FreeUltrafilterNat.finite] simp del: of-nat-Suc)

lemma *hypreal-of-real-not-eq-epsilon*: *hypreal-of-real* $x \neq \varepsilon$
using *not-ex-hypreal-of-real-eq-epsilon* **by auto**

lemma *hypreal-epsilon-not-zero*: $\varepsilon \neq 0$
by (*simp add: epsilon-def star-zero-def star-of-def star-n-eq-iff FreeUltrafilter-Nat.proper*
del: star-of-zero)

lemma *hypreal-epsilon-inverse-omega*: $\varepsilon = \text{inverse } \omega$
by (*simp add: epsilon-def omega-def star-n-inverse*)

lemma *hypreal-epsilon-gt-zero*: $0 < \varepsilon$
by (*simp add: hypreal-epsilon-inverse-omega*)

4.7 Absolute Value Function for the Hyperreals

lemma *hrabs-add-less*: $|x| < r \implies |y| < s \implies |x + y| < r + s$
for $x y r s :: \text{hypreal}$

by (simp add: abs-if split: if-split-asm)

lemma hrabs-less-gt-zero: $|x| < r \implies 0 < r$
 for $x r :: \text{hypreal}$
 by (blast intro!: order-le-less-trans abs-ge-zero)

lemma hrabs-disj: $|x| = x \vee |x| = -x$
 for $x :: 'a::\text{abs-if}$
 by (simp add: abs-if)

lemma hrabs-add-lemma-disj: $y + -x + (y + -z) = |x + -z| \implies y = z \vee x = y$
 for $x y z :: \text{hypreal}$
 by (simp add: abs-if split: if-split-asm)

4.8 Embedding the Naturals into the Hyperreals

abbreviation hypreal-of-nat :: $\text{nat} \Rightarrow \text{hypreal}$
 where hypreal-of-nat \equiv of-nat

lemma SNat-eq: $\text{Nats} = \{n. \exists N. n = \text{hypreal-of-nat } N\}$
 by (simp add: Nats-def image-def)

Naturals embedded in hyperreals: is a hyperreal c.f. NS extension.

lemma hypreal-of-nat: $\text{hypreal-of-nat } m = \text{star-n } (\lambda n. \text{real } m)$
 by (simp add: star-of-def [symmetric])

declaration (

K (Lin-Arith.add-inj-thms [$\text{@}\{thm \text{star-of-le}\} RS \text{iffD2}$,
 $\text{@}\{thm \text{star-of-less}\} RS \text{iffD2}$, $\text{@}\{thm \text{star-of-eq}\} RS \text{iffD2}$]
 $\#>$ Lin-Arith.add-simps [$\text{@}\{thm \text{star-of-zero}\}$, $\text{@}\{thm \text{star-of-one}\}$,
 $\text{@}\{thm \text{star-of-numeral}\}$, $\text{@}\{thm \text{star-of-add}\}$,
 $\text{@}\{thm \text{star-of-minus}\}$, $\text{@}\{thm \text{star-of-diff}\}$, $\text{@}\{thm \text{star-of-mult}\}$]
 $\#>$ Lin-Arith.add-inj-const ($\text{@}\{const-name \text{StarDef.star-of}\}$, $\text{@}\{typ \text{real} \Rightarrow \text{hypreal}\}$))
 \rangle

simproc-setup fast-arith-hypreal (($m::\text{hypreal}$) < n | ($m::\text{hypreal}$) \leq n | ($m::\text{hypreal}$) = n) =
 $\langle K$ Lin-Arith.simproc

4.9 Exponentials on the Hyperreals

lemma hpowr-0 [simp]: $r \wedge 0 = (1::\text{hypreal})$
 for $r :: \text{hypreal}$
 by (rule power-0)

lemma hpowr-Suc [simp]: $r \wedge (\text{Suc } n) = r * (r \wedge n)$
 for $r :: \text{hypreal}$
 by (rule power-Suc)

lemma *hrealpow-two*: $r \hat{\text{Suc}} (\text{Suc } 0) = r * r$
for $r :: \text{hypreal}$
by *simp*

lemma *hrealpow-two-le* [*simp*]: $0 \leq r \hat{\text{Suc}} (\text{Suc } 0)$
for $r :: \text{hypreal}$
by (*auto simp add: zero-le-mult-iff*)

lemma *hrealpow-two-le-add-order* [*simp*]: $0 \leq u \hat{\text{Suc}} (\text{Suc } 0) + v \hat{\text{Suc}} (\text{Suc } 0)$
for $u v :: \text{hypreal}$
by (*simp only: hrealpow-two-le add-nonneg-nonneg*)

lemma *hrealpow-two-le-add-order2* [*simp*]: $0 \leq u \hat{\text{Suc}} (\text{Suc } 0) + v \hat{\text{Suc}} (\text{Suc } 0) + w \hat{\text{Suc}} (\text{Suc } 0)$
for $u v w :: \text{hypreal}$
by (*simp only: hrealpow-two-le add-nonneg-nonneg*)

lemma *hypreal-add-nonneg-eq-0-iff*: $0 \leq x \implies 0 \leq y \implies x + y = 0 \longleftrightarrow x = 0 \wedge y = 0$
for $x y :: \text{hypreal}$
by *arith*

lemma *hypreal-three-squares-add-zero-iff*: $x * x + y * y + z * z = 0 \longleftrightarrow x = 0 \wedge y = 0 \wedge z = 0$
for $x y z :: \text{hypreal}$
by (*simp only: zero-le-square add-nonneg-nonneg hypreal-add-nonneg-eq-0-iff*)
auto

lemma *hrealpow-three-squares-add-zero-iff* [*simp*]:
 $x \hat{\text{Suc}} (\text{Suc } 0) + y \hat{\text{Suc}} (\text{Suc } 0) + z \hat{\text{Suc}} (\text{Suc } 0) = 0 \longleftrightarrow x = 0 \wedge y = 0 \wedge z = 0$
for $x y z :: \text{hypreal}$
by (*simp only: hypreal-three-squares-add-zero-iff hrealpow-two*)

lemma *hrabs-hrealpow-two* [*simp*]: $|x \hat{\text{Suc}} (\text{Suc } 0)| = x \hat{\text{Suc}} (\text{Suc } 0)$
for $x :: \text{hypreal}$
by (*simp add: abs-mult*)

lemma *two-hrealpow-ge-one* [*simp*]: $(1 :: \text{hypreal}) \leq 2 \hat{\text{Suc}} n$
using *power-increasing* [*of 0 n 2::hypreal*] **by** *simp*

lemma *hrealpow*: $\text{star-n } X \hat{\text{Suc}} m = \text{star-n } (\lambda n. (X n :: \text{real}) \hat{\text{Suc}} m)$
by (*induct m*) (*auto simp: star-n-one-num star-n-mult*)

lemma *hrealpow-sum-square-expand*:

$(x + y) \wedge \text{Suc} (\text{Suc } 0) =$
 $x \wedge \text{Suc} (\text{Suc } 0) + y \wedge \text{Suc} (\text{Suc } 0) + (\text{hypreal-of-nat} (\text{Suc} (\text{Suc } 0))) * x * y$
for $x y :: \text{hypreal}$
by (*simp add: distrib-left distrib-right*)

lemma *power-hypreal-of-real-numeral*:

$(\text{numeral } v :: \text{hypreal}) \wedge n = \text{hypreal-of-real} ((\text{numeral } v) \wedge n)$
by *simp*

declare *power-hypreal-of-real-numeral* [*of - numeral w, simp*] **for** w

lemma *power-hypreal-of-real-neg-numeral*:

$(-\text{numeral } v :: \text{hypreal}) \wedge n = \text{hypreal-of-real} ((-\text{numeral } v) \wedge n)$
by *simp*

declare *power-hypreal-of-real-neg-numeral* [*of - numeral w, simp*] **for** w

4.10 Powers with Hypernatural Exponents

Hypernatural powers of hyperreals.

definition *pow* :: $'a::\text{power star} \Rightarrow \text{nat star} \Rightarrow 'a \text{ star}$ (**infixr** *pow* 80)
where *hyperpow-def* [*transfer-unfold*]: $R \text{ pow } N = (*f2* \text{ op } \wedge) R N$

lemma *Standard-hyperpow* [*simp*]: $r \in \text{Standard} \Longrightarrow n \in \text{Standard} \Longrightarrow r \text{ pow } n \in \text{Standard}$

by (*simp add: hyperpow-def*)

lemma *hyperpow*: $\text{star-n } X \text{ pow } \text{star-n } Y = \text{star-n } (\lambda n. X n \wedge Y n)$

by (*simp add: hyperpow-def starfun2-star-n*)

lemma *hyperpow-zero* [*simp*]: $\bigwedge n. (0::'a::\{\text{power, semiring-0}\} \text{ star}) \text{ pow } (n + (1::\text{hypnat})) = 0$

by *transfer simp*

lemma *hyperpow-not-zero*: $\bigwedge r n. r \neq (0::'a::\{\text{field}\} \text{ star}) \Longrightarrow r \text{ pow } n \neq 0$

by *transfer (rule power-not-zero)*

lemma *hyperpow-inverse*: $\bigwedge r n. r \neq (0::'a::\text{field } \text{star}) \Longrightarrow \text{inverse } (r \text{ pow } n) = (\text{inverse } r) \text{ pow } n$

by *transfer (rule power-inverse [symmetric])*

lemma *hyperpow-hrabs*: $\bigwedge r n. |r::'a::\{\text{linordered-idom}\} \text{ star}| \text{ pow } n = |r \text{ pow } n|$

by *transfer (rule power-abs [symmetric])*

lemma *hyperpow-add*: $\bigwedge r n m. (r::'a::\text{monoid-mult } \text{star}) \text{ pow } (n + m) = (r \text{ pow } n) * (r \text{ pow } m)$

by *transfer (rule power-add)*

lemma *hyperpow-one* [*simp*]: $\bigwedge r. (r::'a::\text{monoid-mult } \text{star}) \text{ pow } (1::\text{hypnat}) = r$

by *transfer (rule power-one-right)*

lemma *hyperpow-two*: $\bigwedge r. (r::'a::\text{monoid-mult star}) \text{ pow } (2::\text{hypnat}) = r * r$
by *transfer (rule power2-eq-square)*

lemma *hyperpow-gt-zero*: $\bigwedge r n. (0::'a::\{\text{linordered-semidom}\} \text{ star}) < r \implies 0 < r \text{ pow } n$
by *transfer (rule zero-less-power)*

lemma *hyperpow-ge-zero*: $\bigwedge r n. (0::'a::\{\text{linordered-semidom}\} \text{ star}) \leq r \implies 0 \leq r \text{ pow } n$
by *transfer (rule zero-le-power)*

lemma *hyperpow-le*: $\bigwedge x y n. (0::'a::\{\text{linordered-semidom}\} \text{ star}) < x \implies x \leq y \implies x \text{ pow } n \leq y \text{ pow } n$
by *transfer (rule power-mono [OF - order-less-imp-le])*

lemma *hyperpow-eq-one* [*simp*]: $\bigwedge n. 1 \text{ pow } n = (1::'a::\text{monoid-mult star})$
by *transfer (rule power-one)*

lemma *hrabs-hyperpow-minus* [*simp*]: $\bigwedge (a::'a::\text{linordered-idom star}) n. |(-a) \text{ pow } n| = |a \text{ pow } n|$
by *transfer (rule abs-power-minus)*

lemma *hyperpow-mult*: $\bigwedge r s n. (r * s::'a::\text{comm-monoid-mult star}) \text{ pow } n = (r \text{ pow } n) * (s \text{ pow } n)$
by *transfer (rule power-mult-distrib)*

lemma *hyperpow-two-le* [*simp*]: $\bigwedge r. (0::'a::\{\text{monoid-mult, linordered-ring-strict}\} \text{ star}) \leq r \text{ pow } 2$
by (*auto simp add: hyperpow-two zero-le-mult-iff*)

lemma *hrabs-hyperpow-two* [*simp*]:
 $|x \text{ pow } 2| = (x::'a::\{\text{monoid-mult, linordered-ring-strict}\} \text{ star}) \text{ pow } 2$
by (*simp only: abs-of-nonneg hyperpow-two-le*)

lemma *hyperpow-two-hrabs* [*simp*]: $|x::'a::\text{linordered-idom star}| \text{ pow } 2 = x \text{ pow } 2$
by (*simp add: hyperpow-hrabs*)

The precondition could be weakened to $(0::'a) \leq x$.

lemma *hypreal-mult-less-mono*: $u < v \implies x < y \implies 0 < v \implies 0 < x \implies u * x < v * y$
for $u v x y :: \text{hypreal}$
by (*simp add: mult-strict-mono order-less-imp-le*)

lemma *hyperpow-two-gt-one*: $\bigwedge r::'a::\text{linordered-semidom star}. 1 < r \implies 1 < r \text{ pow } 2$
by *transfer simp*

lemma *hyperpow-two-ge-one*: $\bigwedge r::'a::\text{linordered-semidom star}. 1 \leq r \implies 1 \leq r \text{ pow } 2$

by transfer (rule one-le-power)

lemma two-hyperpow-ge-one [simp]: $(1::\text{hypreal}) \leq 2 \text{ pow } n$
 apply (rule-tac $y = 1 \text{ pow } n$ in order-trans)
 apply (rule-tac [2] hyperpow-le)
 apply auto
 done

lemma hyperpow-minus-one2 [simp]: $\bigwedge n. (-1) \text{ pow } (2 * n) = (1::\text{hypreal})$
 by transfer (rule power-minus1-even)

lemma hyperpow-less-le: $\bigwedge r n N. (0::\text{hypreal}) \leq r \implies r \leq 1 \implies n < N \implies r \text{ pow } N \leq r \text{ pow } n$
 by transfer (rule power-decreasing [OF order-less-imp-le])

lemma hyperpow-SHNat-le:
 $0 \leq r \implies r \leq (1::\text{hypreal}) \implies N \in \text{HNatInfinite} \implies \forall n \in \text{Nats}. r \text{ pow } N \leq r \text{ pow } n$
 by (auto intro!: hyperpow-less-le simp: HNatInfinite-iff)

lemma hyperpow-realpow: $(\text{hypreal-of-real } r) \text{ pow } (\text{hypnat-of-nat } n) = \text{hypreal-of-real } (r \wedge n)$
 by transfer (rule refl)

lemma hyperpow-SReal [simp]: $(\text{hypreal-of-real } r) \text{ pow } (\text{hypnat-of-nat } n) \in \mathbb{R}$
 by (simp add: Reals-eq-Standard)

lemma hyperpow-zero-HNatInfinite [simp]: $N \in \text{HNatInfinite} \implies (0::\text{hypreal}) \text{ pow } N = 0$
 by (drule HNatInfinite-is-Suc, auto)

lemma hyperpow-le-le: $(0::\text{hypreal}) \leq r \implies r \leq 1 \implies n \leq N \implies r \text{ pow } N \leq r \text{ pow } n$
 apply (drule order-le-less [of n, THEN iffD1])
 apply (auto intro: hyperpow-less-le)
 done

lemma hyperpow-Suc-le-self2: $(0::\text{hypreal}) \leq r \implies r < 1 \implies r \text{ pow } (n + (1::\text{hypnat})) \leq r$
 apply (drule-tac $n = (1::\text{hypnat})$ in hyperpow-le-le)
 apply auto
 done

lemma hyperpow-hypnat-of-nat: $\bigwedge x. x \text{ pow } \text{hypnat-of-nat } n = x \wedge n$
 by transfer (rule refl)

lemma of-hypreal-hyperpow:
 $\bigwedge x n. \text{of-hypreal } (x \text{ pow } n) = (\text{of-hypreal } x::'a::\{\text{real-algebra-1}\} \text{ star}) \text{ pow } n$
 by transfer (rule of-real-power)

end

5 Infinite Numbers, Infinitesimals, Infinitely Close Relation

theory NSA

imports HyperDef HOL-Library.Lub-Glb

begin

definition *hnorm* :: 'a::real-normed-vector star \Rightarrow real star

where [transfer-unfold]: *hnorm* = *f* norm

definition *Infinitesimal* :: ('a::real-normed-vector) star set

where *Infinitesimal* = {*x*. $\forall r \in \text{Reals. } 0 < r \longrightarrow \text{hnorm } x < r$ }

definition *HFinite* :: ('a::real-normed-vector) star set

where *HFinite* = {*x*. $\exists r \in \text{Reals. } \text{hnorm } x < r$ }

definition *HInfinite* :: ('a::real-normed-vector) star set

where *HInfinite* = {*x*. $\forall r \in \text{Reals. } r < \text{hnorm } x$ }

definition *approx* :: 'a::real-normed-vector star \Rightarrow 'a star \Rightarrow bool (infixl \approx 50)

where $x \approx y \longleftrightarrow x - y \in \text{Infinitesimal}$

— the “infinitely close” relation

definition *st* :: hypreal \Rightarrow hypreal

where *st* = ($\lambda x. \text{SOME } r. x \in \text{HFinite} \wedge r \in \mathbb{R} \wedge r \approx x$)

— the standard part of a hyperreal

definition *monad* :: 'a::real-normed-vector star \Rightarrow 'a star set

where *monad* *x* = {*y*. $x \approx y$ }

definition *galaxy* :: 'a::real-normed-vector star \Rightarrow 'a star set

where *galaxy* *x* = {*y*. $(x + -y) \in \text{HFinite}$ }

lemma *SReal-def*: $\mathbb{R} \equiv \{x. \exists r. x = \text{hypreal-of-real } r\}$

by (*simp add: Reals-def image-def*)

5.1 Nonstandard Extension of the Norm Function

definition *scaleHR* :: real star \Rightarrow 'a star \Rightarrow 'a::real-normed-vector star

where [transfer-unfold]: *scaleHR* = starfun2 scaleR

lemma *Standard-hnorm* [*simp*]: $x \in \text{Standard} \Longrightarrow \text{hnorm } x \in \text{Standard}$

by (*simp add: hnorm-def*)

lemma *star-of-norm* [*simp*]: $\text{star-of } (\text{norm } x) = \text{hnorm } (\text{star-of } x)$

by *transfer* (rule *refl*)

lemma *hnorm-ge-zero* [*simp*]: $\bigwedge x::'a::\text{real-normed-vector star. } 0 \leq \text{hnorm } x$
by *transfer* (rule *norm-ge-zero*)

lemma *hnorm-eq-zero* [*simp*]: $\bigwedge x::'a::\text{real-normed-vector star. } \text{hnorm } x = 0 \longleftrightarrow x = 0$
by *transfer* (rule *norm-eq-zero*)

lemma *hnorm-triangle-ineq*: $\bigwedge x y::'a::\text{real-normed-vector star. } \text{hnorm } (x + y) \leq \text{hnorm } x + \text{hnorm } y$
by *transfer* (rule *norm-triangle-ineq*)

lemma *hnorm-triangle-ineq3*: $\bigwedge x y::'a::\text{real-normed-vector star. } |\text{hnorm } x - \text{hnorm } y| \leq \text{hnorm } (x - y)$
by *transfer* (rule *norm-triangle-ineq3*)

lemma *hnorm-scaleR*: $\bigwedge x::'a::\text{real-normed-vector star. } \text{hnorm } (a *_{\mathbb{R}} x) = |a| * \text{hnorm } x$
by *transfer* (rule *norm-scaleR*)

lemma *hnorm-scaleHR*: $\bigwedge a (x::'a::\text{real-normed-vector star}). \text{hnorm } (\text{scaleHR } a x) = |a| * \text{hnorm } x$
by *transfer* (rule *norm-scaleR*)

lemma *hnorm-mult-ineq*: $\bigwedge x y::'a::\text{real-normed-algebra star. } \text{hnorm } (x * y) \leq \text{hnorm } x * \text{hnorm } y$
by *transfer* (rule *norm-mult-ineq*)

lemma *hnorm-mult*: $\bigwedge x y::'a::\text{real-normed-div-algebra star. } \text{hnorm } (x * y) = \text{hnorm } x * \text{hnorm } y$
by *transfer* (rule *norm-mult*)

lemma *hnorm-hyperpow*: $\bigwedge (x::'a::\{\text{real-normed-div-algebra}\} \text{ star}) n. \text{hnorm } (x \text{ pow } n) = \text{hnorm } x \text{ pow } n$
by *transfer* (rule *norm-power*)

lemma *hnorm-one* [*simp*]: $\text{hnorm } (1::'a::\text{real-normed-div-algebra star}) = 1$
by *transfer* (rule *norm-one*)

lemma *hnorm-zero* [*simp*]: $\text{hnorm } (0::'a::\text{real-normed-vector star}) = 0$
by *transfer* (rule *norm-zero*)

lemma *zero-less-hnorm-iff* [*simp*]: $\bigwedge x::'a::\text{real-normed-vector star. } 0 < \text{hnorm } x \longleftrightarrow x \neq 0$
by *transfer* (rule *zero-less-norm-iff*)

lemma *hnorm-minus-cancel* [*simp*]: $\bigwedge x::'a::\text{real-normed-vector star. } \text{hnorm } (- x) = \text{hnorm } x$

by *transfer* (*rule norm-minus-cancel*)

lemma *hnorm-minus-commute*: $\bigwedge a b :: 'a :: \text{real-normed-vector star. } \text{hnorm } (a - b) = \text{hnorm } (b - a)$

by *transfer* (*rule norm-minus-commute*)

lemma *hnorm-triangle-ineq2*: $\bigwedge a b :: 'a :: \text{real-normed-vector star. } \text{hnorm } a - \text{hnorm } b \leq \text{hnorm } (a - b)$

by *transfer* (*rule norm-triangle-ineq2*)

lemma *hnorm-triangle-ineq4*: $\bigwedge a b :: 'a :: \text{real-normed-vector star. } \text{hnorm } (a - b) \leq \text{hnorm } a + \text{hnorm } b$

by *transfer* (*rule norm-triangle-ineq4*)

lemma *abs-hnorm-cancel* [*simp*]: $\bigwedge a :: 'a :: \text{real-normed-vector star. } |\text{hnorm } a| = \text{hnorm } a$

by *transfer* (*rule abs-norm-cancel*)

lemma *hnorm-of-hypreal* [*simp*]: $\bigwedge r. \text{hnorm } (\text{of-hypreal } r :: 'a :: \text{real-normed-algebra-1 star}) = |r|$

by *transfer* (*rule norm-of-real*)

lemma *nonzero-hnorm-inverse*:

$\bigwedge a :: 'a :: \text{real-normed-div-algebra star. } a \neq 0 \implies \text{hnorm } (\text{inverse } a) = \text{inverse } (\text{hnorm } a)$

by *transfer* (*rule nonzero-norm-inverse*)

lemma *hnorm-inverse*:

$\bigwedge a :: 'a :: \{\text{real-normed-div-algebra, division-ring}\} \text{ star. } \text{hnorm } (\text{inverse } a) = \text{inverse } (\text{hnorm } a)$

by *transfer* (*rule norm-inverse*)

lemma *hnorm-divide*: $\bigwedge a b :: 'a :: \{\text{real-normed-field, field}\} \text{ star. } \text{hnorm } (a / b) = \text{hnorm } a / \text{hnorm } b$

by *transfer* (*rule norm-divide*)

lemma *hypreal-hnorm-def* [*simp*]: $\bigwedge r :: \text{hypreal. } \text{hnorm } r = |r|$

by *transfer* (*rule real-norm-def*)

lemma *hnorm-add-less*:

$\bigwedge (x :: 'a :: \text{real-normed-vector star}) y r s. \text{hnorm } x < r \implies \text{hnorm } y < s \implies \text{hnorm } (x + y) < r + s$

by *transfer* (*rule norm-add-less*)

lemma *hnorm-mult-less*:

$\bigwedge (x :: 'a :: \text{real-normed-algebra star}) y r s. \text{hnorm } x < r \implies \text{hnorm } y < s \implies \text{hnorm } (x * y) < r * s$

by *transfer* (*rule norm-mult-less*)

lemma *hnorm-scaleHR-less*: $|x| < r \implies \text{hnorm } y < s \implies \text{hnorm } (\text{scaleHR } x \ y) < r * s$

by (*simp only: hnorm-scaleHR*) (*simp add: mult-strict-mono'*)

5.2 Closure Laws for the Standard Reals

lemma *Reals-minus-iff* [*simp*]: $-x \in \mathbb{R} \longleftrightarrow x \in \mathbb{R}$

apply *auto*

apply (*drule Reals-minus*)

apply *auto*

done

lemma *Reals-add-cancel*: $x + y \in \mathbb{R} \implies y \in \mathbb{R} \implies x \in \mathbb{R}$

by (*drule (1) Reals-diff*) *simp*

lemma *SReal-hrabs*: $x \in \mathbb{R} \implies |x| \in \mathbb{R}$

for $x :: \text{hypreal}$

by (*simp add: Reals-eq-Standard*)

lemma *SReal-hypreal-of-real* [*simp*]: *hypreal-of-real* $x \in \mathbb{R}$

by (*simp add: Reals-eq-Standard*)

lemma *SReal-divide-numeral*: $r \in \mathbb{R} \implies r / (\text{numeral } w :: \text{hypreal}) \in \mathbb{R}$

by *simp*

ε is not in Reals because it is an infinitesimal

lemma *SReal-epsilon-not-mem*: $\varepsilon \notin \mathbb{R}$

by (*auto simp: SReal-def hypreal-of-real-not-eq-epsilon* [*symmetric*])

lemma *SReal-omega-not-mem*: $\omega \notin \mathbb{R}$

by (*auto simp: SReal-def hypreal-of-real-not-eq-omega* [*symmetric*])

lemma *SReal-UNIV-real*: $\{x. \text{hypreal-of-real } x \in \mathbb{R}\} = (\text{UNIV} :: \text{real set})$

by *simp*

lemma *SReal-iff*: $x \in \mathbb{R} \longleftrightarrow (\exists y. x = \text{hypreal-of-real } y)$

by (*simp add: SReal-def*)

lemma *hypreal-of-real-image*: *hypreal-of-real* ‘ $(\text{UNIV} :: \text{real set}) = \mathbb{R}$

by (*simp add: Reals-eq-Standard Standard-def*)

lemma *inv-hypreal-of-real-image*: *inv hypreal-of-real* ‘ $\mathbb{R} = \text{UNIV}$

apply (*auto simp add: SReal-def*)

apply (*rule inj-star-of* [*THEN inv-f-f*, *THEN subst*], *blast*)

done

lemma *SReal-hypreal-of-real-image*: $\exists x. x \in P \implies P \subseteq \mathbb{R} \implies \exists Q. P = \text{hypreal-of-real } ‘Q$

unfolding *SReal-def image-def* **by** *blast*

```

lemma SReal-dense:  $x \in \mathbb{R} \implies y \in \mathbb{R} \implies x < y \implies \exists r \in \text{Reals}. x < r \wedge r < y$ 
for  $x\ y :: \text{hypreal}$ 
apply (auto simp: SReal-def)
apply (drule dense)
apply auto
done

```

Completeness of Reals, but both lemmas are unused.

```

lemma SReal-sup-lemma:
 $P \subseteq \mathbb{R} \implies (\exists x \in P. y < x) = (\exists X. \text{hypreal-of-real } X \in P \wedge y < \text{hypreal-of-real } X)$ 
by (blast dest!: SReal-iff [THEN iffD1])

```

```

lemma SReal-sup-lemma2:
 $P \subseteq \mathbb{R} \implies \exists x. x \in P \implies \exists y \in \text{Reals}. \forall x \in P. x < y \implies$ 
 $(\exists X. X \in \{w. \text{hypreal-of-real } w \in P\}) \wedge$ 
 $(\exists Y. \forall X \in \{w. \text{hypreal-of-real } w \in P\}. X < Y)$ 
apply (rule conjI)
apply (fast dest!: SReal-iff [THEN iffD1])
apply (auto, frule subsetD, assumption)
apply (drule SReal-iff [THEN iffD1])
apply (auto, rule-tac x = ya in exI, auto)
done

```

5.3 Set of Finite Elements is a Subring of the Extended Reals

```

lemma HFinite-add:  $x \in \text{HFinite} \implies y \in \text{HFinite} \implies x + y \in \text{HFinite}$ 
unfolding HFinite-def by (blast intro!: Reals-add hnorm-add-less)

```

```

lemma HFinite-mult:  $x \in \text{HFinite} \implies y \in \text{HFinite} \implies x * y \in \text{HFinite}$ 
for  $x\ y :: 'a::\text{real-normed-algebra star}$ 
unfolding HFinite-def by (blast intro!: Reals-mult hnorm-mult-less)

```

```

lemma HFinite-scaleHR:  $x \in \text{HFinite} \implies y \in \text{HFinite} \implies \text{scaleHR } x\ y \in \text{HFinite}$ 
by (auto simp: HFinite-def intro!: Reals-mult hnorm-scaleHR-less)

```

```

lemma HFinite-minus-iff:  $-x \in \text{HFinite} \longleftrightarrow x \in \text{HFinite}$ 
by (simp add: HFinite-def)

```

```

lemma HFinite-star-of [simp]:  $\text{star-of } x \in \text{HFinite}$ 
apply (simp add: HFinite-def)
apply (rule-tac x=star-of (norm x) + 1 in bezI)
apply (transfer, simp)
apply (blast intro: Reals-add SReal-hypreal-of-real Reals-1)
done

```

```

lemma SReal-subset-HFinite:  $(\mathbb{R}::\text{hypreal set}) \subseteq \text{HFinite}$ 

```

by (auto simp add: SReal-def)

lemma *HFfiniteD*: $x \in \text{HFfinite} \implies \exists t \in \text{Reals}. \text{hnorm } x < t$
 by (simp add: HFfinite-def)

lemma *HFfinite-hrabs-iff* [iff]: $|x| \in \text{HFfinite} \longleftrightarrow x \in \text{HFfinite}$
 for $x :: \text{hypreal}$
 by (simp add: HFfinite-def)

lemma *HFfinite-hnorm-iff* [iff]: $\text{hnorm } x \in \text{HFfinite} \longleftrightarrow x \in \text{HFfinite}$
 for $x :: \text{hypreal}$
 by (simp add: HFfinite-def)

lemma *HFfinite-numeral* [simp]: numeral $w \in \text{HFfinite}$
 unfolding star-numeral-def by (rule HFfinite-star-of)

As always with numerals, 0 and 1 are special cases.

lemma *HFfinite-0* [simp]: $0 \in \text{HFfinite}$
 unfolding star-zero-def by (rule HFfinite-star-of)

lemma *HFfinite-1* [simp]: $1 \in \text{HFfinite}$
 unfolding star-one-def by (rule HFfinite-star-of)

lemma *hrealpow-HFfinite*: $x \in \text{HFfinite} \implies x \wedge n \in \text{HFfinite}$
 for $x :: 'a :: \{\text{real-normed-algebra}, \text{monoid-mult}\}$ star
 by (induct n) (auto simp add: power-Suc intro: HFfinite-mult)

lemma *HFfinite-bounded*: $x \in \text{HFfinite} \implies y \leq x \implies 0 \leq y \implies y \in \text{HFfinite}$
 for $x \ y :: \text{hypreal}$
 apply (cases $x \leq 0$)
 apply (drule-tac $y = x$ in order-trans)
 apply (drule-tac [2] order-antisym)
 apply (auto simp add: linorder-not-le)
 apply (auto intro: order-le-less-trans simp add: abs-if HFfinite-def)
 done

5.4 Set of Infinitesimals is a Subring of the Hyperreals

lemma *InfinitesimalI*: $(\bigwedge r. r \in \mathbb{R} \implies 0 < r \implies \text{hnorm } x < r) \implies x \in \text{Infinitesimal}$
 by (simp add: Infinitesimal-def)

lemma *InfinitesimalD*: $x \in \text{Infinitesimal} \implies \forall r \in \text{Reals}. 0 < r \longrightarrow \text{hnorm } x < r$
 by (simp add: Infinitesimal-def)

lemma *InfinitesimalII2*: $(\bigwedge r. 0 < r \implies \text{hnorm } x < \text{star-of } r) \implies x \in \text{Infinitesimal}$
 by (auto simp add: Infinitesimal-def SReal-def)

lemma *InfinesimalD2*: $x \in \text{Infinesimal} \implies 0 < r \implies \text{hnorm } x < \text{star-of } r$
by (*auto simp add: Infinesimal-def SReal-def*)

lemma *Infinesimal-zero* [*iff*]: $0 \in \text{Infinesimal}$
by (*simp add: Infinesimal-def*)

lemma *hypreal-sum-of-halves*: $x / 2 + x / 2 = x$
for $x :: \text{hypreal}$
by *auto*

lemma *Infinesimal-add*: $x \in \text{Infinesimal} \implies y \in \text{Infinesimal} \implies x + y \in \text{Infinesimal}$
apply (*rule InfinesimalI*)
apply (*rule hypreal-sum-of-halves [THEN subst]*)
apply (*drule half-gt-zero*)
apply (*blast intro: hnorm-add-less SReal-divide-numeral dest: InfinesimalD*)
done

lemma *Infinesimal-minus-iff* [*simp*]: $-x \in \text{Infinesimal} \longleftrightarrow x \in \text{Infinesimal}$
by (*simp add: Infinesimal-def*)

lemma *Infinesimal-hnorm-iff*: $\text{hnorm } x \in \text{Infinesimal} \longleftrightarrow x \in \text{Infinesimal}$
by (*simp add: Infinesimal-def*)

lemma *Infinesimal-hrabs-iff* [*iff*]: $|x| \in \text{Infinesimal} \longleftrightarrow x \in \text{Infinesimal}$
for $x :: \text{hypreal}$
by (*simp add: abs-if*)

lemma *Infinesimal-of-hypreal-iff* [*simp*]:
(*of-hypreal* $x :: 'a :: \text{real-normed-algebra-1 star}$) $\in \text{Infinesimal} \longleftrightarrow x \in \text{Infinesimal}$
by (*subst Infinesimal-hnorm-iff [symmetric]*) *simp*

lemma *Infinesimal-diff*: $x \in \text{Infinesimal} \implies y \in \text{Infinesimal} \implies x - y \in \text{Infinesimal}$
using *Infinesimal-add [of $x - y$]* **by** *simp*

lemma *Infinesimal-mult*: $x \in \text{Infinesimal} \implies y \in \text{Infinesimal} \implies x * y \in \text{Infinesimal}$
for $x y :: 'a :: \text{real-normed-algebra star}$
apply (*rule InfinesimalI*)
apply (*subgoal-tac hnorm (x * y) < 1 * r*)
apply (*simp only: mult-1*)
apply (*rule hnorm-mult-less*)
apply (*simp-all add: InfinesimalD*)
done

lemma *Infinesimal-HFinite-mult*: $x \in \text{Infinesimal} \implies y \in \text{HFinite} \implies x * y \in \text{Infinesimal}$
for $x y :: 'a :: \text{real-normed-algebra star}$

```

apply (rule InfinitesimalI)
apply (drule HFiniteD, clarify)
apply (subgoal-tac  $0 < t$ )
apply (subgoal-tac  $\text{hnorm } (x * y) < (r / t) * t$ , simp)
apply (subgoal-tac  $0 < r / t$ )
apply (rule hnorm-mult-less)
apply (simp add: InfinitesimalD)
apply assumption
apply simp
apply (erule order-le-less-trans [OF hnorm-ge-zero])
done

```

lemma *Infinitesimal-HFinite-scaleHR*:

```

 $x \in \text{Infinitesimal} \implies y \in \text{HFinite} \implies \text{scaleHR } x \ y \in \text{Infinitesimal}$ 
apply (rule InfinitesimalI)
apply (drule HFiniteD, clarify)
apply (drule InfinitesimalD)
apply (simp add: hnorm-scaleHR)
apply (subgoal-tac  $0 < t$ )
apply (subgoal-tac  $|x| * \text{hnorm } y < (r / t) * t$ , simp)
apply (subgoal-tac  $0 < r / t$ )
apply (rule mult-strict-mono', simp-all)
apply (erule order-le-less-trans [OF hnorm-ge-zero])
done

```

lemma *Infinitesimal-HFinite-mult2*:

```

 $x \in \text{Infinitesimal} \implies y \in \text{HFinite} \implies y * x \in \text{Infinitesimal}$ 
for  $x \ y :: 'a :: \text{real-normed-algebra star}$ 
apply (rule InfinitesimalI)
apply (drule HFiniteD, clarify)
apply (subgoal-tac  $0 < t$ )
apply (subgoal-tac  $\text{hnorm } (y * x) < t * (r / t)$ , simp)
apply (subgoal-tac  $0 < r / t$ )
apply (rule hnorm-mult-less)
apply assumption
apply (simp add: InfinitesimalD)
apply simp
apply (erule order-le-less-trans [OF hnorm-ge-zero])
done

```

lemma *Infinitesimal-scaleR2*: $x \in \text{Infinitesimal} \implies a *_{\mathbb{R}} x \in \text{Infinitesimal}$

```

apply (case-tac  $a = 0$ , simp)
apply (rule InfinitesimalI)
apply (drule InfinitesimalD)
apply (drule-tac  $x=r / |star-of a|$  in bspec)
apply (simp add: Reals-eq-Standard)
apply simp
apply (simp add: hnorm-scaleR pos-less-divide-eq mult commute)
done

```


lemma *Compl-HFinite: $- HFinite = HInfinite$*
apply (*auto simp add: HInfinite-def HFinite-def linorder-not-less*)
apply (*rule-tac $y=r + 1$ in order-less-le-trans, simp*)
apply *simp*
done

lemma *HInfinite-inverse-Infinitesimal: $x \in HInfinite \implies inverse\ x \in Infinitesimal$*
for $x :: 'a::real-normed-div-algebra\ star$
apply (*rule InfinitesimalI*)
apply (*subgoal-tac $x \neq 0$*)
apply (*rule inverse-less-imp-less*)
apply (*simp add: nonzero-hnorm-inverse*)
apply (*simp add: HInfinite-def Reals-inverse*)
apply *assumption*
apply (*clarify, simp add: Compl-HFinite [symmetric]*)
done

lemma *HInfiniteI: $(\bigwedge r. r \in \mathbb{R} \implies r < hnorm\ x) \implies x \in HInfinite$*
by (*simp add: HInfinite-def*)

lemma *HInfiniteD: $x \in HInfinite \implies r \in \mathbb{R} \implies r < hnorm\ x$*
by (*simp add: HInfinite-def*)

lemma *HInfinite-mult: $x \in HInfinite \implies y \in HInfinite \implies x * y \in HInfinite$*
for $x\ y :: 'a::real-normed-div-algebra\ star$
apply (*rule HInfiniteI, simp only: hnorm-mult*)
apply (*subgoal-tac $r * 1 < hnorm\ x * hnorm\ y$, simp only: mult-1*)
apply (*case-tac $x = 0$, simp add: HInfinite-def*)
apply (*rule mult-strict-mono*)
apply (*simp-all add: HInfiniteD*)
done

lemma *hypreal-add-zero-less-le-mono: $r < x \implies 0 \leq y \implies r < x + y$*
for $r\ x\ y :: hypreal$
by (*auto dest: add-less-le-mono*)

lemma *HInfinite-add-ge-zero: $x \in HInfinite \implies 0 \leq y \implies 0 \leq x \implies x + y \in HInfinite$*
for $x\ y :: hypreal$
by (*auto simp: abs-if add commute HInfinite-def*)

lemma *HInfinite-add-ge-zero2: $x \in HInfinite \implies 0 \leq y \implies 0 \leq x \implies y + x \in HInfinite$*
for $x\ y :: hypreal$
by (*auto intro!: HInfinite-add-ge-zero simp add: add commute*)

lemma *HInfinite-add-gt-zero: $x \in HInfinite \implies 0 < y \implies 0 < x \implies x + y \in HInfinite$*

for $x\ y :: \text{hypreal}$
by (*blast intro: HInfinite-add-ge-zero order-less-imp-le*)

lemma *HInfinite-minus-iff*: $-x \in \text{HInfinite} \longleftrightarrow x \in \text{HInfinite}$
by (*simp add: HInfinite-def*)

lemma *HInfinite-add-le-zero*: $x \in \text{HInfinite} \implies y \leq 0 \implies x \leq 0 \implies x + y \in \text{HInfinite}$

for $x\ y :: \text{hypreal}$
apply (*drule HInfinite-minus-iff [THEN iffD2]*)
apply (*rule HInfinite-minus-iff [THEN iffD1]*)
apply (*simp only: minus-add add commute*)
apply (*rule HInfinite-add-ge-zero*)
apply *simp-all*
done

lemma *HInfinite-add-lt-zero*: $x \in \text{HInfinite} \implies y < 0 \implies x < 0 \implies x + y \in \text{HInfinite}$

for $x\ y :: \text{hypreal}$
by (*blast intro: HInfinite-add-le-zero order-less-imp-le*)

lemma *HFinite-sum-squares*:

$a \in \text{HFinite} \implies b \in \text{HFinite} \implies c \in \text{HFinite} \implies a * a + b * b + c * c \in \text{HFinite}$

for $a\ b\ c :: 'a::\text{real-normed-algebra star}$
by (*auto intro: HFinite-mult HFinite-add*)

lemma *not-Infinitesimal-not-zero*: $x \notin \text{Infinitesimal} \implies x \neq 0$
by *auto*

lemma *not-Infinitesimal-not-zero2*: $x \in \text{HFinite} - \text{Infinitesimal} \implies x \neq 0$
by *auto*

lemma *HFinite-diff-Infinitesimal-hrabs*:

$x \in \text{HFinite} - \text{Infinitesimal} \implies |x| \in \text{HFinite} - \text{Infinitesimal}$
for $x :: \text{hypreal}$
by *blast*

lemma *hnorm-le-Infinitesimal*: $e \in \text{Infinitesimal} \implies \text{hnorm } x \leq e \implies x \in \text{Infinitesimal}$

by (*auto simp: Infinitesimal-def abs-less-iff*)

lemma *hnorm-less-Infinitesimal*: $e \in \text{Infinitesimal} \implies \text{hnorm } x < e \implies x \in \text{Infinitesimal}$

by (*erule hnorm-le-Infinitesimal, erule order-less-imp-le*)

lemma *hrabs-le-Infinitesimal*: $e \in \text{Infinitesimal} \implies |x| \leq e \implies x \in \text{Infinitesimal}$

for $x :: \text{hypreal}$
by (*erule hnorm-le-Infinitesimal*) *simp*

lemma *hrabs-less-Infinitesimal*: $e \in \text{Infinitesimal} \implies |x| < e \implies x \in \text{Infinitesimal}$
for $x :: \text{hypreal}$
by (*erule hnorm-less-Infinitesimal*) *simp*

lemma *Infinitesimal-interval*:
 $e \in \text{Infinitesimal} \implies e' \in \text{Infinitesimal} \implies e' < x \implies x < e \implies x \in \text{Infinitesimal}$
for $x :: \text{hypreal}$
by (*auto simp add: Infinitesimal-def abs-less-iff*)

lemma *Infinitesimal-interval2*:
 $e \in \text{Infinitesimal} \implies e' \in \text{Infinitesimal} \implies e' \leq x \implies x \leq e \implies x \in \text{Infinitesimal}$
for $x :: \text{hypreal}$
by (*auto intro: Infinitesimal-interval simp add: order-le-less*)

lemma *lemma-Infinitesimal-hyperpow*: $x \in \text{Infinitesimal} \implies 0 < N \implies |x \text{ pow } N| \leq |x|$
for $x :: \text{hypreal}$
apply (*unfold Infinitesimal-def*)
apply (*auto intro!: hyperpow-Suc-le-self2*)
simp: hyperpow-hrabs [symmetric] hypnat-gt-zero-iff2 abs-ge-zero
done

lemma *Infinitesimal-hyperpow*: $x \in \text{Infinitesimal} \implies 0 < N \implies x \text{ pow } N \in \text{Infinitesimal}$
for $x :: \text{hypreal}$
apply (*rule hrabs-le-Infinitesimal*)
apply (*rule-tac [2] lemma-Infinitesimal-hyperpow*)
apply *auto*
done

lemma *hrealpow-hyperpow-Infinitesimal-iff*:
 $(x \wedge n \in \text{Infinitesimal}) \longleftrightarrow x \text{ pow } (\text{hypnat-of-nat } n) \in \text{Infinitesimal}$
by (*simp only: hyperpow-hypnat-of-nat*)

lemma *Infinitesimal-hrealpow*: $x \in \text{Infinitesimal} \implies 0 < n \implies x \wedge n \in \text{Infinitesimal}$
for $x :: \text{hypreal}$
by (*simp add: hrealpow-hyperpow-Infinitesimal-iff Infinitesimal-hyperpow*)

lemma *not-Infinitesimal-mult*:
 $x \notin \text{Infinitesimal} \implies y \notin \text{Infinitesimal} \implies x * y \notin \text{Infinitesimal}$
for $x y :: 'a :: \text{real-normed-div-algebra star}$
apply (*unfold Infinitesimal-def, clarify, rename-tac r s*)
apply (*simp only: linorder-not-less hnorm-mult*)
apply (*drule-tac x = r * s in bspec*)
apply (*fast intro: Reals-mult*)

```

apply simp
apply (drule-tac c = s and d = hnorm y and a = r and b = hnorm x in
mult-mono)
apply simp-all
done

```

```

lemma Infinitesimal-mult-disj:  $x * y \in \text{Infinitesimal} \implies x \in \text{Infinitesimal} \vee y \in \text{Infinitesimal}$ 
for x y :: 'a::real-normed-div-algebra star
apply (rule ccontr)
apply (drule de-Morgan-disj [THEN iffD1])
apply (fast dest: not-Infinitesimal-mult)
done

```

```

lemma HFinite-Infinitesimal-not-zero:  $x \in \text{HFinite} - \text{Infinitesimal} \implies x \neq 0$ 
by blast

```

```

lemma HFinite-Infinitesimal-diff-mult:
 $x \in \text{HFinite} - \text{Infinitesimal} \implies y \in \text{HFinite} - \text{Infinitesimal} \implies x * y \in \text{HFinite} - \text{Infinitesimal}$ 
for x y :: 'a::real-normed-div-algebra star
apply clarify
apply (blast dest: HFinite-mult not-Infinitesimal-mult)
done

```

```

lemma Infinitesimal-subset-HFinite:  $\text{Infinitesimal} \subseteq \text{HFinite}$ 
apply (simp add: Infinitesimal-def HFinite-def)
apply auto
apply (rule-tac x = 1 in bexI)
apply auto
done

```

```

lemma Infinitesimal-star-of-mult:  $x \in \text{Infinitesimal} \implies x * \text{star-of } r \in \text{Infinitesimal}$ 
for x :: 'a::real-normed-algebra star
by (erule HFinite-star-of [THEN [2] Infinitesimal-HFinite-mult])

```

```

lemma Infinitesimal-star-of-mult2:  $x \in \text{Infinitesimal} \implies \text{star-of } r * x \in \text{Infinitesimal}$ 
for x :: 'a::real-normed-algebra star
by (erule HFinite-star-of [THEN [2] Infinitesimal-HFinite-mult2])

```

5.5 The Infinitely Close Relation

```

lemma mem-infmal-iff:  $x \in \text{Infinitesimal} \iff x \approx 0$ 
by (simp add: Infinitesimal-def approx-def)

```

```

lemma approx-minus-iff:  $x \approx y \iff x - y \approx 0$ 
by (simp add: approx-def)

```

```

lemma approx-minus-iff2:  $x \approx y \iff -y + x \approx 0$ 

```

by (simp add: approx-def add commute)

lemma approx-refl [iff]: $x \approx x$
by (simp add: approx-def Infinitesimal-def)

lemma hypreal-minus-distrib1: $-(y + -x) = x + -y$
for $x y :: 'a::ab-group-add$
by (simp add: add commute)

lemma approx-sym: $x \approx y \implies y \approx x$
apply (simp add: approx-def)
apply (drule Infinitesimal-minus-iff [THEN iffD2])
apply simp
done

lemma approx-trans: $x \approx y \implies y \approx z \implies x \approx z$
apply (simp add: approx-def)
apply (drule (1) Infinitesimal-add)
apply simp
done

lemma approx-trans2: $r \approx x \implies s \approx x \implies r \approx s$
by (blast intro: approx-sym approx-trans)

lemma approx-trans3: $x \approx r \implies x \approx s \implies r \approx s$
by (blast intro: approx-sym approx-trans)

lemma approx-reorient: $x \approx y \longleftrightarrow y \approx x$
by (blast intro: approx-sym)

Reorientation simplification procedure: reorients (polymorphic) $0 = x$, $1 = x$, $nnn = x$ provided x isn't 0 , 1 or a numeral.

simproc-setup approx-reorient-simproc
($0 \approx x \mid 1 \approx y \mid \text{numeral } w \approx z \mid -1 \approx y \mid - \text{numeral } w \approx r$) =
(
 let val rule = @{thm approx-reorient} RS eq-reflection
 fun proc phi ss ct =
 case Thm.term-of ct of
 - \$ t \$ u => if can HOLogic.dest-number u then NONE
 else if can HOLogic.dest-number t then SOME rule else NONE
 | - => NONE
 in proc end
)

lemma Infinitesimal-approx-minus: $x - y \in \text{Infinitesimal} \longleftrightarrow x \approx y$
by (simp add: approx-minus-iff [symmetric] mem-infmal-iff)

lemma approx-monad-iff: $x \approx y \longleftrightarrow \text{monad } x = \text{monad } y$
by (auto simp add: monad-def dest: approx-sym elim!: approx-trans equalityCE)

lemma *Infinitesimal-approx*: $x \in \text{Infinitesimal} \implies y \in \text{Infinitesimal} \implies x \approx y$
apply (*simp add: mem-infmal-iff*)
apply (*blast intro: approx-trans approx-sym*)
done

lemma *approx-add*: $a \approx b \implies c \approx d \implies a + c \approx b + d$
proof (*unfold approx-def*)
assume *inf*: $a - b \in \text{Infinitesimal}$ $c - d \in \text{Infinitesimal}$
have $a + c - (b + d) = (a - b) + (c - d)$ **by** *simp*
also have $\dots \in \text{Infinitesimal}$
using *inf* **by** (*rule Infinitesimal-add*)
finally show $a + c - (b + d) \in \text{Infinitesimal}$.
qed

lemma *approx-minus*: $a \approx b \implies -a \approx -b$
apply (*rule approx-minus-iff [THEN iffD2, THEN approx-sym]*)
apply (*drule approx-minus-iff [THEN iffD1]*)
apply (*simp add: add.commute*)
done

lemma *approx-minus2*: $-a \approx -b \implies a \approx b$
by (*auto dest: approx-minus*)

lemma *approx-minus-cancel [simp]*: $-a \approx -b \longleftrightarrow a \approx b$
by (*blast intro: approx-minus approx-minus2*)

lemma *approx-add-minus*: $a \approx b \implies c \approx d \implies a + -c \approx b + -d$
by (*blast intro!: approx-add approx-minus*)

lemma *approx-diff*: $a \approx b \implies c \approx d \implies a - c \approx b - d$
using *approx-add [of a b - c - d]* **by** *simp*

lemma *approx-mult1*: $a \approx b \implies c \in \text{HFinite} \implies a * c \approx b * c$
for $a b c :: 'a::\text{real-normed-algebra}$ *star*
by (*simp add: approx-def Infinitesimal-HFinite-mult left-diff-distrib [symmetric]*)

lemma *approx-mult2*: $a \approx b \implies c \in \text{HFinite} \implies c * a \approx c * b$
for $a b c :: 'a::\text{real-normed-algebra}$ *star*
by (*simp add: approx-def Infinitesimal-HFinite-mult2 right-diff-distrib [symmetric]*)

lemma *approx-mult-subst*: $u \approx v * x \implies x \approx y \implies v \in \text{HFinite} \implies u \approx v * y$
for $u v x y :: 'a::\text{real-normed-algebra}$ *star*
by (*blast intro: approx-mult2 approx-trans*)

lemma *approx-mult-subst2*: $u \approx x * v \implies x \approx y \implies v \in \text{HFinite} \implies u \approx y * v$
for $u v x y :: 'a::\text{real-normed-algebra}$ *star*
by (*blast intro: approx-mult1 approx-trans*)

lemma *approx-mult-subst-star-of*: $u \approx x * \text{star-of } v \implies x \approx y \implies u \approx y * \text{star-of } v$

for $u x y :: 'a::\text{real-normed-algebra star}$
by (*auto intro: approx-mult-subst2*)

lemma *approx-eq-imp*: $a = b \implies a \approx b$
by (*simp add: approx-def*)

lemma *Infinitesimal-minus-approx*: $x \in \text{Infinitesimal} \implies -x \approx x$
by (*blast intro: Infinitesimal-minus-iff [THEN iffD2] mem-infmal-iff [THEN iffD1] approx-trans2*)

lemma *bex-Infinitesimal-iff*: $(\exists y \in \text{Infinitesimal}. x - z = y) \longleftrightarrow x \approx z$
by (*simp add: approx-def*)

lemma *bex-Infinitesimal-iff2*: $(\exists y \in \text{Infinitesimal}. x = z + y) \longleftrightarrow x \approx z$
by (*force simp add: bex-Infinitesimal-iff [symmetric]*)

lemma *Infinitesimal-add-approx*: $y \in \text{Infinitesimal} \implies x + y = z \implies x \approx z$
apply (*rule bex-Infinitesimal-iff [THEN iffD1]*)
apply (*drule Infinitesimal-minus-iff [THEN iffD2]*)
apply (*auto simp add: add.assoc [symmetric]*)
done

lemma *Infinitesimal-add-approx-self*: $y \in \text{Infinitesimal} \implies x \approx x + y$
apply (*rule bex-Infinitesimal-iff [THEN iffD1]*)
apply (*drule Infinitesimal-minus-iff [THEN iffD2]*)
apply (*auto simp add: add.assoc [symmetric]*)
done

lemma *Infinitesimal-add-approx-self2*: $y \in \text{Infinitesimal} \implies x \approx y + x$
by (*auto dest: Infinitesimal-add-approx-self simp add: add.commute*)

lemma *Infinitesimal-add-minus-approx-self*: $y \in \text{Infinitesimal} \implies x \approx x + -y$
by (*blast intro!: Infinitesimal-add-approx-self Infinitesimal-minus-iff [THEN iffD2]*)

lemma *Infinitesimal-add-cancel*: $y \in \text{Infinitesimal} \implies x + y \approx z \implies x \approx z$
apply (*drule-tac x = x in Infinitesimal-add-approx-self [THEN approx-sym]*)
apply (*erule approx-trans3 [THEN approx-sym], assumption*)
done

lemma *Infinitesimal-add-right-cancel*: $y \in \text{Infinitesimal} \implies x \approx z + y \implies x \approx z$
apply (*drule-tac x = z in Infinitesimal-add-approx-self2 [THEN approx-sym]*)
apply (*erule approx-trans3 [THEN approx-sym]*)
apply (*simp add: add.commute*)
apply (*erule approx-sym*)
done

lemma *approx-add-left-cancel*: $d + b \approx d + c \implies b \approx c$

```

apply (drule approx-minus-iff [THEN iffD1])
apply (simp add: approx-minus-iff [symmetric] ac-simps)
done

lemma approx-add-right-cancel:  $b + d \approx c + d \implies b \approx c$ 
apply (rule approx-add-left-cancel)
apply (simp add: add.commute)
done

lemma approx-add-mono1:  $b \approx c \implies d + b \approx d + c$ 
apply (rule approx-minus-iff [THEN iffD2])
apply (simp add: approx-minus-iff [symmetric] ac-simps)
done

lemma approx-add-mono2:  $b \approx c \implies b + a \approx c + a$ 
by (simp add: add.commute approx-add-mono1)

lemma approx-add-left-iff [simp]:  $a + b \approx a + c \longleftrightarrow b \approx c$ 
by (fast elim: approx-add-left-cancel approx-add-mono1)

lemma approx-add-right-iff [simp]:  $b + a \approx c + a \longleftrightarrow b \approx c$ 
by (simp add: add.commute)

lemma approx-HFinite:  $x \in \text{HFinite} \implies x \approx y \implies y \in \text{HFinite}$ 
apply (drule be-Infinitesimal-iff2 [THEN iffD2], safe)
apply (drule Infinitesimal-subset-HFinite [THEN subsetD, THEN HFinite-minus-iff
[THEN iffD2]])
apply (drule HFinite-add)
apply (auto simp add: add.assoc)
done

lemma approx-star-of-HFinite:  $x \approx \text{star-of } D \implies x \in \text{HFinite}$ 
by (rule approx-sym [THEN [2] approx-HFinite], auto)

lemma approx-mult-HFinite:  $a \approx b \implies c \approx d \implies b \in \text{HFinite} \implies d \in \text{HFinite}$ 
 $\implies a * c \approx b * d$ 
for  $a b c d :: 'a::\text{real-normed-algebra star}$ 
apply (rule approx-trans)
apply (rule-tac [2] approx-mult2)
apply (rule approx-mult1)
prefer 2 apply (blast intro: approx-HFinite approx-sym, auto)
done

lemma scaleHR-left-diff-distrib:  $\bigwedge a b x. \text{scaleHR } (a - b) x = \text{scaleHR } a x - \text{scaleHR } b x$ 
by transfer (rule scaleR-left-diff-distrib)

lemma approx-scaleR1:  $a \approx \text{star-of } b \implies c \in \text{HFinite} \implies \text{scaleHR } a c \approx b *_R c$ 
apply (unfold approx-def)

```



```

apply (drule (1) Infinitesimal-HFfinite-scaleHR)
apply (simp only: scaleHR-left-diff-distrib)
apply (simp add: scaleHR-def star-scaleR-def [symmetric])
done

```

```

lemma approx-scaleR2:  $a \approx b \implies c *_R a \approx c *_R b$ 
by (simp add: approx-def Infinitesimal-scaleR2 scaleR-right-diff-distrib [symmetric])

```

```

lemma approx-scaleR-HFfinite:  $a \approx \text{star-of } b \implies c \approx d \implies d \in \text{HFfinite} \implies$ 
 $\text{scaleHR } a \ c \approx b *_R d$ 
apply (rule approx-trans)
apply (rule-tac [2] approx-scaleR2)
apply (rule approx-scaleR1)
prefer 2 apply (blast intro: approx-HFfinite approx-sym, auto)
done

```

```

lemma approx-mult-star-of:  $a \approx \text{star-of } b \implies c \approx \text{star-of } d \implies a * c \approx \text{star-of}$ 
 $b * \text{star-of } d$ 
for  $a \ c :: 'a::\text{real-normed-algebra star}$ 
by (blast intro!: approx-mult-HFfinite approx-star-of-HFfinite HFfinite-star-of)

```

```

lemma approx-SReal-mult-cancel-zero:  $a \in \mathbb{R} \implies a \neq 0 \implies a * x \approx 0 \implies x \approx$ 
 $0$ 
for  $a \ x :: \text{hypreal}$ 
apply (drule Reals-inverse [THEN SReal-subset-HFfinite [THEN subsetD]])
apply (auto dest: approx-mult2 simp add: mult.assoc [symmetric])
done

```

```

lemma approx-mult-SReal1:  $a \in \mathbb{R} \implies x \approx 0 \implies x * a \approx 0$ 
for  $a \ x :: \text{hypreal}$ 
by (auto dest: SReal-subset-HFfinite [THEN subsetD] approx-mult1)

```

```

lemma approx-mult-SReal2:  $a \in \mathbb{R} \implies x \approx 0 \implies a * x \approx 0$ 
for  $a \ x :: \text{hypreal}$ 
by (auto dest: SReal-subset-HFfinite [THEN subsetD] approx-mult2)

```

```

lemma approx-mult-SReal-zero-cancel-iff [simp]:  $a \in \mathbb{R} \implies a \neq 0 \implies a * x \approx$ 
 $0 \iff x \approx 0$ 
for  $a \ x :: \text{hypreal}$ 
by (blast intro: approx-SReal-mult-cancel-zero approx-mult-SReal2)

```

```

lemma approx-SReal-mult-cancel:  $a \in \mathbb{R} \implies a \neq 0 \implies a * w \approx a * z \implies w \approx$ 
 $z$ 
for  $a \ w \ z :: \text{hypreal}$ 
apply (drule Reals-inverse [THEN SReal-subset-HFfinite [THEN subsetD]])
apply (auto dest: approx-mult2 simp add: mult.assoc [symmetric])
done

```

```

lemma approx-SReal-mult-cancel-iff1 [simp]:  $a \in \mathbb{R} \implies a \neq 0 \implies a * w \approx a *$ 

```

```

 $z \longleftrightarrow w \approx z$ 
for  $a w z :: \text{hypreal}$ 
by (auto intro!: approx-mult2 SReal-subset-HFinite [THEN subsetD]
     intro: approx-SReal-mult-cancel)

```

```

lemma approx-le-bound:  $z \leq f \implies f \approx g \implies g \leq z \implies f \approx z$ 
for  $z :: \text{hypreal}$ 
apply (simp add: beX-Infinitesimal-iff2 [symmetric], auto)
apply (rule-tac x = g + y - z in beXI)
apply simp
apply (rule Infinitesimal-interval2)
apply (rule-tac [2] Infinitesimal-zero, auto)
done

```

```

lemma approx-hnorm:  $x \approx y \implies \text{hnorm } x \approx \text{hnorm } y$ 
for  $x y :: 'a::\text{real-normed-vector star}$ 
proof (unfold approx-def)
assume  $x - y \in \text{Infinitesimal}$ 
then have  $\text{hnorm } (x - y) \in \text{Infinitesimal}$ 
by (simp only: Infinitesimal-hnorm-iff)
moreover have  $(0::\text{real star}) \in \text{Infinitesimal}$ 
by (rule Infinitesimal-zero)
moreover have  $0 \leq |\text{hnorm } x - \text{hnorm } y|$ 
by (rule abs-ge-zero)
moreover have  $|\text{hnorm } x - \text{hnorm } y| \leq \text{hnorm } (x - y)$ 
by (rule hnorm-triangle-ineq3)
ultimately have  $|\text{hnorm } x - \text{hnorm } y| \in \text{Infinitesimal}$ 
by (rule Infinitesimal-interval2)
then show  $\text{hnorm } x - \text{hnorm } y \in \text{Infinitesimal}$ 
by (simp only: Infinitesimal-hrabs-iff)
qed

```

5.6 Zero is the Only Infinitesimal that is also a Real

```

lemma Infinitesimal-less-SReal:  $x \in \mathbb{R} \implies y \in \text{Infinitesimal} \implies 0 < x \implies y < x$ 
for  $x y :: \text{hypreal}$ 
apply (simp add: Infinitesimal-def)
apply (rule abs-ge-self [THEN order-le-less-trans], auto)
done

```

```

lemma Infinitesimal-less-SReal2:  $y \in \text{Infinitesimal} \implies \forall r \in \text{Reals. } 0 < r \implies y < r$ 
for  $y :: \text{hypreal}$ 
by (blast intro: Infinitesimal-less-SReal)

```

```

lemma SReal-not-Infinitesimal:  $0 < y \implies y \in \mathbb{R} \implies y \notin \text{Infinitesimal}$ 
for  $y :: \text{hypreal}$ 
apply (simp add: Infinitesimal-def)

```

apply (*auto simp add: abs-if*)
done

lemma *SReal-minus-not-Infinitesimal*: $y < 0 \implies y \in \mathbb{R} \implies y \notin \text{Infinitesimal}$
for $y :: \text{hypreal}$
apply (*subst Infinitesimal-minus-iff [symmetric]*)
apply (*rule SReal-not-Infinitesimal, auto*)
done

lemma *SReal-Int-Infinitesimal-zero*: $\mathbb{R} \text{ Int } \text{Infinitesimal} = \{0 :: \text{hypreal}\}$
apply *auto*
apply (*cut-tac x = x and y = 0 in linorder-less-linear*)
apply (*blast dest: SReal-not-Infinitesimal SReal-minus-not-Infinitesimal*)
done

lemma *SReal-Infinitesimal-zero*: $x \in \mathbb{R} \implies x \in \text{Infinitesimal} \implies x = 0$
for $x :: \text{hypreal}$
using *SReal-Int-Infinitesimal-zero* **by** *blast*

lemma *SReal-HFfinite-diff-Infinitesimal*: $x \in \mathbb{R} \implies x \neq 0 \implies x \in \text{HFfinite} - \text{Infinitesimal}$
for $x :: \text{hypreal}$
by (*auto dest: SReal-Infinitesimal-zero SReal-subset-HFfinite [THEN subsetD]*)

lemma *hypreal-of-real-HFfinite-diff-Infinitesimal*:
 $\text{hypreal-of-real } x \neq 0 \implies \text{hypreal-of-real } x \in \text{HFfinite} - \text{Infinitesimal}$
by (*rule SReal-HFfinite-diff-Infinitesimal*) *auto*

lemma *star-of-Infinitesimal-iff-0 [iff]*: $\text{star-of } x \in \text{Infinitesimal} \iff x = 0$
apply (*auto simp add: Infinitesimal-def*)
apply (*drule-tac x=hnorm (star-of x) in bspec*)
apply (*simp add: SReal-def*)
apply (*rule-tac x=norm x in exI, simp*)
apply *simp*
done

lemma *star-of-HFfinite-diff-Infinitesimal*: $x \neq 0 \implies \text{star-of } x \in \text{HFfinite} - \text{Infinitesimal}$
by *simp*

lemma *numeral-not-Infinitesimal [simp]*:
 $\text{numeral } w \neq (0 :: \text{hypreal}) \implies (\text{numeral } w :: \text{hypreal}) \notin \text{Infinitesimal}$
by (*fast dest: Reals-numeral [THEN SReal-Infinitesimal-zero]*)

Again: 1 is a special case, but not 0 this time.

lemma *one-not-Infinitesimal [simp]*:
 $(1 :: 'a :: \{\text{real-normed-vector, zero-neq-one}\} \text{star}) \notin \text{Infinitesimal}$
apply (*simp only: star-one-def star-of-Infinitesimal-iff-0*)
apply *simp*
done

```

lemma approx-SReal-not-zero:  $y \in \mathbb{R} \implies x \approx y \implies y \neq 0 \implies x \neq 0$ 
  for  $x\ y :: \text{hypreal}$ 
  apply (cut-tac  $x = 0$  and  $y = y$  in linorder-less-linear, simp)
  apply (blast dest: approx-sym [THEN mem-infmal-iff [THEN iffD2]]
    SReal-not-Infinitesimal SReal-minus-not-Infinitesimal)
  done

```

```

lemma HFinite-diff-Infinitesimal-approx:
   $x \approx y \implies y \in \text{HFinite} - \text{Infinitesimal} \implies x \in \text{HFinite} - \text{Infinitesimal}$ 
  apply (auto intro: approx-sym [THEN [2] approx-HFinite] simp: mem-infmal-iff)
  apply (drule approx-trans3, assumption)
  apply (blast dest: approx-sym)
  done

```

The premise $y \neq 0$ is essential; otherwise $x / y = 0$ and we lose the *HFinite* premise.

```

lemma Infinitesimal-ratio:
   $y \neq 0 \implies y \in \text{Infinitesimal} \implies x / y \in \text{HFinite} \implies x \in \text{Infinitesimal}$ 
  for  $x\ y :: 'a::\{\text{real-normed-div-algebra,field}\}$  star
  apply (drule Infinitesimal-HFinite-mult2, assumption)
  apply (simp add: divide-inverse mult.assoc)
  done

```

```

lemma Infinitesimal-SReal-divide:  $x \in \text{Infinitesimal} \implies y \in \mathbb{R} \implies x / y \in \text{Infinitesimal}$ 
  for  $x\ y :: \text{hypreal}$ 
  apply (simp add: divide-inverse)
  apply (auto intro!: Infinitesimal-HFinite-mult
    dest!: Reals-inverse [THEN SReal-subset-HFinite [THEN subsetD]])
  done

```

6 Standard Part Theorem

Every finite $x \in R^*$ is infinitely close to a unique real number (i.e. a member of *Reals*).

6.1 Uniqueness: Two Infinitely Close Reals are Equal

```

lemma star-of-approx-iff [simp]: star-of  $x \approx$  star-of  $y \iff x = y$ 
  apply safe
  apply (simp add: approx-def)
  apply (simp only: star-of-diff [symmetric])
  apply (simp only: star-of-Infinitesimal-iff-0)
  apply simp
  done

```

```

lemma SReal-approx-iff:  $x \in \mathbb{R} \implies y \in \mathbb{R} \implies x \approx y \iff x = y$ 

```

```

for  $x y :: \text{hypreal}$ 
apply auto
apply (simp add: approx-def)
apply (drule (1) Reals-diff)
apply (drule (1) SReal-Infinitesimal-zero)
apply simp
done

```

```

lemma numeral-approx-iff [simp]:
  ( $\text{numeral } v \approx (\text{numeral } w :: 'a::\{\text{numeral,real-normed-vector}\} \text{star})) =$ 
   ( $\text{numeral } v = (\text{numeral } w :: 'a)$ )
apply (unfold star-numeral-def)
apply (rule star-of-approx-iff)
done

```

And also for $0 \approx \#nn$ and $1 \approx \#nn$, $\#nn \approx 0$ and $\#nn \approx 1$.

```

lemma [simp]:
  ( $\text{numeral } w \approx (0::'a::\{\text{numeral,real-normed-vector}\} \text{star})) = (\text{numeral } w = (0::'a))$ 
  ( $(0::'a::\{\text{numeral,real-normed-vector}\} \text{star}) \approx \text{numeral } w = (\text{numeral } w = (0::'a))$ 
  ( $\text{numeral } w \approx (1::'b::\{\text{numeral,one,real-normed-vector}\} \text{star})) = (\text{numeral } w =$ 
  ( $1::'b))$ 
  ( $(1::'b::\{\text{numeral,one,real-normed-vector}\} \text{star}) \approx \text{numeral } w = (\text{numeral } w =$ 
  ( $1::'b))$ 
   $\neg (0 \approx (1::'c::\{\text{zero-neq-one,real-normed-vector}\} \text{star}))$ 
   $\neg (1 \approx (0::'c::\{\text{zero-neq-one,real-normed-vector}\} \text{star}))$ 
   apply (unfold star-numeral-def star-zero-def star-one-def)
   apply (unfold star-of-approx-iff)
   apply (auto intro: sym)
done

```

```

lemma star-of-approx-numeral-iff [simp]: star-of k ≈ numeral w ↔ k = numeral
w
by (subst star-of-approx-iff [symmetric]) auto

```

```

lemma star-of-approx-zero-iff [simp]: star-of k ≈ 0 ↔ k = 0
by (simp-all add: star-of-approx-iff [symmetric])

```

```

lemma star-of-approx-one-iff [simp]: star-of k ≈ 1 ↔ k = 1
by (simp-all add: star-of-approx-iff [symmetric])

```

```

lemma approx-unique-real: r ∈ ℝ ⇒ s ∈ ℝ ⇒ r ≈ x ⇒ s ≈ x ⇒ r = s
for  $r s :: \text{hypreal}$ 
by (blast intro: SReal-approx-iff [THEN iffD1] approx-trans2)

```

6.2 Existence of Unique Real Infinitely Close

6.2.1 Lifting of the Ub and Lub Properties

```

lemma hypreal-of-real-isUb-iff: isUb ℝ (hypreal-of-real ‘ Q) (hypreal-of-real Y) =
isUb UNIV Q Y

```

for $Q :: \text{real set}$ **and** $Y :: \text{real}$
by (*simp add: isUb-def settle-def*)

lemma *hypreal-of-real-isLub1*: $\text{isLub } \mathbb{R} (\text{hypreal-of-real } 'Q) (\text{hypreal-of-real } Y)$
 $\implies \text{isLub UNIV } Q Y$

for $Q :: \text{real set}$ **and** $Y :: \text{real}$
apply (*simp add: isLub-def leastP-def*)
apply (*auto intro: hypreal-of-real-isUb-iff [THEN iffD2]*)
simp add: hypreal-of-real-isUb-iff setge-def)
done

lemma *hypreal-of-real-isLub2*: $\text{isLub UNIV } Q Y \implies \text{isLub } \mathbb{R} (\text{hypreal-of-real } 'Q)$
 $(\text{hypreal-of-real } Y)$

for $Q :: \text{real set}$ **and** $Y :: \text{real}$
apply (*auto simp add: isLub-def leastP-def hypreal-of-real-isUb-iff setge-def*)
apply (*metis SReal-iff hypreal-of-real-isUb-iff isUbD2a star-of-le*)
done

lemma *hypreal-of-real-isLub-iff*:

$\text{isLub } \mathbb{R} (\text{hypreal-of-real } 'Q) (\text{hypreal-of-real } Y) = \text{isLub } (\text{UNIV} :: \text{real set}) Q Y$

for $Q :: \text{real set}$ **and** $Y :: \text{real}$
by (*blast intro: hypreal-of-real-isLub1 hypreal-of-real-isLub2*)

lemma *lemma-isUb-hypreal-of-real*: $\text{isUb } \mathbb{R} P Y \implies \exists Yo. \text{isUb } \mathbb{R} P (\text{hypreal-of-real } Yo)$

by (*auto simp add: SReal-iff isUb-def*)

lemma *lemma-isLub-hypreal-of-real*: $\text{isLub } \mathbb{R} P Y \implies \exists Yo. \text{isLub } \mathbb{R} P (\text{hypreal-of-real } Yo)$

by (*auto simp add: isLub-def leastP-def isUb-def SReal-iff*)

lemma *lemma-isLub-hypreal-of-real2*: $\exists Yo. \text{isLub } \mathbb{R} P (\text{hypreal-of-real } Yo) \implies \exists Y. \text{isLub } \mathbb{R} P Y$

by (*auto simp add: isLub-def leastP-def isUb-def*)

lemma *SReal-complete*: $P \subseteq \mathbb{R} \implies \exists x. x \in P \implies \exists Y. \text{isUb } \mathbb{R} P Y \implies \exists t :: \text{hypreal}. \text{isLub } \mathbb{R} P t$

apply (*frule SReal-hypreal-of-real-image*)
apply (*auto, drule lemma-isUb-hypreal-of-real*)
apply (*auto intro!: reals-complete lemma-isLub-hypreal-of-real2*)
simp add: hypreal-of-real-isLub-iff hypreal-of-real-isUb-iff)
done

Lemmas about lubs.

lemma *lemma-st-part-ub*: $x \in \text{HFinite} \implies \exists u. \text{isUb } \mathbb{R} \{s. s \in \mathbb{R} \wedge s < x\} u$

for $x :: \text{hypreal}$
apply (*drule HFiniteD, safe*)
apply (*rule exI, rule isUbI*)

apply (*auto intro: setleI isUbI simp add: abs-less-iff*)
done

lemma *lemma-st-part-nonempty*: $x \in HFinite \implies \exists y. y \in \{s. s \in \mathbb{R} \wedge s < x\}$
for $x :: hypreal$
apply (*drule HFiniteD, safe*)
apply (*drule Reals-minus*)
apply (*rule-tac x = -t in exI*)
apply (*auto simp add: abs-less-iff*)
done

lemma *lemma-st-part-lub*: $x \in HFinite \implies \exists t. isLub \mathbb{R} \{s. s \in \mathbb{R} \wedge s < x\} t$
for $x :: hypreal$
by (*blast intro!: SReal-complete lemma-st-part-ub lemma-st-part-nonempty Collect-restrict*)

lemma *lemma-st-part-le1*:
 $x \in HFinite \implies isLub \mathbb{R} \{s. s \in \mathbb{R} \wedge s < x\} t \implies r \in \mathbb{R} \implies 0 < r \implies x \leq t + r$
for $x r t :: hypreal$
apply (*frule isLubD1a*)
apply (*rule ccontr, drule linorder-not-le [THEN iffD2]*)
apply (*drule (1) Reals-add*)
apply (*drule-tac y = r + t in isLubD1 [THEN setleD], auto*)
done

lemma *hypreal-setle-less-trans*: $S * \leq x \implies x < y \implies S * \leq y$
for $x y :: hypreal$
apply (*simp add: setle-def*)
apply (*auto dest!: bspec order-le-less-trans intro: order-less-imp-le*)
done

lemma *hypreal-gt-isUb*: $isUb R S x \implies x < y \implies y \in R \implies isUb R S y$
for $x y :: hypreal$
apply (*simp add: isUb-def*)
apply (*blast intro: hypreal-setle-less-trans*)
done

lemma *lemma-st-part-gt-ub*: $x \in HFinite \implies x < y \implies y \in \mathbb{R} \implies isUb \mathbb{R} \{s. s \in \mathbb{R} \wedge s < x\} y$
for $x y :: hypreal$
by (*auto dest: order-less-trans intro: order-less-imp-le intro!: isUbI setleI*)

lemma *lemma-minus-le-zero*: $t \leq t + -r \implies r \leq 0$
for $r t :: hypreal$
apply (*drule-tac c = -t in add-left-mono*)
apply (*auto simp add: add.assoc [symmetric]*)
done

lemma *lemma-st-part-le2*:

```

 $x \in HFinite \implies isLub \mathbb{R} \{s. s \in \mathbb{R} \wedge s < x\} t \implies r \in \mathbb{R} \implies 0 < r \implies t + -r \leq x$ 
for  $x r t :: hypreal$ 
apply (frule isLubD1a)
apply (rule ccontr, drule linorder-not-le [THEN iffD1])
apply (drule Reals-minus, drule-tac  $a = t$  in Reals-add, assumption)
apply (drule lemma-st-part-gt-ub, assumption+)
apply (drule isLub-le-isUb, assumption)
apply (drule lemma-minus-le-zero)
apply (auto dest: order-less-le-trans)
done

```

lemma lemma-st-part1a:

```

 $x \in HFinite \implies isLub \mathbb{R} \{s. s \in \mathbb{R} \wedge s < x\} t \implies r \in \mathbb{R} \implies 0 < r \implies x + -t \leq r$ 
for  $x r t :: hypreal$ 
apply (subgoal-tac  $x \leq t + r$ )
apply (auto intro: lemma-st-part-le1)
done

```

lemma lemma-st-part2a:

```

 $x \in HFinite \implies isLub \mathbb{R} \{s. s \in \mathbb{R} \wedge s < x\} t \implies r \in \mathbb{R} \implies 0 < r \implies -(x + -t) \leq r$ 
for  $x r t :: hypreal$ 
apply (subgoal-tac ( $t + -r \leq x$ ))
apply simp
apply (rule lemma-st-part-le2)
apply auto
done

```

lemma lemma-SReal-ub: $x \in \mathbb{R} \implies isUb \mathbb{R} \{s. s \in \mathbb{R} \wedge s < x\} x$

```

for  $x :: hypreal$ 
by (auto intro: isUbI setleI order-less-imp-le)

```

lemma lemma-SReal-lub: $x \in \mathbb{R} \implies isLub \mathbb{R} \{s. s \in \mathbb{R} \wedge s < x\} x$

```

for  $x :: hypreal$ 
apply (auto intro!: isLubI2 lemma-SReal-ub setgeI)
apply (frule isUbD2a)
apply (rule-tac  $x = x$  and  $y = y$  in linorder-cases)
apply (auto intro!: order-less-imp-le)
apply (drule SReal-dense, assumption, safe)
apply (drule-tac  $y = r$  in isUbD)
apply (auto dest: order-less-le-trans)
done

```

lemma lemma-st-part-not-eq1:

```

 $x \in HFinite \implies isLub \mathbb{R} \{s. s \in \mathbb{R} \wedge s < x\} t \implies r \in \mathbb{R} \implies 0 < r \implies x + -t \neq r$ 
for  $x r t :: hypreal$ 

```



```

apply auto
apply (frule isLubD1a [THEN Reals-minus])
using Reals-add-cancel [of x - t] apply simp
apply (drule-tac x = x in lemma-SReal-lub)
apply (drule isLub-unique, assumption, auto)
done

```

lemma *lemma-st-part-not-eq2*:

```

 $x \in HFinite \implies isLub \mathbb{R} \{s. s \in \mathbb{R} \wedge s < x\} t \implies r \in \mathbb{R} \implies 0 < r \implies -(x + -t) \neq r$ 
for  $x r t :: hypreal$ 
apply (auto)
apply (frule isLubD1a)
using Reals-add-cancel [of - x t] apply simp
apply (drule-tac x = x in lemma-SReal-lub)
apply (drule isLub-unique, assumption, auto)
done

```

lemma *lemma-st-part-major*:

```

 $x \in HFinite \implies isLub \mathbb{R} \{s. s \in \mathbb{R} \wedge s < x\} t \implies r \in \mathbb{R} \implies 0 < r \implies |x - t| < r$ 
for  $x r t :: hypreal$ 
apply (frule lemma-st-part1a)
apply (frule-tac [4] lemma-st-part2a, auto)
apply (drule order-le-imp-less-or-eq)+
apply auto
using lemma-st-part-not-eq2 apply fastforce
using lemma-st-part-not-eq1 apply fastforce
done

```

lemma *lemma-st-part-major2*:

```

 $x \in HFinite \implies isLub \mathbb{R} \{s. s \in \mathbb{R} \wedge s < x\} t \implies \forall r \in Reals. 0 < r \longrightarrow |x - t| < r$ 
for  $x t :: hypreal$ 
by (blast dest!: lemma-st-part-major)

```

Existence of real and Standard Part Theorem.

```

lemma lemma-st-part-Ex:  $x \in HFinite \implies \exists t \in Reals. \forall r \in Reals. 0 < r \longrightarrow |x - t| < r$ 
for  $x :: hypreal$ 
apply (frule lemma-st-part-lub, safe)
apply (frule isLubD1a)
apply (blast dest: lemma-st-part-major2)
done

```

lemma *st-part-Ex*: $x \in HFinite \implies \exists t \in Reals. x \approx t$

```

for  $x :: hypreal$ 
apply (simp add: approx-def Infinitesimal-def)
apply (drule lemma-st-part-Ex, auto)

```

done

There is a unique real infinitely close.

lemma *st-part-Ex1*: $x \in \mathit{HFinite} \implies \exists ! t :: \mathit{hypreal}. t \in \mathbb{R} \wedge x \approx t$
apply (*drule st-part-Ex, safe*)
apply (*drule-tac [2] approx-sym, drule-tac [2] approx-sym, drule-tac [2] approx-sym*)
apply (*auto intro!: approx-unique-real*)
done

6.3 Finite, Infinite and Infinitesimal

lemma *HFinite-Int-HInfinite-empty* [*simp*]: $\mathit{HFinite} \ \mathit{Int} \ \mathit{HInfinite} = \{\}$
apply (*simp add: HFinite-def HInfinite-def*)
apply (*auto dest: order-less-trans*)
done

lemma *HFinite-not-HInfinite*:
assumes $x: x \in \mathit{HFinite}$
shows $x \notin \mathit{HInfinite}$
proof
assume $x': x \in \mathit{HInfinite}$
with x **have** $x \in \mathit{HFinite} \cap \mathit{HInfinite}$ **by** *blast*
then show *False* **by** *auto*
qed

lemma *not-HFinite-HInfinite*: $x \notin \mathit{HFinite} \implies x \in \mathit{HInfinite}$
apply (*simp add: HInfinite-def HFinite-def, auto*)
apply (*drule-tac x = r + 1 in bspec*)
apply (*auto*)
done

lemma *HInfinite-HFinite-disj*: $x \in \mathit{HInfinite} \vee x \in \mathit{HFinite}$
by (*blast intro: not-HFinite-HInfinite*)

lemma *HInfinite-HFinite-iff*: $x \in \mathit{HInfinite} \longleftrightarrow x \notin \mathit{HFinite}$
by (*blast dest: HFinite-not-HInfinite not-HFinite-HInfinite*)

lemma *HFinite-HInfinite-iff*: $x \in \mathit{HFinite} \longleftrightarrow x \notin \mathit{HInfinite}$
by (*simp add: HInfinite-HFinite-iff*)

lemma *HInfinite-diff-HFinite-Infinitesimal-disj*:
 $x \notin \mathit{Infinitesimal} \implies x \in \mathit{HInfinite} \vee x \in \mathit{HFinite} - \mathit{Infinitesimal}$
by (*fast intro: not-HFinite-HInfinite*)

lemma *HFinite-inverse*: $x \in \mathit{HFinite} \implies x \notin \mathit{Infinitesimal} \implies \mathit{inverse} \ x \in \mathit{HFinite}$
for $x :: 'a :: \mathit{real-normed-div-algebra} \ \mathit{star}$
apply (*subgoal-tac x \neq 0*)

```

apply (cut-tac x = inverse x in HInfinite-HFinite-disj)
apply (auto dest!: HInfinite-inverse-Infinitesimal simp: nonzero-inverse-inverse-eq)
done

```

```

lemma HFinite-inverse2: x ∈ HFinite – Infinitesimal ⇒ inverse x ∈ HFinite
for x :: 'a::real-normed-div-algebra star
by (blast intro: HFinite-inverse)

```

Stronger statement possible in fact.

```

lemma Infinitesimal-inverse-HFinite: x ∉ Infinitesimal ⇒ inverse x ∈ HFinite
for x :: 'a::real-normed-div-algebra star
apply (drule HInfinite-diff-HFinite-Infinitesimal-disj)
apply (blast intro: HFinite-inverse HInfinite-inverse-Infinitesimal Infinitesimal-subset-HFinite
[THEN subsetD])
done

```

```

lemma HFinite-not-Infinitesimal-inverse:
x ∈ HFinite – Infinitesimal ⇒ inverse x ∈ HFinite – Infinitesimal
for x :: 'a::real-normed-div-algebra star
apply (auto intro: Infinitesimal-inverse-HFinite)
apply (drule Infinitesimal-HFinite-mult2, assumption)
apply (simp add: not-Infinitesimal-not-zero)
done

```

```

lemma approx-inverse: x ≈ y ⇒ y ∈ HFinite – Infinitesimal ⇒ inverse x ≈
inverse y
for x y :: 'a::real-normed-div-algebra star
apply (frule HFinite-diff-Infinitesimal-approx, assumption)
apply (frule not-Infinitesimal-not-zero2)
apply (frule-tac x = x in not-Infinitesimal-not-zero2)
apply (drule HFinite-inverse2)+
apply (drule approx-mult2, assumption, auto)
apply (drule-tac c = inverse x in approx-mult1, assumption)
apply (auto intro: approx-sym simp add: mult.assoc)
done

```

```

lemmas star-of-approx-inverse = star-of-HFinite-diff-Infinitesimal [THEN [2] approx-inverse]
lemmas hypreal-of-real-approx-inverse = hypreal-of-real-HFinite-diff-Infinitesimal
[THEN [2] approx-inverse]

```

```

lemma inverse-add-Infinitesimal-approx:
x ∈ HFinite – Infinitesimal ⇒ h ∈ Infinitesimal ⇒ inverse (x + h) ≈ inverse
x
for x h :: 'a::real-normed-div-algebra star
by (auto intro: approx-inverse approx-sym Infinitesimal-add-approx-self)

```

```

lemma inverse-add-Infinitesimal-approx2:
x ∈ HFinite – Infinitesimal ⇒ h ∈ Infinitesimal ⇒ inverse (h + x) ≈ inverse

```

x

```

for  $x\ h :: 'a::\text{real-normed-div-algebra star}$ 
apply (rule add.commute [THEN subst])
apply (blast intro: inverse-add-Infinitesimal-approx)
done

```

lemma *inverse-add-Infinitesimal-approx-Infinitesimal*:
 $x \in \text{HFinite} - \text{Infinitesimal} \implies h \in \text{Infinitesimal} \implies \text{inverse}(x + h) - \text{inverse}$
 $x \approx h$

```

for  $x\ h :: 'a::\text{real-normed-div-algebra star}$ 
apply (rule approx-trans2)
apply (auto intro: inverse-add-Infinitesimal-approx
  simp add: mem-infmal-iff approx-minus-iff [symmetric])
done

```

lemma *Infinitesimal-square-iff*: $x \in \text{Infinitesimal} \longleftrightarrow x * x \in \text{Infinitesimal}$

```

for  $x :: 'a::\text{real-normed-div-algebra star}$ 
apply (auto intro: Infinitesimal-mult)
apply (rule ccontr, frule Infinitesimal-inverse-HFinite)
apply (frule not-Infinitesimal-not-zero)
apply (auto dest: Infinitesimal-HFinite-mult simp add: mult.assoc)
done

```

declare *Infinitesimal-square-iff* [symmetric, simp]

lemma *HFinite-square-iff* [simp]: $x * x \in \text{HFinite} \longleftrightarrow x \in \text{HFinite}$

```

for  $x :: 'a::\text{real-normed-div-algebra star}$ 
apply (auto intro: HFinite-mult)
apply (auto dest: HInfinite-mult simp add: HFinite-HInfinite-iff)
done

```

lemma *HInfinite-square-iff* [simp]: $x * x \in \text{HInfinite} \longleftrightarrow x \in \text{HInfinite}$

```

for  $x :: 'a::\text{real-normed-div-algebra star}$ 
by (auto simp add: HInfinite-HFinite-iff)

```

lemma *approx-HFinite-mult-cancel*: $a \in \text{HFinite} - \text{Infinitesimal} \implies a * w \approx a * z \implies w \approx z$

```

for  $a\ w\ z :: 'a::\text{real-normed-div-algebra star}$ 
apply safe
apply (frule HFinite-inverse, assumption)
apply (drule not-Infinitesimal-not-zero)
apply (auto dest: approx-mult2 simp add: mult.assoc [symmetric])
done

```

lemma *approx-HFinite-mult-cancel-iff1*: $a \in \text{HFinite} - \text{Infinitesimal} \implies a * w \approx a * z \longleftrightarrow w \approx z$

```

for  $a\ w\ z :: 'a::\text{real-normed-div-algebra star}$ 
by (auto intro: approx-mult2 approx-HFinite-mult-cancel)

```

lemma *HInfinite-HFinite-add-cancel*: $x + y \in \text{HInfinite} \implies y \in \text{HFinite} \implies x$

```

∈ HInfinite
  apply (rule ccontr)
  apply (drule HFinite-HInfinite-iff [THEN iffD2])
  apply (auto dest: HFinite-add simp add: HInfinite-HFinite-iff)
  done

```

```

lemma HInfinite-HFinite-add:  $x \in HInfinite \implies y \in HFinite \implies x + y \in HInfinite$ 
  apply (rule-tac  $y = -y$  in HInfinite-HFinite-add-cancel)
  apply (auto simp add: add.assoc HFinite-minus-iff)
  done

```

```

lemma HInfinite-ge-HInfinite:  $x \in HInfinite \implies x \leq y \implies 0 \leq x \implies y \in HInfinite$ 
  for  $x y :: \text{hypreal}$ 
  by (auto intro: HFinite-bounded simp add: HInfinite-HFinite-iff)

```

```

lemma Infinitesimal-inverse-HInfinite:  $x \in Infinitesimal \implies x \neq 0 \implies \text{inverse } x \in HInfinite$ 
  for  $x :: 'a::\text{real-normed-div-algebra star}$ 
  apply (rule ccontr, drule HFinite-HInfinite-iff [THEN iffD2])
  apply (auto dest: Infinitesimal-HFinite-mult2)
  done

```

```

lemma HInfinite-HFinite-not-Infinitesimal-mult:
 $x \in HInfinite \implies y \in HFinite - Infinitesimal \implies x * y \in HInfinite$ 
  for  $x y :: 'a::\text{real-normed-div-algebra star}$ 
  apply (rule ccontr, drule HFinite-HInfinite-iff [THEN iffD2])
  apply (frule HFinite-Infinitesimal-not-zero)
  apply (drule HFinite-not-Infinitesimal-inverse)
  apply (safe, drule HFinite-mult)
  apply (auto simp add: mult.assoc HFinite-HInfinite-iff)
  done

```

```

lemma HInfinite-HFinite-not-Infinitesimal-mult2:
 $x \in HInfinite \implies y \in HFinite - Infinitesimal \implies y * x \in HInfinite$ 
  for  $x y :: 'a::\text{real-normed-div-algebra star}$ 
  apply (rule ccontr, drule HFinite-HInfinite-iff [THEN iffD2])
  apply (frule HFinite-Infinitesimal-not-zero)
  apply (drule HFinite-not-Infinitesimal-inverse)
  apply (safe, drule-tac  $x = \text{inverse } y$  in HFinite-mult)
  apply assumption
  apply (auto simp add: mult.assoc [symmetric] HFinite-HInfinite-iff)
  done

```

```

lemma HInfinite-gt-SReal:  $x \in HInfinite \implies 0 < x \implies y \in \mathbb{R} \implies y < x$ 
  for  $x y :: \text{hypreal}$ 
  by (auto dest!: bspec simp add: HInfinite-def abs-if order-less-imp-le)

```

lemma *HInfinite-gt-zero-gt-one*: $x \in \text{HInfinite} \implies 0 < x \implies 1 < x$
for $x :: \text{hypreal}$
by (*auto intro: HInfinite-gt-SReal*)

lemma *not-HInfinite-one* [*simp*]: $1 \notin \text{HInfinite}$
by (*simp add: HInfinite-HFinite-iff*)

lemma *approx-hrabs-disj*: $|x| \approx x \vee |x| \approx -x$
for $x :: \text{hypreal}$
using *hrabs-disj [of x] by auto*

6.4 Theorems about Monads

lemma *monad-hrabs-Un-subset*: $\text{monad } |x| \leq \text{monad } x \cup \text{monad } (-x)$
for $x :: \text{hypreal}$
by (*rule hrabs-disj [of x, THEN disjE] auto*)

lemma *Infinitesimal-monad-eq*: $e \in \text{Infinitesimal} \implies \text{monad } (x + e) = \text{monad } x$
by (*fast intro!: Infinitesimal-add-approx-self [THEN approx-sym] approx-monad-iff [THEN iffD1]*)

lemma *mem-monad-iff*: $u \in \text{monad } x \longleftrightarrow -u \in \text{monad } (-x)$
by (*simp add: monad-def*)

lemma *Infinitesimal-monad-zero-iff*: $x \in \text{Infinitesimal} \longleftrightarrow x \in \text{monad } 0$
by (*auto intro: approx-sym simp add: monad-def mem-infmal-iff*)

lemma *monad-zero-minus-iff*: $x \in \text{monad } 0 \longleftrightarrow -x \in \text{monad } 0$
by (*simp add: Infinitesimal-monad-zero-iff [symmetric]*)

lemma *monad-zero-hrabs-iff*: $x \in \text{monad } 0 \longleftrightarrow |x| \in \text{monad } 0$
for $x :: \text{hypreal}$
by (*rule hrabs-disj [of x, THEN disjE] (auto simp: monad-zero-minus-iff [symmetric])*)

lemma *mem-monad-self* [*simp*]: $x \in \text{monad } x$
by (*simp add: monad-def*)

6.5 Proof that $x \approx y$ implies $|x| \approx |y|$

lemma *approx-subset-monad*: $x \approx y \implies \{x, y\} \leq \text{monad } x$
by (*simp (no-asm) (simp add: approx-monad-iff)*)

lemma *approx-subset-monad2*: $x \approx y \implies \{x, y\} \leq \text{monad } y$
apply (*drule approx-sym*)
apply (*fast dest: approx-subset-monad*)
done

lemma *mem-monad-approx*: $u \in \text{monad } x \implies x \approx u$
by (*simp add: monad-def*)

```

lemma approx-mem-monad:  $x \approx u \implies u \in \text{monad } x$ 
  by (simp add: monad-def)

lemma approx-mem-monad2:  $x \approx u \implies x \in \text{monad } u$ 
  apply (simp add: monad-def)
  apply (blast intro!: approx-sym)
  done

lemma approx-mem-monad-zero:  $x \approx y \implies x \in \text{monad } 0 \implies y \in \text{monad } 0$ 
  apply (drule mem-monad-approx)
  apply (fast intro: approx-mem-monad approx-trans)
  done

lemma Infinitesimal-approx-hrabs:  $x \approx y \implies x \in \text{Infinitesimal} \implies |x| \approx |y|$ 
  for  $x y :: \text{hypreal}$ 
  apply (drule Infinitesimal-monad-zero-iff [THEN iffD1])
  apply (blast intro: approx-mem-monad-zero monad-zero-hrabs-iff [THEN iffD1]
    mem-monad-approx approx-trans3)
  done

lemma less-Infinitesimal-less:  $0 < x \implies x \notin \text{Infinitesimal} \implies e \in \text{Infinitesimal}$ 
 $\implies e < x$ 
  for  $x :: \text{hypreal}$ 
  apply (rule ccontr)
  apply (auto intro: Infinitesimal-zero [THEN [2] Infinitesimal-interval]
    dest!: order-le-imp-less-or-eq simp add: linorder-not-less)
  done

lemma Ball-mem-monad-gt-zero:  $0 < x \implies x \notin \text{Infinitesimal} \implies u \in \text{monad } x$ 
 $\implies 0 < u$ 
  for  $u x :: \text{hypreal}$ 
  apply (drule mem-monad-approx [THEN approx-sym])
  apply (erule bex-Infinitesimal-iff2 [THEN iffD2, THEN bexE])
  apply (drule-tac  $e = -x$  in less-Infinitesimal-less, auto)
  done

lemma Ball-mem-monad-less-zero:  $x < 0 \implies x \notin \text{Infinitesimal} \implies u \in \text{monad}$ 
 $x \implies u < 0$ 
  for  $u x :: \text{hypreal}$ 
  apply (drule mem-monad-approx [THEN approx-sym])
  apply (erule bex-Infinitesimal-iff [THEN iffD2, THEN bexE])
  apply (cut-tac  $x = -x$  and  $e = xa$  in less-Infinitesimal-less, auto)
  done

lemma lemma-approx-gt-zero:  $0 < x \implies x \notin \text{Infinitesimal} \implies x \approx y \implies 0 < y$ 
  for  $x y :: \text{hypreal}$ 
  by (blast dest: Ball-mem-monad-gt-zero approx-subset-monad)

```

lemma *lemma-approx-less-zero*: $x < 0 \implies x \notin \text{Infinitesimal} \implies x \approx y \implies y < 0$

for $x\ y :: \text{hypreal}$
by (*blast dest: Ball-mem-monad-less-zero approx-subset-monad*)

lemma *approx-hrabs*: $x \approx y \implies |x| \approx |y|$

for $x\ y :: \text{hypreal}$
by (*drule approx-hnorm simp*)

lemma *approx-hrabs-zero-cancel*: $|x| \approx 0 \implies x \approx 0$

for $x :: \text{hypreal}$
using *hrabs-disj [of x]* **by** (*auto dest: approx-minus*)

lemma *approx-hrabs-add-Infinitesimal*: $e \in \text{Infinitesimal} \implies |x| \approx |x + e|$

for $e\ x :: \text{hypreal}$
by (*fast intro: approx-hrabs Infinitesimal-add-approx-self*)

lemma *approx-hrabs-add-minus-Infinitesimal*: $e \in \text{Infinitesimal} \implies |x| \approx |x - e|$

for $e\ x :: \text{hypreal}$
by (*fast intro: approx-hrabs Infinitesimal-add-minus-approx-self*)

lemma *hrabs-add-Infinitesimal-cancel*:

$e \in \text{Infinitesimal} \implies e' \in \text{Infinitesimal} \implies |x + e| = |y + e'| \implies |x| \approx |y|$
for $e\ e'\ x\ y :: \text{hypreal}$
apply (*drule-tac x = x in approx-hrabs-add-Infinitesimal*)
apply (*drule-tac x = y in approx-hrabs-add-Infinitesimal*)
apply (*auto intro: approx-trans2*)
done

lemma *hrabs-add-minus-Infinitesimal-cancel*:

$e \in \text{Infinitesimal} \implies e' \in \text{Infinitesimal} \implies |x - e| = |y - e'| \implies |x| \approx |y|$
for $e\ e'\ x\ y :: \text{hypreal}$
apply (*drule-tac x = x in approx-hrabs-add-minus-Infinitesimal*)
apply (*drule-tac x = y in approx-hrabs-add-minus-Infinitesimal*)
apply (*auto intro: approx-trans2*)
done

6.6 More HFinite and Infinitesimal Theorems

Interesting slightly counterintuitive theorem: necessary for proving that an open interval is an NS open set.

lemma *Infinitesimal-add-hypreal-of-real-less*:

$x < y \implies u \in \text{Infinitesimal} \implies \text{hypreal-of-real } x + u < \text{hypreal-of-real } y$
apply (*simp add: Infinitesimal-def*)
apply (*drule-tac x = hypreal-of-real y + -hypreal-of-real x in bspec, simp*)
apply (*simp add: abs-less-iff*)
done

lemma *Infininitesimal-add-hrabs-hypreal-of-real-less:*

$x \in \text{Infininitesimal} \implies |\text{hypreal-of-real } r| < \text{hypreal-of-real } y \implies$
 $|\text{hypreal-of-real } r + x| < \text{hypreal-of-real } y$
apply (drule-tac $x = \text{hypreal-of-real } r$ **in** approx-hrabs-add-Infininitesimal)
apply (drule approx-sym [THEN bex-Infininitesimal-iff2 [THEN iffD2]])
apply (auto intro!: *Infininitesimal-add-hypreal-of-real-less*
simp del: star-of-abs simp add: star-of-abs [symmetric])
done

lemma *Infininitesimal-add-hrabs-hypreal-of-real-less2:*

$x \in \text{Infininitesimal} \implies |\text{hypreal-of-real } r| < \text{hypreal-of-real } y \implies$
 $|x + \text{hypreal-of-real } r| < \text{hypreal-of-real } y$
apply (rule add.commute [THEN subst])
apply (erule *Infininitesimal-add-hrabs-hypreal-of-real-less*, assumption)
done

lemma *hypreal-of-real-le-add-Infininitesimal-cancel:*

$u \in \text{Infininitesimal} \implies v \in \text{Infininitesimal} \implies$
 $\text{hypreal-of-real } x + u \leq \text{hypreal-of-real } y + v \implies$
 $\text{hypreal-of-real } x \leq \text{hypreal-of-real } y$
apply (simp add: linorder-not-less [symmetric], auto)
apply (drule-tac $u = v - u$ **in** *Infininitesimal-add-hypreal-of-real-less*)
apply (auto simp add: *Infininitesimal-diff*)
done

lemma *hypreal-of-real-le-add-Infininitesimal-cancel2:*

$u \in \text{Infininitesimal} \implies v \in \text{Infininitesimal} \implies$
 $\text{hypreal-of-real } x + u \leq \text{hypreal-of-real } y + v \implies x \leq y$
by (blast intro: star-of-le [THEN iffD1] intro!: *hypreal-of-real-le-add-Infininitesimal-cancel*)

lemma *hypreal-of-real-less-Infininitesimal-le-zero:*

$\text{hypreal-of-real } x < e \implies e \in \text{Infininitesimal} \implies \text{hypreal-of-real } x \leq 0$
apply (rule linorder-not-less [THEN iffD1], safe)
apply (drule *Infininitesimal-interval*)
apply (drule-tac [4] *SReal-hypreal-of-real* [THEN *SReal-Infininitesimal-zero*],
auto)
done

lemma *Infininitesimal-add-not-zero:* $h \in \text{Infininitesimal} \implies x \neq 0 \implies \text{star-of } x + h \neq 0$

apply auto
apply (subgoal-tac $h = - \text{star-of } x$)
apply (auto intro: minus-unique [symmetric])
done

lemma *Infininitesimal-square-cancel* [simp]: $x * x + y * y \in \text{Infininitesimal} \implies x * x \in \text{Infininitesimal}$

for $x y :: \text{hypreal}$

```

apply (rule Infinitesimal-interval2)
  apply (rule-tac [3] zero-le-square, assumption)
  apply auto
done

```

```

lemma HFinite-square-cancel [simp]:  $x * x + y * y \in HFinite \implies x * x \in HFinite$ 
  for  $x\ y :: \text{hypreal}$ 
  apply (rule HFinite-bounded, assumption)
  apply auto
done

```

```

lemma Infinitesimal-square-cancel2 [simp]:  $x * x + y * y \in Infinitesimal \implies y * y \in Infinitesimal$ 
  for  $x\ y :: \text{hypreal}$ 
  apply (rule Infinitesimal-square-cancel)
  apply (rule add.commute [THEN subst])
  apply simp
done

```

```

lemma HFinite-square-cancel2 [simp]:  $x * x + y * y \in HFinite \implies y * y \in HFinite$ 
  for  $x\ y :: \text{hypreal}$ 
  apply (rule HFinite-square-cancel)
  apply (rule add.commute [THEN subst])
  apply simp
done

```

```

lemma Infinitesimal-sum-square-cancel [simp]:
   $x * x + y * y + z * z \in Infinitesimal \implies x * x \in Infinitesimal$ 
  for  $x\ y\ z :: \text{hypreal}$ 
  apply (rule Infinitesimal-interval2, assumption)
  apply (rule-tac [2] zero-le-square, simp)
  apply (insert zero-le-square [of y])
  apply (insert zero-le-square [of z], simp del:zero-le-square)
done

```

```

lemma HFinite-sum-square-cancel [simp]:  $x * x + y * y + z * z \in HFinite \implies x * x \in HFinite$ 
  for  $x\ y\ z :: \text{hypreal}$ 
  apply (rule HFinite-bounded, assumption)
  apply (rule-tac [2] zero-le-square)
  apply (insert zero-le-square [of y])
  apply (insert zero-le-square [of z], simp del:zero-le-square)
done

```

```

lemma Infinitesimal-sum-square-cancel2 [simp]:
   $y * y + x * x + z * z \in Infinitesimal \implies x * x \in Infinitesimal$ 
  for  $x\ y\ z :: \text{hypreal}$ 
  apply (rule Infinitesimal-sum-square-cancel)

```

```

apply (simp add: ac-simps)
done

```

```

lemma HFfinite-sum-square-cancel2 [simp]:  $y * y + x * x + z * z \in HFfinite \implies$ 
 $x * x \in HFfinite$ 
for  $x\ y\ z :: hypreal$ 
apply (rule HFfinite-sum-square-cancel)
apply (simp add: ac-simps)
done

```

```

lemma Infinitesimal-sum-square-cancel3 [simp]:
 $z * z + y * y + x * x \in Infinitesimal \implies x * x \in Infinitesimal$ 
for  $x\ y\ z :: hypreal$ 
apply (rule Infinitesimal-sum-square-cancel)
apply (simp add: ac-simps)
done

```

```

lemma HFfinite-sum-square-cancel3 [simp]:  $z * z + y * y + x * x \in HFfinite \implies$ 
 $x * x \in HFfinite$ 
for  $x\ y\ z :: hypreal$ 
apply (rule HFfinite-sum-square-cancel)
apply (simp add: ac-simps)
done

```

```

lemma monad-hrabs-less:  $y \in monad\ x \implies 0 < hypreal-of-real\ e \implies |y - x| <$ 
 $hypreal-of-real\ e$ 
apply (drule mem-monad-approx [THEN approx-sym])
apply (drule beX-Infinitesimal-iff [THEN iffD2])
apply (auto dest!: InfinitesimalD)
done

```

```

lemma mem-monad-SReal-HFfinite:  $x \in monad\ (hypreal-of-real\ a) \implies x \in HFfinite$ 
apply (drule mem-monad-approx [THEN approx-sym])
apply (drule beX-Infinitesimal-iff2 [THEN iffD2])
apply (safe dest!: Infinitesimal-subset-HFfinite [THEN subsetD])
apply (erule SReal-hypreal-of-real [THEN SReal-subset-HFfinite [THEN subsetD],
THEN HFfinite-add])
done

```

6.7 Theorems about Standard Part

```

lemma st-approx-self:  $x \in HFfinite \implies st\ x \approx x$ 
apply (simp add: st-def)
apply (frule st-part-Ex, safe)
apply (rule someI2)
apply (auto intro: approx-sym)
done

```

```

lemma st-SReal:  $x \in HFfinite \implies st\ x \in \mathbb{R}$ 

```

```

apply (simp add: st-def)
apply (frule st-part-Ex, safe)
apply (rule someI2)
apply (auto intro: approx-sym)
done

```

```

lemma st-HFinite:  $x \in HFinite \implies st\ x \in HFinite$ 
by (erule st-SReal [THEN SReal-subset-HFinite [THEN subsetD]])

```

```

lemma st-unique:  $r \in \mathbb{R} \implies r \approx x \implies st\ x = r$ 
apply (frule SReal-subset-HFinite [THEN subsetD])
apply (drule (1) approx-HFinite)
apply (unfold st-def)
apply (rule some-equality)
apply (auto intro: approx-unique-real)
done

```

```

lemma st-SReal-eq:  $x \in \mathbb{R} \implies st\ x = x$ 
by (metis approx-refl st-unique)

```

```

lemma st-hypreal-of-real [simp]:  $st\ (hypreal-of-real\ x) = hypreal-of-real\ x$ 
by (rule SReal-hypreal-of-real [THEN st-SReal-eq])

```

```

lemma st-eq-approx:  $x \in HFinite \implies y \in HFinite \implies st\ x = st\ y \implies x \approx y$ 
by (auto dest!: st-approx-self elim!: approx-trans3)

```

```

lemma approx-st-eq:
  assumes  $x: x \in HFinite$  and  $y: y \in HFinite$  and  $xy: x \approx y$ 
  shows  $st\ x = st\ y$ 
proof –
  have  $st\ x \approx x$   $st\ y \approx y$   $st\ x \in \mathbb{R}$   $st\ y \in \mathbb{R}$ 
  by (simp-all add: st-approx-self st-SReal x y)
  with  $xy$  show ?thesis
  by (fast elim: approx-trans approx-trans2 SReal-approx-iff [THEN iffD1])
qed

```

```

lemma st-eq-approx-iff:  $x \in HFinite \implies y \in HFinite \implies x \approx y \iff st\ x = st\ y$ 
by (blast intro: approx-st-eq st-eq-approx)

```

```

lemma st-Infinitesimal-add-SReal:  $x \in \mathbb{R} \implies e \in Infinitesimal \implies st\ (x + e) = x$ 
apply (erule st-unique)
apply (erule Infinitesimal-add-approx-self)
done

```

```

lemma st-Infinitesimal-add-SReal2:  $x \in \mathbb{R} \implies e \in Infinitesimal \implies st\ (e + x) = x$ 
apply (erule st-unique)
apply (erule Infinitesimal-add-approx-self2)

```

done

lemma *HFinite-st-Infinitesimal-add*: $x \in \text{HFinite} \implies \exists e \in \text{Infinitesimal}. x = \text{st}(x) + e$
by (*blast dest!*: *st-approx-self* [*THEN approx-sym*] *bx-Infinitesimal-iff2* [*THEN iffD2*])

lemma *st-add*: $x \in \text{HFinite} \implies y \in \text{HFinite} \implies \text{st}(x + y) = \text{st } x + \text{st } y$
by (*simp add*: *st-unique st-SReal st-approx-self approx-add*)

lemma *st-numeral* [*simp*]: $\text{st}(\text{numeral } w) = \text{numeral } w$
by (*rule Reals-numeral* [*THEN st-SReal-eq*])

lemma *st-neg-numeral* [*simp*]: $\text{st}(- \text{numeral } w) = - \text{numeral } w$

proof –

from *Reals-numeral* **have** $\text{numeral } w \in \mathbb{R}$.

then have $- \text{numeral } w \in \mathbb{R}$ **by** *simp*

with *st-SReal-eq* **show** *?thesis* .

qed

lemma *st-0* [*simp*]: $\text{st } 0 = 0$
by (*simp add*: *st-SReal-eq*)

lemma *st-1* [*simp*]: $\text{st } 1 = 1$
by (*simp add*: *st-SReal-eq*)

lemma *st-neg-1* [*simp*]: $\text{st}(- 1) = - 1$
by (*simp add*: *st-SReal-eq*)

lemma *st-minus*: $x \in \text{HFinite} \implies \text{st}(- x) = - \text{st } x$
by (*simp add*: *st-unique st-SReal st-approx-self approx-minus*)

lemma *st-diff*: $\llbracket x \in \text{HFinite}; y \in \text{HFinite} \rrbracket \implies \text{st}(x - y) = \text{st } x - \text{st } y$
by (*simp add*: *st-unique st-SReal st-approx-self approx-diff*)

lemma *st-mult*: $\llbracket x \in \text{HFinite}; y \in \text{HFinite} \rrbracket \implies \text{st}(x * y) = \text{st } x * \text{st } y$
by (*simp add*: *st-unique st-SReal st-approx-self approx-mult-HFinite*)

lemma *st-Infinitesimal*: $x \in \text{Infinitesimal} \implies \text{st } x = 0$
by (*simp add*: *st-unique mem-infmal-iff*)

lemma *st-not-Infinitesimal*: $\text{st}(x) \neq 0 \implies x \notin \text{Infinitesimal}$
by (*fast intro*: *st-Infinitesimal*)

lemma *st-inverse*: $x \in \text{HFinite} \implies \text{st } x \neq 0 \implies \text{st}(\text{inverse } x) = \text{inverse}(\text{st } x)$
apply (*rule-tac c1 = st x in mult-left-cancel* [*THEN iffD1*])
apply (*auto simp add*: *st-mult* [*symmetric*] *st-not-Infinitesimal HFinite-inverse*)
apply (*subst right-inverse, auto*)
done

lemma *st-divide* [*simp*]: $x \in \mathit{HFinite} \implies y \in \mathit{HFinite} \implies \mathit{st} \ y \neq 0 \implies \mathit{st} \ (x / y) = \mathit{st} \ x / \mathit{st} \ y$
by (*simp add: divide-inverse st-mult st-not-Infinitesimal HFinite-inverse st-inverse*)

lemma *st-idempotent* [*simp*]: $x \in \mathit{HFinite} \implies \mathit{st} \ (\mathit{st} \ x) = \mathit{st} \ x$
by (*blast intro: st-HFinite st-approx-self approx-st-eq*)

lemma *Infinitesimal-add-st-less*:
 $x \in \mathit{HFinite} \implies y \in \mathit{HFinite} \implies u \in \mathit{Infinitesimal} \implies \mathit{st} \ x < \mathit{st} \ y \implies \mathit{st} \ x + u < \mathit{st} \ y$
apply (*drule st-SReal*)
apply (*auto intro!: Infinitesimal-add-hypreal-of-real-less simp add: SReal-iff*)
done

lemma *Infinitesimal-add-st-le-cancel*:
 $x \in \mathit{HFinite} \implies y \in \mathit{HFinite} \implies u \in \mathit{Infinitesimal} \implies \mathit{st} \ x \leq \mathit{st} \ y + u \implies \mathit{st} \ x \leq \mathit{st} \ y$
apply (*simp add: linorder-not-less [symmetric]*)
apply (*auto dest: Infinitesimal-add-st-less*)
done

lemma *st-le*: $x \in \mathit{HFinite} \implies y \in \mathit{HFinite} \implies x \leq y \implies \mathit{st} \ x \leq \mathit{st} \ y$
by (*metis approx-le-bound approx-sym linear st-SReal st-approx-self st-part-Ex1*)

lemma *st-zero-le*: $0 \leq x \implies x \in \mathit{HFinite} \implies 0 \leq \mathit{st} \ x$
apply (*subst st-0 [symmetric]*)
apply (*rule st-le, auto*)
done

lemma *st-zero-ge*: $x \leq 0 \implies x \in \mathit{HFinite} \implies \mathit{st} \ x \leq 0$
apply (*subst st-0 [symmetric]*)
apply (*rule st-le, auto*)
done

lemma *st-hrabs*: $x \in \mathit{HFinite} \implies |\mathit{st} \ x| = \mathit{st} \ |x|$
apply (*simp add: linorder-not-le st-zero-le abs-if st-minus linorder-not-less*)
apply (*auto dest!: st-zero-ge [OF order-less-imp-le]*)
done

6.8 Alternative Definitions using Free Ultrafilter

6.8.1 *HFinite*

lemma *HFinite-FreeUltrafilterNat*:
 $\mathit{star-n} \ X \in \mathit{HFinite} \implies \exists u. \mathit{eventually} \ (\lambda n. \mathit{norm} \ (X \ n) < u) \ \mathcal{U}$
apply (*auto simp add: HFinite-def SReal-def*)
apply (*rule-tac x=r in exI*)
apply (*simp add: hnorm-def star-of-def starfun-star-n*)
apply (*simp add: star-less-def starP2-star-n*)

done

lemma *FreeUltrafilterNat-HFinite:*

$\exists u. \text{eventually } (\lambda n. \text{norm } (X n) < u) \mathcal{U} \implies \text{star-}n X \in \text{HFinite}$

apply (auto simp add: HFinite-def mem-Rep-star-iff)

apply (rule-tac $x = \text{star-of } u$ in *bxI*)

apply (simp add: hnorm-def starfun-star-n star-of-def)

apply (simp add: star-less-def starP2-star-n)

apply (simp add: SReal-def)

done

lemma *HFinite-FreeUltrafilterNat-iff:*

$\text{star-}n X \in \text{HFinite} \iff (\exists u. \text{eventually } (\lambda n. \text{norm } (X n) < u) \mathcal{U})$

by (blast intro!: HFinite-FreeUltrafilterNat FreeUltrafilterNat-HFinite)

6.8.2 HInfinite

lemma *lemma-Compl-eq:* $-\{n. u < \text{norm } (f n)\} = \{n. \text{norm } (f n) \leq u\}$

by auto

lemma *lemma-Compl-eq2:* $-\{n. \text{norm } (f n) < u\} = \{n. u \leq \text{norm } (f n)\}$

by auto

lemma *lemma-Int-eq1:* $\{n. \text{norm } (f n) \leq u\} \text{Int } \{n. u \leq \text{norm } (f n)\} = \{n. \text{norm}(f n) = u\}$

by auto

lemma *lemma-FreeUltrafilterNat-one:* $\{n. \text{norm } (f n) = u\} \leq \{n. \text{norm } (f n) < u + (1::\text{real})\}$

by auto

Exclude this type of sets from free ultrafilter for Infinite numbers!

lemma *FreeUltrafilterNat-const-Finite:*

$\text{eventually } (\lambda n. \text{norm } (X n) = u) \mathcal{U} \implies \text{star-}n X \in \text{HFinite}$

apply (rule FreeUltrafilterNat-HFinite)

apply (rule-tac $x = u + 1$ in *exI*)

apply (auto elim: eventually-mono)

done

lemma *HInfinite-FreeUltrafilterNat:*

$\text{star-}n X \in \text{HInfinite} \implies \forall u. \text{eventually } (\lambda n. u < \text{norm } (X n)) \mathcal{U}$

apply (drule HInfinite-HFinite-iff [THEN iffD1])

apply (simp add: HFinite-FreeUltrafilterNat-iff)

apply (rule allI, drule-tac $x = u + 1$ in *spec*)

apply (simp add: FreeUltrafilterNat.eventually-not-iff [symmetric])

apply (auto elim: eventually-mono)

done

lemma *lemma-Int-HI:* $\{n. \text{norm } (Xa n) < u\} \cap \{n. X n = Xa n\} \subseteq \{n. \text{norm}$

$(X\ n) < u$
for $u :: \text{real}$
by *auto*

lemma *lemma-Int-HIa*: $\{n. u < \text{norm } (X\ n)\} \cap \{n. \text{norm } (X\ n) < u\} = \{\}$
by (*auto intro: order-less-asym*)

lemma *FreeUltrafilterNat-HInfinite*:

$\forall u. \text{eventually } (\lambda n. u < \text{norm } (X\ n))\ \mathcal{U} \implies \text{star-}n\ X \in \text{HInfinite}$
apply (*rule HInfinite-HFinite-iff [THEN iffD2]*)
apply (*safe, drule HFinite-FreeUltrafilterNat, safe*)
apply (*drule-tac x = u in spec*)

proof –

fix u
assume $\forall_{Fn\ \text{in } \mathcal{U}}. \text{norm } (X\ n) < u \ \forall_{Fn\ \text{in } \mathcal{U}}. u < \text{norm } (X\ n)$
then have $\forall_F\ x\ \text{in } \mathcal{U}. \text{False}$
by *eventually-elim auto*
then show *False*
by (*simp add: eventually-False FreeUltrafilterNat.proper*)

qed

lemma *HInfinite-FreeUltrafilterNat-iff*:

$\text{star-}n\ X \in \text{HInfinite} \iff (\forall u. \text{eventually } (\lambda n. u < \text{norm } (X\ n))\ \mathcal{U})$
by (*blast intro!: HInfinite-FreeUltrafilterNat FreeUltrafilterNat-HInfinite*)

6.8.3 Infinitesimal

lemma *ball-SReal-eq*: $(\forall x::\text{hypreal} \in \text{Reals}. P\ x) \iff (\forall x::\text{real}. P\ (\text{star-of } x))$
by (*auto simp: SReal-def*)

lemma *Infinitesimal-FreeUltrafilterNat*:

$\text{star-}n\ X \in \text{Infinitesimal} \implies \forall u > 0. \text{eventually } (\lambda n. \text{norm } (X\ n) < u)\ \mathcal{U}$
apply (*simp add: Infinitesimal-def ball-SReal-eq*)
apply (*simp add: hnorm-def starfun-star-n star-of-def*)
apply (*simp add: star-less-def starP2-star-n*)
done

lemma *FreeUltrafilterNat-Infinitesimal*:

$\forall u > 0. \text{eventually } (\lambda n. \text{norm } (X\ n) < u)\ \mathcal{U} \implies \text{star-}n\ X \in \text{Infinitesimal}$
apply (*simp add: Infinitesimal-def ball-SReal-eq*)
apply (*simp add: hnorm-def starfun-star-n star-of-def*)
apply (*simp add: star-less-def starP2-star-n*)
done

lemma *Infinitesimal-FreeUltrafilterNat-iff*:

$(\text{star-}n\ X \in \text{Infinitesimal}) = (\forall u > 0. \text{eventually } (\lambda n. \text{norm } (X\ n) < u)\ \mathcal{U})$
by (*blast intro!: Infinitesimal-FreeUltrafilterNat FreeUltrafilterNat-Infinitesimal*)

Infinitesimals as smaller than $1/n$ for all $n::\text{nat } (> 0)$.

lemma *lemma-Infinitesimal*: $(\forall r. 0 < r \longrightarrow x < r) \longleftrightarrow (\forall n. x < \text{inverse} (\text{real} (\text{Suc } n)))$

apply (*auto simp del: of-nat-Suc*)
apply (*blast dest!: reals-Archimedean intro: order-less-trans*)
done

lemma *lemma-Infinitesimal2*:

$(\forall r \in \text{Reals}. 0 < r \longrightarrow x < r) \longleftrightarrow (\forall n. x < \text{inverse}(\text{hypreal-of-nat} (\text{Suc } n)))$

apply *safe*
apply (*drule-tac x = inverse (hypreal-of-real (real (Suc n))) in bspec*)
apply *simp-all*
using *less-imp-of-nat-less* **apply** *fastforce*
apply (*auto dest!: reals-Archimedean simp add: SReal-iff simp del: of-nat-Suc*)
apply (*drule star-of-less [THEN iffD2]*)
apply *simp*
apply (*blast intro: order-less-trans*)
done

lemma *Infinitesimal-hypreal-of-nat-iff*:

$\text{Infinitesimal} = \{x. \forall n. \text{hnorm } x < \text{inverse} (\text{hypreal-of-nat} (\text{Suc } n))\}$

apply (*simp add: Infinitesimal-def*)
apply (*auto simp add: lemma-Infinitesimal2*)
done

6.9 Proof that ω is an infinite number

It will follow that ε is an infinitesimal number.

lemma *Suc-Un-eq*: $\{n. n < \text{Suc } m\} = \{n. n < m\} \cup \{n. n = m\}$
by (*auto simp add: less-Suc-eq*)

Prove that any segment is finite and hence cannot belong to \mathcal{U} .

lemma *finite-real-of-nat-segment*: *finite* $\{n::\text{nat}. \text{real } n < \text{real } (m::\text{nat})\}$
by *auto*

lemma *finite-real-of-nat-less-real*: *finite* $\{n::\text{nat}. \text{real } n < u\}$
apply (*cut-tac x = u in reals-Archimedean2, safe*)
apply (*rule finite-real-of-nat-segment [THEN [2] finite-subset]*)
apply (*auto dest: order-less-trans*)
done

lemma *lemma-real-le-Un-eq*: $\{n. f n \leq u\} = \{n. f n < u\} \cup \{n. u = (f n :: \text{real})\}$
by (*auto dest: order-le-imp-less-or-eq simp add: order-less-imp-le*)

lemma *finite-real-of-nat-le-real*: *finite* $\{n::\text{nat}. \text{real } n \leq u\}$
by (*auto simp add: lemma-real-le-Un-eq lemma-finite-omega-set finite-real-of-nat-less-real*)

lemma *finite-rabs-real-of-nat-le-real*: *finite* $\{n::\text{nat}. |\text{real } n| \leq u\}$
by (*simp add: finite-real-of-nat-le-real*)

lemma *rabs-real-of-nat-le-real-FreeUltrafilterNat*:
 \neg eventually $(\lambda n. |real\ n| \leq u)\ \mathcal{U}$
by (*blast intro!*: *FreeUltrafilterNat.finite finite-rabs-real-of-nat-le-real*)

lemma *FreeUltrafilterNat-nat-gt-real*: eventually $(\lambda n. u < real\ n)\ \mathcal{U}$
apply (*rule FreeUltrafilterNat.finite'*)
apply (*subgoal-tac* $\{n::nat. \neg u < real\ n\} = \{n. real\ n \leq u\}$)
apply (*auto simp add: finite-real-of-nat-le-real*)
done

The complement of $\{n. |real\ n| \leq u\} = \{n. u < |real\ n|\}$ is in \mathcal{U} by property of (free) ultrafilters.

lemma *Compl-real-le-eq*: $-\{n::nat. real\ n \leq u\} = \{n. u < real\ n\}$
by (*auto dest!*: *order-le-less-trans simp add: linorder-not-le*)

ω is a member of *HInfinite*.

theorem *HInfinite-omega* [*simp*]: $\omega \in HInfinite$
apply (*simp add: omega-def*)
apply (*rule FreeUltrafilterNat-HInfinite*)
apply *clarify*
apply (*rule-tac* $u1 = u - 1$ **in** *eventually-mono* [*OF FreeUltrafilterNat-nat-gt-real*])
apply *auto*
done

Epsilon is a member of *Infinitesimal*.

lemma *Infinitesimal-epsilon* [*simp*]: $\epsilon \in Infinitesimal$
by (*auto intro!*: *HInfinite-inverse-Infinitesimal HInfinite-omega*
simp add: hypreal-epsilon-inverse-omega)

lemma *HFinite-epsilon* [*simp*]: $\epsilon \in HFinite$
by (*auto intro*: *Infinitesimal-subset-HFinite* [*THEN subsetD*])

lemma *epsilon-approx-zero* [*simp*]: $\epsilon \approx 0$
by (*simp add: mem-infmal-iff* [*symmetric*])

Needed for proof that we define a hyperreal $[<X(n)] \approx hypreal\ of\ real\ a$ given that $\forall n. |X\ n - a| < 1/n$. Used in proof of *NSLIM* \Rightarrow *LIM*.

lemma *real-of-nat-less-inverse-iff*: $0 < u \Rightarrow u < inverse\ (real\ (Suc\ n)) \leftrightarrow real\ (Suc\ n) < inverse\ u$
apply (*simp add: inverse-eq-divide*)
apply (*subst pos-less-divide-eq, assumption*)
apply (*subst pos-less-divide-eq*)
apply *simp*
apply (*simp add: mult commute*)
done

lemma *finite-inverse-real-of-posnat-gt-real*: $0 < u \Rightarrow finite\ \{n. u < inverse\ (real\ (Suc\ n))\}$

proof (*simp only: real-of-nat-less-inverse-iff*)
have $\{n. 1 + \text{real } n < \text{inverse } u\} = \{n. \text{real } n < \text{inverse } u - 1\}$
by *fastforce*
then show *finite* $\{n. \text{real } (\text{Suc } n) < \text{inverse } u\}$
using *finite-real-of-nat-less-real* [*of inverse u - 1*]
by *auto*
qed

lemma *lemma-real-le-Un-eq2*:
 $\{n. u \leq \text{inverse}(\text{real}(\text{Suc } n))\} =$
 $\{n. u < \text{inverse}(\text{real}(\text{Suc } n))\} \cup \{n. u = \text{inverse}(\text{real}(\text{Suc } n))\}$
by (*auto dest: order-le-imp-less-or-eq simp add: order-less-imp-le*)

lemma *finite-inverse-real-of-posnat-ge-real*: $0 < u \implies \text{finite } \{n. u \leq \text{inverse}(\text{real}(\text{Suc } n))\}$
by (*auto simp add: lemma-real-le-Un-eq2 lemma-finite-epsilon-set finite-inverse-real-of-posnat-gt-real simp del: of-nat-Suc*)

lemma *inverse-real-of-posnat-ge-real-FreeUltrafilterNat*:
 $0 < u \implies \neg \text{eventually } (\lambda n. u \leq \text{inverse}(\text{real}(\text{Suc } n))) \mathcal{U}$
by (*blast intro!: FreeUltrafilterNat.finite finite-inverse-real-of-posnat-ge-real*)

The complement of $\{n. u \leq \text{inverse}(\text{real}(\text{Suc } n))\} = \{n. \text{inverse}(\text{real}(\text{Suc } n)) < u\}$ is in \mathcal{U} by property of (free) ultrafilters.

lemma *Compl-le-inverse-eq*: $-\{n. u \leq \text{inverse}(\text{real}(\text{Suc } n))\} = \{n. \text{inverse}(\text{real}(\text{Suc } n)) < u\}$
by (*auto dest!: order-le-less-trans simp add: linorder-not-le*)

lemma *FreeUltrafilterNat-inverse-real-of-posnat*:
 $0 < u \implies \text{eventually } (\lambda n. \text{inverse}(\text{real}(\text{Suc } n)) < u) \mathcal{U}$
by (*drule inverse-real-of-posnat-ge-real-FreeUltrafilterNat*)
(*simp add: FreeUltrafilterNat.eventually-not-iff not-le[symmetric]*)

Example of an hypersequence (i.e. an extended standard sequence) whose term with an hypernatural suffix is an infinitesimal i.e. the n 'th term of the hypersequence is a member of *Infinitesimal*

lemma *SEQ-Infinitesimal*: $(** (\lambda n::\text{nat}. \text{inverse}(\text{real}(\text{Suc } n)))) \text{whn} \in \text{Infinitesimal}$
by (*simp add: hypnat-omega-def starfun-star-n star-n-inverse Infinitesimal-FreeUltrafilterNat-iff FreeUltrafilterNat-inverse-real-of-posnat del: of-nat-Suc*)

Example where we get a hyperreal from a real sequence for which a particular property holds. The theorem is used in proofs about equivalence of nonstandard and standard neighbourhoods. Also used for equivalence of nonstandard and standard definitions of pointwise limit.

$|X(n) - x| < 1/n \implies [\langle X \ n \rangle] - \text{hypreal-of-real } x \in \text{Infinitesimal}$

lemma *real-seq-to-hypreal-Infinitesimal*:

$\forall n. \text{norm } (X\ n - x) < \text{inverse } (\text{real } (\text{Suc } n)) \implies \text{star-}n\ X - \text{star-of } x \in \text{Infinitesimal}$

unfolding *star-n-diff star-of-def Infinitesimal-FreeUltrafilterNat-iff star-n-inverse*
by (*auto dest!: FreeUltrafilterNat-inverse-real-of-posnat*
intro: order-less-trans elim!: eventually-mono)

lemma *real-seq-to-hypreal-approx*:

$\forall n. \text{norm } (X\ n - x) < \text{inverse } (\text{real } (\text{Suc } n)) \implies \text{star-}n\ X \approx \text{star-of } x$
by (*metis bex-Infinitesimal-iff real-seq-to-hypreal-Infinitesimal*)

lemma *real-seq-to-hypreal-approx2*:

$\forall n. \text{norm } (x - X\ n) < \text{inverse}(\text{real}(\text{Suc } n)) \implies \text{star-}n\ X \approx \text{star-of } x$
by (*metis norm-minus-commute real-seq-to-hypreal-approx*)

lemma *real-seq-to-hypreal-Infinitesimal2*:

$\forall n. \text{norm}(X\ n - Y\ n) < \text{inverse}(\text{real}(\text{Suc } n)) \implies \text{star-}n\ X - \text{star-}n\ Y \in \text{Infinitesimal}$

unfolding *Infinitesimal-FreeUltrafilterNat-iff star-n-diff*
by (*auto dest!: FreeUltrafilterNat-inverse-real-of-posnat*
intro: order-less-trans elim!: eventually-mono)

end

7 Nonstandard Complex Numbers

theory *NSComplex*

imports *NSA*

begin

type-synonym *hcomplex = complex star*

abbreviation *hcomplex-of-complex :: complex \Rightarrow complex star*
where *hcomplex-of-complex \equiv star-of*

abbreviation *hcmmod :: complex star \Rightarrow real star*
where *hcmmod \equiv hnorm*

7.0.1 Real and Imaginary parts

definition *hRe :: hcomplex \Rightarrow hypreal*
where *hRe = *f* Re*

definition *hIm :: hcomplex \Rightarrow hypreal*
where *hIm = *f* Im*

7.0.2 Imaginary unit

definition *iii :: hcomplex*
where *iii = star-of i*

7.0.3 Complex conjugate

definition $hcnj :: hcomplex \Rightarrow hcomplex$
where $hcnj = *f* cnj$

7.0.4 Argand

definition $hsgn :: hcomplex \Rightarrow hcomplex$
where $hsgn = *f* sgn$

definition $harg :: hcomplex \Rightarrow hypreal$
where $harg = *f* arg$

definition — abbreviation for $\cos a + i \sin a$
 $hcis :: hypreal \Rightarrow hcomplex$
where $hcis = *f* cis$

7.0.5 Injection from hyperreals

abbreviation $hcomplex-of-hypreal :: hypreal \Rightarrow hcomplex$
where $hcomplex-of-hypreal \equiv of-hypreal$

definition — abbreviation for $r * (\cos a + i \sin a)$
 $hrcis :: hypreal \Rightarrow hypreal \Rightarrow hcomplex$
where $hrcis = *f2* rcis$

7.0.6 $e^{\wedge}(x + iy)$

definition $hExp :: hcomplex \Rightarrow hcomplex$
where $hExp = *f* exp$

definition $HComplex :: hypreal \Rightarrow hypreal \Rightarrow hcomplex$
where $HComplex = *f2* Complex$

lemmas $hcomplex-defs$ [transfer-unfold] =
 $hRe-def$ $hIm-def$ $iii-def$ $hcnj-def$ $hsgn-def$ $harg-def$ $hcis-def$
 $hrcis-def$ $hExp-def$ $HComplex-def$

lemma $Standard-hRe$ [simp]: $x \in Standard \Longrightarrow hRe x \in Standard$
by (simp add: $hcomplex-defs$)

lemma $Standard-hIm$ [simp]: $x \in Standard \Longrightarrow hIm x \in Standard$
by (simp add: $hcomplex-defs$)

lemma $Standard-iii$ [simp]: $iii \in Standard$
by (simp add: $hcomplex-defs$)

lemma $Standard-hcnj$ [simp]: $x \in Standard \Longrightarrow hcnj x \in Standard$
by (simp add: $hcomplex-defs$)

lemma *Standard-hsgn* [simp]: $x \in \text{Standard} \implies \text{hsgn } x \in \text{Standard}$
by (simp add: hcomplex-defs)

lemma *Standard-harg* [simp]: $x \in \text{Standard} \implies \text{harg } x \in \text{Standard}$
by (simp add: hcomplex-defs)

lemma *Standard-hcis* [simp]: $r \in \text{Standard} \implies \text{hcis } r \in \text{Standard}$
by (simp add: hcomplex-defs)

lemma *Standard-hExp* [simp]: $x \in \text{Standard} \implies \text{hExp } x \in \text{Standard}$
by (simp add: hcomplex-defs)

lemma *Standard-hrcis* [simp]: $r \in \text{Standard} \implies s \in \text{Standard} \implies \text{hrcis } r \ s \in \text{Standard}$
by (simp add: hcomplex-defs)

lemma *Standard-HComplex* [simp]: $r \in \text{Standard} \implies s \in \text{Standard} \implies \text{HComplex } r \ s \in \text{Standard}$
by (simp add: hcomplex-defs)

lemma *hcmmod-def*: $\text{hcmmod} = *f* \ \text{cmmod}$
by (rule hnorm-def)

7.1 Properties of Nonstandard Real and Imaginary Parts

lemma *hcomplex-hRe-hIm-cancel-iff*: $\bigwedge w \ z. \ w = z \iff \text{hRe } w = \text{hRe } z \wedge \text{hIm } w = \text{hIm } z$
by transfer (rule complex-Re-Im-cancel-iff)

lemma *hcomplex-equality* [intro?]: $\bigwedge z \ w. \ \text{hRe } z = \text{hRe } w \implies \text{hIm } z = \text{hIm } w \implies z = w$
by transfer (rule complex-equality)

lemma *hcomplex-hRe-zero* [simp]: $\text{hRe } 0 = 0$
by transfer simp

lemma *hcomplex-hIm-zero* [simp]: $\text{hIm } 0 = 0$
by transfer simp

lemma *hcomplex-hRe-one* [simp]: $\text{hRe } 1 = 1$
by transfer simp

lemma *hcomplex-hIm-one* [simp]: $\text{hIm } 1 = 0$
by transfer simp

7.2 Addition for Nonstandard Complex Numbers

lemma *hRe-add*: $\bigwedge x \ y. \ \text{hRe } (x + y) = \text{hRe } x + \text{hRe } y$
by transfer simp

lemma *hIm-add*: $\bigwedge x y. \text{hIm } (x + y) = \text{hIm } x + \text{hIm } y$
 by *transfer simp*

7.3 More Minus Laws

lemma *hRe-minus*: $\bigwedge z. \text{hRe } (-z) = - \text{hRe } z$
 by *transfer (rule uminus-complex.sel)*

lemma *hIm-minus*: $\bigwedge z. \text{hIm } (-z) = - \text{hIm } z$
 by *transfer (rule uminus-complex.sel)*

lemma *hcomplex-add-minus-eq-minus*: $x + y = 0 \implies x = -y$
 for $x y :: \text{hcomplex}$
 apply (*drule minus-unique*)
 apply (*simp add: minus-equation-iff [of x y]*)
 done

lemma *hcomplex-i-mult-eq [simp]*: $iii * iii = -1$
 by *transfer (rule i-squared)*

lemma *hcomplex-i-mult-left [simp]*: $\bigwedge z. iii * (iii * z) = -z$
 by *transfer (rule complex-i-mult-minus)*

lemma *hcomplex-i-not-zero [simp]*: $iii \neq 0$
 by *transfer (rule complex-i-not-zero)*

7.4 More Multiplication Laws

lemma *hcomplex-mult-minus-one*: $-1 * z = -z$
 for $z :: \text{hcomplex}$
 by *simp*

lemma *hcomplex-mult-minus-one-right*: $z * -1 = -z$
 for $z :: \text{hcomplex}$
 by *simp*

lemma *hcomplex-mult-left-cancel*: $c \neq 0 \implies c * a = c * b \longleftrightarrow a = b$
 for $a b c :: \text{hcomplex}$
 by *simp*

lemma *hcomplex-mult-right-cancel*: $c \neq 0 \implies a * c = b * c \longleftrightarrow a = b$
 for $a b c :: \text{hcomplex}$
 by *simp*

7.5 Subtraction and Division

lemma *hcomplex-diff-eq-eq [simp]*: $x - y = z \longleftrightarrow x = z + y$
 for $x y z :: \text{hcomplex}$
 by (*rule diff-eq-eq*)

7.6 Embedding Properties for $hcomplex$ -of- $hypreal$ Map

lemma hRe - $hcomplex$ -of- $hypreal$ [simp]: $\bigwedge z. hRe (hcomplex\text{-of-}hypreal\ z) = z$
 by transfer (rule Re - $complex$ -of- $real$)

lemma hIm - $hcomplex$ -of- $hypreal$ [simp]: $\bigwedge z. hIm (hcomplex\text{-of-}hypreal\ z) = 0$
 by transfer (rule Im - $complex$ -of- $real$)

lemma $hcomplex$ -of- $hypreal$ - ϵ -not-zero [simp]: $hcomplex\text{-of-}hypreal\ \epsilon \neq 0$
 by (simp add: $hypreal$ - ϵ -not-zero)

7.7 $HComplex$ theorems

lemma hRe - $HComplex$ [simp]: $\bigwedge x\ y. hRe (HComplex\ x\ y) = x$
 by transfer simp

lemma hIm - $HComplex$ [simp]: $\bigwedge x\ y. hIm (HComplex\ x\ y) = y$
 by transfer simp

lemma $hcomplex$ -surj [simp]: $\bigwedge z. HComplex (hRe\ z) (hIm\ z) = z$
 by transfer (rule $complex$ -surj)

lemma $hcomplex$ -induct [case-names rect]:
 $(\bigwedge x\ y. P (HComplex\ x\ y)) \implies P\ z$
 by (rule $hcomplex$ -surj [THEN subst]) blast

7.8 Modulus (Absolute Value) of Nonstandard Complex Number

lemma $hcomplex$ -of- $hypreal$ -abs:
 $hcomplex\text{-of-}hypreal\ |x| = hcomplex\text{-of-}hypreal (hmod (hcomplex\text{-of-}hypreal\ x))$
 by simp

lemma $HComplex$ -inject [simp]: $\bigwedge x\ y\ x'\ y'. HComplex\ x\ y = HComplex\ x'\ y' \iff$
 $x = x' \wedge y = y'$
 by transfer (rule $complex$.inject)

lemma $HComplex$ -add [simp]:
 $\bigwedge x1\ y1\ x2\ y2. HComplex\ x1\ y1 + HComplex\ x2\ y2 = HComplex (x1 + x2) (y1 + y2)$
 by transfer (rule $complex$ -add)

lemma $HComplex$ -minus [simp]: $\bigwedge x\ y. - HComplex\ x\ y = HComplex (- x) (- y)$
 by transfer (rule $complex$ -minus)

lemma $HComplex$ -diff [simp]:
 $\bigwedge x1\ y1\ x2\ y2. HComplex\ x1\ y1 - HComplex\ x2\ y2 = HComplex (x1 - x2) (y1 - y2)$
 by transfer (rule $complex$ -diff)

lemma *HComplex-mult* [simp]:

$\bigwedge x1\ y1\ x2\ y2. HComplex\ x1\ y1 * HComplex\ x2\ y2 = HComplex\ (x1*x2 - y1*y2)$
 $(x1*y2 + y1*x2)$
 by transfer (rule complex-mult)

HComplex-inverse is proved below.

lemma *hcomplex-of-hypreal-eq*: $\bigwedge r. hcomplex-of-hypreal\ r = HComplex\ r\ 0$
 by transfer (rule complex-of-real-def)

lemma *HComplex-add-hcomplex-of-hypreal* [simp]:

$\bigwedge x\ y\ r. HComplex\ x\ y + hcomplex-of-hypreal\ r = HComplex\ (x + r)\ y$
 by transfer (rule Complex-add-complex-of-real)

lemma *hcomplex-of-hypreal-add-HComplex* [simp]:

$\bigwedge r\ x\ y. hcomplex-of-hypreal\ r + HComplex\ x\ y = HComplex\ (r + x)\ y$
 by transfer (rule complex-of-real-add-Complex)

lemma *HComplex-mult-hcomplex-of-hypreal*:

$\bigwedge x\ y\ r. HComplex\ x\ y * hcomplex-of-hypreal\ r = HComplex\ (x * r)\ (y * r)$
 by transfer (rule Complex-mult-complex-of-real)

lemma *hcomplex-of-hypreal-mult-HComplex*:

$\bigwedge r\ x\ y. hcomplex-of-hypreal\ r * HComplex\ x\ y = HComplex\ (r * x)\ (r * y)$
 by transfer (rule complex-of-real-mult-Complex)

lemma *i-hcomplex-of-hypreal* [simp]: $\bigwedge r. iii * hcomplex-of-hypreal\ r = HComplex\ 0\ r$

by transfer (rule i-complex-of-real)

lemma *hcomplex-of-hypreal-i* [simp]: $\bigwedge r. hcomplex-of-hypreal\ r * iii = HComplex\ 0\ r$

by transfer (rule complex-of-real-i)

7.9 Conjugation

lemma *hcomplex-hcnj-cancel-iff* [iff]: $\bigwedge x\ y. hcnj\ x = hcnj\ y \longleftrightarrow x = y$
 by transfer (rule complex-cnj-cancel-iff)

lemma *hcomplex-hcnj-hcnj* [simp]: $\bigwedge z. hcnj\ (hcnj\ z) = z$
 by transfer (rule complex-cnj-cnj)

lemma *hcomplex-hcnj-hcomplex-of-hypreal* [simp]:

$\bigwedge x. hcnj\ (hcomplex-of-hypreal\ x) = hcomplex-of-hypreal\ x$
 by transfer (rule complex-cnj-complex-of-real)

lemma *hcomplex-hmod-hcnj* [simp]: $\bigwedge z. hmod\ (hcnj\ z) = hmod\ z$
 by transfer (rule complex-mod-cnj)

lemma *hcomplex-hcnj-minus*: $\bigwedge z. \text{hcnj } (-z) = - \text{hcnj } z$
by *transfer (rule complex-cnj-minus)*

lemma *hcomplex-hcnj-inverse*: $\bigwedge z. \text{hcnj } (\text{inverse } z) = \text{inverse } (\text{hcnj } z)$
by *transfer (rule complex-cnj-inverse)*

lemma *hcomplex-hcnj-add*: $\bigwedge w z. \text{hcnj } (w + z) = \text{hcnj } w + \text{hcnj } z$
by *transfer (rule complex-cnj-add)*

lemma *hcomplex-hcnj-diff*: $\bigwedge w z. \text{hcnj } (w - z) = \text{hcnj } w - \text{hcnj } z$
by *transfer (rule complex-cnj-diff)*

lemma *hcomplex-hcnj-mult*: $\bigwedge w z. \text{hcnj } (w * z) = \text{hcnj } w * \text{hcnj } z$
by *transfer (rule complex-cnj-mult)*

lemma *hcomplex-hcnj-divide*: $\bigwedge w z. \text{hcnj } (w / z) = \text{hcnj } w / \text{hcnj } z$
by *transfer (rule complex-cnj-divide)*

lemma *hcnj-one [simp]*: $\text{hcnj } 1 = 1$
by *transfer (rule complex-cnj-one)*

lemma *hcomplex-hcnj-zero [simp]*: $\text{hcnj } 0 = 0$
by *transfer (rule complex-cnj-zero)*

lemma *hcomplex-hcnj-zero-iff [iff]*: $\bigwedge z. \text{hcnj } z = 0 \longleftrightarrow z = 0$
by *transfer (rule complex-cnj-zero-iff)*

lemma *hcomplex-mult-hcnj*: $\bigwedge z. z * \text{hcnj } z = \text{hcomplex-of-hypreal } ((\text{hRe } z)^2 + (\text{hIm } z)^2)$
by *transfer (rule complex-mult-cnj)*

7.10 More Theorems about the Function *hcm*

lemma *hcm-hcomplex-of-hypreal-of-nat [simp]*:
 $\text{hcm } (\text{hcomplex-of-hypreal } (\text{hypreal-of-nat } n)) = \text{hypreal-of-nat } n$
by *simp*

lemma *hcm-hcomplex-of-hypreal-of-hypnat [simp]*:
 $\text{hcm } (\text{hcomplex-of-hypreal}(\text{hypreal-of-hypnat } n)) = \text{hypreal-of-hypnat } n$
by *simp*

lemma *hcm-mult-hcnj*: $\bigwedge z. \text{hcm } (z * \text{hcnj } z) = (\text{hcm } z)^2$
by *transfer (rule complex-mod-mult-cnj)*

lemma *hcm-triangle-ineq2 [simp]*: $\bigwedge a b. \text{hcm } (b + a) - \text{hcm } b \leq \text{hcm } a$
by *transfer (rule complex-mod-triangle-ineq2)*

lemma *hcm-diff-ineq [simp]*: $\bigwedge a b. \text{hcm } a - \text{hcm } b \leq \text{hcm } (a + b)$
by *transfer (rule norm-diff-ineq)*

7.11 Exponentiation

lemma *hcomplexpow-0* [*simp*]: $z \wedge 0 = 1$
for $z :: \text{hcomplex}$
by (*rule power-0*)

lemma *hcomplexpow-Suc* [*simp*]: $z \wedge (\text{Suc } n) = z * (z \wedge n)$
for $z :: \text{hcomplex}$
by (*rule power-Suc*)

lemma *hcomplexpow-i-squared* [*simp*]: $iii^2 = -1$
by *transfer* (*rule power2-i*)

lemma *hcomplex-of-hypreal-pow*: $\wedge x. \text{hcomplex-of-hypreal } (x \wedge n) = \text{hcomplex-of-hypreal } x \wedge n$
by *transfer* (*rule of-real-power*)

lemma *hcomplex-hcnj-pow*: $\wedge z. \text{hcnj } (z \wedge n) = \text{hcnj } z \wedge n$
by *transfer* (*rule complex-cnj-power*)

lemma *hcmmod-hcomplexpow*: $\wedge x. \text{hcmmod } (x \wedge n) = \text{hcmmod } x \wedge n$
by *transfer* (*rule norm-power*)

lemma *hcpow-minus*:
 $\wedge x n. (-x :: \text{hcomplex}) \text{ pow } n = (\text{if } (*p* \text{ even}) \text{ then } (x \text{ pow } n) \text{ else } -(x \text{ pow } n))$
by *transfer simp*

lemma *hcpow-mult*: $(r * s) \text{ pow } n = (r \text{ pow } n) * (s \text{ pow } n)$
for $r s :: \text{hcomplex}$
by (*fact hyperpow-mult*)

lemma *hcpow-zero2* [*simp*]: $\wedge n. 0 \text{ pow } (\text{hSuc } n) = (0::'a::\text{semiring-1 star})$
by *transfer* (*rule power-0-Suc*)

lemma *hcpow-not-zero* [*simp,intro*]: $\wedge r n. r \neq 0 \implies r \text{ pow } n \neq (0::\text{hcomplex})$
by (*fact hyperpow-not-zero*)

lemma *hcpow-zero-zero*: $r \text{ pow } n = 0 \implies r = 0$
for $r :: \text{hcomplex}$
by (*blast intro: ccontr dest: hcpow-not-zero*)

7.12 The Function *hsgn*

lemma *hsgn-zero* [*simp*]: $\text{hsgn } 0 = 0$
by *transfer* (*rule sgn-zero*)

lemma *hsgn-one* [*simp*]: $\text{hsgn } 1 = 1$
by *transfer* (*rule sgn-one*)

lemma *hsgn-minus*: $\bigwedge z. \text{hsgn } (-z) = - \text{hsgn } z$
by *transfer (rule sgn-minus)*

lemma *hsgn-eq*: $\bigwedge z. \text{hsgn } z = z / \text{hcomplex-of-hypreal } (\text{hmod } z)$
by *transfer (rule sgn-eq)*

lemma *hmod-i*: $\bigwedge x y. \text{hmod } (\text{HComplex } x y) = (*f* \text{sqrt}) (x^2 + y^2)$
by *transfer (rule complex-norm)*

lemma *hcomplex-eq-cancel-iff1 [simp]*:
 $\text{hcomplex-of-hypreal } xa = \text{HComplex } x y \longleftrightarrow xa = x \wedge y = 0$
by *(simp add: hcomplex-of-hypreal-eq)*

lemma *hcomplex-eq-cancel-iff2 [simp]*:
 $\text{HComplex } x y = \text{hcomplex-of-hypreal } xa \longleftrightarrow x = xa \wedge y = 0$
by *(simp add: hcomplex-of-hypreal-eq)*

lemma *HComplex-eq-0 [simp]*: $\bigwedge x y. \text{HComplex } x y = 0 \longleftrightarrow x = 0 \wedge y = 0$
by *transfer (rule Complex-eq-0)*

lemma *HComplex-eq-1 [simp]*: $\bigwedge x y. \text{HComplex } x y = 1 \longleftrightarrow x = 1 \wedge y = 0$
by *transfer (rule Complex-eq-1)*

lemma *i-eq-HComplex-0-1*: $iii = \text{HComplex } 0 1$
by *transfer (simp add: complex-eq-iff)*

lemma *HComplex-eq-i [simp]*: $\bigwedge x y. \text{HComplex } x y = iii \longleftrightarrow x = 0 \wedge y = 1$
by *transfer (rule Complex-eq-i)*

lemma *hRe-hsgn [simp]*: $\bigwedge z. \text{hRe } (\text{hsgn } z) = \text{hRe } z / \text{hmod } z$
by *transfer (rule Re-sgn)*

lemma *hIm-hsgn [simp]*: $\bigwedge z. \text{hIm } (\text{hsgn } z) = \text{hIm } z / \text{hmod } z$
by *transfer (rule Im-sgn)*

lemma *HComplex-inverse*: $\bigwedge x y. \text{inverse } (\text{HComplex } x y) = \text{HComplex } (x / (x^2 + y^2)) (-y / (x^2 + y^2))$
by *transfer (rule complex-inverse)*

lemma *hRe-mult-i-eq [simp]*: $\bigwedge y. \text{hRe } (iii * \text{hcomplex-of-hypreal } y) = 0$
by *transfer simp*

lemma *hIm-mult-i-eq [simp]*: $\bigwedge y. \text{hIm } (iii * \text{hcomplex-of-hypreal } y) = y$
by *transfer simp*

lemma *hmod-mult-i [simp]*: $\bigwedge y. \text{hmod } (iii * \text{hcomplex-of-hypreal } y) = |y|$
by *transfer (simp add: norm-complex-def)*

lemma *hmod-mult-i2 [simp]*: $\bigwedge y. \text{hmod } (\text{hcomplex-of-hypreal } y * iii) = |y|$

by transfer (simp add: norm-complex-def)

7.12.1 harg

lemma *cos-harg-i-mult-zero* [simp]: $\bigwedge y. y \neq 0 \implies (*f* \cos) (harg (HComplex 0 y)) = 0$

by transfer (simp add: Complex-eq)

7.13 Polar Form for Nonstandard Complex Numbers

lemma *complex-split-polar2*: $\forall n. \exists r a. (z n) = \text{complex-of-real } r * \text{Complex } (\cos a) (\sin a)$

unfolding *Complex-eq* by (auto intro: complex-split-polar)

lemma *hcomplex-split-polar*:

$\bigwedge z. \exists r a. z = \text{hcomplex-of-hypreal } r * (HComplex ((*f* \cos) a) ((*f* \sin) a))$
by transfer (simp add: Complex-eq complex-split-polar)

lemma *hcis-eq*:

$\bigwedge a. \text{hcis } a = \text{hcomplex-of-hypreal } ((*f* \cos) a) + iii * \text{hcomplex-of-hypreal } ((*f* \sin) a)$

by transfer (simp add: complex-eq-iff)

lemma *hrcis-Ex*: $\bigwedge z. \exists r a. z = \text{hrcis } r a$

by transfer (rule rcis-Ex)

lemma *hRe-hcomplex-polar* [simp]:

$\bigwedge r a. \text{hRe } (\text{hcomplex-of-hypreal } r * HComplex ((*f* \cos) a) ((*f* \sin) a)) = r * (*f* \cos) a$

by transfer simp

lemma *hRe-hrcis* [simp]: $\bigwedge r a. \text{hRe } (\text{hrcis } r a) = r * (*f* \cos) a$

by transfer (rule Re-rcis)

lemma *hIm-hcomplex-polar* [simp]:

$\bigwedge r a. \text{hIm } (\text{hcomplex-of-hypreal } r * HComplex ((*f* \cos) a) ((*f* \sin) a)) = r * (*f* \sin) a$

by transfer simp

lemma *hIm-hrcis* [simp]: $\bigwedge r a. \text{hIm } (\text{hrcis } r a) = r * (*f* \sin) a$

by transfer (rule Im-rcis)

lemma *hcmmod-unit-one* [simp]: $\bigwedge a. \text{hcmmod } (HComplex ((*f* \cos) a) ((*f* \sin) a)) = 1$

by transfer (simp add: cmmod-unit-one)

lemma *hcmmod-complex-polar* [simp]:

$\bigwedge r a. \text{hcmmod } (\text{hcomplex-of-hypreal } r * HComplex ((*f* \cos) a) ((*f* \sin) a)) = |r|$

by transfer (simp add: Complex-eq cmmod-complex-polar)

lemma *hcmmod-hrcis* [*simp*]: $\bigwedge r a. \text{hcmmod}(\text{hrcis } r a) = |r|$
by *transfer* (*rule complex-mod-rcis*)

$$(r1 * \text{hrcis } a) * (r2 * \text{hrcis } b) = r1 * r2 * \text{hrcis } (a + b)$$

lemma *hcis-hrcis-eq*: $\bigwedge a. \text{hcis } a = \text{hrcis } 1 a$
by *transfer* (*rule cis-rcis-eq*)

declare *hcis-hrcis-eq* [*symmetric, simp*]

lemma *hrcis-mult*: $\bigwedge a b r1 r2. \text{hrcis } r1 a * \text{hrcis } r2 b = \text{hrcis } (r1 * r2) (a + b)$
by *transfer* (*rule rcis-mult*)

lemma *hcis-mult*: $\bigwedge a b. \text{hcis } a * \text{hcis } b = \text{hcis } (a + b)$
by *transfer* (*rule cis-mult*)

lemma *hcis-zero* [*simp*]: $\text{hcis } 0 = 1$
by *transfer* (*rule cis-zero*)

lemma *hrcis-zero-mod* [*simp*]: $\bigwedge a. \text{hrcis } 0 a = 0$
by *transfer* (*rule rcis-zero-mod*)

lemma *hrcis-zero-arg* [*simp*]: $\bigwedge r. \text{hrcis } r 0 = \text{hcomplex-of-hypreal } r$
by *transfer* (*rule rcis-zero-arg*)

lemma *hcomplex-i-mult-minus* [*simp*]: $\bigwedge x. iiii * (iii * x) = - x$
by *transfer* (*rule complex-i-mult-minus*)

lemma *hcomplex-i-mult-minus2* [*simp*]: $iii * iii * x = - x$
by *simp*

lemma *hcis-hypreal-of-nat-Suc-mult*:
 $\bigwedge a n. \text{hcis } (\text{hypreal-of-nat } (\text{Suc } n) * a) = \text{hcis } a * \text{hcis } (\text{hypreal-of-nat } n * a)$
by *transfer* (*simp add: distrib-right cis-mult*)

lemma *NSDeMoivre*: $\bigwedge a. (\text{hcis } a) ^ n = \text{hcis } (\text{hypreal-of-nat } n * a)$
by *transfer* (*rule DeMoivre*)

lemma *hcis-hypreal-of-hypnat-Suc-mult*:
 $\bigwedge a n. \text{hcis } (\text{hypreal-of-hypnat } (n + 1) * a) = \text{hcis } a * \text{hcis } (\text{hypreal-of-hypnat } n * a)$
by *transfer* (*simp add: distrib-right cis-mult*)

lemma *NSDeMoivre-ext*: $\bigwedge a n. (\text{hcis } a) \text{ pow } n = \text{hcis } (\text{hypreal-of-hypnat } n * a)$
by *transfer* (*rule DeMoivre*)

lemma *NSDeMoivre2*: $\bigwedge a r. (\text{hrcis } r a) ^ n = \text{hrcis } (r ^ n) (\text{hypreal-of-nat } n * a)$
by *transfer* (*rule DeMoivre2*)

lemma *DeMoiivre2-ext*: $\bigwedge a r n. (\text{hrcis } r a) \text{ pow } n = \text{hrcis } (r \text{ pow } n) (\text{hypreal-of-hypnat } n * a)$

by *transfer (rule DeMoiivre2)*

lemma *hcis-inverse [simp]*: $\bigwedge a. \text{inverse } (\text{hcis } a) = \text{hcis } (- a)$

by *transfer (rule cis-inverse)*

lemma *hrcis-inverse*: $\bigwedge a r. \text{inverse } (\text{hrcis } r a) = \text{hrcis } (\text{inverse } r) (- a)$

by *transfer (simp add: rcis-inverse inverse-eq-divide [symmetric])*

lemma *hRe-hcis [simp]*: $\bigwedge a. \text{hRe } (\text{hcis } a) = (* * \text{cos}) a$

by *transfer simp*

lemma *hIm-hcis [simp]*: $\bigwedge a. \text{hIm } (\text{hcis } a) = (* * \text{sin}) a$

by *transfer simp*

lemma *cos-n-hRe-hcis-pow-n*: $(* * \text{cos}) (\text{hypreal-of-nat } n * a) = \text{hRe } (\text{hcis } a ^ n)$

by *(simp add: NSDeMoiivre)*

lemma *sin-n-hIm-hcis-pow-n*: $(* * \text{sin}) (\text{hypreal-of-nat } n * a) = \text{hIm } (\text{hcis } a ^ n)$

by *(simp add: NSDeMoiivre)*

lemma *cos-n-hRe-hcis-hcpow-n*: $(* * \text{cos}) (\text{hypreal-of-hypnat } n * a) = \text{hRe } (\text{hcis } a \text{ pow } n)$

by *(simp add: NSDeMoiivre-ext)*

lemma *sin-n-hIm-hcis-hcpow-n*: $(* * \text{sin}) (\text{hypreal-of-hypnat } n * a) = \text{hIm } (\text{hcis } a \text{ pow } n)$

by *(simp add: NSDeMoiivre-ext)*

lemma *hExp-add*: $\bigwedge a b. \text{hExp } (a + b) = \text{hExp } a * \text{hExp } b$

by *transfer (rule exp-add)*

7.14 *hcomplex-of-complex*: the Injection from type *complex* to *hcomplex*

lemma *hcomplex-of-complex-i*: $\text{iii} = \text{hcomplex-of-complex } i$

by *(rule iii-def)*

lemma *hRe-hcomplex-of-complex*: $\text{hRe } (\text{hcomplex-of-complex } z) = \text{hypreal-of-real } (\text{Re } z)$

by *transfer (rule refl)*

lemma *hIm-hcomplex-of-complex*: $\text{hIm } (\text{hcomplex-of-complex } z) = \text{hypreal-of-real } (\text{Im } z)$

by *transfer (rule refl)*

lemma *hcmmod-hcomplex-of-complex*: $hcmmod (hcomplex-of-complex\ x) = hypreal-of-real (cmmod\ x)$
 by *transfer* (rule *refl*)

7.15 Numerals and Arithmetic

lemma *hcomplex-of-hypreal-eq-hcomplex-of-complex*:
 $hcomplex-of-hypreal (hypreal-of-real\ x) = hcomplex-of-complex (complex-of-real\ x)$
 by *transfer* (rule *refl*)

lemma *hcomplex-hypreal-numeral*:
 $hcomplex-of-complex (numeral\ w) = hcomplex-of-hypreal (numeral\ w)$
 by *transfer* (rule *of-real-numeral* [*symmetric*])

lemma *hcomplex-hypreal-neg-numeral*:
 $hcomplex-of-complex (-\ numeral\ w) = hcomplex-of-hypreal (-\ numeral\ w)$
 by *transfer* (rule *of-real-neg-numeral* [*symmetric*])

lemma *hcomplex-numeral-hcnj* [*simp*]: $hcnj (numeral\ v :: hcomplex) = numeral\ v$
 by *transfer* (rule *complex-cnj-numeral*)

lemma *hcomplex-numeral-hcmmod* [*simp*]: $hcmmod (numeral\ v :: hcomplex) = (numeral\ v :: hypreal)$
 by *transfer* (rule *norm-numeral*)

lemma *hcomplex-neg-numeral-hcmmod* [*simp*]: $hcmmod (-\ numeral\ v :: hcomplex) = (numeral\ v :: hypreal)$
 by *transfer* (rule *norm-neg-numeral*)

lemma *hcomplex-numeral-hRe* [*simp*]: $hRe (numeral\ v :: hcomplex) = numeral\ v$
 by *transfer* (rule *complex-Re-numeral*)

lemma *hcomplex-numeral-hIm* [*simp*]: $hIm (numeral\ v :: hcomplex) = 0$
 by *transfer* (rule *complex-Im-numeral*)

end

8 Star-Transforms in Non-Standard Analysis

theory *Star*
 imports *NSA*
 begin

definition — internal sets
 $starset-n :: (nat \Rightarrow 'a\ set) \Rightarrow 'a\ star\ set$ (**sn** - [80] 80)
 where **sn** $As = Iset (star-n\ As)$

definition *InternalSets* :: $'a\ star\ set\ set$

where $InternalSets = \{X. \exists As. X = *sn* As\}$

definition — nonstandard extension of function

$is-starext :: ('a star \Rightarrow 'a star) \Rightarrow ('a \Rightarrow 'a) \Rightarrow bool$

where $is-starext F f \longleftrightarrow$

$(\forall x y. \exists X \in Rep-star x. \exists Y \in Rep-star y. y = F x \longleftrightarrow eventually (\lambda n. Y n = f(X n)) U)$

definition — internal functions

$starfun-n :: (nat \Rightarrow 'a \Rightarrow 'b) \Rightarrow 'a star \Rightarrow 'b star \quad (*fn* - [80] 80)$

where $*fn* F = Ifun (star-n F)$

definition $InternalFuns :: ('a star \Rightarrow 'b star) set$

where $InternalFuns = \{X. \exists F. X = *fn* F\}$

8.1 Preamble - Pulling \exists over \forall

This proof does not need AC and was suggested by the referee for the JCM Paper: let $f x$ be least y such that $Q x y$.

lemma *no-choice*: $\forall x. \exists y. Q x y \implies \exists f :: 'a \Rightarrow nat. \forall x. Q x (f x)$

by (*rule exI* [**where** $x = \lambda x. LEAST y. Q x y$]) (*blast intro: LeastI*)

8.2 Properties of the Star-transform Applied to Sets of Reals

lemma *STAR-star-of-image-subset*: $star-of ' A \subseteq *s* A$

by *auto*

lemma *STAR-hypreal-of-real-Int*: $*s* X \cap \mathbb{R} = hypreal-of-real ' X$

by (*auto simp add: SReal-def*)

lemma *STAR-star-of-Int*: $*s* X \cap Standard = star-of ' X$

by (*auto simp add: Standard-def*)

lemma *lemma-not-hyprealA*: $x \notin hypreal-of-real ' A \implies \forall y \in A. x \neq hypreal-of-real y$

by *auto*

lemma *lemma-not-starA*: $x \notin star-of ' A \implies \forall y \in A. x \neq star-of y$

by *auto*

lemma *lemma-Compl-eq*: $-\{n. X n = xa\} = \{n. X n \neq xa\}$

by *auto*

lemma *STAR-real-seq-to-hypreal*: $\forall n. (X n) \notin M \implies star-n X \notin *s* M$

by (*simp add: starset-def star-of-def Iset-star-n FreeUltrafilterNat.proper*)

lemma *STAR-singleton*: $*s* \{x\} = \{star-of x\}$

by *simp*

lemma *STAR-not-mem*: $x \notin F \implies \text{star-of } x \notin ** F$
by *transfer*

lemma *STAR-subset-closed*: $x \in ** A \implies A \subseteq B \implies x \in ** B$
by (*erule rev-subsetD*) *simp*

Nonstandard extension of a set (defined using a constant sequence) as a special case of an internal set.

lemma *starset-n-starset*: $\forall n. As\ n = A \implies *sn* As = ** A$
by (*drule fun-eq-iff [THEN iffD2]*) (*simp add: starset-n-def starset-def star-of-def*)

8.3 Theorems about nonstandard extensions of functions

Nonstandard extension of a function (defined using a constant sequence) as a special case of an internal function.

lemma *starfun-n-starfun*: $\forall n. F\ n = f \implies *fn* F = ** f$
apply (*drule fun-eq-iff [THEN iffD2]*)
apply (*simp add: starfun-n-def starfun-def star-of-def*)
done

Prove that *abs* for hypreal is a nonstandard extension of *abs* for real w/o use of congruence property (proved after this for general nonstandard extensions of real valued functions).

Proof now Uses the ultrafilter tactic!

lemma *hrabs-is-starext-rabs*: *is-starext abs abs*
apply (*simp add: is-starext-def, safe*)
apply (*rule-tac x=x in star-cases*)
apply (*rule-tac x=y in star-cases*)
apply (*unfold star-n-def, auto*)
apply (*rule bezI, rule-tac [2] lemma-starrel-refl*)
apply (*rule bezI, rule-tac [2] lemma-starrel-refl*)
apply (*fold star-n-def*)
apply (*unfold star-abs-def starfun-def star-of-def*)
apply (*simp add: Ifun-star-n star-n-eq-iff*)
done

Nonstandard extension of functions.

lemma *starfun*: $(** f) (star-n\ X) = star-n (\lambda n. f (X\ n))$
by (*rule starfun-star-n*)

lemma *starfun-if-eq*: $\bigwedge w. w \neq \text{star-of } x \implies (** (\lambda z. \text{if } z = x \text{ then } a \text{ else } g\ z))$
 $w = (** g) w$
by *transfer simp*

Multiplication: $(** f) x (** g) = *(f\ x\ g)$

lemma *starfun-mult*: $\bigwedge x. (** f) x (** g) x = (** (\lambda x. f\ x\ * g\ x)) x$
by *transfer (rule refl)*

declare *starfun-mult* [*symmetric, simp*]

Addition: $(*f) + (*g) = *(f + g)$

lemma *starfun-add*: $\bigwedge x. (**f) x + (**g) x = (**(\lambda x. f x + g x)) x$
by *transfer (rule refl)*

declare *starfun-add* [*symmetric, simp*]

Subtraction: $(*f) + -(*g) = *(f + -g)$

lemma *starfun-minus*: $\bigwedge x. - (**f) x = (**(\lambda x. - f x)) x$
by *transfer (rule refl)*

declare *starfun-minus* [*symmetric, simp*]

lemma *starfun-add-minus*: $\bigwedge x. (**f) x + -(**g) x = (**(\lambda x. f x + -g x)) x$

by *transfer (rule refl)*

declare *starfun-add-minus* [*symmetric, simp*]

lemma *starfun-diff*: $\bigwedge x. (**f) x - (**g) x = (**(\lambda x. f x - g x)) x$
by *transfer (rule refl)*

declare *starfun-diff* [*symmetric, simp*]

Composition: $(*f) \circ (*g) = *(f \circ g)$

lemma *starfun-o2*: $(\lambda x. (**f) ((**g) x)) = **(\lambda x. f (g x))$
by *transfer (rule refl)*

lemma *starfun-o*: $(**f) \circ (**g) = (**(\lambda x. f (g x)))$
by *(transfer o-def) (rule refl)*

NS extension of constant function.

lemma *starfun-const-fun* [*simp*]: $\bigwedge x. (**(\lambda x. k)) x = \text{star-of } k$
by *transfer (rule refl)*

The NS extension of the identity function.

lemma *starfun-Id* [*simp*]: $\bigwedge x. (**(\lambda x. x)) x = x$
by *transfer (rule refl)*

This is trivial, given *starfun-Id*.

lemma *starfun-Idfun-approx*: $x \approx \text{star-of } a \implies (**(\lambda x. x)) x \approx \text{star-of } a$
by *(simp only: starfun-Id)*

The Star-function is a (nonstandard) extension of the function.

lemma *is-starext-starfun*: *is-starext* $(**f) f$
apply *(auto simp: is-starext-def)*
apply *(rule-tac x = x in star-cases)*
apply *(rule-tac x = y in star-cases)*
apply *(auto intro!: bexI [OF - Rep-star-star-n] simp: starfun star-n-eq-iff)*
done

Any nonstandard extension is in fact the Star-function.

lemma *is-starfun-starext*: $is-starext F f \implies F = *f* f$
apply (*simp add: is-starext-def*)
apply (*rule ext*)
apply (*rule-tac x = x in star-cases*)
apply (*drule-tac x = x in spec*)
apply (*drule-tac x = (*f* f) x in spec*)
apply (*auto simp add: starfun-star-n*)
apply (*simp add: star-n-eq-iff [symmetric]*)
apply (*simp add: starfun-star-n [of f, symmetric]*)
done

lemma *is-starext-starfun-iff*: $is-starext F f \iff F = *f* f$
by (*blast intro: is-starfun-starext is-starext-starfun*)

Extended function has same solution as its standard version for real arguments. i.e they are the same for all real arguments.

lemma *starfun-eq*: $(*f* f) (star-of a) = star-of (f a)$
by (*rule starfun-star-of*)

lemma *starfun-approx*: $(*f* f) (star-of a) \approx star-of (f a)$
by *simp*

Useful for NS definition of derivatives.

lemma *starfun-lambda-cancel*: $\bigwedge x'. (*f* (\lambda h. f (x + h))) x' = (*f* f) (star-of x + x')$
by *transfer (rule refl)*

lemma *starfun-lambda-cancel2*: $(*f* (\lambda h. f (g (x + h)))) x' = (*f* (f \circ g)) (star-of x + x')$
unfolding *o-def* **by** (*rule starfun-lambda-cancel*)

lemma *starfun-mult-HFinite-approx*:
 $(*f* f) x \approx l \implies (*f* g) x \approx m \implies l \in HFinite \implies m \in HFinite \implies$
 $(*f* (\lambda x. f x * g x)) x \approx l * m$
for $l m :: 'a::real-normed-algebra$ *star*
apply (*drule (3) approx-mult-HFinite*)
apply (*auto intro: approx-HFinite [OF - approx-sym]*)
done

lemma *starfun-add-approx*: $(*f* f) x \approx l \implies (*f* g) x \approx m \implies (*f* (\%x. f x + g x)) x \approx l + m$
by (*auto intro: approx-add*)

Examples: *hrabs* is nonstandard extension of *rabs*, *inverse* is nonstandard extension of *inverse*.

Can be proved easily using theorem *starfun* and properties of ultrafilter as for *inverse* below we use the theorem we proved above instead.

lemma *starfun-rabs-hrabs*: $*f* \text{ abs} = \text{abs}$
by (*simp only: star-abs-def*)

lemma *starfun-inverse-inverse* [*simp*]: $(*f* \text{ inverse}) x = \text{inverse } x$
by (*simp only: star-inverse-def*)

lemma *starfun-inverse*: $\bigwedge x. \text{inverse } ((*f* f) x) = (*f* (\lambda x. \text{inverse } (f x))) x$
by *transfer (rule refl)*
declare *starfun-inverse* [*symmetric, simp*]

lemma *starfun-divide*: $\bigwedge x. (*f* f) x / (*f* g) x = (*f* (\lambda x. f x / g x)) x$
by *transfer (rule refl)*
declare *starfun-divide* [*symmetric, simp*]

lemma *starfun-inverse2*: $\bigwedge x. \text{inverse } ((*f* f) x) = (*f* (\lambda x. \text{inverse } (f x))) x$
by *transfer (rule refl)*

General lemma/theorem needed for proofs in elementary topology of the reals.

lemma *starfun-mem-starset*: $\bigwedge x. (*f* f) x : *s* A \implies x \in *s* \{x. f x \in A\}$
by *transfer simp*

Alternative definition for *hrabs* with *rabs* function applied entrywise to equivalence class representative. This is easily proved using *starfun* and *ns extension thm*.

lemma *hypreal-hrabs*: $|\text{star-}n X| = \text{star-}n (\lambda n. |X n|)$
by (*simp only: starfun-rabs-hrabs [symmetric] starfun*)

Nonstandard extension of set through nonstandard extension of *rabs* function i.e. *hrabs*. A more general result should be where we replace *rabs* by some arbitrary function *f* and *hrabs* by its NS extension. See second NS set extension below.

lemma *STAR-rabs-add-minus*: $*s* \{x. |x + - y| < r\} = \{x. |x + -\text{star-of } y| < \text{star-of } r\}$
by *transfer (rule refl)*

lemma *STAR-starfun-rabs-add-minus*:
 $*s* \{x. |f x + - y| < r\} = \{x. |(*f* f) x + -\text{star-of } y| < \text{star-of } r\}$
by *transfer (rule refl)*

Another characterization of Infinitesimal and one of \approx relation. In this theory since *hypreal-hrabs* proved here. Maybe move both theorems??

lemma *Infinitesimal-FreeUltrafilterNat-iff2*:
 $\text{star-}n X \in \text{Infinitesimal} \iff (\forall m. \text{eventually } (\lambda n. \text{norm } (X n) < \text{inverse } (\text{real } (\text{Suc } m)))) \mathcal{U}$
by (*simp add: Infinitesimal-hypreal-of-nat-iff star-of-def hnorm-def star-of-nat-def starfun-star-n star-n-inverse star-n-less*)

```

lemma HNatInfinite-inverse-Infinitesimal [simp]:
   $n \in \text{HNatInfinite} \implies \text{inverse} (\text{hypreal-of-hypnat } n) \in \text{Infinitesimal}$ 
  apply (cases  $n$ )
  apply (auto simp: of-hypnat-def starfun-star-n star-n-inverse
    HNatInfinite-FreeUltrafilterNat-iff Infinitesimal-FreeUltrafilterNat-iff2)
  apply (drule-tac  $x = \text{Suc } m$  in spec)
  apply (auto elim!: eventually-mono)
  done

lemma approx-FreeUltrafilterNat-iff:
   $\text{star-n } X \approx \text{star-n } Y \iff (\forall r > 0. \text{eventually } (\lambda n. \text{norm } (X \ n - Y \ n) < r) \ \mathcal{U})$ 
  apply (subst approx-minus-iff)
  apply (rule mem-infmal-iff [THEN subst])
  apply (simp add: star-n-diff)
  apply (simp add: Infinitesimal-FreeUltrafilterNat-iff)
  done

lemma approx-FreeUltrafilterNat-iff2:
   $\text{star-n } X \approx \text{star-n } Y \iff (\forall m. \text{eventually } (\lambda n. \text{norm } (X \ n - Y \ n) < \text{inverse} (\text{real } (\text{Suc } m))) \ \mathcal{U})$ 
  apply (subst approx-minus-iff)
  apply (rule mem-infmal-iff [THEN subst])
  apply (simp add: star-n-diff)
  apply (simp add: Infinitesimal-FreeUltrafilterNat-iff2)
  done

lemma inj-starfun: inj starfun
  apply (rule inj-onI)
  apply (rule ext, rule ccontr)
  apply (drule-tac  $x = \text{star-n } (\lambda n. xa)$  in fun-cong)
  apply (auto simp add: starfun star-n-eq-iff FreeUltrafilterNat.proper)
  done

```

end

9 Star-transforms for the Hypernaturals

```

theory NatStar
  imports Star
begin

```

```

lemma star-n-eq-starfun-whn: star-n X = (*f* X) whn
  by (simp add: hypnat-omega-def starfun-def star-of-def Ifun-star-n)

```

```

lemma starset-n-Un: *sn* (λn. (A n) ∪ (B n)) = *sn* A ∪ *sn* B
  apply (simp add: starset-n-def star-n-eq-starfun-whn Un-def)
  apply (rule-tac  $x = \text{whn}$  in spec, transfer, simp)
  done

```

lemma *InternalSets-Un*: $X \in \text{InternalSets} \implies Y \in \text{InternalSets} \implies X \cup Y \in \text{InternalSets}$

by (*auto simp add: InternalSets-def starset-n-Un [symmetric]*)

lemma *starset-n-Int*: $*sn* (\lambda n. A \ n \ \cap \ B \ n) = *sn* A \ \cap \ *sn* B$

apply (*simp add: starset-n-def star-n-eq-starfun-whn Int-def*)

apply (*rule-tac x=whn in spec, transfer, simp*)

done

lemma *InternalSets-Int*: $X \in \text{InternalSets} \implies Y \in \text{InternalSets} \implies X \cap Y \in \text{InternalSets}$

by (*auto simp add: InternalSets-def starset-n-Int [symmetric]*)

lemma *starset-n-Compl*: $*sn* ((\lambda n. - A \ n)) = - (*sn* A)$

apply (*simp add: starset-n-def star-n-eq-starfun-whn Compl-eq*)

apply (*rule-tac x=whn in spec, transfer, simp*)

done

lemma *InternalSets-Compl*: $X \in \text{InternalSets} \implies - X \in \text{InternalSets}$

by (*auto simp add: InternalSets-def starset-n-Compl [symmetric]*)

lemma *starset-n-diff*: $*sn* (\lambda n. (A \ n) - (B \ n)) = *sn* A - *sn* B$

apply (*simp add: starset-n-def star-n-eq-starfun-whn set-diff-eq*)

apply (*rule-tac x=whn in spec, transfer, simp*)

done

lemma *InternalSets-diff*: $X \in \text{InternalSets} \implies Y \in \text{InternalSets} \implies X - Y \in \text{InternalSets}$

by (*auto simp add: InternalSets-def starset-n-diff [symmetric]*)

lemma *NatStar-SHNat-subset*: $\text{Nats} \leq *s* (\text{UNIV}:: \text{nat set})$

by *simp*

lemma *NatStar-hypreal-of-real-Int*: $*s* X \ \text{Int} \ \text{Nats} = \text{hypnat-of-nat} \ 'X$

by (*auto simp add: SHNat-eq*)

lemma *starset-starset-n-eq*: $*s* X = *sn* (\lambda n. X)$

by (*simp add: starset-n-starset*)

lemma *InternalSets-starset-n [simp]*: $(*s* X) \in \text{InternalSets}$

by (*auto simp add: InternalSets-def starset-starset-n-eq*)

lemma *InternalSets-UNIV-diff*: $X \in \text{InternalSets} \implies \text{UNIV} - X \in \text{InternalSets}$

apply (*subgoal-tac UNIV - X = - X*)

apply (*auto intro: InternalSets-Compl*)

done

9.1 Nonstandard Extensions of Functions

Example of transfer of a property from reals to hyperreals — used for limit comparison of sequences.

lemma *starfun-le-mono*: $\forall n. N \leq n \longrightarrow f\ n \leq g\ n \implies$
 $\forall n. \text{hypnat-of-nat } N \leq n \longrightarrow (*f* f)\ n \leq (*f* g)\ n$
by *transfer*

And another:

lemma *starfun-less-mono*:
 $\forall n. N \leq n \longrightarrow f\ n < g\ n \implies \forall n. \text{hypnat-of-nat } N \leq n \longrightarrow (*f* f)\ n < (*f* g)\ n$
by *transfer*

Nonstandard extension when we increment the argument by one.

lemma *starfun-shift-one*: $\bigwedge N. (*f* (\lambda n. f\ (Suc\ n)))\ N = (*f* f)\ (N + (1::\text{hypnat}))$
by *transfer simp*

Nonstandard extension with absolute value.

lemma *starfun-abs*: $\bigwedge N. (*f* (\lambda n. |f\ n|))\ N = |(*f* f)\ N|$
by *transfer (rule refl)*

The *hyperpow* function as a nonstandard extension of *realpow*.

lemma *starfun-pow*: $\bigwedge N. (*f* (\lambda n. r\ ^\ n))\ N = \text{hypreal-of-real } r\ \text{pow } N$
by *transfer (rule refl)*

lemma *starfun-pow2*: $\bigwedge N. (*f* (\lambda n. X\ n\ ^\ m))\ N = (*f* X)\ N\ \text{pow } \text{hypnat-of-nat } m$
by *transfer (rule refl)*

lemma *starfun-pow3*: $\bigwedge R. (*f* (\lambda r. r\ ^\ n))\ R = R\ \text{pow } \text{hypnat-of-nat } n$
by *transfer (rule refl)*

The *hypreal-of-hypnat* function as a nonstandard extension of *real*.

lemma *starfunNat-real-of-nat*: $(*f* \text{real}) = \text{hypreal-of-hypnat}$
by *transfer (simp add: fun-eq-iff)*

lemma *starfun-inverse-real-of-nat-eq*:
 $N \in \text{HNatInfinite} \implies (*f* (\lambda x::\text{nat. inverse } (\text{real } x)))\ N = \text{inverse } (\text{hypreal-of-hypnat } N)$
apply (*rule-tac f1 = inverse in starfun-o2 [THEN subst]*)
apply (*subgoal-tac hypreal-of-hypnat N \neq 0*)
apply (*simp-all add: zero-less-HNatInfinite starfunNat-real-of-nat*)
done

Internal functions – some redundancy with **f** now.

lemma *starfun-n*: $(*fn* f)\ (\text{star-n } X) = \text{star-n } (\lambda n. f\ n\ (X\ n))$

by (*simp add: starfun-n-def Ifun-star-n*)

Multiplication: $(*fn) x (*gn) = *(fn x gn)$

lemma *starfun-n-mult*: $(*fn* f) z * (*fn* g) z = (*fn* (\lambda i x. f i x * g i x)) z$
by (*cases z*) (*simp add: starfun-n star-n-mult*)

Addition: $(*fn) + (*gn) = *(fn + gn)$

lemma *starfun-n-add*: $(*fn* f) z + (*fn* g) z = (*fn* (\lambda i x. f i x + g i x)) z$
by (*cases z*) (*simp add: starfun-n star-n-add*)

Subtraction: $(*fn) - (*gn) = *(fn + - gn)$

lemma *starfun-n-add-minus*: $(*fn* f) z + - (*fn* g) z = (*fn* (\lambda i x. f i x + -g i x)) z$
by (*cases z*) (*simp add: starfun-n star-n-minus star-n-add*)

Composition: $(*fn) \circ (*gn) = *(fn \circ gn)$

lemma *starfun-n-const-fun [simp]*: $(*fn* (\lambda i x. k)) z = \text{star-of } k$
by (*cases z*) (*simp add: starfun-n star-of-def*)

lemma *starfun-n-minus*: $- (*fn* f) x = (*fn* (\lambda i x. - (f i) x)) x$
by (*cases x*) (*simp add: starfun-n star-n-minus*)

lemma *starfun-n-eq [simp]*: $(*fn* f) (\text{star-of } n) = \text{star-n } (\lambda i. f i n)$
by (*simp add: starfun-n star-of-def*)

lemma *starfun-eq-iff*: $((*f* f) = (*f* g)) \longleftrightarrow f = g$
by *transfer (rule refl)*

lemma *starfunNat-inverse-real-of-nat-Infinitesimal [simp]*:
 $N \in \text{HNatInfinite} \implies (*f* (\%x. \text{inverse } (\text{real } x))) N \in \text{Infinitesimal}$
apply (*rule-tac f1 = inverse in starfun-o2 [THEN subst]*)
apply (*subgoal-tac hypreal-of-hypnat N $\sim = 0$*)
apply (*simp-all add: zero-less-HNatInfinite starfunNat-real-of-nat*)
done

9.2 Nonstandard Characterization of Induction

lemma *hypnat-induct-obj*:

$\bigwedge n. ((*p* P) (0::\text{hypnat}) \wedge (\forall n. (*p* P) n \longrightarrow (*p* P) (n + 1))) \longrightarrow (*p* P) n$
by *transfer (induct-tac n, auto)*

lemma *hypnat-induct*:

$\bigwedge n. (*p* P) (0::\text{hypnat}) \implies (\bigwedge n. (*p* P) n \implies (*p* P) (n + 1)) \implies (*p* P) n$
by *transfer (induct-tac n, auto)*

lemma *starP2-eq-iff*: $(*p2* (op =)) = (op =)$

by *transfer (rule refl)*

lemma *starP2-eq-iff2*: ($*p2*$ ($\lambda x y. x = y$)) $X Y \longleftrightarrow X = Y$
by (*simp add: starP2-eq-iff*)

lemma *nonempty-nat-set-Least-mem*: $c \in S \implies (LEAST n. n \in S) \in S$
for $S :: nat\ set$
by (*erule LeastI*)

lemma *nonempty-set-star-has-least*:
 $\bigwedge S::nat\ set\ star. Iset\ S \neq \{\}$ $\implies \exists n \in Iset\ S. \forall m \in Iset\ S. n \leq m$
apply (*transfer empty-def*)
apply (*rule-tac x=LEAST n. n \in S in bestI*)
apply (*simp add: Least-le*)
apply (*rule LeastI-ex, auto*)
done

lemma *nonempty-InternalNatSet-has-least*: $S \in InternalSets \implies S \neq \{\} \implies \exists n \in S. \forall m \in S. n \leq m$
for $S :: hypnat\ set$
apply (*clarsimp simp add: InternalSets-def starset-n-def*)
apply (*erule nonempty-set-star-has-least*)
done

Goldblatt, page 129 Thm 11.3.2.

lemma *internal-induct-lemma*:
 $\bigwedge X::nat\ set\ star.$
 $(0::hypnat) \in Iset\ X \implies \forall n. n \in Iset\ X \longrightarrow n + 1 \in Iset\ X \implies Iset\ X =$
 $(UNIV::hypnat\ set)$
apply (*transfer UNIV-def*)
apply (*rule equalityI [OF subset-UNIV subsetI]*)
apply (*induct-tac x, auto*)
done

lemma *internal-induct*:
 $X \in InternalSets \implies (0::hypnat) \in X \implies \forall n. n \in X \longrightarrow n + 1 \in X \implies X =$
 $(UNIV::hypnat\ set)$
apply (*clarsimp simp add: InternalSets-def starset-n-def*)
apply (*erule (1) internal-induct-lemma*)
done

end

10 Sequences and Convergence (Nonstandard)

theory *HSEQ*
imports *HOL.Limits NatStar*
abbrevs $----> = \longrightarrow_{NS}$
begin

definition *NSLIMSEQ* :: (nat ⇒ 'a::real-normed-vector) ⇒ 'a ⇒ bool
 (((-)/ —————_{NS} (-)) [60, 60] 60) **where**
 — Nonstandard definition of convergence of sequence
 $X \longrightarrow_{NS} L \iff (\forall N \in \mathit{HNatInfinite}. (*f* X) N \approx \mathit{star-of} L)$

definition *nslim* :: (nat ⇒ 'a::real-normed-vector) ⇒ 'a
where *nslim* X = (THE L. X —————_{NS} L)
 — Nonstandard definition of limit using choice operator

definition *NSconvergent* :: (nat ⇒ 'a::real-normed-vector) ⇒ bool
where *NSconvergent* X ⇔ (∃ L. X —————_{NS} L)
 — Nonstandard definition of convergence

definition *NSBseq* :: (nat ⇒ 'a::real-normed-vector) ⇒ bool
where *NSBseq* X ⇔ (∀ N ∈ *HNatInfinite*. (*f* X) N ∈ *HFinite*)
 — Nonstandard definition for bounded sequence

definition *NSCauchy* :: (nat ⇒ 'a::real-normed-vector) ⇒ bool
where *NSCauchy* X ⇔ (∀ M ∈ *HNatInfinite*. ∀ N ∈ *HNatInfinite*. (*f* X) N
 M ≈ (*f* X) N)
 — Nonstandard definition

10.1 Limits of Sequences

lemma *NSLIMSEQ-iff*: (X —————_{NS} L) ⇔ (∀ N ∈ *HNatInfinite*. (*f* X) N
 ≈ *star-of* L)
by (*simp add: NSLIMSEQ-def*)

lemma *NSLIMSEQ-I*: (∧ N. N ∈ *HNatInfinite* ⇒ *starfun* X N ≈ *star-of* L) ⇒
 X —————_{NS} L
by (*simp add: NSLIMSEQ-def*)

lemma *NSLIMSEQ-D*: X —————_{NS} L ⇒ N ∈ *HNatInfinite* ⇒ *starfun* X N ≈
star-of L
by (*simp add: NSLIMSEQ-def*)

lemma *NSLIMSEQ-const*: (λ n. k) —————_{NS} k
by (*simp add: NSLIMSEQ-def*)

lemma *NSLIMSEQ-add*: X —————_{NS} a ⇒ Y —————_{NS} b ⇒ (λ n. X n + Y
 n) —————_{NS} a + b
by (*auto intro: approx-add simp add: NSLIMSEQ-def*)

lemma *NSLIMSEQ-add-const*: f —————_{NS} a ⇒ (λ n. f n + b) —————_{NS} a + b
by (*simp only: NSLIMSEQ-add NSLIMSEQ-const*)

lemma *NSLIMSEQ-mult*: X —————_{NS} a ⇒ Y —————_{NS} b ⇒ (λ n. X n * Y

$n) \longrightarrow_{NS} a * b$
for $a b :: 'a::real-normed-algebra$
by (*auto intro!*: *approx-mult-HFinite simp add: NSLIMSEQ-def*)

lemma *NSLIMSEQ-minus*: $X \longrightarrow_{NS} a \implies (\lambda n. - X n) \longrightarrow_{NS} - a$
by (*auto simp add: NSLIMSEQ-def*)

lemma *NSLIMSEQ-minus-cancel*: $(\lambda n. - X n) \longrightarrow_{NS} - a \implies X \longrightarrow_{NS} a$
by (*drule NSLIMSEQ-minus*) *simp*

lemma *NSLIMSEQ-diff*: $X \longrightarrow_{NS} a \implies Y \longrightarrow_{NS} b \implies (\lambda n. X n - Y n) \longrightarrow_{NS} a - b$
using *NSLIMSEQ-add* [*of X a - Y - b*] **by** (*simp add: NSLIMSEQ-minus fun-Compl-def*)

lemma *NSLIMSEQ-add-minus*: $X \longrightarrow_{NS} a \implies Y \longrightarrow_{NS} b \implies (\lambda n. X n + - Y n) \longrightarrow_{NS} a + - b$
by (*simp add: NSLIMSEQ-diff*)

lemma *NSLIMSEQ-diff-const*: $f \longrightarrow_{NS} a \implies (\lambda n. f n - b) \longrightarrow_{NS} a - b$
by (*simp add: NSLIMSEQ-diff NSLIMSEQ-const*)

lemma *NSLIMSEQ-inverse*: $X \longrightarrow_{NS} a \implies a \neq 0 \implies (\lambda n. inverse (X n)) \longrightarrow_{NS} inverse a$
for $a :: 'a::real-normed-div-algebra$
by (*simp add: NSLIMSEQ-def star-of-approx-inverse*)

lemma *NSLIMSEQ-mult-inverse*: $X \longrightarrow_{NS} a \implies Y \longrightarrow_{NS} b \implies b \neq 0 \implies (\lambda n. X n / Y n) \longrightarrow_{NS} a / b$
for $a b :: 'a::real-normed-field$
by (*simp add: NSLIMSEQ-mult NSLIMSEQ-inverse divide-inverse*)

lemma *starfun-hnorm*: $\bigwedge x. hnorm ((*f* f) x) = (*f* (\lambda x. norm (f x))) x$
by *transfer simp*

lemma *NSLIMSEQ-norm*: $X \longrightarrow_{NS} a \implies (\lambda n. norm (X n)) \longrightarrow_{NS} norm a$
by (*simp add: NSLIMSEQ-def starfun-hnorm [symmetric] approx-hnorm*)

Uniqueness of limit.

lemma *NSLIMSEQ-unique*: $X \longrightarrow_{NS} a \implies X \longrightarrow_{NS} b \implies a = b$
apply (*simp add: NSLIMSEQ-def*)
apply (*drule HNatInfinite-whn [THEN [2] bspec]*)
apply (*auto dest: approx-trans3*)
done

lemma *NSLIMSEQ-pow* [*rule-format*]: $(X \longrightarrow_{NS} a) \longrightarrow ((\lambda n. (X n) ^ m) \longrightarrow_{NS} a ^ m)$

for $a :: 'a::\{\text{real-normed-algebra},\text{power}\}$
by (*induct m*) (*auto intro: NSLIMSEQ-mult NSLIMSEQ-const*)

We can now try and derive a few properties of sequences, starting with the limit comparison property for sequences.

lemma *NSLIMSEQ-le*: $f \longrightarrow_{NS} l \implies g \longrightarrow_{NS} m \implies \exists N. \forall n \geq N. f\ n \leq g\ n \implies l \leq m$

for $l\ m :: \text{real}$
apply (*simp add: NSLIMSEQ-def, safe*)
apply (*drule starfun-le-mono*)
apply (*drule HNatInfinite-whn [THEN [2] bspec]*)
apply (*drule-tac x = whn in spec*)
apply (*drule beX-Infininitesimal-iff2 [THEN iffD2]*)
apply *clarify*
apply (*auto intro: hypreal-of-real-le-add-Infininitesimal-cancel2*)
done

lemma *NSLIMSEQ-le-const*: $X \longrightarrow_{NS} r \implies \forall n. a \leq X\ n \implies a \leq r$
for $a\ r :: \text{real}$
by (*erule NSLIMSEQ-le [OF NSLIMSEQ-const]*) *auto*

lemma *NSLIMSEQ-le-const2*: $X \longrightarrow_{NS} r \implies \forall n. X\ n \leq a \implies r \leq a$
for $a\ r :: \text{real}$
by (*erule NSLIMSEQ-le [OF - NSLIMSEQ-const]*) *auto*

Shift a convergent series by 1: By the equivalence between Cauchiness and convergence and because the successor of an infinite hypernatural is also infinite.

lemma *NSLIMSEQ-Suc*: $f \longrightarrow_{NS} l \implies (\lambda n. f(Suc\ n)) \longrightarrow_{NS} l$
apply (*unfold NSLIMSEQ-def*)
apply *safe*
apply (*drule-tac x=N + 1 in bspec*)
apply (*erule HNatInfinite-add*)
apply (*simp add: starfun-shift-one*)
done

lemma *NSLIMSEQ-imp-Suc*: $(\lambda n. f(Suc\ n)) \longrightarrow_{NS} l \implies f \longrightarrow_{NS} l$
apply (*unfold NSLIMSEQ-def*)
apply *safe*
apply (*drule-tac x=N - 1 in bspec*)
apply (*erule Nats-1 [THEN [2] HNatInfinite-diff]*)
apply (*simp add: starfun-shift-one one-le-HNatInfinite*)
done

lemma *NSLIMSEQ-Suc-iff*: $(\lambda n. f(Suc\ n)) \longrightarrow_{NS} l \iff f \longrightarrow_{NS} l$
by (*blast intro: NSLIMSEQ-imp-Suc NSLIMSEQ-Suc*)

10.1.1 Equivalence of *LIMSEQ* and *NSLIMSEQ*

lemma *LIMSEQ-NSLIMSEQ*:

assumes $X: X \longrightarrow L$

shows $X \longrightarrow_{NS} L$

proof (rule *NSLIMSEQ-I*)

fix N

assume $N: N \in \text{HNatInfinite}$

have $\text{starfun } X \ N - \text{star-of } L \in \text{Infinitesimal}$

proof (rule *InfinitesimalI2*)

fix $r :: \text{real}$

assume $r: 0 < r$

from *LIMSEQ-D* [*OF* X r] obtain no where $\forall n \geq no. \text{norm } (X \ n - L) < r ..$

then have $\forall n \geq \text{star-of } no. \text{hnorm } (\text{starfun } X \ n - \text{star-of } L) < \text{star-of } r$

by *transfer*

then show $\text{hnorm } (\text{starfun } X \ N - \text{star-of } L) < \text{star-of } r$

using N by (*simp add: star-of-le-HNatInfinite*)

qed

then show $\text{starfun } X \ N \approx \text{star-of } L$

by (*simp only: approx-def*)

qed

lemma *NSLIMSEQ-LIMSEQ*:

assumes $X: X \longrightarrow_{NS} L$

shows $X \longrightarrow L$

proof (rule *LIMSEQ-I*)

fix $r :: \text{real}$

assume $r: 0 < r$

have $\exists no. \forall n \geq no. \text{hnorm } (\text{starfun } X \ n - \text{star-of } L) < \text{star-of } r$

proof (*intro exI allI impI*)

fix n

assume $whn \leq n$

with *HNatInfinite-whn* have $n \in \text{HNatInfinite}$

by (*rule HNatInfinite-upward-closed*)

with X have $\text{starfun } X \ n \approx \text{star-of } L$

by (*rule NSLIMSEQ-D*)

then have $\text{starfun } X \ n - \text{star-of } L \in \text{Infinitesimal}$

by (*simp only: approx-def*)

then show $\text{hnorm } (\text{starfun } X \ n - \text{star-of } L) < \text{star-of } r$

using r by (*rule InfinitesimalD2*)

qed

then show $\exists no. \forall n \geq no. \text{norm } (X \ n - L) < r$

by *transfer*

qed

theorem *LIMSEQ-NSLIMSEQ-iff*: $f \longrightarrow L \iff f \longrightarrow_{NS} L$

by (*blast intro: LIMSEQ-NSLIMSEQ NSLIMSEQ-LIMSEQ*)

10.1.2 Derived theorems about *NSLIMSEQ*

We prove the NS version from the standard one, since the NS proof seems more complicated than the standard one above!

lemma *NSLIMSEQ-norm-zero*: $(\lambda n. \text{norm } (X \ n)) \longrightarrow_{NS} 0 \longleftrightarrow X \longrightarrow_{NS} 0$

by (*simp add: LIMSEQ-NSLIMSEQ-iff [symmetric] tendsto-norm-zero-iff*)

lemma *NSLIMSEQ-rabs-zero*: $(\lambda n. |f \ n|) \longrightarrow_{NS} 0 \longleftrightarrow f \longrightarrow_{NS} (0::\text{real})$

by (*simp add: LIMSEQ-NSLIMSEQ-iff [symmetric] tendsto-rabs-zero-iff*)

Generalization to other limits.

lemma *NSLIMSEQ-imp-rabs*: $f \longrightarrow_{NS} l \implies (\lambda n. |f \ n|) \longrightarrow_{NS} |l|$

for $l :: \text{real}$

by (*simp add: NSLIMSEQ-def (auto intro: approx-hrabs simp add: starfun-abs)*)

lemma *NSLIMSEQ-inverse-zero*: $\forall y::\text{real}. \exists N. \forall n \geq N. y < f \ n \implies (\lambda n. \text{inverse } (f \ n)) \longrightarrow_{NS} 0$

by (*simp add: LIMSEQ-NSLIMSEQ-iff [symmetric] LIMSEQ-inverse-zero*)

lemma *NSLIMSEQ-inverse-real-of-nat*: $(\lambda n. \text{inverse } (\text{real } (\text{Suc } n))) \longrightarrow_{NS} 0$

by (*simp add: LIMSEQ-NSLIMSEQ-iff [symmetric] LIMSEQ-inverse-real-of-nat del: of-nat-Suc*)

lemma *NSLIMSEQ-inverse-real-of-nat-add*: $(\lambda n. r + \text{inverse } (\text{real } (\text{Suc } n))) \longrightarrow_{NS} r$

by (*simp add: LIMSEQ-NSLIMSEQ-iff [symmetric] LIMSEQ-inverse-real-of-nat-add del: of-nat-Suc*)

lemma *NSLIMSEQ-inverse-real-of-nat-add-minus*: $(\lambda n. r + - \text{inverse } (\text{real } (\text{Suc } n))) \longrightarrow_{NS} r$

using *LIMSEQ-inverse-real-of-nat-add-minus* **by** (*simp add: LIMSEQ-NSLIMSEQ-iff [symmetric]*)

lemma *NSLIMSEQ-inverse-real-of-nat-add-minus-mult*:

$(\lambda n. r * (1 + - \text{inverse } (\text{real } (\text{Suc } n)))) \longrightarrow_{NS} r$

using *LIMSEQ-inverse-real-of-nat-add-minus-mult*

by (*simp add: LIMSEQ-NSLIMSEQ-iff [symmetric]*)

10.2 Convergence

lemma *nslimI*: $X \longrightarrow_{NS} L \implies \text{nslim } X = L$

by (*simp add: nslim-def (blast intro: NSLIMSEQ-unique)*)

lemma *lim-nslim-iff*: $\text{lim } X = \text{nslim } X$

by (*simp add: lim-def nslim-def LIMSEQ-NSLIMSEQ-iff*)

lemma *NSconvergentD*: *NSconvergent* $X \implies \exists L. X \longrightarrow_{NS} L$

by (*simp add: NSconvergent-def*)

lemma *NSconvergentI*: $X \longrightarrow_{NS} L \implies NSconvergent X$
by (*auto simp add: NSconvergent-def*)

lemma *convergent-NSconvergent-iff*: $convergent X = NSconvergent X$
by (*simp add: convergent-def NSconvergent-def LIMSEQ-NSLIMSEQ-iff*)

lemma *NSconvergent-NSLIMSEQ-iff*: $NSconvergent X \longleftrightarrow X \longrightarrow_{NS} nslim X$
by (*auto intro: theI NSLIMSEQ-unique simp add: NSconvergent-def nslim-def*)

10.3 Bounded Monotonic Sequences

lemma *NSBseqD*: $NSBseq X \implies N \in HNatInfinite \implies (*f* X) N \in HFinite$
by (*simp add: NSBseq-def*)

lemma *Standard-subset-HFinite*: $Standard \subseteq HFinite$
by (*auto simp: Standard-def*)

lemma *NSBseqD2*: $NSBseq X \implies (*f* X) N \in HFinite$
apply (*cases N \in HNatInfinite*)
apply (*erule (1) NSBseqD*)
apply (*rule subsetD [OF Standard-subset-HFinite]*)
apply (*simp add: HNatInfinite-def Nats-eq-Standard*)
done

lemma *NSBseqI*: $\forall N \in HNatInfinite. (*f* X) N \in HFinite \implies NSBseq X$
by (*simp add: NSBseq-def*)

The standard definition implies the nonstandard definition.

lemma *Bseq-NSBseq*: $Bseq X \implies NSBseq X$
unfolding *NSBseq-def*

proof *safe*

assume $X: Bseq X$

fix N

assume $N: N \in HNatInfinite$

from *BseqD [OF X]* **obtain** K **where** $\forall n. norm (X n) \leq K$

by *fast*

then have $\forall N. hnorm (starfun X N) \leq star-of K$

by *transfer*

then have $hnorm (starfun X N) \leq star-of K$

by *simp*

also have $star-of K < star-of (K + 1)$

by *simp*

finally have $\exists x \in Reals. hnorm (starfun X N) < x$

by (*rule bexI*) *simp*

then show $starfun X N \in HFinite$

by (*simp add: HFinite-def*)

qed

The nonstandard definition implies the standard definition.

lemma *SReal-less-omega*: $r \in \mathbb{R} \implies r < \omega$
using *HInfinite-omega*
by (*simp add: HInfinite-def*) (*simp add: order-less-imp-le*)

lemma *NSBseq-Bseq*: $NSBseq\ X \implies Bseq\ X$
proof (*rule ccontr*)
let $?n = \lambda K. LEAST\ n. K < norm\ (X\ n)$
assume $NSBseq\ X$
then have $finite: (*f* X) ((*f* ?n) \omega) \in HFinite$
by (*rule NSBseqD2*)
assume $\neg Bseq\ X$
then have $\forall K > 0. \exists n. K < norm\ (X\ n)$
by (*simp add: Bseq-def linorder-not-le*)
then have $\forall K > 0. K < norm\ (X\ (?n\ K))$
by (*auto intro: LeastI-ex*)
then have $\forall K > 0. K < hnorm\ ((*f* X) ((*f* ?n) K))$
by *transfer*
then have $\omega < hnorm\ ((*f* X) ((*f* ?n) \omega))$
by *simp*
then have $\forall r \in \mathbb{R}. r < hnorm\ ((*f* X) ((*f* ?n) \omega))$
by (*simp add: order-less-trans [OF SReal-less-omega]*)
then have $(*f* X) ((*f* ?n) \omega) \in HInfinite$
by (*simp add: HInfinite-def*)
with *finite* **show** *False*
by (*simp add: HFinite-HInfinite-iff*)
qed

Equivalence of nonstandard and standard definitions for a bounded sequence.

lemma *Bseq-NSBseq-iff*: $Bseq\ X = NSBseq\ X$
by (*blast intro!: NSBseq-Bseq Bseq-NSBseq*)

A convergent sequence is bounded: Boundedness as a necessary condition for convergence. The nonstandard version has no existential, as usual.

lemma *NSconvergent-NSBseq*: $NSconvergent\ X \implies NSBseq\ X$
by (*simp add: NSconvergent-def NSBseq-def NSLIMSEQ-def*)
(blast intro: HFinite-star-of approx-sym approx-HFinite)

Standard Version: easily now proved using equivalence of NS and standard definitions.

lemma *convergent-Bseq*: $convergent\ X \implies Bseq\ X$
for $X :: nat \Rightarrow 'b::real-normed-vector$
by (*simp add: NSconvergent-NSBseq convergent-NSconvergent-iff Bseq-NSBseq-iff*)

10.3.1 Upper Bounds and Lubs of Bounded Sequences

lemma *NSBseq-isUb*: $NSBseq\ X \implies \exists U::real. isUb\ UNIV\ \{x. \exists n. X\ n = x\}\ U$
by (*simp add: Bseq-NSBseq-iff [symmetric] Bseq-isUb*)

lemma *NSBseq-isLub*: $NSBseq\ X \implies \exists U::real. isLub\ UNIV\ \{x. \exists n. X\ n = x\}$
 U
by (*simp add: Bseq-NSBseq-iff [symmetric] Bseq-isLub*)

10.3.2 A Bounded and Monotonic Sequence Converges

The best of both worlds: Easier to prove this result as a standard theorem and then use equivalence to ”transfer” it into the equivalent nonstandard form if needed!

lemma *Bmonoseq-NSLIMSEQ*: $\forall n \geq m. X\ n = X\ m \implies \exists L. X \longrightarrow_{NS} L$
by (*auto dest!: Bmonoseq-LIMSEQ simp add: LIMSEQ-NSLIMSEQ-iff*)

lemma *NSBseq-mono-NSconvergent*: $NSBseq\ X \implies \forall m. \forall n \geq m. X\ m \leq X\ n$
 $\implies NSconvergent\ X$
for $X :: nat \Rightarrow real$
by (*auto intro: Bseq-mono-convergent*
simp: convergent-NSconvergent-iff [symmetric] Bseq-NSBseq-iff [symmetric])

10.4 Cauchy Sequences

lemma *NSCauchyI*:
 $(\bigwedge M\ N. M \in HNatInfinite \implies N \in HNatInfinite \implies starfun\ X\ M \approx starfun\ X\ N) \implies NSCauchy\ X$
by (*simp add: NSCauchy-def*)

lemma *NSCauchyD*:
 $NSCauchy\ X \implies M \in HNatInfinite \implies N \in HNatInfinite \implies starfun\ X\ M \approx starfun\ X\ N$
by (*simp add: NSCauchy-def*)

10.4.1 Equivalence Between NS and Standard

lemma *Cauchy-NSCauchy*:
assumes $X: Cauchy\ X$
shows $NSCauchy\ X$
proof (*rule NSCauchyI*)
fix M
assume $M: M \in HNatInfinite$
fix N
assume $N: N \in HNatInfinite$
have $starfun\ X\ M - starfun\ X\ N \in Infinitesimal$
proof (*rule InfinitesimalI2*)
fix $r :: real$
assume $r: 0 < r$
from *CauchyD [OF X r]* **obtain** k **where** $\forall m \geq k. \forall n \geq k. norm\ (X\ m - X\ n) < r ..$
then have $\forall m \geq star-of\ k. \forall n \geq star-of\ k. hnorm\ (starfun\ X\ m - starfun\ X\ n) < star-of\ r$

```

    by transfer
  then show  $hnorm (starfun X M - starfun X N) < star-of r$ 
    using  $M N$  by (simp add: star-of-le-HNatInfinite)
qed
then show  $starfun X M \approx starfun X N$ 
  by (simp only: approx-def)
qed

lemma NSCauchy-Cauchy:
  assumes  $X: NSCauchy X$ 
  shows  $Cauchy X$ 
proof (rule CauchyI)
  fix  $r :: real$ 
  assume  $r: 0 < r$ 
  have  $\exists k. \forall m \geq k. \forall n \geq k. hnorm (starfun X m - starfun X n) < star-of r$ 
  proof (intro exI allI impI)
    fix  $M$ 
    assume  $whn \leq M$ 
    with  $HNatInfinite-whn$  have  $M: M \in HNatInfinite$ 
      by (rule HNatInfinite-upward-closed)
    fix  $N$ 
    assume  $whn \leq N$ 
    with  $HNatInfinite-whn$  have  $N: N \in HNatInfinite$ 
      by (rule HNatInfinite-upward-closed)
    from  $X M N$  have  $starfun X M \approx starfun X N$ 
      by (rule NSCauchyD)
    then have  $starfun X M - starfun X N \in Infinitesimal$ 
      by (simp only: approx-def)
    then show  $hnorm (starfun X M - starfun X N) < star-of r$ 
      using  $r$  by (rule InfinitesimalD2)
  qed
  then show  $\exists k. \forall m \geq k. \forall n \geq k. norm (X m - X n) < r$ 
    by transfer
qed

```

```

theorem NSCauchy-Cauchy-iff:  $NSCauchy X = Cauchy X$ 
  by (blast intro!: NSCauchy-Cauchy Cauchy-NSCauchy)

```

10.4.2 Cauchy Sequences are Bounded

A Cauchy sequence is bounded – nonstandard version.

```

lemma NSCauchy-NSBseq:  $NSCauchy X \implies NSBseq X$ 
  by (simp add: Cauchy-Bseq Bseq-NSBseq-iff [symmetric] NSCauchy-Cauchy-iff)

```

10.4.3 Cauchy Sequences are Convergent

Equivalence of Cauchy criterion and convergence: We will prove this using our NS formulation which provides a much easier proof than using the stan-

dard definition. We do not need to use properties of subsequences such as boundedness, monotonicity etc... Compare with Harrison’s corresponding proof in HOL which is much longer and more complicated. Of course, we do not have problems which he encountered with guessing the right instantiations for his ‘epsilon-delta’ proof(s) in this case since the NS formulations do not involve existential quantifiers.

lemma *NSconvergent-NSCauchy*: $NSconvergent\ X \implies NSCauchy\ X$
by (*simp add: NSconvergent-def NSLIMSEQ-def NSCauchy-def*) (*auto intro: approx-trans2*)

lemma *real-NSCauchy-NSconvergent*: $NSCauchy\ X \implies NSconvergent\ X$
for $X :: nat \Rightarrow real$
apply (*simp add: NSconvergent-def NSLIMSEQ-def*)
apply (*frule NSCauchy-NSBseq*)
apply (*simp add: NSBseq-def NSCauchy-def*)
apply (*drule HNatInfinite-whn [THEN [2] bspec]*)
apply (*drule HNatInfinite-whn [THEN [2] bspec]*)
apply (*auto dest!: st-part-Ex simp add: SReal-iff*)
apply (*blast intro: approx-trans3*)
done

lemma *NSCauchy-NSconvergent*: $NSCauchy\ X \implies NSconvergent\ X$
for $X :: nat \Rightarrow 'a::banach$
apply (*drule NSCauchy-Cauchy [THEN Cauchy-convergent]*)
apply (*erule convergent-NSconvergent-iff [THEN iffD1]*)
done

lemma *NSCauchy-NSconvergent-iff*: $NSCauchy\ X = NSconvergent\ X$
for $X :: nat \Rightarrow 'a::banach$
by (*fast intro: NSCauchy-NSconvergent NSconvergent-NSCauchy*)

10.5 Power Sequences

The sequence x^n tends to 0 if $(0::'a) \leq x$ and $x < (1::'a)$. Proof will use (NS) Cauchy equivalence for convergence and also fact that bounded and monotonic sequence converges.

We now use NS criterion to bring proof of theorem through.

lemma *NSLIMSEQ-realpow-zero*: $0 \leq x \implies x < 1 \implies (\lambda n. x ^ n) \longrightarrow_{NS} 0$
for $x :: real$
apply (*simp add: NSLIMSEQ-def*)
apply (*auto dest!: convergent-realpow simp add: convergent-NSconvergent-iff*)
apply (*frule NSconvergentD*)
apply (*auto simp add: NSLIMSEQ-def NSCauchy-NSconvergent-iff [symmetric]*
NSCauchy-def starfun-pow)
apply (*frule HNatInfinite-add-one*)
apply (*drule bspec, assumption*)
apply (*drule bspec, assumption*)

```

apply (drule-tac x = N + 1 in bspec, assumption)
apply (simp add: hyperpow-add)
apply (drule approx-mult-subst-star-of, assumption)
apply (drule approx-trans3, assumption)
apply (auto simp del: star-of-mult simp add: star-of-mult [symmetric])
done

```

```

lemma NSLIMSEQ-rabs-realpow-zero: |c| < 1  $\implies$  ( $\lambda n. |c|^n$ )  $\longrightarrow_{NS}$  0
for c :: real
by (simp add: LIMSEQ-rabs-realpow-zero LIMSEQ-NSLIMSEQ-iff [symmetric])

```

```

lemma NSLIMSEQ-rabs-realpow-zero2: |c| < 1  $\implies$  ( $\lambda n. c^n$ )  $\longrightarrow_{NS}$  0
for c :: real
by (simp add: LIMSEQ-rabs-realpow-zero2 LIMSEQ-NSLIMSEQ-iff [symmetric])

```

end

11 Finite Summation and Infinite Series for Hyperreals

```

theory HSeries
imports HSEQ
begin

```

```

definition sumhr :: hypnat  $\times$  hypnat  $\times$  (nat  $\Rightarrow$  real)  $\Rightarrow$  hypreal
where sumhr = ( $\lambda(M,N,f). \text{starfun2 } (\lambda m n. \text{sum } f \{m..<n\}) M N$ )

```

```

definition NSsums :: (nat  $\Rightarrow$  real)  $\Rightarrow$  real  $\Rightarrow$  bool (infixr NSsums 80)
where f NSsums s = ( $\lambda n. \text{sum } f \{..<n\}$ )  $\longrightarrow_{NS}$  s

```

```

definition NSsummable :: (nat  $\Rightarrow$  real)  $\Rightarrow$  bool
where NSsummable f  $\longleftrightarrow$  ( $\exists s. f \text{ NSsums } s$ )

```

```

definition NSsuminf :: (nat  $\Rightarrow$  real)  $\Rightarrow$  real
where NSsuminf f = (THE s. f NSsums s)

```

```

lemma sumhr-app: sumhr (M, N, f) = (*f2* ( $\lambda m n. \text{sum } f \{m..<n\}$ )) M N
by (simp add: sumhr-def)

```

Base case in definition of *sumr*.

```

lemma sumhr-zero [simp]:  $\bigwedge m. \text{sumhr } (m, 0, f) = 0$ 
unfolding sumhr-app by transfer simp

```

Recursive case in definition of *sumr*.

```

lemma sumhr-if:

```

$\bigwedge m n. \text{sumhr } (m, n + 1, f) = (\text{if } n + 1 \leq m \text{ then } 0 \text{ else } \text{sumhr } (m, n, f) + (*f* f) n)$

unfolding *sumhr-app* **by** *transfer simp*

lemma *sumhr-Suc-zero* [*simp*]: $\bigwedge n. \text{sumhr } (n + 1, n, f) = 0$

unfolding *sumhr-app* **by** *transfer simp*

lemma *sumhr-eq-bounds* [*simp*]: $\bigwedge n. \text{sumhr } (n, n, f) = 0$

unfolding *sumhr-app* **by** *transfer simp*

lemma *sumhr-Suc* [*simp*]: $\bigwedge m. \text{sumhr } (m, m + 1, f) = (*f* f) m$

unfolding *sumhr-app* **by** *transfer simp*

lemma *sumhr-add-lbound-zero* [*simp*]: $\bigwedge k m. \text{sumhr } (m + k, k, f) = 0$

unfolding *sumhr-app* **by** *transfer simp*

lemma *sumhr-add*: $\bigwedge m n. \text{sumhr } (m, n, f) + \text{sumhr } (m, n, g) = \text{sumhr } (m, n, \lambda i. f i + g i)$

unfolding *sumhr-app* **by** *transfer (rule sum.distrib [symmetric])*

lemma *sumhr-mult*: $\bigwedge m n. \text{hypreal-of-real } r * \text{sumhr } (m, n, f) = \text{sumhr } (m, n, \lambda n. r * f n)$

unfolding *sumhr-app* **by** *transfer (rule sum-distrib-left)*

lemma *sumhr-split-add*: $\bigwedge n p. n < p \implies \text{sumhr } (0, n, f) + \text{sumhr } (n, p, f) = \text{sumhr } (0, p, f)$

unfolding *sumhr-app* **by** *transfer (simp add: sum-add-nat-ivl)*

lemma *sumhr-split-diff*: $n < p \implies \text{sumhr } (0, p, f) - \text{sumhr } (0, n, f) = \text{sumhr } (n, p, f)$

by (*drule sumhr-split-add [symmetric, where f = f]*) *simp*

lemma *sumhr-hrabs*: $\bigwedge m n. |\text{sumhr } (m, n, f)| \leq \text{sumhr } (m, n, \lambda i. |f i|)$

unfolding *sumhr-app* **by** *transfer (rule sum-abs)*

Other general version also needed.

lemma *sumhr-fun-hypnat-eq*:

$(\forall r. m \leq r \wedge r < n \implies f r = g r) \implies$
 $\text{sumhr } (\text{hypnat-of-nat } m, \text{hypnat-of-nat } n, f) =$
 $\text{sumhr } (\text{hypnat-of-nat } m, \text{hypnat-of-nat } n, g)$

unfolding *sumhr-app* **by** *transfer simp*

lemma *sumhr-const*: $\bigwedge n. \text{sumhr } (0, n, \lambda i. r) = \text{hypreal-of-hypnat } n * \text{hypreal-of-real } r$

unfolding *sumhr-app* **by** *transfer simp*

lemma *sumhr-less-bounds-zero* [*simp*]: $\bigwedge m n. n < m \implies \text{sumhr } (m, n, f) = 0$

unfolding *sumhr-app* **by** *transfer simp*

lemma *sumhr-minus*: $\bigwedge m n. \text{sumhr } (m, n, \lambda i. - f i) = - \text{sumhr } (m, n, f)$
unfolding *sumhr-app* **by** *transfer* (rule *sum-negf*)

lemma *sumhr-shift-bounds*:
 $\bigwedge m n. \text{sumhr } (m + \text{hypnat-of-nat } k, n + \text{hypnat-of-nat } k, f) =$
 $\text{sumhr } (m, n, \lambda i. f (i + k))$
unfolding *sumhr-app* **by** *transfer* (rule *sum-shift-bounds-nat-ivl*)

11.1 Nonstandard Sums

Infinite sums are obtained by summing to some infinite hypernatural (such as *whn*).

lemma *sumhr-hypreal-of-hypnat-omega*: $\text{sumhr } (0, \text{whn}, \lambda i. 1) = \text{hypreal-of-hypnat } \text{whn}$
by (*simp add: sumhr-const*)

lemma *sumhr-hypreal-omega-minus-one*: $\text{sumhr}(0, \text{whn}, \lambda i. 1) = \omega - 1$
apply (*simp add: sumhr-const*)

apply (*unfold star-class-defs omega-def hypnat-omega-def of-hypnat-def star-of-def*)
apply (*simp add: starfun-star-n starfun2-star-n*)
done

lemma *sumhr-minus-one-realpow-zero* [*simp*]: $\bigwedge N. \text{sumhr } (0, N + N, \lambda i. (-1)^{(i + 1)}) = 0$
unfolding *sumhr-app*
apply *transfer*
apply (*simp del: power-Suc add: mult-2 [symmetric]*)
apply (*induct-tac N*)
apply *simp-all*
done

lemma *sumhr-interval-const*:
 $(\forall n. m \leq \text{Suc } n \longrightarrow f n = r) \wedge m \leq na \implies$
 $\text{sumhr } (\text{hypnat-of-nat } m, \text{hypnat-of-nat } na, f) = \text{hypreal-of-nat } (na - m) * \text{hypreal-of-real } r$
unfolding *sumhr-app* **by** *transfer simp*

lemma *starfunNat-sumr*: $\bigwedge N. (*f* (\lambda n. \text{sum } f \{0..<n\})) N = \text{sumhr } (0, N, f)$
unfolding *sumhr-app* **by** *transfer* (rule *refl*)

lemma *sumhr-hrabs-approx* [*simp*]: $\text{sumhr } (0, M, f) \approx \text{sumhr } (0, N, f) \implies$
 $|\text{sumhr } (M, N, f)| \approx 0$
using *linorder-less-linear* [**where** $x = M$ **and** $y = N$]
apply *auto*
apply (*drule approx-sym [THEN approx-minus-iff [THEN iffD1]]*)
apply (*auto dest: approx-hrabs simp add: sumhr-split-diff*)
done

11.2 Infinite sums: Standard and NS theorems

lemma *sums-NSsums-iff*: $f \text{ sums } l \iff f \text{ NSsums } l$

by (*simp add: sums-def NSsums-def LIMSEQ-NSLIMSEQ-iff*)

lemma *summable-NSsummable-iff*: $\text{summable } f \iff \text{NSsummable } f$

by (*simp add: summable-def NSsummable-def sums-NSsums-iff*)

lemma *suminf-NSsuminf-iff*: $\text{suminf } f = \text{NSsuminf } f$

by (*simp add: suminf-def NSsuminf-def sums-NSsums-iff*)

lemma *NSsums-NSsummable*: $f \text{ NSsums } l \implies \text{NSsummable } f$

unfolding *NSsums-def NSsummable-def* **by** *blast*

lemma *NSsummable-NSsums*: $\text{NSsummable } f \implies f \text{ NSsums } (\text{NSsuminf } f)$

unfolding *NSsummable-def NSsuminf-def NSsums-def*

by (*blast intro: theI NSLIMSEQ-unique*)

lemma *NSsums-unique*: $f \text{ NSsums } s \implies s = \text{NSsuminf } f$

by (*simp add: suminf-NSsuminf-iff [symmetric] sums-NSsums-iff sums-unique*)

lemma *NSseries-zero*: $\forall m. n \leq \text{Suc } m \longrightarrow f m = 0 \implies f \text{ NSsums } (\text{sum } f \{..<n\})$

by (*auto simp add: sums-NSsums-iff [symmetric] not-le[symmetric] intro!: sums-finite*)

lemma *NSsummable-NSCauchy*:

$\text{NSsummable } f \iff (\forall M \in \text{HNatInfinite}. \forall N \in \text{HNatInfinite}. |\text{sumhr } (M, N, f)| \approx 0)$

apply (*auto simp add: summable-NSsummable-iff [symmetric]*)

summable-iff-convergent convergent-NSconvergent-iff atLeast0LessThan[symmetric]

NSCauchy-NSconvergent-iff [symmetric] NSCauchy-def starfunNat-sumr)

apply (*cut-tac x = M and y = N in linorder-less-linear*)

apply *auto*

apply (*rule approx-minus-iff [THEN iffD2, THEN approx-sym]*)

apply (*rule-tac [2] approx-minus-iff [THEN iffD2]*)

apply (*auto dest: approx-hrabs-zero-cancel simp: sumhr-split-diff atLeast0LessThan[symmetric]*)

done

Terms of a convergent series tend to zero.

lemma *NSsummable-NSLIMSEQ-zero*: $\text{NSsummable } f \implies f \longrightarrow_{NS} 0$

apply (*auto simp add: NSLIMSEQ-def NSsummable-NSCauchy*)

apply (*drule bspec*)

apply *auto*

apply (*drule-tac x = N + 1 in bspec*)

apply (*auto intro: HNatInfinite-add-one approx-hrabs-zero-cancel*)

done

Nonstandard comparison test.

lemma *NSsummable-comparison-test*: $\exists N. \forall n. N \leq n \longrightarrow |f n| \leq g n \implies \text{NSsummable } g \implies \text{NSsummable } f$

apply (*fold summable-NSsummable-iff*)

apply (*rule summable-comparison-test, simp, assumption*)
done

lemma *NSsummable-rabs-comparison-test*:
 $\exists N. \forall n. N \leq n \longrightarrow |f\ n| \leq g\ n \Longrightarrow NSsummable\ g \Longrightarrow NSsummable\ (\lambda k. |f\ k|)$
by (*rule NSsummable-comparison-test*) *auto*

end

12 Limits and Continuity (Nonstandard)

theory *HLim*
imports *Star*
abbrevs $----> = -\square\rightarrow_{NS}$
begin

Nonstandard Definitions.

definition *NSLIM* :: (*'a::real-normed-vector* \Rightarrow *'b::real-normed-vector*) \Rightarrow *'a* \Rightarrow *'b* \Rightarrow *bool*
 $(((-)/ -(-)/\rightarrow_{NS} (-)) [60, 0, 60] 60)$
where $f -a\rightarrow_{NS} L \longleftrightarrow (\forall x. x \neq \text{star-of } a \wedge x \approx \text{star-of } a \longrightarrow (*f* f) x \approx \text{star-of } L)$

definition *isNSCont* :: (*'a::real-normed-vector* \Rightarrow *'b::real-normed-vector*) \Rightarrow *'a* \Rightarrow *bool*
where — NS definition dispenses with limit notions
 $isNSCont\ f\ a \longleftrightarrow (\forall y. y \approx \text{star-of } a \longrightarrow (*f* f) y \approx \text{star-of } (f\ a))$

definition *isNSUCont* :: (*'a::real-normed-vector* \Rightarrow *'b::real-normed-vector*) \Rightarrow *bool*
where $isNSUCont\ f \longleftrightarrow (\forall x\ y. x \approx y \longrightarrow (*f* f) x \approx (*f* f) y)$

12.1 Limits of Functions

lemma *NSLIM-I*: $(\bigwedge x. x \neq \text{star-of } a \Longrightarrow x \approx \text{star-of } a \Longrightarrow \text{starfun } f\ x \approx \text{star-of } L) \Longrightarrow f -a\rightarrow_{NS} L$
by (*simp add: NSLIM-def*)

lemma *NSLIM-D*: $f -a\rightarrow_{NS} L \Longrightarrow x \neq \text{star-of } a \Longrightarrow x \approx \text{star-of } a \Longrightarrow \text{starfun } f\ x \approx \text{star-of } L$
by (*simp add: NSLIM-def*)

Proving properties of limits using nonstandard definition. The properties hold for standard limits as well!

lemma *NSLIM-mult*: $f -x\rightarrow_{NS} l \Longrightarrow g -x\rightarrow_{NS} m \Longrightarrow (\lambda x. f\ x * g\ x) -x\rightarrow_{NS} (l * m)$
for $l\ m :: 'a::\text{real-normed-algebra}$
by (*auto simp add: NSLIM-def intro!: approx-mult-HFinite*)

lemma *starfun-scaleR* [*simp*]: $\text{starfun } (\lambda x. f x *_R g x) = (\lambda x. \text{scaleHR } (\text{starfun } f x) (\text{starfun } g x))$

by *transfer* (*rule refl*)

lemma *NSLIM-scaleR*: $f -x \rightarrow_{NS} l \implies g -x \rightarrow_{NS} m \implies (\lambda x. f x *_R g x) -x \rightarrow_{NS} (l *_R m)$

by (*auto simp add: NSLIM-def intro!: approx-scaleR-HFfinite*)

lemma *NSLIM-add*: $f -x \rightarrow_{NS} l \implies g -x \rightarrow_{NS} m \implies (\lambda x. f x + g x) -x \rightarrow_{NS} (l + m)$

by (*auto simp add: NSLIM-def intro!: approx-add*)

lemma *NSLIM-const* [*simp*]: $(\lambda x. k) -x \rightarrow_{NS} k$

by (*simp add: NSLIM-def*)

lemma *NSLIM-minus*: $f -a \rightarrow_{NS} L \implies (\lambda x. - f x) -a \rightarrow_{NS} -L$

by (*simp add: NSLIM-def*)

lemma *NSLIM-diff*: $f -x \rightarrow_{NS} l \implies g -x \rightarrow_{NS} m \implies (\lambda x. f x - g x) -x \rightarrow_{NS} (l - m)$

by (*simp only: NSLIM-add NSLIM-minus diff-conv-add-uminus*)

lemma *NSLIM-add-minus*: $f -x \rightarrow_{NS} l \implies g -x \rightarrow_{NS} m \implies (\lambda x. f x + - g x) -x \rightarrow_{NS} (l + -m)$

by (*simp only: NSLIM-add NSLIM-minus*)

lemma *NSLIM-inverse*: $f -a \rightarrow_{NS} L \implies L \neq 0 \implies (\lambda x. \text{inverse } (f x)) -a \rightarrow_{NS} (\text{inverse } L)$

for $L :: 'a::\text{real-normed-div-algebra}$

apply (*simp add: NSLIM-def, clarify*)

apply (*drule spec*)

apply (*auto simp add: star-of-approx-inverse*)

done

lemma *NSLIM-zero*:

assumes $f: f -a \rightarrow_{NS} l$

shows $(\lambda x. f(x) - l) -a \rightarrow_{NS} 0$

proof –

have $(\lambda x. f x - l) -a \rightarrow_{NS} l - l$

by (*rule NSLIM-diff [OF f NSLIM-const]*)

then show *?thesis* **by** *simp*

qed

lemma *NSLIM-zero-cancel*: $(\lambda x. f x - l) -x \rightarrow_{NS} 0 \implies f -x \rightarrow_{NS} l$

apply (*drule-tac g = $\lambda x. l$ and $m = l$ in NSLIM-add*)

apply (*auto simp add: add.assoc*)

done

lemma *NSLIM-const-not-eq*: $k \neq L \implies \neg (\lambda x. k) -a \rightarrow_{NS} L$

```

for  $a :: 'a::\text{real-normed-algebra-1}$ 
apply (simp add: NSLIM-def)
apply (rule-tac  $x=\text{star-of } a + \text{of-hypreal } \varepsilon$  in  $exI$ )
apply (simp add: hypreal-epsilon-not-zero approx-def)
done

```

```

lemma NSLIM-not-zero:  $k \neq 0 \implies \neg (\lambda x. k) -a \rightarrow_{NS} 0$ 
for  $a :: 'a::\text{real-normed-algebra-1}$ 
by (rule NSLIM-const-not-eq)

```

```

lemma NSLIM-const-eq:  $(\lambda x. k) -a \rightarrow_{NS} L \implies k = L$ 
for  $a :: 'a::\text{real-normed-algebra-1}$ 
by (rule ccontr) (blast dest: NSLIM-const-not-eq)

```

```

lemma NSLIM-unique:  $f -a \rightarrow_{NS} L \implies f -a \rightarrow_{NS} M \implies L = M$ 
for  $a :: 'a::\text{real-normed-algebra-1}$ 
by (drule (1) NSLIM-diff) (auto dest!: NSLIM-const-eq)

```

```

lemma NSLIM-mult-zero:  $f -x \rightarrow_{NS} 0 \implies g -x \rightarrow_{NS} 0 \implies (\lambda x. f x * g x) -x \rightarrow_{NS} 0$ 
for  $f g :: 'a::\text{real-normed-vector} \Rightarrow 'b::\text{real-normed-algebra}$ 
by (drule NSLIM-mult) auto

```

```

lemma NSLIM-self:  $(\lambda x. x) -a \rightarrow_{NS} a$ 
by (simp add: NSLIM-def)

```

12.1.1 Equivalence of filterlim and NSLIM

```

lemma LIM-NSLIM:
  assumes  $f: f -a \rightarrow L$ 
  shows  $f -a \rightarrow_{NS} L$ 

```

```

proof (rule NSLIM-I)

```

```

  fix  $x$ 

```

```

  assume  $neg: x \neq \text{star-of } a$ 

```

```

  assume  $approx: x \approx \text{star-of } a$ 

```

```

  have  $\text{starfun } f x - \text{star-of } L \in \text{Infinitesimal}$ 

```

```

  proof (rule InfinitesimalI2)

```

```

    fix  $r :: \text{real}$ 

```

```

    assume  $r: 0 < r$ 

```

```

    from LIM-D [OF  $f r$ ] obtain  $s$ 

```

```

    where  $s: 0 < s$  and  $\text{less-r}: \bigwedge x. x \neq a \implies \text{norm } (x - a) < s \implies \text{norm } (f x - L) < r$ 

```

```

    by fast

```

```

    from less-r have less-r':

```

```

       $\bigwedge x. x \neq \text{star-of } a \implies \text{hnorm } (x - \text{star-of } a) < \text{star-of } s \implies$ 

```

```

       $\text{hnorm } (\text{starfun } f x - \text{star-of } L) < \text{star-of } r$ 

```

```

    by transfer

```

```

    from approx have  $x - \text{star-of } a \in \text{Infinitesimal}$ 

```

```

    by (simp only: approx-def)

```

then have $hnorm (x - star-of a) < star-of s$
using s **by** (rule *InfinitesimalD2*)
with neq **show** $hnorm (starfun f x - star-of L) < star-of r$
by (rule *less-r'*)
qed
then show $starfun f x \approx star-of L$
by (unfold *approx-def*)
qed

lemma *NSLIM-LIM*:

assumes $f: f -a \rightarrow_{NS} L$
shows $f -a \rightarrow L$
proof (rule *LIM-I*)
fix $r :: real$
assume $r: 0 < r$
have $\exists s > 0. \forall x. x \neq star-of a \wedge hnorm (x - star-of a) < s \longrightarrow$
 $hnorm (starfun f x - star-of L) < star-of r$
proof (rule *exI, safe*)
show $0 < \varepsilon$
by (rule *hypreal-epsilon-gt-zero*)
next
fix x
assume $neq: x \neq star-of a$
assume $hnorm (x - star-of a) < \varepsilon$
with *Infinitesimal-epsilon* **have** $x - star-of a \in Infinitesimal$
by (rule *hnorm-less-Infinitesimal*)
then have $x \approx star-of a$
by (unfold *approx-def*)
with $f neq$ **have** $starfun f x \approx star-of L$
by (rule *NSLIM-D*)
then have $starfun f x - star-of L \in Infinitesimal$
by (unfold *approx-def*)
then show $hnorm (starfun f x - star-of L) < star-of r$
using r **by** (rule *InfinitesimalD2*)
qed
then show $\exists s > 0. \forall x. x \neq a \wedge norm (x - a) < s \longrightarrow norm (f x - L) < r$
by *transfer*
qed

theorem *LIM-NSLIM-iff*: $f -x \rightarrow L \longleftrightarrow f -x \rightarrow_{NS} L$
by (blast *intro: LIM-NSLIM NSLIM-LIM*)

12.2 Continuity

lemma *isNSContD*: $isNSCont f a \Longrightarrow y \approx star-of a \Longrightarrow (*f* f) y \approx star-of (f a)$
by (*simp add: isNSCont-def*)

lemma *isNSCont-NSLIM*: $isNSCont f a \Longrightarrow f -a \rightarrow_{NS} (f a)$

by (simp add: isNSCont-def NSLIM-def)

lemma *NSLIM-isNSCont*: $f -a \rightarrow_{NS} (f a) \implies isNSCont f a$
apply (auto simp add: isNSCont-def NSLIM-def)
apply (case-tac $y = star-of a$)
apply auto
done

NS continuity can be defined using NS Limit in similar fashion to standard definition of continuity.

lemma *isNSCont-NSLIM-iff*: $isNSCont f a \longleftrightarrow f -a \rightarrow_{NS} (f a)$
by (blast intro: isNSCont-NSLIM NSLIM-isNSCont)

Hence, NS continuity can be given in terms of standard limit.

lemma *isNSCont-LIM-iff*: $(isNSCont f a) = (f -a \rightarrow (f a))$
by (simp add: LIM-NSLIM-iff isNSCont-NSLIM-iff)

Moreover, it’s trivial now that NS continuity is equivalent to standard continuity.

lemma *isNSCont-isCont-iff*: $isNSCont f a \longleftrightarrow isCont f a$
by (simp add: isCont-def) (rule isNSCont-LIM-iff)

Standard continuity \implies NS continuity.

lemma *isCont-isNSCont*: $isCont f a \implies isNSCont f a$
by (erule isNSCont-isCont-iff [THEN iffD2])

NS continuity \implies Standard continuity.

lemma *isNSCont-isCont*: $isNSCont f a \implies isCont f a$
by (erule isNSCont-isCont-iff [THEN iffD1])

Alternative definition of continuity.

Prove equivalence between NS limits – seems easier than using standard definition.

lemma *NSLIM-h-iff*: $f -a \rightarrow_{NS} L \longleftrightarrow (\lambda h. f (a + h)) -0 \rightarrow_{NS} L$
apply (simp add: NSLIM-def, auto)
apply (drule-tac $x = star-of a + x$ **in** spec)
apply (drule-tac [2] $x = - star-of a + x$ **in** spec, safe, simp)
apply (erule mem-infmal-iff [THEN iffD2, THEN Infinitesimal-add-approx-self [THEN approx-sym]])
apply (erule-tac [3] approx-minus-iff2 [THEN iffD1])
prefer 2 **apply** (simp add: add commute)
apply (rule-tac $x = x$ **in** star-cases)
apply (rule-tac [2] $x = x$ **in** star-cases)
apply (auto simp add: starfun star-of-def star-n-minus star-n-add add.assoc star-n-zero-num)
done

lemma *NSLIM-isCont-iff*: $f - a \rightarrow_{NS} f a \iff (\lambda h. f (a + h)) - 0 \rightarrow_{NS} f a$
by (*fact NSLIM-h-iff*)

lemma *isNSCont-minus*: $isNSCont f a \implies isNSCont (\lambda x. - f x) a$
by (*simp add: isNSCont-def*)

lemma *isNSCont-inverse*: $isNSCont f x \implies f x \neq 0 \implies isNSCont (\lambda x. inverse (f x)) x$
for $f :: 'a::real-normed-vector \Rightarrow 'b::real-normed-div-algebra$
by (*auto intro: isCont-inverse simp add: isNSCont-isCont-iff*)

lemma *isNSCont-const [simp]*: $isNSCont (\lambda x. k) a$
by (*simp add: isNSCont-def*)

lemma *isNSCont-abs [simp]*: $isNSCont abs a$
for $a :: real$
by (*auto simp: isNSCont-def intro: approx-hrabs simp: starfun-rabs-hrabs*)

12.3 Uniform Continuity

lemma *isNSUContD*: $isNSUCont f \implies x \approx y \implies (*f* f) x \approx (*f* f) y$
by (*simp add: isNSUCont-def*)

lemma *isUCont-isNSUCont*:
fixes $f :: 'a::real-normed-vector \Rightarrow 'b::real-normed-vector$
assumes $f: isUCont f$
shows $isNSUCont f$
unfolding *isNSUCont-def*

proof *safe*

fix $x y :: 'a star$

assume *approx*: $x \approx y$

have *starfun* $f x - starfun f y \in Infinitesimal$

proof (*rule InfinitesimalI2*)

fix $r :: real$

assume $r: 0 < r$

with f **obtain** s **where** $0 < s$

and *less-r*: $\bigwedge x y. norm (x - y) < s \implies norm (f x - f y) < r$

by (*auto simp add: isUCont-def dist-norm*)

from *less-r* **have** *less-r'*:

$\bigwedge x y. hnorm (x - y) < star-of s \implies hnorm (starfun f x - starfun f y) < star-of r$

by *transfer*

from *approx* **have** $x - y \in Infinitesimal$

by (*unfold approx-def*)

then **have** $hnorm (x - y) < star-of s$

using s **by** (*rule InfinitesimalD2*)

then **show** $hnorm (starfun f x - starfun f y) < star-of r$

by (*rule less-r'*)

qed

```

then show starfun f x  $\approx$  starfun f y
  by (unfold approx-def)
qed

```

```

lemma isNSUCont-isUCont:

```

```

  fixes f :: 'a::real-normed-vector  $\Rightarrow$  'b::real-normed-vector
  assumes f: isNSUCont f
  shows isUCont f
  unfolding isUCont-def dist-norm
proof safe
  fix r :: real
  assume r: 0 < r
  have  $\exists s > 0. \forall x y. \text{hnorm } (x - y) < s \longrightarrow \text{hnorm } (\text{starfun } f x - \text{starfun } f y) <$ 
    star-of r
  proof (rule exI, safe)
    show 0 <  $\varepsilon$ 
      by (rule hypreal-epsilon-gt-zero)
  next
    fix x y :: 'a star
    assume hnrm (x - y) <  $\varepsilon$ 
    with Infinitesimal-epsilon have x - y  $\in$  Infinitesimal
      by (rule hnrm-less-Infinitesimal)
    then have x  $\approx$  y
      by (unfold approx-def)
    with f have starfun f x  $\approx$  starfun f y
      by (simp add: isNSUCont-def)
    then have starfun f x - starfun f y  $\in$  Infinitesimal
      by (unfold approx-def)
    then show hnrm (starfun f x - starfun f y) < star-of r
      using r by (rule InfinitesimalD2)
  qed
  then show  $\exists s > 0. \forall x y. \text{norm } (x - y) < s \longrightarrow \text{norm } (f x - f y) < r$ 
    by transfer
qed

```

```

end

```

13 Differentiation (Nonstandard)

```

theory HDeriv
  imports HLim
begin

```

Nonstandard Definitions.

```

definition nsderiv :: ['a::real-normed-field  $\Rightarrow$  'a, 'a, 'a]  $\Rightarrow$  bool
  ((NSDERIV (-)/ (-) /> (-)) [1000, 1000, 60] 60)
  where NSDERIV f x /> D  $\longleftrightarrow$ 
    ( $\forall h \in$  Infinitesimal - {0}. (( $*$ f* f)(star-of x + h) - star-of (f x)) / h  $\approx$ 
    star-of D)

```

definition *NSdifferentiable* :: [*'a*::*real-normed-field* \Rightarrow *'a*, *'a*] \Rightarrow *bool*
 (**infixl** *NSdifferentiable* 60)
where *f NSdifferentiable x* \longleftrightarrow ($\exists D. NSDERIV f x \text{ :> } D$)

definition *increment* :: (*real* \Rightarrow *real*) \Rightarrow *real* \Rightarrow *hypreal* \Rightarrow *hypreal*
where *increment f x h* =
 (*SOME inc. f NSdifferentiable x* \wedge *inc* = (**f* f*) (*hypreal-of-real x + h*) -
hypreal-of-real (f x))

13.1 Derivatives

lemma *DERIV-NS-iff*: (*DERIV f x :> D*) \longleftrightarrow ($\lambda h. (f (x + h) - f x) / h$)
 $-0 \rightarrow_{NS} D$
by (*simp add: DERIV-def LIM-NSLIM-iff*)

lemma *NS-DERIV-D*: *DERIV f x :> D* \Longrightarrow ($\lambda h. (f (x + h) - f x) / h$) $-0 \rightarrow_{NS} D$
by (*simp add: DERIV-def LIM-NSLIM-iff*)

lemma *hnorm-of-hypreal*: $\bigwedge r. hnorm ((*f* of-real) r :: 'a :: real-normed-div-algebra star) = |r|$
by *transfer (rule norm-of-real)*

lemma *Infinitesimal-of-hypreal*:
x \in *Infinitesimal* \Longrightarrow ((**f* of-real*) *x* :: *'a*::*real-normed-div-algebra star*) \in *Infinitesimal*
apply (*rule InfinitesimalI2*)
apply (*drule (1) InfinitesimalD2*)
apply (*simp add: hnorm-of-hypreal*)
done

lemma *of-hypreal-eq-0-iff*: $\bigwedge x. ((*f* of-real) x = (0 :: 'a :: real-algebra-1 star)) =$
 $(x = 0)$
by *transfer (rule of-real-eq-0-iff)*

lemma *NSDeriv-unique*: *NSDERIV f x :> D* \Longrightarrow *NSDERIV f x :> E* \Longrightarrow *D = E*
apply (*subgoal-tac (*f* of-real) $\varepsilon \in$ Infinitesimal - {0 :: 'a star}*)
apply (*simp only: nsderiv-def*)
apply (*drule (1) bspec*)
apply (*drule (1) approx-trans3*)
apply *simp*
apply (*simp add: Infinitesimal-of-hypreal*)
apply (*simp add: of-hypreal-eq-0-iff hypreal-epsilon-not-zero*)
done

First *NSDERIV* in terms of *NSLIM*.

First equivalence.

lemma *NSDERIV-NSLIM-iff*: (*NSDERIV f x :> D*) \longleftrightarrow ($\lambda h. (f (x + h) - f x)$)


```

/ h) - 0 →NS D
apply (auto simp add: nsderiv-def NSLIM-def)
apply (drule-tac x = xa in bspec)
apply (rule-tac [3] ccontr)
apply (drule-tac [3] x = h in spec)
apply (auto simp add: mem-infmal-iff starfun-lambda-cancel)
done

```

Second equivalence.

lemma *NSDERIV-NSLIM-iff2*: $(NSDERIV f x :> D) \longleftrightarrow (\lambda z. (f z - f x) / (z - x)) -x \rightarrow_{NS} D$
by (simp add: NSDERIV-NSLIM-iff DERIV-LIM-iff LIM-NSLIM-iff [symmetric])

While we’re at it!

lemma *NSDERIV-iff2*:
 $(NSDERIV f x :> D) \longleftrightarrow$
 $(\forall w. w \neq \text{star-of } x \wedge w \approx \text{star-of } x \longrightarrow (*f* (\lambda z. (f z - f x) / (z - x))) w \approx \text{star-of } D)$
by (simp add: NSDERIV-NSLIM-iff2 NSLIM-def)

lemma *hypreal-not-eq-minus-iff*: $x \neq a \longleftrightarrow x - a \neq (0::'a::\text{ab-group-add})$
by auto

lemma *NSDERIVD5*:
 $(NSDERIV f x :> D) \implies$
 $(\forall u. u \approx \text{hypreal-of-real } x \longrightarrow$
 $(*f* (\lambda z. f z - f x)) u \approx \text{hypreal-of-real } D * (u - \text{hypreal-of-real } x))$
apply (auto simp add: NSDERIV-iff2)
apply (case-tac u = hypreal-of-real x, auto)
apply (drule-tac x = u **in** spec, auto)
apply (drule-tac c = u - hypreal-of-real x **and** b = hypreal-of-real D **in** approx-mult1)
apply (drule-tac [!] hypreal-not-eq-minus-iff [THEN iffD1])
apply (subgoal-tac [2] (*f* (\lambda z. z - x)) u ≠ (0::hypreal))
apply (auto simp: approx-minus-iff [THEN iffD1, THEN mem-infmal-iff [THEN iffD2]])
Infinitesimal-subset-HFfinite [THEN subsetD])
done

lemma *NSDERIVD4*:
 $(NSDERIV f x :> D) \implies$
 $(\forall h \in \text{Infinitesimal.}$
 $(*f* f)(\text{hypreal-of-real } x + h) - \text{hypreal-of-real } (f x) \approx \text{hypreal-of-real } D * h)$
apply (auto simp add: nsderiv-def)
apply (case-tac h = 0)
apply auto
apply (drule-tac x = h **in** bspec)
apply (drule-tac [2] c = h **in** approx-mult1)

```

apply (auto intro: Infinitesimal-subset-HFinite [THEN subsetD])
done

```

lemma NSDERIVD3:

```

(NSDERIV f x :> D) ==>
  ∀ h ∈ Infinitesimal - {0}.
    (( *f* f) (hypreal-of-real x + h) - hypreal-of-real (f x)) ≈ hypreal-of-real D
* h
apply (auto simp add: nsderiv-def)
apply (rule ccontr, drule-tac x = h in bspec)
apply (drule-tac [2] c = h in approx-mult1)
apply (auto intro: Infinitesimal-subset-HFinite [THEN subsetD] simp add:
mult.assoc)
done

```

Differentiability implies continuity nice and simple ”algebraic” proof.

```

lemma NSDERIV-isNSCont: NSDERIV f x :> D ==> isNSCont f x
apply (auto simp add: nsderiv-def isNSCont-NSLIM-iff NSLIM-def)
apply (drule approx-minus-iff [THEN iffD1])
apply (drule hypreal-not-eq-minus-iff [THEN iffD1])
apply (drule-tac x = xa - star-of x in bspec)
prefer 2 apply (simp add: add.assoc [symmetric])
apply (auto simp add: mem-infmal-iff [symmetric] add.commute)
apply (drule-tac c = xa - star-of x in approx-mult1)
apply (auto intro: Infinitesimal-subset-HFinite [THEN subsetD] simp add:
mult.assoc)
apply (drule-tac x3=D in
  HFinite-star-of [THEN [2] Infinitesimal-HFinite-mult, THEN mem-infmal-iff
[THEN iffD1]])
apply (auto simp add: mult.commute intro: approx-trans approx-minus-iff [THEN
iffD2])
done

```

Differentiation rules for combinations of functions follow from clear, straight-forward, algebraic manipulations.

Constant function.

```

lemma NSDERIV-const [simp]: NSDERIV (λx. k) x :> 0
by (simp add: NSDERIV-NSLIM-iff)

```

Sum of functions- proved easily.

```

lemma NSDERIV-add:
  NSDERIV f x :> Da ==> NSDERIV g x :> Db ==> NSDERIV (λx. f x + g x)
x :> Da + Db
apply (auto simp add: NSDERIV-NSLIM-iff NSLIM-def)
apply (auto simp add: add-divide-distrib diff-divide-distrib dest!: spec)
apply (drule-tac b = star-of Da and d = star-of Db in approx-add)
apply (auto simp add: ac-simps algebra-simps)
done

```

Product of functions - Proof is trivial but tedious and long due to rearrangement of terms.

lemma *lemma-nsderiv1*: $(a * b) - (c * d) = (b * (a - c)) + (c * (b - d))$
for $a b c d :: 'a::comm-ring star$
by (*simp add: right-diff-distrib ac-simps*)

lemma *lemma-nsderiv2*: $(x - y) / z = star-of D + yb \implies z \neq 0 \implies$
 $z \in Infinitesimal \implies yb \in Infinitesimal \implies x - y \approx 0$
for $x y z :: 'a::real-normed-field star$
apply (*simp add: nonzero-divide-eq-eq*)
apply (*auto intro!: Infinitesimal-HFinite-mult2 HFinite-add*
simp add: mult.assoc mem-infmal-iff [symmetric])
apply (*erule Infinitesimal-subset-HFinite [THEN subsetD]*)
done

lemma *NSDERIV-mult*:

NSDERIV f x :=> Da \implies *NSDERIV g x :=> Db* \implies
NSDERIV $(\lambda x. f x * g x) x :=> (Da * g x) + (Db * f x)$
apply (*auto simp add: NSDERIV-NSLIM-iff NSLIM-def*)
apply (*auto dest!: spec simp add: starfun-lambda-cancel lemma-nsderiv1*)
apply (*simp (no-asm) add: add-divide-distrib diff-divide-distrib*)
apply (*drule bex-Infinitesimal-iff2 [THEN iffD2]*)
apply (*auto simp add: times-divide-eq-right [symmetric]*
simp del: times-divide-eq-right times-divide-eq-left)
apply (*drule-tac D = Db in lemma-nsderiv2, assumption+*)
apply (*drule-tac approx-minus-iff [THEN iffD2, THEN bex-Infinitesimal-iff2*
[THEN iffD2]])
apply (*auto intro!: approx-add-mono1 simp: distrib-right distrib-left mult.commute*
add.assoc)
apply (*rule-tac b1 = star-of Db * star-of (f x) in add.commute [THEN subst]*)
apply (*auto intro!: Infinitesimal-add-approx-self2 [THEN approx-sym]*
Infinitesimal-add Infinitesimal-mult Infinitesimal-star-of-mult Infinitesimal-star-of-mult2
simp add: add.assoc [symmetric])
done

Multiplying by a constant.

lemma *NSDERIV-cmult*: *NSDERIV f x :=> D* \implies *NSDERIV* $(\lambda x. c * f x) x :=>$
 $c * D$
apply (*simp only: times-divide-eq-right [symmetric] NSDERIV-NSLIM-iff*
minus-mult-right right-diff-distrib [symmetric])
apply (*erule NSLIM-const [THEN NSLIM-mult]*)
done

Negation of function.

lemma *NSDERIV-minus*: *NSDERIV f x :=> D* \implies *NSDERIV* $(\lambda x. - f x) x :=>$
 $- D$
proof (*simp add: NSDERIV-NSLIM-iff*)
assume $(\lambda h. (f (x + h) - f x) / h) - 0 \rightarrow_{NS} D$

```

then have deriv: ( $\lambda h. - ((f(x+h) - f x) / h) - 0 \rightarrow_{NS} - D$ )
  by (rule NSLIM-minus)
have  $\forall h. - ((f(x+h) - f x) / h) = (- f(x+h) + f x) / h$ 
  by (simp add: minus-divide-left)
with deriv have ( $\lambda h. (- f(x+h) + f x) / h - 0 \rightarrow_{NS} - D$ )
  by simp
then show ( $\lambda h. (f(x+h) - f x) / h - 0 \rightarrow_{NS} D \implies (\lambda h. (f x - f(x+h)) / h) - 0 \rightarrow_{NS} - D$ )
  by simp
qed

```

Subtraction.

lemma NSDERIV-add-minus:

```

NSDERIV f x :> Da  $\implies$  NSDERIV g x :> Db  $\implies$  NSDERIV ( $\lambda x. f x + - g x$ ) x :> Da + - Db
by (blast dest: NSDERIV-add NSDERIV-minus)

```

lemma NSDERIV-diff:

```

NSDERIV f x :> Da  $\implies$  NSDERIV g x :> Db  $\implies$  NSDERIV ( $\lambda x. f x - g x$ ) x :> Da - Db
using NSDERIV-add-minus [of f x Da g Db] by simp

```

Similarly to the above, the chain rule admits an entirely straightforward derivation. Compare this with Harrison’s HOL proof of the chain rule, which proved to be trickier and required an alternative characterisation of differentiability- the so-called Carathedory derivative. Our main problem is manipulation of terms.

13.2 Lemmas

lemma NSDERIV-zero:

```

NSDERIV g x :> D  $\implies$  (*f* g) (star-of x + xa) = star-of (g x)  $\implies$ 
  xa  $\in$  Infinitesimal  $\implies$  xa  $\neq$  0  $\implies$  D = 0
apply (simp add: nsderiv-def)
apply (drule bspec)
apply auto
done

```

Can be proved differently using NSLIM-isCont-iff.

lemma NSDERIV-approx:

```

NSDERIV f x :> D  $\implies$  h  $\in$  Infinitesimal  $\implies$  h  $\neq$  0  $\implies$ 
  (*f* f) (star-of x + h) - star-of (f x)  $\approx$  0
apply (simp add: nsderiv-def)
apply (simp add: mem-infmal-iff [symmetric])
apply (rule Infinitesimal-ratio)
apply (rule-tac [3] approx-star-of-HFinite, auto)
done

```

From one version of differentiability

$$f x - f a \text{ -----} \approx Db x - a$$

lemma *NSDERIVD1*: [| *NSDERIV* $f (g x) \text{ :> } Da$;
 ($*f* g$) (*star-of*(x) + xa) \neq *star-of* ($g x$);
 ($*f* g$) (*star-of*(x) + xa) \approx *star-of* ($g x$)
 |] ==> (($*f* f$) (($*f* g$) (*star-of*(x) + xa))
 - *star-of* ($f (g x)$))
 / (($*f* g$) (*star-of*(x) + xa) - *star-of* ($g x$))
 \approx *star-of*(Da)
by (*auto simp add: NSDERIV-NSLIM-iff2 NSLIM-def*)

From other version of differentiability

$$f (x + h) - f x \text{ -----} \approx Db h$$

lemma *NSDERIVD2*: [| *NSDERIV* $g x \text{ :> } Db$; $xa \in \text{Infinitesimal}$; $xa \neq 0$ |]
 ==> (($*f* g$) (*star-of*(x) + xa) - *star-of*($g x$)) / xa
 \approx *star-of*(Db)
by (*auto simp add: NSDERIV-NSLIM-iff NSLIM-def mem-infmal-iff starfun-lambda-cancel*)

lemma *lemma-chain*: $z \neq 0 \implies x * y = (x * \text{inverse } z) * (z * y)$
for $x y z :: 'a::\text{real-normed-field}$ *star*

proof -
assume $z: z \neq 0$
have $x * y = x * (\text{inverse } z * z) * y$ **by** (*simp add: z*)
then show *?thesis* **by** (*simp add: mult.assoc*)
qed

This proof uses both definitions of differentiability.

lemma *NSDERIV-chain*:

NSDERIV $f (g x) \text{ :> } Da \implies \text{NSDERIV } g x \text{ :> } Db \implies \text{NSDERIV } (f \circ g) x \text{ :> } Da * Db$
apply (*simp (no-asm-simp) add: NSDERIV-NSLIM-iff NSLIM-def mem-infmal-iff [symmetric]*)
apply *clarify*
apply (*frule-tac* $f = g$ **in** *NSDERIV-approx*)
apply (*auto simp add: starfun-lambda-cancel2 starfun-o [symmetric]*)
apply (*case-tac* ($*f* g$) (*star-of* (x) + xa) = *star-of* ($g x$))
apply (*drule-tac* $g = g$ **in** *NSDERIV-zero*)
apply (*auto simp add: divide-inverse*)
apply (*rule-tac* $z1 = (*f* g) (\text{star-of } (x) + xa) - \text{star-of } (g x)$ **and** $y1 = \text{inverse } xa$
in *lemma-chain [THEN ssubst]*)
apply (*erule hypreal-not-eq-minus-iff [THEN iffD1]*)
apply (*rule approx-mult-star-of*)
apply (*simp-all add: divide-inverse [symmetric]*)
apply (*blast intro: NSDERIVD1 approx-minus-iff [THEN iffD2]*)
apply (*blast intro: NSDERIVD2*)
done

Differentiation of natural number powers.

```

lemma NSDERIV-Id [simp]: NSDERIV ( $\lambda x. x$ )  $x$  :> 1
  by (simp add: NSDERIV-NSLIM-iff NSLIM-def del: divide-self-if)

lemma NSDERIV-cmult-Id [simp]: NSDERIV ( $op * c$ )  $x$  :>  $c$ 
  using NSDERIV-Id [THEN NSDERIV-cmult] by simp

lemma NSDERIV-inverse:
  fixes  $x :: 'a::real-normed-field$ 
  assumes  $x \neq 0$  — can't get rid of  $x \neq (0::'a)$  because it isn't continuous at zero

  shows NSDERIV ( $\lambda x. inverse\ x$ )  $x$  :>  $-(inverse\ x \hat{\ } Suc\ (Suc\ 0))$ 
proof –
  {
    fix  $h :: 'a\ star$ 
    assume  $h-Inf$ :  $h \in Infinitesimal$ 
    from this assms have not-0:  $star-of\ x + h \neq 0$ 
      by (rule Infinitesimal-add-not-zero)
    assume  $h \neq 0$ 
    from  $h-Inf$  have  $h * star-of\ x \in Infinitesimal$ 
      by (rule Infinitesimal-HFinite-mult) simp
    with assms have  $inverse\ (-\ (h * star-of\ x) + -\ (star-of\ x * star-of\ x)) \approx$ 
       $inverse\ (-\ (star-of\ x * star-of\ x))$ 
    apply –
    apply (rule inverse-add-Infinitesimal-approx2)
    apply (auto dest!: hypreal-of-real-HFinite-diff-Infinitesimal
      simp add: inverse-minus-eq [symmetric] HFinite-minus-iff)
    done
    moreover from not-0  $\langle h \neq 0 \rangle$  assms
    have  $inverse\ (-\ (h * star-of\ x) + -\ (star-of\ x * star-of\ x)) =$ 
       $(inverse\ (star-of\ x + h) - inverse\ (star-of\ x)) / h$ 
    apply (simp add: division-ring-inverse-diff nonzero-inverse-mult-distrib [symmetric]
      nonzero-inverse-minus-eq [symmetric] ac-simps ring-distrib)
    apply (subst nonzero-inverse-minus-eq [symmetric])
    using distrib-right [symmetric, of  $h\ star-of\ x\ star-of\ x$ ] apply simp
    apply (simp add: field-simps)
    done
    ultimately have  $(inverse\ (star-of\ x + h) - inverse\ (star-of\ x)) / h \approx$ 
       $-(inverse\ (star-of\ x) * inverse\ (star-of\ x))$ 
    using assms by simp
  }
  then show ?thesis by (simp add: nsderiv-def)
qed

```

13.2.1 Equivalence of NS and Standard definitions

```

lemma divideR-eq-divide:  $x /_R y = x / y$ 
  by (simp add: divide-inverse mult.commute)

```

Now equivalence between *NSDERIV* and *DERIV*.

lemma *NSDERIV-DERIV-iff*: $NSDERIV f x \text{ :> } D \longleftrightarrow DERIV f x \text{ :> } D$
by (*simp add: DERIV-def NSDERIV-NSLIM-iff LIM-NSLIM-iff*)

NS version.

lemma *NSDERIV-pow*: $NSDERIV (\lambda x. x \wedge n) x \text{ :> } real\ n * (x \wedge (n - Suc\ 0))$
by (*simp add: NSDERIV-DERIV-iff DERIV-pow*)

Derivative of inverse.

lemma *NSDERIV-inverse-fun*:
 $NSDERIV f x \text{ :> } d \implies f x \neq 0 \implies$
 $NSDERIV (\lambda x. inverse (f x)) x \text{ :> } -(d * inverse (f x \wedge Suc (Suc\ 0)))$
for $x :: 'a :: \{real-normed-field\}$
by (*simp add: NSDERIV-DERIV-iff DERIV-inverse-fun del: power-Suc*)

Derivative of quotient.

lemma *NSDERIV-quotient*:
fixes $x :: 'a :: real-normed-field$
shows $NSDERIV f x \text{ :> } d \implies NSDERIV g x \text{ :> } e \implies g x \neq 0 \implies$
 $NSDERIV (\lambda y. f y / g y) x \text{ :> } (d * g x - (e * f x)) / (g x \wedge Suc (Suc\ 0))$
by (*simp add: NSDERIV-DERIV-iff DERIV-quotient del: power-Suc*)

lemma *CARAT-NSDERIV*:

$NSDERIV f x \text{ :> } l \implies \exists g. (\forall z. f z - f x = g z * (z - x)) \wedge isNSCont\ g\ x \wedge g\ x = l$
by (*auto simp add: NSDERIV-DERIV-iff isNSCont-isCont-iff CARAT-DERIV mult.commute*)

lemma *hypreal-eq-minus-iff3*: $x = y + z \longleftrightarrow x + - z = y$
for $x\ y\ z :: hypreal$
by *auto*

lemma *CARAT-DERIVD*:

assumes *all*: $\forall z. f z - f x = g z * (z - x)$
and *nsc*: $isNSCont\ g\ x$
shows $NSDERIV f x \text{ :> } g\ x$

proof –

from *nsc* **have** $\forall w. w \neq star-of\ x \wedge w \approx star-of\ x \longrightarrow$
 $(*f* g) w * (w - star-of\ x) / (w - star-of\ x) \approx star-of (g x)$

by (*simp add: isNSCont-def*)

with *all* **show** *?thesis*

by (*simp add: NSDERIV-iff2 starfun-if-eq cong: if-cong*)

qed

13.2.2 Differentiability predicate

lemma *NSdifferentiableD*: $f\ NSdifferentiable\ x \implies \exists D. NSDERIV f x \text{ :> } D$
by (*simp add: NSdifferentiable-def*)

lemma *NSdifferentiableI*: $NSDERIV f x \text{ :> } D \implies f\ NSdifferentiable\ x$

by (force simp add: NSdifferentiable-def)

13.3 (NS) Increment

lemma *incrementI*:

f NSdifferentiable *x* \implies
 $\text{increment } f x h = (*f* f) (\text{hypreal-of-real } x + h) - \text{hypreal-of-real } (f x)$
 by (simp add: increment-def)

lemma *incrementI2*:

NSDERIV *f x* $:> D \implies$
 $\text{increment } f x h = (*f* f) (\text{hypreal-of-real } x + h) - \text{hypreal-of-real } (f x)$
 by (erule NSdifferentiableI [THEN incrementI])

The Increment theorem – Keisler p. 65.

lemma *increment-thm*:

NSDERIV *f x* $:> D \implies h \in \text{Infinitesimal} \implies h \neq 0 \implies$
 $\exists e \in \text{Infinitesimal}. \text{increment } f x h = \text{hypreal-of-real } D * h + e * h$
 apply (frule-tac $h = h$ in incrementI2, simp add: nsderiv-def)
 apply (drule bspec, auto)
 apply (drule be-Infinitesimal-iff2 [THEN iffD2], clarify)
 apply (frule-tac $b1 = \text{hypreal-of-real } D + y$ in mult-right-cancel [THEN iffD2])
 apply (erule-tac [2])
 $V = ((*f* f) (\text{hypreal-of-real } x + h) - \text{hypreal-of-real } (f x)) / h =$
 $\text{hypreal-of-real } (D) + y$
 in thin-rl)
 apply assumption
 apply (simp add: times-divide-eq-right [symmetric])
 apply (auto simp add: distrib-right)
 done

lemma *increment-thm2*:

NSDERIV *f x* $:> D \implies h \approx 0 \implies h \neq 0 \implies$
 $\exists e \in \text{Infinitesimal}. \text{increment } f x h = \text{hypreal-of-real } D * h + e * h$
 by (blast dest!: mem-infmal-iff [THEN iffD2] intro!: increment-thm)

lemma *increment-approx-zero*: NSDERIV *f x* $:> D \implies h \approx 0 \implies h \neq 0 \implies$
 $\text{increment } f x h \approx 0$

apply (drule increment-thm2)
 apply (auto intro!: Infinitesimal-HFinite-mult2 HFinite-add
 simp add: distrib-right [symmetric] mem-infmal-iff [symmetric])
 apply (erule Infinitesimal-subset-HFinite [THEN subsetD])
 done

end

14 Nonstandard Extensions of Transcendental Functions

```
theory HTranscendental
imports HOL.Transcendental HSeries HDeriv
begin
```

definition

```
exphr :: real => hypreal where
  — define exponential function using standard part
  exphr x = st(sumhr (0, whn, %n. inverse (fact n) * (x ^ n)))
```

definition

```
sinhr :: real => hypreal where
  sinhr x = st(sumhr (0, whn, %n. sin-coeff n * x ^ n))
```

definition

```
coshhr :: real => hypreal where
  coshhr x = st(sumhr (0, whn, %n. cos-coeff n * x ^ n))
```

14.1 Nonstandard Extension of Square Root Function

```
lemma STAR-sqrt-zero [simp]: (*f* sqrt) 0 = 0
by (simp add: starfun star-n-zero-num)
```

```
lemma STAR-sqrt-one [simp]: (*f* sqrt) 1 = 1
by (simp add: starfun star-n-one-num)
```

```
lemma hypreal-sqrt-pow2-iff: (( *f* sqrt)(x) ^ 2 = x) = (0 ≤ x)
apply (cases x)
apply (auto simp add: star-n-le star-n-zero-num starfun hrealpow star-n-eq-iff
  simp del: hpowr-Suc power-Suc)
done
```

```
lemma hypreal-sqrt-gt-zero-pow2: !!x. 0 < x ==> (*f* sqrt) (x) ^ 2 = x
by (transfer, simp)
```

```
lemma hypreal-sqrt-pow2-gt-zero: 0 < x ==> 0 < (*f* sqrt) (x) ^ 2
by (frule hypreal-sqrt-gt-zero-pow2, auto)
```

```
lemma hypreal-sqrt-not-zero: 0 < x ==> (*f* sqrt) (x) ≠ 0
apply (frule hypreal-sqrt-pow2-gt-zero)
apply (auto simp add: numeral-2-eq-2)
done
```

```
lemma hypreal-inverse-sqrt-pow2:
  0 < x ==> inverse (( *f* sqrt)(x)) ^ 2 = inverse x
apply (cut-tac n = 2 and a = (*f* sqrt) x in power-inverse [symmetric])
apply (auto dest: hypreal-sqrt-gt-zero-pow2)
```

done

lemma *hypreal-sqrt-mult-distrib*:

!!x y. [[0 < x; 0 < y]] ==>

(**f* sqrt*)(x*y) = (**f* sqrt*)(x) * (**f* sqrt*)(y)

apply *transfer*

apply (*auto intro: real-sqrt-mult-distrib*)

done

lemma *hypreal-sqrt-mult-distrib2*:

[[0 ≤ x; 0 ≤ y]] ==>

(**f* sqrt*)(x*y) = (**f* sqrt*)(x) * (**f* sqrt*)(y)

by (*auto intro: hypreal-sqrt-mult-distrib simp add: order-le-less*)

lemma *hypreal-sqrt-approx-zero [simp]*:

0 < x ==> ((**f* sqrt*)(x) ≈ 0) = (x ≈ 0)

apply (*auto simp add: mem-infmal-iff [symmetric]*)

apply (*rule hypreal-sqrt-gt-zero-pow2 [THEN subst]*)

apply (*auto intro: Infinitesimal-mult*

dest!: hypreal-sqrt-gt-zero-pow2 [THEN ssubst]

simp add: numeral-2-eq-2)

done

lemma *hypreal-sqrt-approx-zero2 [simp]*:

0 ≤ x ==> ((**f* sqrt*)(x) ≈ 0) = (x ≈ 0)

by (*auto simp add: order-le-less*)

lemma *hypreal-sqrt-sum-squares [simp]*:

((**f* sqrt*)(x*x + y*y + z*z) ≈ 0) = (x*x + y*y + z*z ≈ 0)

apply (*rule hypreal-sqrt-approx-zero2*)

apply (*rule add-nonneg-nonneg*)

apply (*auto*)

done

lemma *hypreal-sqrt-sum-squares2 [simp]*:

((**f* sqrt*)(x*x + y*y) ≈ 0) = (x*x + y*y ≈ 0)

apply (*rule hypreal-sqrt-approx-zero2*)

apply (*rule add-nonneg-nonneg*)

apply (*auto*)

done

lemma *hypreal-sqrt-gt-zero*: !!x. 0 < x ==> 0 < (**f* sqrt*)(x)

apply *transfer*

apply (*auto intro: real-sqrt-gt-zero*)

done

lemma *hypreal-sqrt-ge-zero*: 0 ≤ x ==> 0 ≤ (**f* sqrt*)(x)

by (*auto intro: hypreal-sqrt-gt-zero simp add: order-le-less*)

lemma *hypreal-sqrt-hrabs* [simp]: $\forall x. (*f* \text{ sqrt})(x^2) = |x|$
by (*transfer*, *simp*)

lemma *hypreal-sqrt-hrabs2* [simp]: $\forall x. (*f* \text{ sqrt})(x*x) = |x|$
by (*transfer*, *simp*)

lemma *hypreal-sqrt-hyperpow-hrabs* [simp]:
 $\forall x. (*f* \text{ sqrt})(x \text{ pow } (\text{hypnat-of-nat } 2)) = |x|$
by (*transfer*, *simp*)

lemma *star-sqrt-HFinite*: $\llbracket x \in \text{HFinite}; 0 \leq x \rrbracket \implies (*f* \text{ sqrt}) x \in \text{HFinite}$
apply (*rule HFinite-square-iff [THEN iffD1]*)
apply (*simp only: hypreal-sqrt-mult-distrib2 [symmetric], simp*)
done

lemma *st-hypreal-sqrt*:
 $\llbracket x \in \text{HFinite}; 0 \leq x \rrbracket \implies \text{st}((*f* \text{ sqrt}) x) = (*f* \text{ sqrt})(\text{st } x)$
apply (*rule power-inject-base [where n=1]*)
apply (*auto intro!: st-zero-le hypreal-sqrt-ge-zero*)
apply (*rule st-mult [THEN subst]*)
apply (*rule-tac [3] hypreal-sqrt-mult-distrib2 [THEN subst]*)
apply (*rule-tac [5] hypreal-sqrt-mult-distrib2 [THEN subst]*)
apply (*auto simp add: st-hrabs st-zero-le star-sqrt-HFinite*)
done

lemma *hypreal-sqrt-sum-squares-ge1* [simp]: $\forall x y. x \leq (*f* \text{ sqrt})(x^2 + y^2)$
by *transfer* (*rule real-sqrt-sum-squares-ge1*)

lemma *HFinite-hypreal-sqrt*:
 $\llbracket 0 \leq x; x \in \text{HFinite} \rrbracket \implies (*f* \text{ sqrt}) x \in \text{HFinite}$
apply (*auto simp add: order-le-less*)
apply (*rule HFinite-square-iff [THEN iffD1]*)
apply (*drule hypreal-sqrt-gt-zero-pow2*)
apply (*simp add: numeral-2-eq-2*)
done

lemma *HFinite-hypreal-sqrt-imp-HFinite*:
 $\llbracket 0 \leq x; (*f* \text{ sqrt}) x \in \text{HFinite} \rrbracket \implies x \in \text{HFinite}$
apply (*auto simp add: order-le-less*)
apply (*drule HFinite-square-iff [THEN iffD2]*)
apply (*drule hypreal-sqrt-gt-zero-pow2*)
apply (*simp add: numeral-2-eq-2 del: HFinite-square-iff*)
done

lemma *HFinite-hypreal-sqrt-iff* [simp]:
 $0 \leq x \implies ((*f* \text{ sqrt}) x \in \text{HFinite}) = (x \in \text{HFinite})$
by (*blast intro: HFinite-hypreal-sqrt HFinite-hypreal-sqrt-imp-HFinite*)

lemma *HFinite-sqrt-sum-squares* [simp]:

```

    (( *f* sqrt)(x*x + y*y) ∈ HFinite) = (x*x + y*y ∈ HFinite)
  apply (rule HFinite-hypreal-sqrt-iff)
  apply (rule add-nonneg-nonneg)
  apply (auto)
  done

```

```

lemma Infinitesimal-hypreal-sqrt:
  [| 0 ≤ x; x ∈ Infinitesimal |] ==> (*f* sqrt) x ∈ Infinitesimal
  apply (auto simp add: order-le-less)
  apply (rule Infinitesimal-square-iff [THEN iffD2])
  apply (drule hypreal-sqrt-gt-zero-pow2)
  apply (simp add: numeral-2-eq-2)
  done

```

```

lemma Infinitesimal-hypreal-sqrt-imp-Infinitesimal:
  [| 0 ≤ x; (*f* sqrt) x ∈ Infinitesimal |] ==> x ∈ Infinitesimal
  apply (auto simp add: order-le-less)
  apply (drule Infinitesimal-square-iff [THEN iffD1])
  apply (drule hypreal-sqrt-gt-zero-pow2)
  apply (simp add: numeral-2-eq-2 del: Infinitesimal-square-iff [symmetric])
  done

```

```

lemma Infinitesimal-hypreal-sqrt-iff [simp]:
  0 ≤ x ==> ((*f* sqrt) x ∈ Infinitesimal) = (x ∈ Infinitesimal)
  by (blast intro: Infinitesimal-hypreal-sqrt-imp-Infinitesimal Infinitesimal-hypreal-sqrt)

```

```

lemma Infinitesimal-sqrt-sum-squares [simp]:
  ((*f* sqrt)(x*x + y*y) ∈ Infinitesimal) = (x*x + y*y ∈ Infinitesimal)
  apply (rule Infinitesimal-hypreal-sqrt-iff)
  apply (rule add-nonneg-nonneg)
  apply (auto)
  done

```

```

lemma HInfinite-hypreal-sqrt:
  [| 0 ≤ x; x ∈ HInfinite |] ==> (*f* sqrt) x ∈ HInfinite
  apply (auto simp add: order-le-less)
  apply (rule HInfinite-square-iff [THEN iffD1])
  apply (drule hypreal-sqrt-gt-zero-pow2)
  apply (simp add: numeral-2-eq-2)
  done

```

```

lemma HInfinite-hypreal-sqrt-imp-HInfinite:
  [| 0 ≤ x; (*f* sqrt) x ∈ HInfinite |] ==> x ∈ HInfinite
  apply (auto simp add: order-le-less)
  apply (drule HInfinite-square-iff [THEN iffD2])
  apply (drule hypreal-sqrt-gt-zero-pow2)
  apply (simp add: numeral-2-eq-2 del: HInfinite-square-iff)
  done

```

lemma *HInfinite-hypreal-sqrt-iff* [simp]:
 $0 \leq x \implies ((*f* \text{ sqrt}) x \in \text{HInfinite}) = (x \in \text{HInfinite})$
by (blast intro: *HInfinite-hypreal-sqrt HInfinite-hypreal-sqrt-imp-HInfinite*)

lemma *HInfinite-sqrt-sum-squares* [simp]:
 $((*f* \text{ sqrt})(x*x + y*y) \in \text{HInfinite}) = (x*x + y*y \in \text{HInfinite})$
apply (rule *HInfinite-hypreal-sqrt-iff*)
apply (rule *add-nonneg-nonneg*)
apply (auto)
done

lemma *HFinite-exp* [simp]:
 $\text{sumhr } (0, \text{whn}, \%n. \text{inverse } (\text{fact } n) * x ^ n) \in \text{HFinite}$
unfolding *sumhr-app*
apply (simp only: *star-zero-def starfun2-star-of atLeast0LessThan*)
apply (rule *NSBseqD2*)
apply (rule *NSconvergent-NSBseq*)
apply (rule *convergent-NSconvergent-iff [THEN iffD1]*)
apply (rule *summable-iff-convergent [THEN iffD1]*)
apply (rule *summable-exp*)
done

lemma *exphr-zero* [simp]: $\text{exphr } 0 = 1$
apply (simp add: *exphr-def sumhr-split-add [OF hypnat-one-less-hypnat-omega, symmetric]*)
apply (rule *st-unique, simp*)
apply (rule *subst [where P= $\lambda x. 1 \approx x$, OF - approx-refl]*)
apply (rule *rev-mp [OF hypnat-one-less-hypnat-omega]*)
apply (rule *tac x=whn in spec*)
apply (unfold *sumhr-app, transfer, simp add: power-0-left*)
done

lemma *coshr-zero* [simp]: $\text{coshr } 0 = 1$
apply (simp add: *coshr-def sumhr-split-add [OF hypnat-one-less-hypnat-omega, symmetric]*)
apply (rule *st-unique, simp*)
apply (rule *subst [where P= $\lambda x. 1 \approx x$, OF - approx-refl]*)
apply (rule *rev-mp [OF hypnat-one-less-hypnat-omega]*)
apply (rule *tac x=whn in spec*)
apply (unfold *sumhr-app, transfer, simp add: cos-coeff-def power-0-left*)
done

lemma *STAR-exp-zero-approx-one* [simp]: $(*f* \text{ exp}) (0::\text{hypreal}) \approx 1$
apply (subgoal-tac (**f* exp*) (0::hypreal) = 1, simp)
apply (transfer, simp)
done

lemma *STAR-exp-Infinitesimal*: $x \in \text{Infinitesimal} \implies (*f* \text{ exp}) (x::\text{hypreal}) \approx 1$

```

apply (case-tac x = 0)
apply (cut-tac [2] x = 0 in DERIV-exp)
apply (auto simp add: NSDERIV-DERIV-iff [symmetric] nsderiv-def)
apply (drule-tac x = x in bspec, auto)
apply (drule-tac c = x in approx-mult1)
apply (auto intro: Infinitesimal-subset-HFinite [THEN subsetD]
      simp add: mult.assoc)
apply (rule approx-add-right-cancel [where d=-1])
apply (rule approx-sym [THEN [2] approx-trans2])
apply (auto simp add: mem-infmal-iff)
done

```

```

lemma STAR-exp-epsilon [simp]: (*f* exp) ε ≈ 1
by (auto intro: STAR-exp-Infinitesimal)

```

```

lemma STAR-exp-add:

```

```

  !!(x::'a:: {banach,real-normed-field}) star y. (*f* exp)(x + y) = (*f* exp) x *
  (*f* exp) y
by transfer (rule exp-add)

```

```

lemma exphr-hypreal-of-real-exp-eq: exphr x = hypreal-of-real (exp x)

```

```

apply (simp add: exphr-def)
apply (rule st-unique, simp)
apply (subst starfunNat-sumr [symmetric])
unfolding atLeast0LessThan
apply (rule NSLIMSEQ-D [THEN approx-sym])
apply (rule LIMSEQ-NSLIMSEQ)
apply (subst sums-def [symmetric])
apply (cut-tac exp-converges [where x=x], simp)
apply (rule HNatInfinite-whn)
done

```

```

lemma starfun-exp-ge-add-one-self [simp]: !!x::hypreal. 0 ≤ x ==> (1 + x) ≤ (
*f* exp) x
by transfer (rule exp-ge-add-one-self-aux)

```

```

lemma starfun-exp-HInfinite:

```

```

  [| x ∈ HInfinite; 0 ≤ x |] ==> (*f* exp) (x::hypreal) ∈ HInfinite
apply (frule starfun-exp-ge-add-one-self)
apply (rule HInfinite-ge-HInfinite, assumption)
apply (rule order-trans [of - 1+x], auto)
done

```

```

lemma starfun-exp-minus:

```

```

  !!x::'a:: {banach,real-normed-field} star. (*f* exp) (-x) = inverse((*f* exp) x)
by transfer (rule exp-minus)

```

lemma *starfun-exp-Infinitesimal*:

$[[x \in HInfinite; x \leq 0]] \implies (*f* \text{ exp}) (x::hypreal) \in Infinitesimal$
apply (*subgoal-tac* $\exists y. x = -y$)
apply (*rule-tac* [2] $x = -x$ **in** *exI*)
apply (*auto intro!*: *HInfinite-inverse-Infinitesimal starfun-exp-HInfinite*
simp add: starfun-exp-minus HInfinite-minus-iff)
done

lemma *starfun-exp-gt-one* [*simp*]: $!!x::hypreal. 0 < x \implies 1 < (*f* \text{ exp}) x$
by transfer (*rule exp-gt-one*)

abbreviation *real-ln* :: *real* \Rightarrow *real* **where**
real-ln \equiv *ln*

lemma *starfun-ln-exp* [*simp*]: $!!x. (*f* \text{ real-ln}) ((*f* \text{ exp}) x) = x$
by transfer (*rule ln-exp*)

lemma *starfun-exp-ln-iff* [*simp*]: $!!x. ((*f* \text{ exp})((*f* \text{ real-ln}) x) = x) = (0 < x)$
by transfer (*rule exp-ln-iff*)

lemma *starfun-exp-ln-eq*: $!!u x. (*f* \text{ exp}) u = x \implies (*f* \text{ real-ln}) x = u$
by transfer (*rule ln-unique*)

lemma *starfun-ln-less-self* [*simp*]: $!!x. 0 < x \implies (*f* \text{ real-ln}) x < x$
by transfer (*rule ln-less-self*)

lemma *starfun-ln-ge-zero* [*simp*]: $!!x. 1 \leq x \implies 0 \leq (*f* \text{ real-ln}) x$
by transfer (*rule ln-ge-zero*)

lemma *starfun-ln-gt-zero* [*simp*]: $!!x. 1 < x \implies 0 < (*f* \text{ real-ln}) x$
by transfer (*rule ln-gt-zero*)

lemma *starfun-ln-not-eq-zero* [*simp*]: $!!x. [[0 < x; x \neq 1]] \implies (*f* \text{ real-ln}) x \neq 0$
by transfer simp

lemma *starfun-ln-HFinite*: $[[x \in HFinite; 1 \leq x]] \implies (*f* \text{ real-ln}) x \in HFinite$
apply (*rule HFinite-bounded*)
apply *assumption*
apply (*simp-all add: starfun-ln-less-self order-less-imp-le*)
done

lemma *starfun-ln-inverse*: $!!x. 0 < x \implies (*f* \text{ real-ln}) (\text{inverse } x) = -(*f* \text{ ln}) x$
by transfer (*rule ln-inverse*)

lemma *starfun-abs-exp-cancel*: $\bigwedge x. |(*f* \text{ exp}) (x::hypreal)| = (*f* \text{ exp}) x$
by transfer (*rule abs-exp-cancel*)

lemma *starfun-exp-less-mono*: $\bigwedge x y :: \text{hypreal}. x < y \implies (*f* \text{exp}) x < (*f* \text{exp}) y$
by *transfer* (*rule exp-less-mono*)

lemma *starfun-exp-HFinite*: $x \in \text{HFinite} \implies (*f* \text{exp}) (x :: \text{hypreal}) \in \text{HFinite}$
apply (*auto simp add: HFinite-def, rename-tac u*)
apply (*rule-tac x = (*f* exp) u in rev-bexI*)
apply (*simp add: Reals-eq-Standard*)
apply (*simp add: starfun-abs-exp-cancel*)
apply (*simp add: starfun-exp-less-mono*)
done

lemma *starfun-exp-add-HFinite-Infinitesimal-approx*:
 $[[x \in \text{Infinitesimal}; z \in \text{HFinite}]] \implies (*f* \text{exp}) (z + x :: \text{hypreal}) \approx (*f* \text{exp}) z$
apply (*simp add: STAR-exp-add*)
apply (*frule STAR-exp-Infinitesimal*)
apply (*drule approx-mult2*)
apply (*auto intro: starfun-exp-HFinite*)
done

lemma *starfun-ln-HInfinite*:
 $[[x \in \text{HInfinite}; 0 < x]] \implies (*f* \text{real-ln}) x \in \text{HInfinite}$
apply (*rule ccontr, drule HFinite-HInfinite-iff [THEN iffD2]*)
apply (*drule starfun-exp-HFinite*)
apply (*simp add: starfun-exp-ln-iff [THEN iffD2] HFinite-HInfinite-iff*)
done

lemma *starfun-exp-HInfinite-Infinitesimal-disj*:
 $x \in \text{HInfinite} \implies (*f* \text{exp}) x \in \text{HInfinite} \mid (*f* \text{exp}) (x :: \text{hypreal}) \in \text{Infinitesimal}$
apply (*insert linorder-linear [of x 0]*)
apply (*auto intro: starfun-exp-HInfinite starfun-exp-Infinitesimal*)
done

lemma *starfun-ln-HFinite-not-Infinitesimal*:
 $[[x \in \text{HFinite} - \text{Infinitesimal}; 0 < x]] \implies (*f* \text{real-ln}) x \in \text{HFinite}$
apply (*rule ccontr, drule HInfinite-HFinite-iff [THEN iffD2]*)
apply (*drule starfun-exp-HInfinite-Infinitesimal-disj*)
apply (*simp add: starfun-exp-ln-iff [symmetric] HInfinite-HFinite-iff*
 $\text{del: starfun-exp-ln-iff}$)
done

lemma *starfun-ln-Infinitesimal-HInfinite*:
 $[[x \in \text{Infinitesimal}; 0 < x]] \implies (*f* \text{real-ln}) x \in \text{HInfinite}$
apply (*drule Infinitesimal-inverse-HInfinite*)
apply (*frule positive-imp-inverse-positive*)


```

apply (drule-tac [2] starfun-ln-HInfinite)
apply (auto simp add: starfun-ln-inverse HInfinite-minus-iff)
done

```

```

lemma starfun-ln-less-zero: !!x. [| 0 < x; x < 1 |] ==> (*f* real-ln) x < 0
by transfer (rule ln-less-zero)

```

```

lemma starfun-ln-Infinitesimal-less-zero:
  [| x ∈ Infinitesimal; 0 < x |] ==> (*f* real-ln) x < 0
by (auto intro!: starfun-ln-less-zero simp add: Infinitesimal-def)

```

```

lemma starfun-ln-HInfinite-gt-zero:
  [| x ∈ HInfinite; 0 < x |] ==> 0 < (*f* real-ln) x
by (auto intro!: starfun-ln-gt-zero simp add: HInfinite-def)

```

```

lemma HFinite-sin [simp]: sumhr (0, whn, %n. sin-coeff n * x ^ n) ∈ HFinite
unfolding sumhr-app
apply (simp only: star-zero-def starfun2-star-of atLeast0LessThan)
apply (rule NSBseqD2)
apply (rule NSconvergent-NSBseq)
apply (rule convergent-NSconvergent-iff [THEN iffD1])
apply (rule summable-iff-convergent [THEN iffD1])
using summable-norm-sin [of x]
apply (simp add: summable-rabs-cancel)
done

```

```

lemma STAR-sin-zero [simp]: (*f* sin) 0 = 0
by transfer (rule sin-zero)

```

```

lemma STAR-sin-Infinitesimal [simp]:
  fixes x :: 'a::{real-normed-field,banach} star
  shows x ∈ Infinitesimal ==> (*f* sin) x ≈ x
apply (case-tac x = 0)
apply (cut-tac [2] x = 0 in DERIV-sin)
apply (auto simp add: NSDERIV-DERIV-iff [symmetric] nsderiv-def)
apply (drule bspec [where x = x], auto)
apply (drule approx-mult1 [where c = x])
apply (auto intro: Infinitesimal-subset-HFinite [THEN subsetD]
  simp add: mult.assoc)
done

```

```

lemma HFinite-cos [simp]: sumhr (0, whn, %n. cos-coeff n * x ^ n) ∈ HFinite
unfolding sumhr-app
apply (simp only: star-zero-def starfun2-star-of atLeast0LessThan)
apply (rule NSBseqD2)
apply (rule NSconvergent-NSBseq)

```

```

apply (rule convergent-NSconvergent-iff [THEN iffD1])
apply (rule summable-iff-convergent [THEN iffD1])
using summable-norm-cos [of x]
apply (simp add: summable-rabs-cancel)
done

```

```

lemma STAR-cos-zero [simp]: ( *f* cos) 0 = 1
by transfer (rule cos-zero)

```

```

lemma STAR-cos-Infinitesimal [simp]:
  fixes x :: 'a::{real-normed-field,banach} star
  shows x ∈ Infinitesimal ==> ( *f* cos) x ≈ 1
apply (case-tac x = 0)
apply (cut-tac [2] x = 0 in DERIV-cos)
apply (auto simp add: NSDERIV-DERIV-iff [symmetric] nsderiv-def)
apply (drule bspec [where x = x])
apply auto
apply (drule approx-mult1 [where c = x])
apply (auto intro: Infinitesimal-subset-HFinite [THEN subsetD]
  simp add: mult.assoc)
apply (rule approx-add-right-cancel [where d = -1])
apply simp
done

```

```

lemma STAR-tan-zero [simp]: ( *f* tan) 0 = 0
by transfer (rule tan-zero)

```

```

lemma STAR-tan-Infinitesimal: x ∈ Infinitesimal ==> ( *f* tan) x ≈ x
apply (case-tac x = 0)
apply (cut-tac [2] x = 0 in DERIV-tan)
apply (auto simp add: NSDERIV-DERIV-iff [symmetric] nsderiv-def)
apply (drule bspec [where x = x], auto)
apply (drule approx-mult1 [where c = x])
apply (auto intro: Infinitesimal-subset-HFinite [THEN subsetD]
  simp add: mult.assoc)
done

```

```

lemma STAR-sin-cos-Infinitesimal-mult:
  fixes x :: 'a::{real-normed-field,banach} star
  shows x ∈ Infinitesimal ==> ( *f* sin) x * ( *f* cos) x ≈ x
using approx-mult-HFinite [of ( *f* sin) x - ( *f* cos) x 1]
by (simp add: Infinitesimal-subset-HFinite [THEN subsetD])

```

```

lemma HFinite-pi: hypreal-of-real pi ∈ HFinite
by simp

```

```

lemma lemma-split-hypreal-of-real:

```

$N \in \text{HNatInfinite}$
 $\implies \text{hypreal-of-real } a =$
 $\text{hypreal-of-hypnat } N * (\text{inverse}(\text{hypreal-of-hypnat } N) * \text{hypreal-of-real } a)$
by (*simp add: mult.assoc [symmetric] zero-less-HNatInfinite*)

lemma *STAR-sin-Infinitesimal-divide:*

fixes $x :: 'a :: \{\text{real-normed-field}, \text{banach}\}$ *star*
shows $[|x \in \text{Infinitesimal}; x \neq 0|] \implies (*f* \text{ sin}) x/x \approx 1$
using *DERIV-sin [of 0::'a]*
by (*simp add: NSDERIV-DERIV-iff [symmetric] nsderiv-def*)

lemma *lemma-sin-pi:*

$n \in \text{HNatInfinite}$
 $\implies (*f* \text{ sin}) (\text{inverse}(\text{hypreal-of-hypnat } n)) / (\text{inverse}(\text{hypreal-of-hypnat } n)) \approx 1$
apply (*rule STAR-sin-Infinitesimal-divide*)
apply (*auto simp add: zero-less-HNatInfinite*)
done

lemma *STAR-sin-inverse-HNatInfinite:*

$n \in \text{HNatInfinite}$
 $\implies (*f* \text{ sin}) (\text{inverse}(\text{hypreal-of-hypnat } n)) * \text{hypreal-of-hypnat } n \approx 1$
apply (*frule lemma-sin-pi*)
apply (*simp add: divide-inverse*)
done

lemma *Infinitesimal-pi-divide-HNatInfinite:*

$N \in \text{HNatInfinite}$
 $\implies \text{hypreal-of-real } \pi / (\text{hypreal-of-hypnat } N) \in \text{Infinitesimal}$
apply (*simp add: divide-inverse*)
apply (*auto intro: Infinitesimal-HFinite-mult2*)
done

lemma *pi-divide-HNatInfinite-not-zero [simp]:*

$N \in \text{HNatInfinite} \implies \text{hypreal-of-real } \pi / (\text{hypreal-of-hypnat } N) \neq 0$
by (*simp add: zero-less-HNatInfinite*)

lemma *STAR-sin-pi-divide-HNatInfinite-approx-pi:*

$n \in \text{HNatInfinite}$
 $\implies (*f* \text{ sin}) (\text{hypreal-of-real } \pi / (\text{hypreal-of-hypnat } n)) * \text{hypreal-of-hypnat } n$
 $\approx \text{hypreal-of-real } \pi$
apply (*frule STAR-sin-Infinitesimal-divide*
 $[OF \text{ Infinitesimal-pi-divide-HNatInfinite}$
 $\text{ pi-divide-HNatInfinite-not-zero}]$)

```

apply (auto)
apply (rule approx-SReal-mult-cancel [of inverse (hypreal-of-real pi)])
apply (auto intro: Reals-inverse simp add: divide-inverse ac-simps)
done

```

```

lemma STAR-sin-pi-divide-HNatInfinite-approx-pi2:
  n ∈ HNatInfinite
  ==> hypreal-of-hypnat n *
    (*f* sin) (hypreal-of-real pi / (hypreal-of-hypnat n))
    ≈ hypreal-of-real pi
apply (rule mult.commute [THEN subst])
apply (erule STAR-sin-pi-divide-HNatInfinite-approx-pi)
done

```

```

lemma starfunNat-pi-divide-n-Infinitesimal:
  N ∈ HNatInfinite ==> (*f* (%x. pi / real x)) N ∈ Infinitesimal
by (auto intro!: Infinitesimal-HFinite-mult2
  simp add: starfun-mult [symmetric] divide-inverse
  starfun-inverse [symmetric] starfunNat-real-of-nat)

```

```

lemma STAR-sin-pi-divide-n-approx:
  N ∈ HNatInfinite ==>
    (*f* sin) (( *f* (%x. pi / real x)) N) ≈
    hypreal-of-real pi / (hypreal-of-hypnat N)
apply (simp add: starfunNat-real-of-nat [symmetric])
apply (rule STAR-sin-Infinitesimal)
apply (simp add: divide-inverse)
apply (rule Infinitesimal-HFinite-mult2)
apply (subst starfun-inverse)
apply (erule starfunNat-inverse-real-of-nat-Infinitesimal)
apply simp
done

```

```

lemma NSLIMSEQ-sin-pi: (%n. real n * sin (pi / real n)) →NS pi
apply (auto simp add: NSLIMSEQ-def starfun-mult [symmetric] starfunNat-real-of-nat)
apply (rule-tac f1 = sin in starfun-o2 [THEN subst])
apply (auto simp add: starfun-mult [symmetric] starfunNat-real-of-nat divide-inverse)
apply (rule-tac f1 = inverse in starfun-o2 [THEN subst])
apply (auto dest: STAR-sin-pi-divide-HNatInfinite-approx-pi
  simp add: starfunNat-real-of-nat mult.commute divide-inverse)
done

```

```

lemma NSLIMSEQ-cos-one: (%n. cos (pi / real n)) →NS 1
apply (simp add: NSLIMSEQ-def, auto)
apply (rule-tac f1 = cos in starfun-o2 [THEN subst])
apply (rule STAR-cos-Infinitesimal)
apply (auto intro!: Infinitesimal-HFinite-mult2
  simp add: starfun-mult [symmetric] divide-inverse
  starfun-inverse [symmetric] starfunNat-real-of-nat)

```

done

lemma *NSLIMSEQ-sin-cos-pi*:

(%n. real n * sin (pi / real n) * cos (pi / real n)) \longrightarrow_{NS} pi

by (insert NSLIMSEQ-mult [OF NSLIMSEQ-sin-pi NSLIMSEQ-cos-one], simp)

A familiar approximation to $\cos x$ when x is small

lemma *STAR-cos-Infinitesimal-approx*:

fixes $x :: 'a :: \{\text{real-normed-field}, \text{banach}\}$ star

shows $x \in \text{Infinitesimal} \implies (*f* \cos) x \approx 1 - x^2$

apply (rule STAR-cos-Infinitesimal [THEN approx-trans])

apply (auto simp add: Infinitesimal-approx-minus [symmetric]

add.assoc [symmetric] numeral-2-eq-2)

done

lemma *STAR-cos-Infinitesimal-approx2*:

fixes $x :: \text{hypreal}$ — perhaps could be generalised, like many other hypreal results

shows $x \in \text{Infinitesimal} \implies (*f* \cos) x \approx 1 - (x^2)/2$

apply (rule STAR-cos-Infinitesimal [THEN approx-trans])

apply (auto intro: Infinitesimal-SReal-divide Infinitesimal-mult

simp add: Infinitesimal-approx-minus [symmetric] numeral-2-eq-2)

done

end

15 Non-Standard Complex Analysis

theory NSCA

imports NSComplex HTranscendental

begin

abbreviation

SComplex :: *hcomplex set* **where**

SComplex \equiv *Standard*

definition — standard part map

stc :: *hcomplex* \implies *hcomplex* **where**

stc $x = (\text{SOME } r. x \in \text{HFinite} \ \& \ r : \text{SComplex} \ \& \ r \approx x)$

15.1 Closure Laws for SComplex, the Standard Complex Numbers

lemma *SComplex-minus-iff* [simp]: $(-x \in \text{SComplex}) = (x \in \text{SComplex})$

by (auto, drule Standard-minus, auto)

lemma *SComplex-add-cancel*:

$[[x + y \in \text{SComplex}; y \in \text{SComplex}]] \implies x \in \text{SComplex}$

by (drule (1) Standard-diff, simp)

lemma *SReal-hcmod-hcomplex-of-complex* [simp]:

$hcmod (hcomplex-of-complex r) \in \mathbb{R}$

by (simp add: Reals-eq-Standard)

lemma *SReal-hcmod-numeral* [simp]: $hcmod (numeral w :: hcomplex) \in \mathbb{R}$

by (simp add: Reals-eq-Standard)

lemma *SReal-hcmod-SComplex*: $x \in SComplex \implies hcmod x \in \mathbb{R}$

by (simp add: Reals-eq-Standard)

lemma *SComplex-divide-numeral*:

$r \in SComplex \implies r / (numeral w :: hcomplex) \in SComplex$

by simp

lemma *SComplex-UNIV-complex*:

$\{x. hcomplex-of-complex x \in SComplex\} = (UNIV :: complex set)$

by simp

lemma *SComplex-iff*: $(x \in SComplex) = (\exists y. x = hcomplex-of-complex y)$

by (simp add: Standard-def image-def)

lemma *hcomplex-of-complex-image*:

$hcomplex-of-complex \text{ `}(UNIV :: complex set) = SComplex$

by (simp add: Standard-def)

lemma *inv-hcomplex-of-complex-image*: $inv \text{ `} hcomplex-of-complex \text{ `}(SComplex = UNIV$

by (auto simp add: Standard-def image-def) (metis inj-star-of inv-f-f)

lemma *SComplex-hcomplex-of-complex-image*:

$[\exists x. x: P; P \leq SComplex] \implies \exists Q. P = hcomplex-of-complex \text{ `} Q$

apply (simp add: Standard-def, blast)

done

lemma *SComplex-SReal-dense*:

$[\exists x \in SComplex; y \in SComplex; hcmod x < hcmod y$

$] \implies \exists r \in Reals. hcmod x < r \ \& \ r < hcmod y$

apply (auto intro: SReal-dense simp add: SReal-hcmod-SComplex)

done

15.2 The Finite Elements form a Subring

lemma *HFinite-hcmod-hcomplex-of-complex* [simp]:

$hcmod (hcomplex-of-complex r) \in HFinite$

by (auto intro!: SReal-subset-HFinite [THEN subsetD])

lemma *HFinite-hcmod-iff*: $(x \in HFinite) = (hcmod x \in HFinite)$

by (simp add: HFinite-def)

lemma *HFinite-bounded-hcmod*:

$\llbracket x \in \text{HFinite}; y \leq \text{hcmod } x; 0 \leq y \rrbracket \implies y: \text{HFinite}$

by (*auto intro: HFinite-bounded simp add: HFinite-hcmod-iff*)

15.3 The Complex Infinitesimals form a Subring

lemma *hcomplex-sum-of-halves*: $x/(2::\text{hcomplex}) + x/(2::\text{hcomplex}) = x$

by *auto*

lemma *Infinitesimal-hcmod-iff*:

$(z \in \text{Infinitesimal}) = (\text{hcmod } z \in \text{Infinitesimal})$

by (*simp add: Infinitesimal-def*)

lemma *HInfinite-hcmod-iff*: $(z \in \text{HInfinite}) = (\text{hcmod } z \in \text{HInfinite})$

by (*simp add: HInfinite-def*)

lemma *HFinite-diff-Infinitesimal-hcmod*:

$x \in \text{HFinite} - \text{Infinitesimal} \implies \text{hcmod } x \in \text{HFinite} - \text{Infinitesimal}$

by (*simp add: HFinite-hcmod-iff Infinitesimal-hcmod-iff*)

lemma *hcmod-less-Infinitesimal*:

$\llbracket e \in \text{Infinitesimal}; \text{hcmod } x < \text{hcmod } e \rrbracket \implies x \in \text{Infinitesimal}$

by (*auto elim: hrabs-less-Infinitesimal simp add: Infinitesimal-hcmod-iff*)

lemma *hcmod-le-Infinitesimal*:

$\llbracket e \in \text{Infinitesimal}; \text{hcmod } x \leq \text{hcmod } e \rrbracket \implies x \in \text{Infinitesimal}$

by (*auto elim: hrabs-le-Infinitesimal simp add: Infinitesimal-hcmod-iff*)

lemma *Infinitesimal-interval-hcmod*:

$\llbracket e \in \text{Infinitesimal};$
 $e' \in \text{Infinitesimal};$
 $\text{hcmod } e' < \text{hcmod } x; \text{hcmod } x < \text{hcmod } e$
 $\rrbracket \implies x \in \text{Infinitesimal}$

by (*auto intro: Infinitesimal-interval simp add: Infinitesimal-hcmod-iff*)

lemma *Infinitesimal-interval2-hcmod*:

$\llbracket e \in \text{Infinitesimal};$
 $e' \in \text{Infinitesimal};$
 $\text{hcmod } e' \leq \text{hcmod } x; \text{hcmod } x \leq \text{hcmod } e$
 $\rrbracket \implies x \in \text{Infinitesimal}$

by (*auto intro: Infinitesimal-interval2 simp add: Infinitesimal-hcmod-iff*)

15.4 The “Infinitely Close” Relation

lemma *approx-SComplex-mult-cancel-zero*:

$\llbracket a \in \text{SComplex}; a \neq 0; a * x \approx 0 \rrbracket \implies x \approx 0$

apply (*drule Standard-inverse [THEN Standard-subset-HFinite [THEN subsetD]]*)

apply (*auto dest: approx-mult2 simp add: mult.assoc [symmetric]*)

done

lemma *approx-mult-SComplex1*: $[[a \in SComplex; x \approx 0]] \implies x*a \approx 0$
by (*auto dest: Standard-subset-HFfinite [THEN subsetD] approx-mult1*)

lemma *approx-mult-SComplex2*: $[[a \in SComplex; x \approx 0]] \implies a*x \approx 0$
by (*auto dest: Standard-subset-HFfinite [THEN subsetD] approx-mult2*)

lemma *approx-mult-SComplex-zero-cancel-iff [simp]*:
 $[[a \in SComplex; a \neq 0]] \implies (a*x \approx 0) = (x \approx 0)$
by (*blast intro: approx-SComplex-mult-cancel-zero approx-mult-SComplex2*)

lemma *approx-SComplex-mult-cancel*:
 $[[a \in SComplex; a \neq 0; a*w \approx a*z]] \implies w \approx z$
apply (*drule Standard-inverse [THEN Standard-subset-HFfinite [THEN subsetD]]*)
apply (*auto dest: approx-mult2 simp add: mult.assoc [symmetric]*)
done

lemma *approx-SComplex-mult-cancel-iff1 [simp]*:
 $[[a \in SComplex; a \neq 0]] \implies (a*w \approx a*z) = (w \approx z)$
by (*auto intro!: approx-mult2 Standard-subset-HFfinite [THEN subsetD]*
intro: approx-SComplex-mult-cancel)

lemma *approx-hcmod-approx-zero*: $(x \approx y) = (hcmod (y - x) \approx 0)$
apply (*subst hnorm-minus-commute*)
apply (*simp add: approx-def Infinitesimal-hcmod-iff*)
done

lemma *approx-approx-zero-iff*: $(x \approx 0) = (hcmod x \approx 0)$
by (*simp add: approx-hcmod-approx-zero*)

lemma *approx-minus-zero-cancel-iff [simp]*: $(-x \approx 0) = (x \approx 0)$
by (*simp add: approx-def*)

lemma *Infinitesimal-hcmod-add-diff*:
 $u \approx 0 \implies hcmod(x + u) - hcmod x \in Infinitesimal$
apply (*drule approx-approx-zero-iff [THEN iffD1]*)
apply (*rule-tac e = hcmod u and e' = - hcmod u in Infinitesimal-interval2*)
apply (*auto simp add: mem-infmal-iff [symmetric]*)
apply (*rule-tac c1 = hcmod x in add-le-cancel-left [THEN iffD1]*)
apply *auto*
done

lemma *approx-hcmod-add-hcmod*: $u \approx 0 \implies hcmod(x + u) \approx hcmod x$
apply (*rule approx-minus-iff [THEN iffD2]*)
apply (*auto intro: Infinitesimal-hcmod-add-diff simp add: mem-infmal-iff [symmetric]*)
done

15.5 Zero is the Only Infinitesimal Complex Number

lemma *Infinitesimal-less-SComplex*:

$[[x \in SComplex; y \in Infinitesimal; 0 < hmod x]] \implies hmod y < hmod x$
by (*auto intro: Infinitesimal-less-SReal SReal-hmod-SComplex simp add: Infinitesimal-hmod-iff*)

lemma *SComplex-Int-Infinitesimal-zero*: $SComplex \text{ Int } Infinitesimal = \{0\}$

by (*auto simp add: Standard-def Infinitesimal-hmod-iff*)

lemma *SComplex-Infinitesimal-zero*:

$[[x \in SComplex; x \in Infinitesimal]] \implies x = 0$
by (*cut-tac SComplex-Int-Infinitesimal-zero, blast*)

lemma *SComplex-HFinite-diff-Infinitesimal*:

$[[x \in SComplex; x \neq 0]] \implies x \in HFinite - Infinitesimal$
by (*auto dest: SComplex-Infinitesimal-zero Standard-subset-HFinite [THEN subsetD]*)

lemma *hcomplex-of-complex-HFinite-diff-Infinitesimal*:

$hcomplex\text{-of-complex } x \neq 0$
 $\implies hcomplex\text{-of-complex } x \in HFinite - Infinitesimal$
by (*rule SComplex-HFinite-diff-Infinitesimal, auto*)

lemma *numeral-not-Infinitesimal [simp]*:

$numeral\ w \neq (0::hcomplex) \implies (numeral\ w::hcomplex) \notin Infinitesimal$
by (*fast dest: Standard-numeral [THEN SComplex-Infinitesimal-zero]*)

lemma *approx-SComplex-not-zero*:

$[[y \in SComplex; x \approx y; y \neq 0]] \implies x \neq 0$
by (*auto dest: SComplex-Infinitesimal-zero approx-sym [THEN mem-infmal-iff [THEN iffD2]]*)

lemma *SComplex-approx-iff*:

$[[x \in SComplex; y \in SComplex]] \implies (x \approx y) = (x = y)$
by (*auto simp add: Standard-def*)

lemma *numeral-Infinitesimal-iff [simp]*:

$((numeral\ w :: hcomplex) \in Infinitesimal) =$
 $(numeral\ w = (0::hcomplex))$

apply (*rule iffI*)

apply (*fast dest: Standard-numeral [THEN SComplex-Infinitesimal-zero]*)

apply (*simp (no-asm-simp)*)

done

lemma *approx-unique-complex*:

$[[r \in SComplex; s \in SComplex; r \approx x; s \approx x]] \implies r = s$
by (*blast intro: SComplex-approx-iff [THEN iffD1] approx-trans2*)

15.6 Properties of hRe , hIm and $HComplex$

lemma *abs-hRe-le-hcmod*: $\bigwedge x. |hRe\ x| \leq hcmod\ x$
by *transfer* (rule *abs-Re-le-cmod*)

lemma *abs-hIm-le-hcmod*: $\bigwedge x. |hIm\ x| \leq hcmod\ x$
by *transfer* (rule *abs-Im-le-cmod*)

lemma *Infinesimal-hRe*: $x \in \text{Infinesimal} \implies hRe\ x \in \text{Infinesimal}$
apply (rule *InfinesimalI2*, *simp*)
apply (rule *order-le-less-trans* [*OF* *abs-hRe-le-hcmod*])
apply (erule (1) *InfinesimalD2*)
done

lemma *Infinesimal-hIm*: $x \in \text{Infinesimal} \implies hIm\ x \in \text{Infinesimal}$
apply (rule *InfinesimalI2*, *simp*)
apply (rule *order-le-less-trans* [*OF* *abs-hIm-le-hcmod*])
apply (erule (1) *InfinesimalD2*)
done

lemma *real-sqrt-lessI*: $\llbracket 0 < u; x < u^2 \rrbracket \implies \text{sqrt}\ x < u$

by (frule *real-sqrt-less-mono*) *simp*

lemma *hypreal-sqrt-lessI*:
 $\bigwedge x\ u. \llbracket 0 < u; x < u^2 \rrbracket \implies (*f*\ \text{sqrt})\ x < u$
by *transfer* (rule *real-sqrt-lessI*)

lemma *hypreal-sqrt-ge-zero*: $\bigwedge x. 0 \leq x \implies 0 \leq (*f*\ \text{sqrt})\ x$
by *transfer* (rule *real-sqrt-ge-zero*)

lemma *Infinesimal-sqrt*:
 $\llbracket x \in \text{Infinesimal}; 0 \leq x \rrbracket \implies (*f*\ \text{sqrt})\ x \in \text{Infinesimal}$
apply (rule *InfinesimalI2*)
apply (drule-tac $r=r^2$ **in** *InfinesimalD2*, *simp*)
apply (*simp* *add*: *hypreal-sqrt-ge-zero*)
apply (rule *hypreal-sqrt-lessI*, *simp-all*)
done

lemma *Infinesimal-HComplex*:
 $\llbracket x \in \text{Infinesimal}; y \in \text{Infinesimal} \rrbracket \implies HComplex\ x\ y \in \text{Infinesimal}$
apply (rule *Infinesimal-hcmod-iff* [*THEN* *iffD2*])
apply (*simp* *add*: *hcmod-i*)
apply (rule *Infinesimal-add*)
apply (erule *Infinesimal-hrealpow*, *simp*)
apply (erule *Infinesimal-hrealpow*, *simp*)
done

lemma *hcomplex-Infinesimal-iff*:
 $(x \in \text{Infinesimal}) = (hRe\ x \in \text{Infinesimal} \wedge hIm\ x \in \text{Infinesimal})$

apply (*safe intro!*: *Infinitesimal-hRe Infinitesimal-hIm*)
apply (*drule* (1) *Infinitesimal-HComplex, simp*)
done

lemma *hRe-diff* [*simp*]: $\bigwedge x y. hRe (x - y) = hRe x - hRe y$
by *transfer simp*

lemma *hIm-diff* [*simp*]: $\bigwedge x y. hIm (x - y) = hIm x - hIm y$
by *transfer simp*

lemma *approx-hRe*: $x \approx y \implies hRe x \approx hRe y$
unfolding *approx-def* **by** (*drule Infinitesimal-hRe*) *simp*

lemma *approx-hIm*: $x \approx y \implies hIm x \approx hIm y$
unfolding *approx-def* **by** (*drule Infinitesimal-hIm*) *simp*

lemma *approx-HComplex*:
 $\llbracket a \approx b; c \approx d \rrbracket \implies HComplex a c \approx HComplex b d$
unfolding *approx-def* **by** (*simp add: Infinitesimal-HComplex*)

lemma *hcomplex-approx-iff*:
 $(x \approx y) = (hRe x \approx hRe y \wedge hIm x \approx hIm y)$
unfolding *approx-def* **by** (*simp add: hcomplex-Infinitesimal-iff*)

lemma *HFinite-hRe*: $x \in HFinite \implies hRe x \in HFinite$
apply (*auto simp add: HFinite-def SReal-def*)
apply (*rule-tac x=star-of r in exI, simp*)
apply (*erule order-le-less-trans [OF abs-hRe-le-hcmod]*)
done

lemma *HFinite-hIm*: $x \in HFinite \implies hIm x \in HFinite$
apply (*auto simp add: HFinite-def SReal-def*)
apply (*rule-tac x=star-of r in exI, simp*)
apply (*erule order-le-less-trans [OF abs-hIm-le-hcmod]*)
done

lemma *HFinite-HComplex*:
 $\llbracket x \in HFinite; y \in HFinite \rrbracket \implies HComplex x y \in HFinite$
apply (*subgoal-tac HComplex x 0 + HComplex 0 y \in HFinite, simp*)
apply (*rule HFinite-add*)
apply (*simp add: HFinite-hcmod-iff hcmod-i*)
apply (*simp add: HFinite-hcmod-iff hcmod-i*)
done

lemma *hcomplex-HFinite-iff*:
 $(x \in HFinite) = (hRe x \in HFinite \wedge hIm x \in HFinite)$
apply (*safe intro!: HFinite-hRe HFinite-hIm*)
apply (*drule* (1) *HFinite-HComplex, simp*)
done

lemma *hcomplex-HInfinite-iff*:
 $(x \in HInfinite) = (hRe\ x \in HInfinite \vee hIm\ x \in HInfinite)$
by (*simp add: HInfinite-HFinite-iff hcomplex-HFinite-iff*)

lemma *hcomplex-of-hypreal-approx-iff* [*simp*]:
 $(hcomplex\ of\ hypreal\ x \approx hcomplex\ of\ hypreal\ z) = (x \approx z)$
by (*simp add: hcomplex-approx-iff*)

lemma *Standard-HComplex*:
 $\llbracket x \in Standard; y \in Standard \rrbracket \implies HComplex\ x\ y \in Standard$
by (*simp add: HComplex-def*)

lemma *stc-part-Ex*: $x: HFinite \implies \exists t \in SComplex. x \approx t$
apply (*simp add: hcomplex-HFinite-iff hcomplex-approx-iff*)
apply (*rule-tac x = HComplex (st (hRe x)) (st (hIm x)) in bexI*)
apply (*simp add: st-approx-self [THEN approx-sym]*)
apply (*simp add: Standard-HComplex st-SReal [unfolded Reals-eq-Standard]*)
done

lemma *stc-part-Ex1*: $x: HFinite \implies \exists! t. t \in SComplex \ \& \ x \approx t$
apply (*drule stc-part-Ex, safe*)
apply (*drule-tac [2] approx-sym, drule-tac [2] approx-sym, drule-tac [2] approx-sym*)
apply (*auto intro!: approx-unique-complex*)
done

lemmas *hcomplex-of-complex-approx-inverse =*
hcomplex-of-complex-HFinite-diff-Infinitesimal [THEN [2] approx-inverse]

15.7 Theorems About Monads

lemma *monad-zero-hcmod-iff*: $(x \in monad\ 0) = (hcmod\ x: monad\ 0)$
by (*simp add: Infinitesimal-monad-zero-iff [symmetric] Infinitesimal-hcmod-iff*)

15.8 Theorems About Standard Part

lemma *stc-approx-self*: $x \in HFinite \implies stc\ x \approx x$
apply (*simp add: stc-def*)
apply (*frule stc-part-Ex, safe*)
apply (*rule someI2*)
apply (*auto intro: approx-sym*)
done

lemma *stc-SComplex*: $x \in HFinite \implies stc\ x \in SComplex$
apply (*simp add: stc-def*)
apply (*frule stc-part-Ex, safe*)
apply (*rule someI2*)
apply (*auto intro: approx-sym*)
done

lemma *stc-HFinite*: $x \in \mathit{HFinite} \implies \mathit{stc} x \in \mathit{HFinite}$
by (*erule stc-SComplex* [THEN *Standard-subset-HFinite* [THEN *subsetD*]])

lemma *stc-unique*: $\llbracket y \in \mathit{SComplex}; y \approx x \rrbracket \implies \mathit{stc} x = y$
apply (*frule Standard-subset-HFinite* [THEN *subsetD*])
apply (*drule* (1) *approx-HFinite*)
apply (*unfold stc-def*)
apply (*rule some-equality*)
apply (*auto intro: approx-unique-complex*)
done

lemma *stc-SComplex-eq* [*simp*]: $x \in \mathit{SComplex} \implies \mathit{stc} x = x$
apply (*erule stc-unique*)
apply (*rule approx-refl*)
done

lemma *stc-hcomplex-of-complex*:
 $\mathit{stc} (\mathit{hcomplex-of-complex} x) = \mathit{hcomplex-of-complex} x$
by *auto*

lemma *stc-eq-approx*:
 $\llbracket x \in \mathit{HFinite}; y \in \mathit{HFinite}; \mathit{stc} x = \mathit{stc} y \rrbracket \implies x \approx y$
by (*auto dest!: stc-approx-self elim!: approx-trans3*)

lemma *approx-stc-eq*:
 $\llbracket x \in \mathit{HFinite}; y \in \mathit{HFinite}; x \approx y \rrbracket \implies \mathit{stc} x = \mathit{stc} y$
by (*blast intro: approx-trans approx-trans2 SComplex-approx-iff* [THEN *iffD1*]
dest: stc-approx-self stc-SComplex)

lemma *stc-eq-approx-iff*:
 $\llbracket x \in \mathit{HFinite}; y \in \mathit{HFinite} \rrbracket \implies (x \approx y) = (\mathit{stc} x = \mathit{stc} y)$
by (*blast intro: approx-stc-eq stc-eq-approx*)

lemma *stc-Infinitesimal-add-SComplex*:
 $\llbracket x \in \mathit{SComplex}; e \in \mathit{Infinitesimal} \rrbracket \implies \mathit{stc}(x + e) = x$
apply (*erule stc-unique*)
apply (*erule Infinitesimal-add-approx-self*)
done

lemma *stc-Infinitesimal-add-SComplex2*:
 $\llbracket x \in \mathit{SComplex}; e \in \mathit{Infinitesimal} \rrbracket \implies \mathit{stc}(e + x) = x$
apply (*erule stc-unique*)
apply (*erule Infinitesimal-add-approx-self2*)
done

lemma *HFinite-stc-Infinitesimal-add*:
 $x \in \mathit{HFinite} \implies \exists e \in \mathit{Infinitesimal}. x = \mathit{stc}(x) + e$
by (*blast dest!: stc-approx-self* [THEN *approx-sym*] *bex-Infinitesimal-iff2* [THEN

iffD2)

lemma *stc-add*:

$[[x \in \mathit{HFinite}; y \in \mathit{HFinite}]] \implies \mathit{stc} (x + y) = \mathit{stc}(x) + \mathit{stc}(y)$
by (*simp add: stc-unique stc-SComplex stc-approx-self approx-add*)

lemma *stc-numeral* [*simp*]: $\mathit{stc} (\mathit{numeral} w) = \mathit{numeral} w$

by (*rule Standard-numeral [THEN stc-SComplex-eq]*)

lemma *stc-zero* [*simp*]: $\mathit{stc} 0 = 0$

by *simp*

lemma *stc-one* [*simp*]: $\mathit{stc} 1 = 1$

by *simp*

lemma *stc-minus*: $y \in \mathit{HFinite} \implies \mathit{stc}(-y) = -\mathit{stc}(y)$

by (*simp add: stc-unique stc-SComplex stc-approx-self approx-minus*)

lemma *stc-diff*:

$[[x \in \mathit{HFinite}; y \in \mathit{HFinite}]] \implies \mathit{stc} (x - y) = \mathit{stc}(x) - \mathit{stc}(y)$
by (*simp add: stc-unique stc-SComplex stc-approx-self approx-diff*)

lemma *stc-mult*:

$[[x \in \mathit{HFinite}; y \in \mathit{HFinite}]]$
 $\implies \mathit{stc} (x * y) = \mathit{stc}(x) * \mathit{stc}(y)$
by (*simp add: stc-unique stc-SComplex stc-approx-self approx-mult-HFinite*)

lemma *stc-Infinitesimal*: $x \in \mathit{Infinitesimal} \implies \mathit{stc} x = 0$

by (*simp add: stc-unique mem-infmal-iff*)

lemma *stc-not-Infinitesimal*: $\mathit{stc}(x) \neq 0 \implies x \notin \mathit{Infinitesimal}$

by (*fast intro: stc-Infinitesimal*)

lemma *stc-inverse*:

$[[x \in \mathit{HFinite}; \mathit{stc} x \neq 0]]$
 $\implies \mathit{stc}(\mathit{inverse} x) = \mathit{inverse} (\mathit{stc} x)$
apply (*drule stc-not-Infinitesimal*)
apply (*simp add: stc-unique stc-SComplex stc-approx-self approx-inverse*)
done

lemma *stc-divide* [*simp*]:

$[[x \in \mathit{HFinite}; y \in \mathit{HFinite}; \mathit{stc} y \neq 0]]$
 $\implies \mathit{stc}(x/y) = (\mathit{stc} x) / (\mathit{stc} y)$
by (*simp add: divide-inverse stc-mult stc-not-Infinitesimal HFinite-inverse stc-inverse*)

lemma *stc-idempotent* [*simp*]: $x \in \mathit{HFinite} \implies \mathit{stc}(\mathit{stc}(x)) = \mathit{stc}(x)$

by (*blast intro: stc-HFinite stc-approx-self approx-stc-eq*)

lemma *HFinite-HFinite-hcomplex-of-hypreal*:

$z \in HFinite \implies hcomplex\text{-of-hypreal } z \in HFinite$
by (*simp add: hcomplex-HFinite-iff*)

lemma *SComplex-SReal-hcomplex-of-hypreal*:
 $x \in \mathbb{R} \implies hcomplex\text{-of-hypreal } x \in SComplex$
apply (*rule Standard-of-hypreal*)
apply (*simp add: Reals-eq-Standard*)
done

lemma *stc-hcomplex-of-hypreal*:
 $z \in HFinite \implies stc(hcomplex\text{-of-hypreal } z) = hcomplex\text{-of-hypreal } (st z)$
apply (*rule stc-unique*)
apply (*rule SComplex-SReal-hcomplex-of-hypreal*)
apply (*erule st-SReal*)
apply (*simp add: hcomplex-of-hypreal-approx-iff st-approx-self*)
done

lemma *Infinitesimal-hcnj-iff [simp]*:
 $(hcnj z \in Infinitesimal) = (z \in Infinitesimal)$
by (*simp add: Infinitesimal-hcmod-iff*)

lemma *Infinitesimal-hcomplex-of-hypreal-epsilon [simp]*:
 $hcomplex\text{-of-hypreal } \varepsilon \in Infinitesimal$
by (*simp add: Infinitesimal-hcmod-iff*)

end

16 Star-transforms in NSA, Extending Sets of Complex Numbers and Complex Functions

theory *CStar*
imports *NSCA*
begin

16.1 Properties of the *-Transform Applied to Sets of Reals

lemma *STARC-hcomplex-of-complex-Int: ** X \cap SComplex = hcomplex-of-complex ‘ X*
by (*auto simp: Standard-def*)

lemma *lemma-not-hcomplexA: $x \notin hcomplex\text{-of-complex } ‘ A \implies \forall y \in A. x \neq hcomplex\text{-of-complex } y$*
by *auto*

16.2 Theorems about Nonstandard Extensions of Functions

lemma *starfunC-hcpow: $\bigwedge Z. (*f* (\lambda z. z \hat{ } n)) Z = Z \text{ pow hypnat-of-nat } n$*

by *transfer (rule refl)*

lemma *starfunCR-cmod*: $*f* \text{ cmod} = \text{hcmmod}$
 by *transfer (rule refl)*

16.3 Internal Functions - Some Redundancy With $*f*$ Now

lemma *starfun-Re*: $(*f* (\lambda x. \text{Re } (f x))) = (\lambda x. \text{hRe } ((*f* f) x))$
 by *transfer (rule refl)*

lemma *starfun-Im*: $(*f* (\lambda x. \text{Im } (f x))) = (\lambda x. \text{hIm } ((*f* f) x))$
 by *transfer (rule refl)*

lemma *starfunC-eq-Re-Im-iff*:
 $(*f* f) x = z \iff (*f* (\lambda x. \text{Re } (f x))) x = \text{hRe } z \wedge (*f* (\lambda x. \text{Im } (f x))) x = \text{hIm } z$
 by (*simp add: hcomplex-hRe-hIm-cancel-iff starfun-Re starfun-Im*)

lemma *starfunC-approx-Re-Im-iff*:
 $(*f* f) x \approx z \iff (*f* (\lambda x. \text{Re } (f x))) x \approx \text{hRe } z \wedge (*f* (\lambda x. \text{Im } (f x))) x \approx \text{hIm } z$
 by (*simp add: hcomplex-approx-iff starfun-Re starfun-Im*)

end

17 Limits, Continuity and Differentiation for Complex Functions

theory *CLim*
 imports *CStar*
 begin

declare *hypreal-epsilon-not-zero* [*simp*]

lemma *lemma-complex-mult-inverse-squared* [*simp*]: $x \neq 0 \implies x * (\text{inverse } x)^2 = \text{inverse } x$
 for $x :: \text{complex}$
 by (*simp add: numeral-2-eq-2*)

Changing the quantified variable. Install earlier?

lemma *all-shift*: $(\forall x::'a::\text{comm-ring-1}. P x) \iff (\forall x. P (x - a))$
 apply *auto*
 apply (*drule-tac x = x + a in spec*)
 apply (*simp add: add.assoc*)
 done

lemma *complex-add-minus-iff* [*simp*]: $x + - a = 0 \longleftrightarrow x = a$
for $x a :: \text{complex}$
by (*simp add: diff-eq-eq*)

lemma *complex-add-eq-0-iff* [*iff*]: $x + y = 0 \longleftrightarrow y = - x$
for $x y :: \text{complex}$
apply *auto*
apply (*drule sym [THEN diff-eq-eq [THEN iffD2]]*)
apply *auto*
done

17.1 Limit of Complex to Complex Function

lemma *NSLIM-Re*: $f -a \rightarrow_{NS} L \implies (\lambda x. \text{Re } (f x)) -a \rightarrow_{NS} \text{Re } L$
by (*simp add: NSLIM-def starfunC-approx-Re-Im-iff hRe-hcomplex-of-complex*)

lemma *NSLIM-Im*: $f -a \rightarrow_{NS} L \implies (\lambda x. \text{Im } (f x)) -a \rightarrow_{NS} \text{Im } L$
by (*simp add: NSLIM-def starfunC-approx-Re-Im-iff hIm-hcomplex-of-complex*)

lemma *LIM-Re*: $f -a \rightarrow L \implies (\lambda x. \text{Re } (f x)) -a \rightarrow \text{Re } L$
for $f :: 'a :: \text{real-normed-vector} \Rightarrow \text{complex}$
by (*simp add: LIM-NSLIM-iff NSLIM-Re*)

lemma *LIM-Im*: $f -a \rightarrow L \implies (\lambda x. \text{Im } (f x)) -a \rightarrow \text{Im } L$
for $f :: 'a :: \text{real-normed-vector} \Rightarrow \text{complex}$
by (*simp add: LIM-NSLIM-iff NSLIM-Im*)

lemma *LIM-cnj*: $f -a \rightarrow L \implies (\lambda x. \text{cnj } (f x)) -a \rightarrow \text{cnj } L$
for $f :: 'a :: \text{real-normed-vector} \Rightarrow \text{complex}$
by (*simp add: LIM-eq complex-cnj-diff [symmetric] del: complex-cnj-diff*)

lemma *LIM-cnj-iff*: $((\lambda x. \text{cnj } (f x)) -a \rightarrow \text{cnj } L) \longleftrightarrow f -a \rightarrow L$
for $f :: 'a :: \text{real-normed-vector} \Rightarrow \text{complex}$
by (*simp add: LIM-eq complex-cnj-diff [symmetric] del: complex-cnj-diff*)

lemma *starfun-norm*: $(*f* (\lambda x. \text{norm } (f x))) = (\lambda x. \text{hnorm } ((*f* f) x))$
by *transfer (rule refl)*

lemma *star-of-Re* [*simp*]: $\text{star-of } (\text{Re } x) = \text{hRe } (\text{star-of } x)$
by *transfer (rule refl)*

lemma *star-of-Im* [*simp*]: $\text{star-of } (\text{Im } x) = \text{hIm } (\text{star-of } x)$
by *transfer (rule refl)*

Another equivalence result.

lemma *NSCLIM-NSCRLIM-iff*: $f -x \rightarrow_{NS} L \longleftrightarrow (\lambda y. \text{cmod } (f y - L)) -x \rightarrow_{NS} 0$
by (*simp add: NSLIM-def starfun-norm approx-approx-zero-iff [symmetric] approx-minus-iff [symmetric]*)

Much, much easier standard proof.

lemma *CLIM-CRLIM-iff*: $f -x \rightarrow L \longleftrightarrow (\lambda y. \text{cmod } (f y - L)) -x \rightarrow 0$
for $f :: 'a::\text{real-normed-vector} \Rightarrow \text{complex}$
by (*simp add: LIM-eq*)

So this is nicer nonstandard proof.

lemma *NSCLIM-NSCRLIM-iff2*: $f -x \rightarrow_{NS} L \longleftrightarrow (\lambda y. \text{cmod } (f y - L)) -x \rightarrow_{NS} 0$
by (*simp add: LIM-NSLIM-iff [symmetric] CLIM-CRLIM-iff*)

lemma *NSLIM-NSCRLIM-Re-Im-iff*:
 $f -a \rightarrow_{NS} L \longleftrightarrow (\lambda x. \text{Re } (f x)) -a \rightarrow_{NS} \text{Re } L \wedge (\lambda x. \text{Im } (f x)) -a \rightarrow_{NS} \text{Im } L$
apply (*auto intro: NSLIM-Re NSLIM-Im*)
apply (*auto simp add: NSLIM-def starfun-Re starfun-Im*)
apply (*auto dest!: spec*)
apply (*simp add: hcomplex-approx-iff*)
done

lemma *LIM-CRLIM-Re-Im-iff*: $f -a \rightarrow L \longleftrightarrow (\lambda x. \text{Re } (f x)) -a \rightarrow \text{Re } L \wedge (\lambda x. \text{Im } (f x)) -a \rightarrow \text{Im } L$
for $f :: 'a::\text{real-normed-vector} \Rightarrow \text{complex}$
by (*simp add: LIM-NSLIM-iff NSLIM-NSCRLIM-Re-Im-iff*)

17.2 Continuity

lemma *NSLIM-isContc-iff*: $f -a \rightarrow_{NS} f a \longleftrightarrow (\lambda h. f (a + h)) -0 \rightarrow_{NS} f a$
by (*rule NSLIM-h-iff*)

17.3 Functions from Complex to Reals

lemma *isNSContCR-cmod [simp]*: $\text{isNSCont cmod } a$
by (*auto intro: approx-hnorm*)
simp: starfunCR-cmod hmod-hcomplex-of-complex [symmetric] isNSCont-def)

lemma *isContCR-cmod [simp]*: $\text{isCont cmod } a$
by (*simp add: isNSCont-isCont-iff [symmetric]*)

lemma *isCont-Re*: $\text{isCont } f a \Longrightarrow \text{isCont } (\lambda x. \text{Re } (f x)) a$
for $f :: 'a::\text{real-normed-vector} \Rightarrow \text{complex}$
by (*simp add: isCont-def LIM-Re*)

lemma *isCont-Im*: $\text{isCont } f a \Longrightarrow \text{isCont } (\lambda x. \text{Im } (f x)) a$
for $f :: 'a::\text{real-normed-vector} \Rightarrow \text{complex}$
by (*simp add: isCont-def LIM-Im*)

17.4 Differentiation of Natural Number Powers

lemma *CDERIV-pow [simp]*: $\text{DERIV } (\lambda x. x \wedge n) x :> \text{complex-of-real } (\text{real } n) * (x \wedge (n - \text{Suc } 0))$

```

apply (induct n)
apply (drule-tac [2] DERIV-ident [THEN DERIV-mult])
apply (auto simp add: distrib-right of-nat-Suc)
apply (case-tac n)
apply (auto simp add: ac-simps)
done

```

Nonstandard version.

```

lemma NSCDERIV-pow: NSDERIV ( $\lambda x. x ^ n$ )  $x$  :=> complex-of-real (real n) *
( $x ^ (n - 1)$ )
by (metis CDERIV-pow NSDERIV-DERIV-iff One-nat-def)

```

Can't relax the premise $x \neq (0::'a)$: it isn't continuous at zero.

```

lemma NSCDERIV-inverse:  $x \neq 0 \implies$  NSDERIV ( $\lambda x. \text{inverse } x$ )  $x$  :=> - (inverse
 $x$ )2
for  $x :: \text{complex}$ 
unfolding numeral-2-eq-2 by (rule NSDERIV-inverse)

```

```

lemma CDERIV-inverse:  $x \neq 0 \implies$  DERIV ( $\lambda x. \text{inverse } x$ )  $x$  :=> - (inverse  $x$ )2
for  $x :: \text{complex}$ 
unfolding numeral-2-eq-2 by (rule DERIV-inverse)

```

17.5 Derivative of Reciprocals (Function *inverse*)

```

lemma CDERIV-inverse-fun:
DERIV  $f x$  :=>  $d \implies f x \neq 0 \implies$  DERIV ( $\lambda x. \text{inverse } (f x)$ )  $x$  :=> - (d * inverse
 $((f x)^2)$ )
for  $x :: \text{complex}$ 
unfolding numeral-2-eq-2 by (rule DERIV-inverse-fun)

```

```

lemma NSCDERIV-inverse-fun:
NSDERIV  $f x$  :=>  $d \implies f x \neq 0 \implies$  NSDERIV ( $\lambda x. \text{inverse } (f x)$ )  $x$  :=> - (d *
inverse  $((f x)^2)$ )
for  $x :: \text{complex}$ 
unfolding numeral-2-eq-2 by (rule NSDERIV-inverse-fun)

```

17.6 Derivative of Quotient

```

lemma CDERIV-quotient:
DERIV  $f x$  :=>  $d \implies$  DERIV  $g x$  :=>  $e \implies g(x) \neq 0 \implies$ 
DERIV ( $\lambda y. f y / g y$ )  $x$  :=> (d *  $g x - (e * f x) / (g x)^2$ )
for  $x :: \text{complex}$ 
unfolding numeral-2-eq-2 by (rule DERIV-quotient)

```

```

lemma NSCDERIV-quotient:
NSDERIV  $f x$  :=>  $d \implies$  NSDERIV  $g x$  :=>  $e \implies g x \neq (0::\text{complex}) \implies$ 
NSDERIV ( $\lambda y. f y / g y$ )  $x$  :=> (d *  $g x - (e * f x) / (g x)^2$ )
unfolding numeral-2-eq-2 by (rule NSDERIV-quotient)

```

17.7 Caratheodory Formulation of Derivative at a Point: Standard Proof

lemma *CARAT-CDERIVD*:

$(\forall z. f z - f x = g z * (z - x)) \wedge isNSCont\ g\ x \wedge g\ x = l \implies NSDERIV\ f\ x\ :>$
 l

by *clarify* (rule *CARAT-DERIVD*)

end

18 Logarithms: Non-Standard Version

theory *HLog*

imports *HTranscendental*

begin

lemma *epsilon-ge-zero* [*simp*]: $0 \leq \varepsilon$

by (*simp add: epsilon-def star-n-zero-num star-n-le*)

lemma *hypfinite-witness*: $\varepsilon \in \{x. 0 \leq x \wedge x \in HFinite\}$

by *auto*

definition *powhr* :: $hypreal \Rightarrow hypreal \Rightarrow hypreal$ (**infixr** *powhr* 80)

where [*transfer-unfold*]: $x\ powhr\ a = starfun2\ (op\ powr)\ x\ a$

definition *hlog* :: $hypreal \Rightarrow hypreal \Rightarrow hypreal$

where [*transfer-unfold*]: $hlog\ a\ x = starfun2\ log\ a\ x$

lemma *powhr*: $(star-n\ X)\ powhr\ (star-n\ Y) = star-n\ (\lambda n. (X\ n)\ powr\ (Y\ n))$

by (*simp add: powhr-def starfun2-star-n*)

lemma *powhr-one-eq-one* [*simp*]: $\bigwedge a. 1\ powhr\ a = 1$

by *transfer simp*

lemma *powhr-mult*: $\bigwedge a\ x\ y. 0 < x \implies 0 < y \implies (x * y)\ powhr\ a = (x\ powhr\ a) * (y\ powhr\ a)$

by *transfer (simp add: powr-mult)*

lemma *powhr-gt-zero* [*simp*]: $\bigwedge a\ x. 0 < x\ powhr\ a \longleftrightarrow x \neq 0$

by *transfer simp*

lemma *powhr-not-zero* [*simp*]: $\bigwedge a\ x. x\ powhr\ a \neq 0 \longleftrightarrow x \neq 0$

by *transfer simp*

lemma *powhr-divide*: $\bigwedge a\ x\ y. 0 < x \implies 0 < y \implies (x / y)\ powhr\ a = (x\ powhr\ a) / (y\ powhr\ a)$

by *transfer* (rule *powr-divide*)

lemma *powhr-add*: $\bigwedge a b x. x \text{ powhr } (a + b) = (x \text{ powhr } a) * (x \text{ powhr } b)$
by *transfer* (rule *powr-add*)

lemma *powhr-powhr*: $\bigwedge a b x. (x \text{ powhr } a) \text{ powhr } b = x \text{ powhr } (a * b)$
by *transfer* (rule *powr-powr*)

lemma *powhr-powhr-swap*: $\bigwedge a b x. (x \text{ powhr } a) \text{ powhr } b = (x \text{ powhr } b) \text{ powhr } a$
by *transfer* (rule *powr-powr-swap*)

lemma *powhr-minus*: $\bigwedge a x. x \text{ powhr } (- a) = \text{inverse } (x \text{ powhr } a)$
by *transfer* (rule *powr-minus*)

lemma *powhr-minus-divide*: $x \text{ powhr } (- a) = 1 / (x \text{ powhr } a)$
by (*simp add: divide-inverse powhr-minus*)

lemma *powhr-less-mono*: $\bigwedge a b x. a < b \implies 1 < x \implies x \text{ powhr } a < x \text{ powhr } b$
by *transfer simp*

lemma *powhr-less-cancel*: $\bigwedge a b x. x \text{ powhr } a < x \text{ powhr } b \implies 1 < x \implies a < b$
by *transfer simp*

lemma *powhr-less-cancel-iff* [*simp*]: $1 < x \implies x \text{ powhr } a < x \text{ powhr } b \longleftrightarrow a < b$
by (*blast intro: powhr-less-cancel powhr-less-mono*)

lemma *powhr-le-cancel-iff* [*simp*]: $1 < x \implies x \text{ powhr } a \leq x \text{ powhr } b \longleftrightarrow a \leq b$
by (*simp add: linorder-not-less [symmetric]*)

lemma *hlog*: $\text{hlog } (\text{star-}n \ X) (\text{star-}n \ Y) = \text{star-}n \ (\lambda n. \text{log } (X \ n) \ (Y \ n))$
by (*simp add: hlog-def starfun2-star-n*)

lemma *hlog-starfun-ln*: $\bigwedge x. (*f* \ \ln) \ x = \text{hlog } ((*f* \ \exp) \ 1) \ x$
by *transfer* (rule *log-ln*)

lemma *powhr-hlog-cancel* [*simp*]: $\bigwedge a x. 0 < a \implies a \neq 1 \implies 0 < x \implies a \text{ powhr } (\text{hlog } a \ x) = x$
by *transfer simp*

lemma *hlog-powhr-cancel* [*simp*]: $\bigwedge a y. 0 < a \implies a \neq 1 \implies \text{hlog } a \ (a \text{ powhr } y) = y$
by *transfer simp*

lemma *hlog-mult*:
 $\bigwedge a x y. 0 < a \implies a \neq 1 \implies 0 < x \implies 0 < y \implies \text{hlog } a \ (x * y) = \text{hlog } a \ x + \text{hlog } a \ y$
by *transfer* (rule *log-mult*)

lemma *hlog-as-starfun*: $\bigwedge a x. 0 < a \implies a \neq 1 \implies \text{hlog } a \ x = (*f* \ \ln) \ x / ($

$*f* \ln) a$

by *transfer (simp add: log-def)*

lemma *hlog-eq-div-starfun-ln-mult-hlog*:

$\bigwedge a b x. 0 < a \implies a \neq 1 \implies 0 < b \implies b \neq 1 \implies 0 < x \implies$
 $\text{hlog } a \ x = ((*f* \ln) b / (*f* \ln) a) * \text{hlog } b \ x$

by *transfer (rule log-eq-div-ln-mult-log)*

lemma *powhr-as-starfun*: $\bigwedge a x. x \text{ powhr } a = (\text{if } x = 0 \text{ then } 0 \text{ else } (*f* \text{ exp}) (a * (*f* \text{ real-ln}) x))$

by *transfer (simp add: powr-def)*

lemma *HInfinite-powhr*:

$x \in \text{HInfinite} \implies 0 < x \implies a \in \text{HFinite} - \text{Infinitesimal} \implies 0 < a \implies x \text{ powhr } a \in \text{HInfinite}$

by (*auto intro!*: *starfun-ln-ge-zero starfun-ln-HInfinite*

HInfinite-HFinite-not-Infinitesimal-mult2 starfun-exp-HInfinite

simp add: order-less-imp-le HInfinite-gt-zero-gt-one powhr-as-starfun zero-le-mult-iff)

lemma *hlog-hrabs-HInfinite-Infinitesimal*:

$x \in \text{HFinite} - \text{Infinitesimal} \implies a \in \text{HInfinite} \implies 0 < a \implies \text{hlog } a \ |x| \in \text{Infinitesimal}$

apply (*frule HInfinite-gt-zero-gt-one*)

apply (*auto intro!*: *starfun-ln-HFinite-not-Infinitesimal*

HInfinite-inverse-Infinitesimal Infinitesimal-HFinite-mult2

simp add: starfun-ln-HInfinite not-Infinitesimal-not-zero

hlog-as-starfun divide-inverse)

done

lemma *hlog-HInfinite-as-starfun*: $a \in \text{HInfinite} \implies 0 < a \implies \text{hlog } a \ x = (*f* \ln) x / (*f* \ln) a$

by (*rule hlog-as-starfun auto*)

lemma *hlog-one [simp]*: $\bigwedge a. \text{hlog } a \ 1 = 0$

by *transfer simp*

lemma *hlog-eq-one [simp]*: $\bigwedge a. 0 < a \implies a \neq 1 \implies \text{hlog } a \ a = 1$

by *transfer (rule log-eq-one)*

lemma *hlog-inverse*: $0 < a \implies a \neq 1 \implies 0 < x \implies \text{hlog } a \ (\text{inverse } x) = - \text{hlog } a \ x$

by (*rule add-left-cancel [of hlog a x, THEN iffD1]*) (*simp add: hlog-mult [symmetric]*)

lemma *hlog-divide*: $0 < a \implies a \neq 1 \implies 0 < x \implies 0 < y \implies \text{hlog } a \ (x / y) = \text{hlog } a \ x - \text{hlog } a \ y$

by (*simp add: hlog-mult hlog-inverse divide-inverse*)

lemma *hlog-less-cancel-iff [simp]*:

$\bigwedge a x y. 1 < a \implies 0 < x \implies 0 < y \implies \text{hlog } a \ x < \text{hlog } a \ y \longleftrightarrow x < y$

by *transfer simp*

lemma *hlog-le-cancel-iff* [*simp*]: $1 < a \implies 0 < x \implies 0 < y \implies \text{hlog } a \ x \leq \text{hlog } a \ y \iff x \leq y$
by (*simp add: linorder-not-less* [*symmetric*])

end

theory *Hyperreal*
imports *HLog*
begin

end
theory *Hypercomplex*
imports *CLim Hyperreal*
begin

end

theory *Nonstandard-Analysis*
imports *Hypercomplex*
begin

end