

Machine Words in Isabelle/HOL

Jeremy Dawson, Paul Graunke, Brian Huffman, Gerwin Klein, and John Matthews

October 8, 2017

Abstract

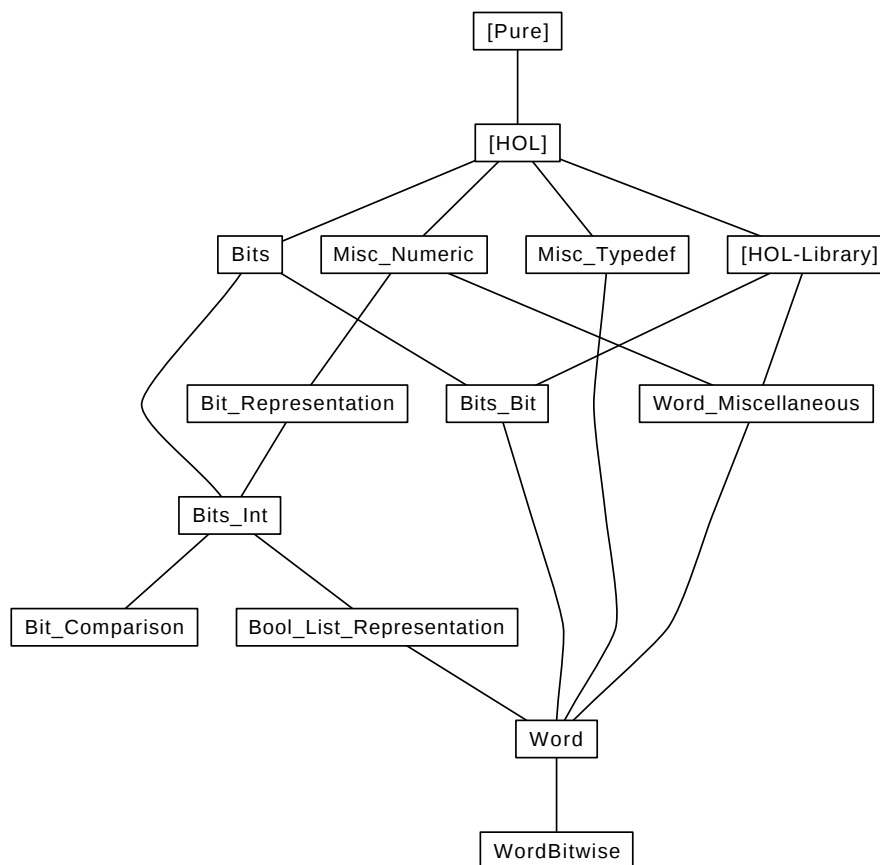
A formalisation of generic, fixed size machine words in Isabelle/HOL.
An earlier version of this formalisation is described in [1].

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1 Syntactic classes for bitwise operations

```

theory Bits
  imports Main
begin

class bit =
  fixes bitNOT :: 'a ⇒ 'a (NOT - [70] 71)
    and bitAND :: 'a ⇒ 'a ⇒ 'a (infixr AND 64)
    and bitOR :: 'a ⇒ 'a ⇒ 'a (infixr OR 59)
    and bitXOR :: 'a ⇒ 'a ⇒ 'a (infixr XOR 59)

```

We want the bitwise operations to bind slightly weaker than $+$ and $-$, but $\sim\sim$ to bind slightly stronger than $*$.

Testing and shifting operations.

```

class bits = bit +
  fixes test-bit :: 'a ⇒ nat ⇒ bool (infixl !! 100)
    and lsb :: 'a ⇒ bool
    and set-bit :: 'a ⇒ nat ⇒ bool ⇒ 'a
    and set-bits :: (nat ⇒ bool) ⇒ 'a (binder BITS 10)
    and shiftl :: 'a ⇒ nat ⇒ 'a (infixl << 55)
    and shiftr :: 'a ⇒ nat ⇒ 'a (infixl >> 55)

class bitss = bits +
  fixes msb :: 'a ⇒ bool

end

```

2 Bit operations in $\mathcal{Z}_\epsilon$

```

theory Bits-Bit
  imports Bits HOL-Library.Bit
begin

```

```

instantiation bit :: bit
begin

```

```

primrec bitNOT-bit
  where
    NOT 0 = (1::bit)
  | NOT 1 = (0::bit)

```

```

primrec bitAND-bit
  where
    0 AND y = (0::bit)
  | 1 AND y = (y::bit)

```

```

primrec bitOR-bit

```

```

where
  0 OR y = (y::bit)
| 1 OR y = (1::bit)

primrec bitXOR-bit
where
  0 XOR y = (y::bit)
| 1 XOR y = (NOT y :: bit)

instance ..

end

lemmas bit-simps =
  bitNOT-bit.simps bitAND-bit.simps bitOR-bit.simps bitXOR-bit.simps

lemma bit-extra-simps [simp]:
  x AND 0 = 0
  x AND 1 = x
  x OR 1 = 1
  x OR 0 = x
  x XOR 1 = NOT x
  x XOR 0 = x
for x :: bit
by (cases x, auto)+

lemma bit-ops-comm:
  x AND y = y AND x
  x OR y = y OR x
  x XOR y = y XOR x
for x :: bit
by (cases y, auto)+

lemma bit-ops-same [simp]:
  x AND x = x
  x OR x = x
  x XOR x = 0
for x :: bit
by (cases x, auto)+

lemma bit-not-not [simp]: NOT (NOT x) = x
for x :: bit
by (cases x) auto

lemma bit-or-def: b OR c = NOT (NOT b AND NOT c)
for b c :: bit
by (induct b) simp-all

lemma bit-xor-def: b XOR c = (b AND NOT c) OR (NOT b AND c)

```

```

for  $b\ c :: \text{bit}$ 
by (induct b) simp-all

lemma bit-NOT-eq-1-iff [simp]:  $\text{NOT } b = 1 \longleftrightarrow b = 0$ 
for  $b :: \text{bit}$ 
by (induct b) simp-all

lemma bit-AND-eq-1-iff [simp]:  $a \text{ AND } b = 1 \longleftrightarrow a = 1 \wedge b = 1$ 
for  $a\ b :: \text{bit}$ 
by (induct a) simp-all

end

```

3 Useful Numerical Lemmas

```

theory Misc-Numeric
imports Main
begin

lemma one-mod-exp-eq-one [simp]:  $1 \bmod (2 * 2 ^ n) = (1 :: \text{int})$ 
by (smt mod-pos-pos-trivial zero-less-power)

lemma mod-2-neq-1-eq-eq-0:  $k \bmod 2 \neq 1 \longleftrightarrow k \bmod 2 = 0$ 
for  $k :: \text{int}$ 
by (fact not-mod-2-eq-1-eq-0)

lemma z1pmod2:  $(2 * b + 1) \bmod 2 = (1 :: \text{int})$ 
for  $b :: \text{int}$ 
by arith

lemma diff-le-eq':  $a - b \leq c \longleftrightarrow a \leq b + c$ 
for  $a\ b\ c :: \text{int}$ 
by arith

lemma emep1:  $\text{even } n \implies \text{even } d \implies 0 \leq d \implies (n + 1) \bmod d = (n \bmod d) + 1$ 
for  $n\ d :: \text{int}$ 
by (auto simp add: pos-zmod-mult-2 add.commute dvd-def)

lemma int-mod-ge:  $a < n \implies 0 < n \implies a \leq a \bmod n$ 
for  $a\ n :: \text{int}$ 
by (metis dual-order.trans le-cases mod-pos-pos-trivial pos-mod-conj)

lemma int-mod-ge':  $b < 0 \implies 0 < n \implies b + n \leq b \bmod n$ 
for  $b\ n :: \text{int}$ 
by (metis add-less-same-cancel2 int-mod-ge mod-add-self2)

lemma int-mod-le':  $0 \leq b - n \implies b \bmod n \leq b - n$ 
for  $b\ n :: \text{int}$ 

```

```

    by (metis minus-mod-self2 zmod-le-nonneg-dividend)

lemma zless2:  $0 < (2 :: int)$ 
  by (fact zero-less-numeral)

lemma zless2p:  $0 < (2 ^ n :: int)$ 
  by arith

lemma zle2p:  $0 \leq (2 ^ n :: int)$ 
  by arith

lemma m1mod2k:  $-1 \bmod 2 ^ n = (2 ^ n - 1 :: int)$ 
  using zless2p by (rule zmod-minus1)

lemma p1mod22k':  $(1 + 2 * b) \bmod (2 * 2 ^ n) = 1 + 2 * (b \bmod 2 ^ n)$ 
  for  $b :: int$ 
  using zle2p by (rule pos-zmod-mult-2)

lemma p1mod22k:  $(2 * b + 1) \bmod (2 * 2 ^ n) = 2 * (b \bmod 2 ^ n) + 1$ 
  for  $b :: int$ 
  by (simp add: p1mod22k' add.commute)

lemma int-mod-lem:  $0 < n \implies 0 \leq b \wedge b < n \longleftrightarrow b \bmod n = b$ 
  for  $b n :: int$ 
  apply safe
  apply (erule (1) mod-pos-pos-trivial)
  apply (erule tac [!] subst)
  apply auto
  done

end

```

4 Integers as implicit bit strings

```

theory Bit-Representation
  imports Misc-Numeric
begin

```

4.1 Constructors and destructors for binary integers

```

definition Bit ::  $int \Rightarrow bool \Rightarrow int$  (infixl BIT 90)
  where  $k \text{ BIT } b = (if\ b\ then\ 1\ else\ 0) + k + k$ 

```

```

lemma Bit-B0:  $k \text{ BIT } False = k + k$ 
  by (simp add: Bit-def)

```

```

lemma Bit-B1:  $k \text{ BIT } True = k + k + 1$ 
  by (simp add: Bit-def)

```

lemma *Bit-B0-2t*: $k \text{ BIT False} = 2 * k$
by (*rule trans*, *rule Bit-B0*) *simp*

lemma *Bit-B1-2t*: $k \text{ BIT True} = 2 * k + 1$
by (*rule trans*, *rule Bit-B1*) *simp*

definition *bin-last* :: $\text{int} \Rightarrow \text{bool}$
where *bin-last* $w \longleftrightarrow w \bmod 2 = 1$

lemma *bin-last-odd*: $\text{bin-last} = \text{odd}$
by (*rule ext*) (*simp add: bin-last-def even-iff-mod-2-eq-zero*)

definition *bin-rest* :: $\text{int} \Rightarrow \text{int}$
where *bin-rest* $w = w \text{ div } 2$

lemma *bin-rl-simp* [*simp*]: $\text{bin-rest } w \text{ BIT bin-last } w = w$
unfolding *bin-rest-def bin-last-def Bit-def*
by (*cases* $w \bmod 2 = 0$) (*use div-mult-mod-eq [of w 2] in simp-all*)

lemma *bin-rest-BIT* [*simp*]: $\text{bin-rest } (x \text{ BIT } b) = x$
unfolding *bin-rest-def Bit-def*
by (*cases* b) *simp-all*

lemma *bin-last-BIT* [*simp*]: $\text{bin-last } (x \text{ BIT } b) = b$
unfolding *bin-last-def Bit-def*
by (*cases* b) *simp-all*

lemma *BIT-eq-iff* [*iff*]: $u \text{ BIT } b = v \text{ BIT } c \longleftrightarrow u = v \wedge b = c$
by (*auto simp: Bit-def*) *arith+*

lemma *BIT-bin-simps* [*simp*]:
 $\text{numeral } k \text{ BIT False} = \text{numeral } (\text{Num.Bit0 } k)$
 $\text{numeral } k \text{ BIT True} = \text{numeral } (\text{Num.Bit1 } k)$
 $(- \text{numeral } k) \text{ BIT False} = - \text{numeral } (\text{Num.Bit0 } k)$
 $(- \text{numeral } k) \text{ BIT True} = - \text{numeral } (\text{Num.BitM } k)$
unfolding *numeral.simps numeral-BitM*
by (*simp-all add: Bit-def del: arith-simps add-numeral-special diff-numeral-special*)

lemma *BIT-special-simps* [*simp*]:
shows $0 \text{ BIT False} = 0$
and $0 \text{ BIT True} = 1$
and $1 \text{ BIT False} = 2$
and $1 \text{ BIT True} = 3$
and $(- 1) \text{ BIT False} = - 2$
and $(- 1) \text{ BIT True} = - 1$
by (*simp-all add: Bit-def*)

lemma *Bit-eq-0-iff*: $w \text{ BIT } b = 0 \longleftrightarrow w = 0 \wedge \neg b$
by (*auto simp: Bit-def*) *arith*

lemma *Bit-eq-m1-iff*: $w \text{ BIT } b = -1 \longleftrightarrow w = -1 \wedge b$
by (*auto simp: Bit-def*) *arith*

lemma *BitM-inc*: $\text{Num.BitM } (\text{Num.inc } w) = \text{Num.Bit1 } w$
by (*induct w*) *simp-all*

lemma *expand-BIT*:
 $\text{numeral } (\text{Num.Bit0 } w) = \text{numeral } w \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit1 } w) = \text{numeral } w \text{ BIT True}$
 $-\text{numeral } (\text{Num.Bit0 } w) = (-\text{numeral } w) \text{ BIT False}$
 $-\text{numeral } (\text{Num.Bit1 } w) = (-\text{numeral } (w + \text{Num.One})) \text{ BIT True}$
by (*simp-all add: add-One BitM-inc*)

lemma *bin-last-numeral-simps* [*simp*]:
 $\neg \text{bin-last } 0$
 $\text{bin-last } 1$
 $\text{bin-last } (-1)$
 bin-last Numeral1
 $\neg \text{bin-last } (\text{numeral } (\text{Num.Bit0 } w))$
 $\text{bin-last } (\text{numeral } (\text{Num.Bit1 } w))$
 $\neg \text{bin-last } (-\text{numeral } (\text{Num.Bit0 } w))$
 $\text{bin-last } (-\text{numeral } (\text{Num.Bit1 } w))$
by (*simp-all add: bin-last-def zmod-zminus1-eq-if*) (*auto simp add: divmod-def*)

lemma *bin-rest-numeral-simps* [*simp*]:
 $\text{bin-rest } 0 = 0$
 $\text{bin-rest } 1 = 0$
 $\text{bin-rest } (-1) = -1$
 $\text{bin-rest Numeral1} = 0$
 $\text{bin-rest } (\text{numeral } (\text{Num.Bit0 } w)) = \text{numeral } w$
 $\text{bin-rest } (\text{numeral } (\text{Num.Bit1 } w)) = \text{numeral } w$
 $\text{bin-rest } (-\text{numeral } (\text{Num.Bit0 } w)) = -\text{numeral } w$
 $\text{bin-rest } (-\text{numeral } (\text{Num.Bit1 } w)) = -\text{numeral } (w + \text{Num.One})$
by (*simp-all add: bin-rest-def zdiv-zminus1-eq-if*) (*auto simp add: divmod-def*)

lemma *less-Bits*: $v \text{ BIT } b < w \text{ BIT } c \longleftrightarrow v < w \vee v \leq w \wedge \neg b \wedge c$
by (*auto simp: Bit-def*)

lemma *le-Bits*: $v \text{ BIT } b \leq w \text{ BIT } c \longleftrightarrow v < w \vee v \leq w \wedge (\neg b \vee c)$
by (*auto simp: Bit-def*)

lemma *pred-BIT-simps* [*simp*]:
 $x \text{ BIT False} - 1 = (x - 1) \text{ BIT True}$
 $x \text{ BIT True} - 1 = x \text{ BIT False}$
by (*simp-all add: Bit-B0-2t Bit-B1-2t*)

lemma *succ-BIT-simps* [*simp*]:
 $x \text{ BIT False} + 1 = x \text{ BIT True}$

$x \text{ BIT True} + 1 = (x + 1) \text{ BIT False}$
by (*simp-all add: Bit-B0-2t Bit-B1-2t*)

lemma *add-BIT-simps* [*simp*]:
 $x \text{ BIT False} + y \text{ BIT False} = (x + y) \text{ BIT False}$
 $x \text{ BIT False} + y \text{ BIT True} = (x + y) \text{ BIT True}$
 $x \text{ BIT True} + y \text{ BIT False} = (x + y) \text{ BIT True}$
 $x \text{ BIT True} + y \text{ BIT True} = (x + y + 1) \text{ BIT False}$
by (*simp-all add: Bit-B0-2t Bit-B1-2t*)

lemma *mult-BIT-simps* [*simp*]:
 $x \text{ BIT False} * y = (x * y) \text{ BIT False}$
 $x * y \text{ BIT False} = (x * y) \text{ BIT False}$
 $x \text{ BIT True} * y = (x * y) \text{ BIT False} + y$
by (*simp-all add: Bit-B0-2t Bit-B1-2t algebra-simps*)

lemma *B-mod-2'*: $X = 2 \implies (w \text{ BIT True}) \bmod X = 1 \wedge (w \text{ BIT False}) \bmod X = 0$
by (*simp add: Bit-B0 Bit-B1*)

lemma *bin-ex-rl*: $\exists w b. w \text{ BIT } b = \text{bin}$
by (*metis bin-rl-simp*)

lemma *bin-exhaust*:
assumes *that*: $\bigwedge x b. \text{bin} = x \text{ BIT } b \implies Q$
shows Q
apply (*insert bin-ex-rl [of bin]*)
apply (*erule exE*)
apply (*rule that*)
apply *force*
done

primrec *bin-nth*
where
 $Z: \text{bin-nth } w \ 0 \longleftrightarrow \text{bin-last } w$
 $| \text{Suc}: \text{bin-nth } w \ (\text{Suc } n) \longleftrightarrow \text{bin-nth } (\text{bin-rest } w) \ n$

lemma *bin-abs-lem*: $\text{bin} = (w \text{ BIT } b) \implies \text{bin} \neq -1 \longrightarrow \text{bin} \neq 0 \longrightarrow \text{nat } |w| < \text{nat } |\text{bin}|$
apply *clarsimp*
apply (*unfold Bit-def*)
apply (*cases b*)
apply (*clarsimp, arith*)
apply (*clarsimp, arith*)
done

lemma *bin-induct*:
assumes *PPls*: $P \ 0$
and *PMin*: $P \ (-1)$

```

    and PBit:  $\bigwedge bin\ bit. P\ bin \implies P\ (bin\ BIT\ bit)$ 
  shows  $P\ bin$ 
    apply (rule-tac  $P=P$  and  $a=bin$  and  $f1=nat \circ abs$  in  $wf-measure$  [THEN
wf-induct])
    apply (simp add: measure-def inv-image-def)
    apply (case-tac  $x$  rule: bin-exhaust)
    apply (frule bin-abs-lem)
    apply (auto simp add : PPls PMin PBit)
  done

```

```

lemma Bit-div2 [simp]:  $(w\ BIT\ b)\ div\ 2 = w$ 
  unfolding bin-rest-def [symmetric] by (rule bin-rest-BIT)

```

```

lemma bin-nth-eq-iff:  $bin-nth\ x = bin-nth\ y \longleftrightarrow x = y$ 
proof -
  have bin-nth-lem [rule-format]:  $\forall y. bin-nth\ x = bin-nth\ y \longrightarrow x = y$ 
  apply (induct  $x$  rule: bin-induct)
  apply safe
  apply (erule rev-mp)
  apply (induct-tac  $y$  rule: bin-induct)
  apply safe
  apply (drule-tac  $x=0$  in fun-cong, force)
  apply (erule notE, rule ext, drule-tac  $x=Suc\ x$  in fun-cong, force)
  apply (drule-tac  $x=0$  in fun-cong, force)
  apply (erule rev-mp)
  apply (induct-tac  $y$  rule: bin-induct)
  apply safe
  apply (drule-tac  $x=0$  in fun-cong, force)
  apply (erule notE, rule ext, drule-tac  $x=Suc\ x$  in fun-cong, force)
  apply (metis Bit-eq-m1-iff Z bin-last-BIT)
  apply (case-tac  $y$  rule: bin-exhaust)
  apply clarify
  apply (erule allE)
  apply (erule impE)
  prefer 2
  apply (erule conjI)
  apply (drule-tac  $x=0$  in fun-cong, force)
  apply (rule ext)
  apply (drule-tac  $x=Suc\ x$  for  $x$  in fun-cong, force)
  done
  show ?thesis
    by (auto elim: bin-nth-lem)
qed

```

```

lemmas bin-eqI = ext [THEN bin-nth-eq-iff [THEN iffD1]]

```

```

lemma bin-eq-iff:  $x = y \longleftrightarrow (\forall n. bin-nth\ x\ n = bin-nth\ y\ n)$ 
  using bin-nth-eq-iff by auto

```

lemma *bin-nth-zero* [simp]: $\neg \text{bin-nth } 0 \ n$
by (induct *n*) auto

lemma *bin-nth-1* [simp]: $\text{bin-nth } 1 \ n \longleftrightarrow n = 0$
by (cases *n*) simp-all

lemma *bin-nth-minus1* [simp]: $\text{bin-nth } (-1) \ n$
by (induct *n*) auto

lemma *bin-nth-0-BIT*: $\text{bin-nth } (w \text{ BIT } b) \ 0 \longleftrightarrow b$
by auto

lemma *bin-nth-Suc-BIT*: $\text{bin-nth } (w \text{ BIT } b) \ (\text{Suc } n) = \text{bin-nth } w \ n$
by auto

lemma *bin-nth-minus* [simp]: $0 < n \implies \text{bin-nth } (w \text{ BIT } b) \ n = \text{bin-nth } w \ (n - 1)$
by (cases *n*) auto

lemma *bin-nth-numeral*: $\text{bin-rest } x = y \implies \text{bin-nth } x \ (\text{numeral } n) = \text{bin-nth } y \ (\text{pred-numeral } n)$
by (simp add: numeral-eq-Suc)

lemmas *bin-nth-numeral-simps* [simp] =
 $\text{bin-nth-numeral } [\text{OF } \text{bin-rest-numeral-simps}(2)]$
 $\text{bin-nth-numeral } [\text{OF } \text{bin-rest-numeral-simps}(5)]$
 $\text{bin-nth-numeral } [\text{OF } \text{bin-rest-numeral-simps}(6)]$
 $\text{bin-nth-numeral } [\text{OF } \text{bin-rest-numeral-simps}(7)]$
 $\text{bin-nth-numeral } [\text{OF } \text{bin-rest-numeral-simps}(8)]$

lemmas *bin-nth-simps* =
 $\text{bin-nth.Z } \text{bin-nth.Suc } \text{bin-nth-zero } \text{bin-nth-minus1}$
 $\text{bin-nth-numeral-simps}$

4.2 Truncating binary integers

definition *bin-sign* :: $\text{int} \Rightarrow \text{int}$
where $\text{bin-sign } k = (\text{if } k \geq 0 \text{ then } 0 \text{ else } -1)$

lemma *bin-sign-simps* [simp]:
 $\text{bin-sign } 0 = 0$
 $\text{bin-sign } 1 = 0$
 $\text{bin-sign } (-1) = -1$
 $\text{bin-sign } (\text{numeral } k) = 0$
 $\text{bin-sign } (- \text{numeral } k) = -1$
 $\text{bin-sign } (w \text{ BIT } b) = \text{bin-sign } w$
by (simp-all add: bin-sign-def Bit-def)

lemma *bin-sign-rest* [simp]: $\text{bin-sign } (\text{bin-rest } w) = \text{bin-sign } w$

```

by (cases w rule: bin-exhaust) auto

primrec bintrunc :: nat  $\Rightarrow$  int  $\Rightarrow$  int
  where
    Z : bintrunc 0 bin = 0
    | Suc : bintrunc (Suc n) bin = bintrunc n (bin-rest bin) BIT (bin-last bin)

primrec sbintrunc :: nat  $\Rightarrow$  int  $\Rightarrow$  int
  where
    Z : sbintrunc 0 bin = (if bin-last bin then -1 else 0)
    | Suc : sbintrunc (Suc n) bin = sbintrunc n (bin-rest bin) BIT (bin-last bin)

lemma sign-bintr: bin-sign (bintrunc n w) = 0
  by (induct n arbitrary: w) auto

lemma bintrunc-mod2p: bintrunc n w = w mod 2 ^ n
  by (induct n arbitrary: w) (auto simp add: bin-last-def bin-rest-def Bit-def zmod-zmult2-eq)

lemma sbintrunc-mod2p: sbintrunc n w = (w + 2 ^ n) mod 2 ^ (Suc n) - 2 ^ n
  apply (induct n arbitrary: w)
  apply (auto simp add: bin-last-odd bin-rest-def Bit-def elim!: evenE oddE)
  apply (smt pos-zmod-mult-2 zle2p)
  apply (simp add: mult-mod-right)
  done

### 4.3 Simplifications for (s)bintrunc

lemma bintrunc-n-0 [simp]: bintrunc n 0 = 0
  by (induct n) auto

lemma sbintrunc-n-0 [simp]: sbintrunc n 0 = 0
  by (induct n) auto

lemma sbintrunc-n-minus1 [simp]: sbintrunc n (- 1) = -1
  by (induct n) auto

lemma bintrunc-Suc-numeral:
  bintrunc (Suc n) 1 = 1
  bintrunc (Suc n) (- 1) = bintrunc n (- 1) BIT True
  bintrunc (Suc n) (numeral (Num.Bit0 w)) = bintrunc n (numeral w) BIT False
  bintrunc (Suc n) (numeral (Num.Bit1 w)) = bintrunc n (numeral w) BIT True
  bintrunc (Suc n) (- numeral (Num.Bit0 w)) = bintrunc n (- numeral w) BIT
  False
  bintrunc (Suc n) (- numeral (Num.Bit1 w)) = bintrunc n (- numeral (w +
  Num.One)) BIT True
  by simp-all

lemma sbintrunc-0-numeral [simp]:
  sbintrunc 0 1 = -1

```

$\text{sbintrunc } 0 \text{ (numeral (Num.Bit0 } w)) = 0$
 $\text{sbintrunc } 0 \text{ (numeral (Num.Bit1 } w)) = -1$
 $\text{sbintrunc } 0 \text{ (- numeral (Num.Bit0 } w)) = 0$
 $\text{sbintrunc } 0 \text{ (- numeral (Num.Bit1 } w)) = -1$
by *simp-all*

lemma *sbintrunc-Suc-numeral*:

$\text{sbintrunc (Suc } n) 1 = 1$
 $\text{sbintrunc (Suc } n) \text{ (numeral (Num.Bit0 } w)) = \text{sbintrunc } n \text{ (numeral } w) \text{ BIT False}$
 $\text{sbintrunc (Suc } n) \text{ (numeral (Num.Bit1 } w)) = \text{sbintrunc } n \text{ (numeral } w) \text{ BIT True}$
 $\text{sbintrunc (Suc } n) \text{ (- numeral (Num.Bit0 } w)) = \text{sbintrunc } n \text{ (- numeral } w) \text{ BIT False}$
 $\text{sbintrunc (Suc } n) \text{ (- numeral (Num.Bit1 } w)) = \text{sbintrunc } n \text{ (- numeral } (w + \text{Num.One})) \text{ BIT True}$
by *simp-all*

lemma *bin-sign-lem*: $(\text{bin-sign } (\text{sbintrunc } n \text{ bin}) = -1) = \text{bin-nth bin } n$

apply (*induct n arbitrary: bin*)
apply (*case-tac bin rule: bin-exhaust, case-tac b, auto*)
done

lemma *nth-bintr*: $\text{bin-nth (bintrunc } m \text{ } w) \text{ } n \longleftrightarrow n < m \wedge \text{bin-nth } w \text{ } n$

apply (*induct n arbitrary: w m*)
apply (*case-tac m, auto*)[1]
apply (*case-tac m, auto*)[1]
done

lemma *nth-sbintr*: $\text{bin-nth (sbintrunc } m \text{ } w) \text{ } n = (\text{if } n < m \text{ then bin-nth } w \text{ } n \text{ else bin-nth } w \text{ } m)$

apply (*induct n arbitrary: w m*)
apply (*case-tac m*)
apply *simp-all*
apply (*case-tac m*)
apply *simp-all*
done

lemma *bin-nth-Bit*: $\text{bin-nth (} w \text{ BIT } b) \text{ } n \longleftrightarrow n = 0 \wedge b \vee (\exists m. n = \text{Suc } m \wedge \text{bin-nth } w \text{ } m)$

by (*cases n*) *auto*

lemma *bin-nth-Bit0*:

$\text{bin-nth (numeral (Num.Bit0 } w)) \text{ } n \longleftrightarrow$
 $(\exists m. n = \text{Suc } m \wedge \text{bin-nth (numeral } w) \text{ } m)$
using *bin-nth-Bit* [**where** $w = \text{numeral } w$ **and** $b = \text{False}$] **by** *simp*

lemma *bin-nth-Bit1*:

$\text{bin-nth (numeral (Num.Bit1 } w)) \text{ } n \longleftrightarrow$
 $n = 0 \vee (\exists m. n = \text{Suc } m \wedge \text{bin-nth (numeral } w) \text{ } m)$
using *bin-nth-Bit* [**where** $w = \text{numeral } w$ **and** $b = \text{True}$] **by** *simp*

lemma *bintrunc-bintrunc-l*: $n \leq m \implies \text{bintrunc } m (\text{bintrunc } n w) = \text{bintrunc } n w$

by (*rule bin-eqI*) (*auto simp: nth-bintr*)

lemma *sbintrunc-sbintrunc-l*: $n \leq m \implies \text{sbintrunc } m (\text{sbintrunc } n w) = \text{sbintrunc } n w$

by (*rule bin-eqI*) (*auto simp: nth-sbintr*)

lemma *bintrunc-bintrunc-ge*: $n \leq m \implies \text{bintrunc } n (\text{bintrunc } m w) = \text{bintrunc } n w$

by (*rule bin-eqI*) (*auto simp: nth-bintr*)

lemma *bintrunc-bintrunc-min* [*simp*]: $\text{bintrunc } m (\text{bintrunc } n w) = \text{bintrunc } (\min m n) w$

by (*rule bin-eqI*) (*auto simp: nth-bintr*)

lemma *sbintrunc-sbintrunc-min* [*simp*]: $\text{sbintrunc } m (\text{sbintrunc } n w) = \text{sbintrunc } (\min m n) w$

by (*rule bin-eqI*) (*auto simp: nth-sbintr min.absorb1 min.absorb2*)

lemmas *bintrunc-Pls* =

bintrunc.Suc [**where** *bin=0*, *simplified bin-last-numeral-simps bin-rest-numeral-simps*]

lemmas *bintrunc-Min* [*simp*] =

bintrunc.Suc [**where** *bin=-1*, *simplified bin-last-numeral-simps bin-rest-numeral-simps*]

lemmas *bintrunc-BIT* [*simp*] =

bintrunc.Suc [**where** *bin=w BIT b*, *simplified bin-last-BIT bin-rest-BIT*] **for** *w b*

lemmas *bintrunc-Sucs* = *bintrunc-Pls bintrunc-Min bintrunc-BIT*

bintrunc-Suc-numeral

lemmas *sbintrunc-Suc-Pls* =

sbintrunc.Suc [**where** *bin=0*, *simplified bin-last-numeral-simps bin-rest-numeral-simps*]

lemmas *sbintrunc-Suc-Min* =

sbintrunc.Suc [**where** *bin=-1*, *simplified bin-last-numeral-simps bin-rest-numeral-simps*]

lemmas *sbintrunc-Suc-BIT* [*simp*] =

sbintrunc.Suc [**where** *bin=w BIT b*, *simplified bin-last-BIT bin-rest-BIT*] **for** *w b*

lemmas *sbintrunc-Sucs* = *sbintrunc-Suc-Pls sbintrunc-Suc-Min sbintrunc-Suc-BIT*

sbintrunc-Suc-numeral

lemmas *sbintrunc-Pls* =

sbintrunc.Z [**where** *bin=0*, *simplified bin-last-numeral-simps bin-rest-numeral-simps*]

lemmas *sbintrunc-Min* =
sbintrunc.Z [**where** *bin* = -1, *simplified bin-last-numeral-simps bin-rest-numeral-simps*]

lemmas *sbintrunc-0-BIT-B0* [*simp*] =
sbintrunc.Z [**where** *bin* = *w BIT False*, *simplified bin-last-numeral-simps bin-rest-numeral-simps*]
for *w*

lemmas *sbintrunc-0-BIT-B1* [*simp*] =
sbintrunc.Z [**where** *bin* = *w BIT True*, *simplified bin-last-BIT bin-rest-numeral-simps*]
for *w*

lemmas *sbintrunc-0-simps* =
sbintrunc-Pls sbintrunc-Min sbintrunc-0-BIT-B0 sbintrunc-0-BIT-B1

lemmas *bintrunc-simps* = *bintrunc.Z bintrunc-Sucs*
lemmas *sbintrunc-simps* = *sbintrunc-0-simps sbintrunc-Sucs*

lemma *bintrunc-minus*: $0 < n \implies \text{bintrunc } (\text{Suc } (n - 1)) \ w = \text{bintrunc } n \ w$
by *auto*

lemma *sbintrunc-minus*: $0 < n \implies \text{sbintrunc } (\text{Suc } (n - 1)) \ w = \text{sbintrunc } n \ w$
by *auto*

lemmas *bintrunc-minus-simps* =
bintrunc-Sucs [*THEN* [2] *bintrunc-minus* [*symmetric*, *THEN trans*]]
lemmas *sbintrunc-minus-simps* =
sbintrunc-Sucs [*THEN* [2] *sbintrunc-minus* [*symmetric*, *THEN trans*]]

lemmas *thobini1* = *arg-cong* [**where** *f* = $\lambda w. w \text{ BIT } b$] **for** *b*

lemmas *bintrunc-BIT-I* = *trans* [*OF bintrunc-BIT thobini1*]
lemmas *bintrunc-Min-I* = *trans* [*OF bintrunc-Min thobini1*]

lemmas *bmsts* = *bintrunc-minus-simps*(1-3) [*THEN thobini1* [*THEN* [2] *trans*]]
lemmas *bintrunc-Pls-minus-I* = *bmsts*(1)
lemmas *bintrunc-Min-minus-I* = *bmsts*(2)
lemmas *bintrunc-BIT-minus-I* = *bmsts*(3)

lemma *bintrunc-Suc-lem*: $\text{bintrunc } (\text{Suc } n) \ x = y \implies m = \text{Suc } n \implies \text{bintrunc } m \ x = y$
by *auto*

lemmas *bintrunc-Suc-Ialts* =
bintrunc-Min-I [*THEN bintrunc-Suc-lem*]
bintrunc-BIT-I [*THEN bintrunc-Suc-lem*]

lemmas *sbintrunc-BIT-I* = *trans* [*OF sbintrunc-Suc-BIT thobini1*]

lemmas *sbintrunc-Suc-Is* =

sbintrunc-Sucs(1-3) [THEN thobini1 [THEN [2] trans]]

lemmas *sbintrunc-Suc-minus-Is* =
sbintrunc-minus-simps(1-3) [THEN thobini1 [THEN [2] trans]]

lemma *sbintrunc-Suc-lem*: *sbintrunc* (Suc *n*) *x* = *y* $\implies$ *m* = Suc *n* $\implies$ *sbintrunc* *m* *x* = *y*
by *auto*

lemmas *sbintrunc-Suc-Ialts* =
sbintrunc-Suc-Is [THEN *sbintrunc-Suc-lem*]

lemma *sbintrunc-bintrunc-lt*: *m* > *n* $\implies$ *sbintrunc* *n* (*bintrunc* *m* *w*) = *sbintrunc* *n* *w*
by (rule *bin-eqI*) (auto simp: *nth-sbintr* *nth-bintr*)

lemma *bintrunc-sbintrunc-le*: *m* $\leq$ Suc *n* $\implies$ *bintrunc* *m* (*sbintrunc* *n* *w*) = *bintrunc* *m* *w*
apply (rule *bin-eqI*)
apply (auto simp: *nth-sbintr* *nth-bintr*)
apply (subgoal-tac *x=n*, *safe*, *arith+*)[1]
apply (subgoal-tac *x=n*, *safe*, *arith+*)[1]
done

lemmas *bintrunc-sbintrunc* [simp] = *order-refl* [THEN *bintrunc-sbintrunc-le*]
lemmas *sbintrunc-bintrunc* [simp] = *lessI* [THEN *sbintrunc-bintrunc-lt*]
lemmas *bintrunc-bintrunc* [simp] = *order-refl* [THEN *bintrunc-bintrunc-l*]
lemmas *sbintrunc-sbintrunc* [simp] = *order-refl* [THEN *sbintrunc-sbintrunc-l*]

lemma *bintrunc-sbintrunc'* [simp]: 0 < *n* $\implies$ *bintrunc* *n* (*sbintrunc* (*n* - 1) *w*) = *bintrunc* *n* *w*
by (cases *n*) (auto simp del: *bintrunc.Suc*)

lemma *sbintrunc-bintrunc'* [simp]: 0 < *n* $\implies$ *sbintrunc* (*n* - 1) (*bintrunc* *n* *w*) = *sbintrunc* (*n* - 1) *w*
by (cases *n*) (auto simp del: *bintrunc.Suc*)

lemma *bin-sbin-eq-iff*: *bintrunc* (Suc *n*) *x* = *bintrunc* (Suc *n*) *y* $\longleftrightarrow$ *sbintrunc* *n* *x* = *sbintrunc* *n* *y*
apply (rule *iffI*)
apply (rule *box-equals* [OF - *sbintrunc-bintrunc* *sbintrunc-bintrunc*])
apply *simp*
apply (rule *box-equals* [OF - *bintrunc-sbintrunc* *bintrunc-sbintrunc*])
apply *simp*
done

lemma *bin-sbin-eq-iff'*:
0 < *n* $\implies$ *bintrunc* *n* *x* = *bintrunc* *n* *y* $\longleftrightarrow$ *sbintrunc* (*n* - 1) *x* = *sbintrunc* (*n* - 1) *y*

by (*cases n*) (*simp-all add: bin-sbin-eq-iff del: bintrunc.Suc*)

lemmas *bintrunc-sbintruncS0* [*simp*] = *bintrunc-sbintrunc'* [*unfolded One-nat-def*]
lemmas *sbintrunc-bintruncS0* [*simp*] = *sbintrunc-bintrunc'* [*unfolded One-nat-def*]

lemmas *bintrunc-bintrunc-l'* = *le-add1* [*THEN bintrunc-bintrunc-l*]
lemmas *sbintrunc-sbintrunc-l'* = *le-add1* [*THEN sbintrunc-sbintrunc-l*]

lemmas *nat-non0-gr* =
trans [*OF iszero-def* [*THEN Not-eq-iff* [*THEN iffD2*]] *refl*]

lemma *bintrunc-numeral*:
bintrunc (*numeral k*) *x* = *bintrunc* (*pred-numeral k*) (*bin-rest x*) *BIT bin-last x*
by (*simp add: numeral-eq-Suc*)

lemma *sbintrunc-numeral*:
sbintrunc (*numeral k*) *x* = *sbintrunc* (*pred-numeral k*) (*bin-rest x*) *BIT bin-last x*
by (*simp add: numeral-eq-Suc*)

lemma *bintrunc-numeral-simps* [*simp*]:
bintrunc (*numeral k*) (*numeral (Num.Bit0 w)*) = *bintrunc* (*pred-numeral k*)
(*numeral w*) *BIT False*
bintrunc (*numeral k*) (*numeral (Num.Bit1 w)*) = *bintrunc* (*pred-numeral k*)
(*numeral w*) *BIT True*
bintrunc (*numeral k*) (*− numeral (Num.Bit0 w)*) = *bintrunc* (*pred-numeral k*)
(*− numeral w*) *BIT False*
bintrunc (*numeral k*) (*− numeral (Num.Bit1 w)*) =
bintrunc (*pred-numeral k*) (*− numeral (w + Num.One)*) *BIT True*
bintrunc (*numeral k*) *1* = *1*
by (*simp-all add: bintrunc-numeral*)

lemma *sbintrunc-numeral-simps* [*simp*]:
sbintrunc (*numeral k*) (*numeral (Num.Bit0 w)*) = *sbintrunc* (*pred-numeral k*)
(*numeral w*) *BIT False*
sbintrunc (*numeral k*) (*numeral (Num.Bit1 w)*) = *sbintrunc* (*pred-numeral k*)
(*numeral w*) *BIT True*
sbintrunc (*numeral k*) (*− numeral (Num.Bit0 w)*) =
sbintrunc (*pred-numeral k*) (*− numeral w*) *BIT False*
sbintrunc (*numeral k*) (*− numeral (Num.Bit1 w)*) =
sbintrunc (*pred-numeral k*) (*− numeral (w + Num.One)*) *BIT True*
sbintrunc (*numeral k*) *1* = *1*
by (*simp-all add: sbintrunc-numeral*)

lemma *no-bintr-alt1*: *bintrunc n* = ($\lambda w. w \bmod 2^n :: int$)
by (*rule ext*) (*rule bintrunc-mod2p*)

```

lemma range-bintrunc: range (bintrunc n) = {i. 0 ≤ i ∧ i < 2 ^ n}
  apply (unfold no-bintr-alt1)
  apply (auto simp add: image-iff)
  apply (rule exI)
  apply (auto intro: int-mod-lem [THEN iffD1, symmetric])
  done

lemma no-sbintr-alt2: sbintrunc n = (λw. (w + 2 ^ n) mod 2 ^ Suc n - 2 ^ n ::
  int)
  by (rule ext) (simp add : sbintrunc-mod2p)

lemma range-sbintrunc: range (sbintrunc n) = {i. -(2 ^ n) ≤ i ∧ i < 2 ^ n}
  apply (unfold no-sbintr-alt2)
  apply (auto simp add: image-iff eq-diff-eq)
  apply (rule exI)
  apply (auto intro: int-mod-lem [THEN iffD1, symmetric])
  done

lemma sb-inc-lem: a + 2 ^ k < 0 ⟹ a + 2 ^ k + 2 ^ (Suc k) ≤ (a + 2 ^ k) mod
  2 ^ (Suc k)
  for a :: int
  apply (erule int-mod-ge' [where n = 2 ^ (Suc k) and b = a + 2 ^ k, simplified
  zless2p])
  apply (rule TrueI)
  done

lemma sb-inc-lem': a < -(2 ^ k) ⟹ a + 2 ^ k + 2 ^ (Suc k) ≤ (a + 2 ^ k) mod
  2 ^ (Suc k)
  for a :: int
  by (rule sb-inc-lem) simp

lemma sbintrunc-inc: x < -(2 ^ n) ⟹ x + 2 ^ (Suc n) ≤ sbintrunc n x
  unfolding no-sbintr-alt2 by (drule sb-inc-lem') simp

lemma sb-dec-lem: 0 ≤ -(2 ^ k) + a ⟹ (a + 2 ^ k) mod (2 * 2 ^ k) ≤ -(2
  ^ k) + a
  for a :: int
  using int-mod-le' [where n = 2 ^ (Suc k) and b = a + 2 ^ k] by simp

lemma sb-dec-lem': 2 ^ k ≤ a ⟹ (a + 2 ^ k) mod (2 * 2 ^ k) ≤ -(2 ^ k) + a
  for a :: int
  by (rule sb-dec-lem) simp

lemma sbintrunc-dec: x ≥ (2 ^ n) ⟹ x - 2 ^ (Suc n) ≥ sbintrunc n x
  unfolding no-sbintr-alt2 by (drule sb-dec-lem') simp

lemmas m2pths = pos-mod-sign pos-mod-bound [OF zless2p]

lemma bintr-ge0: 0 ≤ bintrunc n w

```

```

by (simp add: bintrunc-mod2p)

lemma bintr-lt2p: bintrunc n w < 2 ^ n
by (simp add: bintrunc-mod2p)

lemma bintr-Min: bintrunc n (- 1) = 2 ^ n - 1
by (simp add: bintrunc-mod2p m1mod2k)

lemma sbintr-ge: - (2 ^ n) ≤ sbintrunc n w
by (simp add: sbintrunc-mod2p)

lemma sbintr-lt: sbintrunc n w < 2 ^ n
by (simp add: sbintrunc-mod2p)

lemma sign-Pls-ge-0: bin-sign bin = 0 ⟷ bin ≥ 0
for bin :: int
by (simp add: bin-sign-def)

lemma sign-Min-lt-0: bin-sign bin = -1 ⟷ bin < 0
for bin :: int
by (simp add: bin-sign-def)

lemma bin-rest-trunc: bin-rest (bintrunc n bin) = bintrunc (n - 1) (bin-rest bin)
by (induct n arbitrary: bin) auto

lemma bin-rest-power-trunc:
  (bin-rest ^ k) (bintrunc n bin) = bintrunc (n - k) ((bin-rest ^ k) bin)
by (induct k) (auto simp: bin-rest-trunc)

lemma bin-rest-trunc-i: bintrunc n (bin-rest bin) = bin-rest (bintrunc (Suc n) bin)
by auto

lemma bin-rest-strunc: bin-rest (sbintrunc (Suc n) bin) = sbintrunc n (bin-rest bin)
by (induct n arbitrary: bin) auto

lemma bintrunc-rest [simp]: bintrunc n (bin-rest (bintrunc n bin)) = bin-rest (bintrunc n bin)
apply (induct n arbitrary: bin)
apply simp
apply (case-tac bin rule: bin-exhaust)
apply (auto simp: bintrunc-bintrunc-l)
done

lemma sbintrunc-rest [simp]: sbintrunc n (bin-rest (sbintrunc n bin)) = bin-rest (sbintrunc n bin)
apply (induct n arbitrary: bin)
apply simp

```

```

apply (case-tac bin rule: bin-exhaust)
apply (auto simp: bintrunc-bintrunc-l split: bool.splits)
done

lemma bintrunc-rest': bintrunc n  $\circ$  bin-rest  $\circ$  bintrunc n = bin-rest  $\circ$  bintrunc n
by (rule ext) auto

lemma sbintrunc-rest': sbintrunc n  $\circ$  bin-rest  $\circ$  sbintrunc n = bin-rest  $\circ$  sbintrunc
n
by (rule ext) auto

lemma rco-lem:  $f \circ g \circ f = g \circ f \implies f \circ (g \circ f) \wedge n = g \wedge n \circ f$ 
apply (rule ext)
apply (induct-tac n)
apply (simp-all (no-asm))
apply (drule fun-cong)
apply (unfold o-def)
apply (erule trans)
apply simp
done

lemmas rco-bintr = bintrunc-rest'
[THEN rco-lem [THEN fun-cong], unfolded o-def]
lemmas rco-sbintr = sbintrunc-rest'
[THEN rco-lem [THEN fun-cong], unfolded o-def]

```

4.4 Splitting and concatenation

```

primrec bin-split :: nat  $\Rightarrow$  int  $\Rightarrow$  int  $\times$  int
where
  Z: bin-split 0 w = (w, 0)
| Suc: bin-split (Suc n) w =
  (let (w1, w2) = bin-split n (bin-rest w)
   in (w1, w2 BIT bin-last w))

lemma [code]:
  bin-split (Suc n) w = (let (w1, w2) = bin-split n (bin-rest w) in (w1, w2 BIT
bin-last w))
  bin-split 0 w = (w, 0)
by simp-all

primrec bin-cat :: int  $\Rightarrow$  nat  $\Rightarrow$  int  $\Rightarrow$  int
where
  Z: bin-cat w 0 v = w
| Suc: bin-cat w (Suc n) v = bin-cat w n (bin-rest v) BIT bin-last v

end

```

5 Bitwise Operations on Binary Integers

```
theory Bits-Int
  imports Bits Bit-Representation
begin
```

5.1 Logical operations

bit-wise logical operations on the int type

```
instantiation int :: bit
begin
```

```
definition int-not-def: bitNOT = ( $\lambda x::int. - x - 1$ )
```

```
function bitAND-int
  where bitAND-int x y =
    (if x = 0 then 0 else if x = -1 then y
     else (bin-rest x AND bin-rest y) BIT (bin-last x  $\wedge$  bin-last y))
  by pat-completeness simp
```

```
termination
  by (relation measure (nat  $\circ$  abs  $\circ$  fst), simp-all add: bin-rest-def)
```

```
declare bitAND-int.simps [simp del]
```

```
definition int-or-def:
  bitOR = ( $\lambda x y::int. NOT (NOT x AND NOT y)$ )
```

```
definition int-xor-def:
  bitXOR = ( $\lambda x y::int. (x AND NOT y) OR (NOT x AND y)$ )
```

```
instance ..
```

```
end
```

5.1.1 Basic simplification rules

```
lemma int-not-BIT [simp]:
  NOT (w BIT b) = (NOT w) BIT ( $\neg$  b)
  unfolding int-not-def Bit-def by (cases b, simp-all)
```

```
lemma int-not-simps [simp]:
  NOT (0::int) = -1
  NOT (1::int) = -2
  NOT (- 1::int) = 0
  NOT (numeral w::int) = - numeral (w + Num.One)
  NOT (- numeral (Num.Bit0 w)::int) = numeral (Num.BitM w)
  NOT (- numeral (Num.Bit1 w)::int) = numeral (Num.Bit0 w)
  unfolding int-not-def by simp-all
```

lemma *int-not-not* [simp]: $\text{NOT } (\text{NOT } (x::\text{int})) = x$
unfolding *int-not-def* **by** *simp*

lemma *int-and-0* [simp]: $(0::\text{int}) \text{ AND } x = 0$
by (*simp add: bitAND-int.simps*)

lemma *int-and-m1* [simp]: $(-1::\text{int}) \text{ AND } x = x$
by (*simp add: bitAND-int.simps*)

lemma *int-and-Bits* [simp]:
 $(x \text{ BIT } b) \text{ AND } (y \text{ BIT } c) = (x \text{ AND } y) \text{ BIT } (b \wedge c)$
by (*subst bitAND-int.simps, simp add: Bit-eq-0-iff Bit-eq-m1-iff*)

lemma *int-or-zero* [simp]: $(0::\text{int}) \text{ OR } x = x$
unfolding *int-or-def* **by** *simp*

lemma *int-or-minus1* [simp]: $(-1::\text{int}) \text{ OR } x = -1$
unfolding *int-or-def* **by** *simp*

lemma *int-or-Bits* [simp]:
 $(x \text{ BIT } b) \text{ OR } (y \text{ BIT } c) = (x \text{ OR } y) \text{ BIT } (b \vee c)$
unfolding *int-or-def* **by** *simp*

lemma *int-xor-zero* [simp]: $(0::\text{int}) \text{ XOR } x = x$
unfolding *int-xor-def* **by** *simp*

lemma *int-xor-Bits* [simp]:
 $(x \text{ BIT } b) \text{ XOR } (y \text{ BIT } c) = (x \text{ XOR } y) \text{ BIT } ((b \vee c) \wedge \neg (b \wedge c))$
unfolding *int-xor-def* **by** *auto*

5.1.2 Binary destructors

lemma *bin-rest-NOT* [simp]: $\text{bin-rest } (\text{NOT } x) = \text{NOT } (\text{bin-rest } x)$
by (*cases x rule: bin-exhaust, simp*)

lemma *bin-last-NOT* [simp]: $\text{bin-last } (\text{NOT } x) \longleftrightarrow \neg \text{bin-last } x$
by (*cases x rule: bin-exhaust, simp*)

lemma *bin-rest-AND* [simp]: $\text{bin-rest } (x \text{ AND } y) = \text{bin-rest } x \text{ AND } \text{bin-rest } y$
by (*cases x rule: bin-exhaust, cases y rule: bin-exhaust, simp*)

lemma *bin-last-AND* [simp]: $\text{bin-last } (x \text{ AND } y) \longleftrightarrow \text{bin-last } x \wedge \text{bin-last } y$
by (*cases x rule: bin-exhaust, cases y rule: bin-exhaust, simp*)

lemma *bin-rest-OR* [simp]: $\text{bin-rest } (x \text{ OR } y) = \text{bin-rest } x \text{ OR } \text{bin-rest } y$
by (*cases x rule: bin-exhaust, cases y rule: bin-exhaust, simp*)

lemma *bin-last-OR* [simp]: $\text{bin-last } (x \text{ OR } y) \longleftrightarrow \text{bin-last } x \vee \text{bin-last } y$

by (cases x rule: *bin-exhaust*, cases y rule: *bin-exhaust*, *simp*)

lemma *bin-rest-XOR* [*simp*]: $\text{bin-rest } (x \text{ XOR } y) = \text{bin-rest } x \text{ XOR } \text{bin-rest } y$
by (cases x rule: *bin-exhaust*, cases y rule: *bin-exhaust*, *simp*)

lemma *bin-last-XOR* [*simp*]: $\text{bin-last } (x \text{ XOR } y) \longleftrightarrow (\text{bin-last } x \vee \text{bin-last } y) \wedge \neg (\text{bin-last } x \wedge \text{bin-last } y)$
by (cases x rule: *bin-exhaust*, cases y rule: *bin-exhaust*, *simp*)

lemma *bin-nth-ops*:
 $\forall x y. \text{bin-nth } (x \text{ AND } y) n = (\text{bin-nth } x n \ \& \ \text{bin-nth } y n)$
 $\forall x y. \text{bin-nth } (x \text{ OR } y) n = (\text{bin-nth } x n \ | \ \text{bin-nth } y n)$
 $\forall x y. \text{bin-nth } (x \text{ XOR } y) n = (\text{bin-nth } x n \ \sim \ \text{bin-nth } y n)$
 $\forall x. \text{bin-nth } (\text{NOT } x) n = (\sim \ \text{bin-nth } x n)$
by (*induct n*) *auto*

5.1.3 Derived properties

lemma *int-xor-minus1* [*simp*]: $(-1::\text{int}) \text{ XOR } x = \text{NOT } x$
by (*auto simp add: bin-eq-iff bin-nth-ops*)

lemma *int-xor-extra-simps* [*simp*]:
 $w \text{ XOR } (0::\text{int}) = w$
 $w \text{ XOR } (-1::\text{int}) = \text{NOT } w$
by (*auto simp add: bin-eq-iff bin-nth-ops*)

lemma *int-or-extra-simps* [*simp*]:
 $w \text{ OR } (0::\text{int}) = w$
 $w \text{ OR } (-1::\text{int}) = -1$
by (*auto simp add: bin-eq-iff bin-nth-ops*)

lemma *int-and-extra-simps* [*simp*]:
 $w \text{ AND } (0::\text{int}) = 0$
 $w \text{ AND } (-1::\text{int}) = w$
by (*auto simp add: bin-eq-iff bin-nth-ops*)

lemma *bin-ops-comm*:
shows
int-and-comm: $\forall y::\text{int}. x \text{ AND } y = y \text{ AND } x$ **and**
int-or-comm: $\forall y::\text{int}. x \text{ OR } y = y \text{ OR } x$ **and**
int-xor-comm: $\forall y::\text{int}. x \text{ XOR } y = y \text{ XOR } x$
by (*auto simp add: bin-eq-iff bin-nth-ops*)

lemma *bin-ops-same* [*simp*]:
 $(x::\text{int}) \text{ AND } x = x$
 $(x::\text{int}) \text{ OR } x = x$
 $(x::\text{int}) \text{ XOR } x = 0$
by (*auto simp add: bin-eq-iff bin-nth-ops*)

lemmas *bin-log-esimps* =
int-and-extra-simps int-or-extra-simps int-xor-extra-simps
int-and-0 int-and-m1 int-or-zero int-or-minus1 int-xor-zero int-xor-minus1

lemma *bbw-ao-absorb*:
 $!!y::int. x \text{ AND } (y \text{ OR } x) = x \ \& \ x \text{ OR } (y \text{ AND } x) = x$
by (*auto simp add: bin-eq-iff bin-nth-ops*)

lemma *bbw-ao-absorbs-other*:
 $x \text{ AND } (x \text{ OR } y) = x \wedge (y \text{ AND } x) \text{ OR } x = (x::int)$
 $(y \text{ OR } x) \text{ AND } x = x \wedge x \text{ OR } (x \text{ AND } y) = (x::int)$
 $(x \text{ OR } y) \text{ AND } x = x \wedge (x \text{ AND } y) \text{ OR } x = (x::int)$
by (*auto simp add: bin-eq-iff bin-nth-ops*)

lemmas *bbw-ao-absorbs* [*simp*] = *bbw-ao-absorb bbw-ao-absorbs-other*

lemma *int-xor-not*:
 $!!y::int. (\text{NOT } x) \text{ XOR } y = \text{NOT } (x \text{ XOR } y) \ \& \$
 $x \text{ XOR } (\text{NOT } y) = \text{NOT } (x \text{ XOR } y)$
by (*auto simp add: bin-eq-iff bin-nth-ops*)

lemma *int-and-assoc*:
 $(x \text{ AND } y) \text{ AND } (z::int) = x \text{ AND } (y \text{ AND } z)$
by (*auto simp add: bin-eq-iff bin-nth-ops*)

lemma *int-or-assoc*:
 $(x \text{ OR } y) \text{ OR } (z::int) = x \text{ OR } (y \text{ OR } z)$
by (*auto simp add: bin-eq-iff bin-nth-ops*)

lemma *int-xor-assoc*:
 $(x \text{ XOR } y) \text{ XOR } (z::int) = x \text{ XOR } (y \text{ XOR } z)$
by (*auto simp add: bin-eq-iff bin-nth-ops*)

lemmas *bbw-assocs* = *int-and-assoc int-or-assoc int-xor-assoc*

lemma *bbw-lcs* [*simp*]:
 $(y::int) \text{ AND } (x \text{ AND } z) = x \text{ AND } (y \text{ AND } z)$
 $(y::int) \text{ OR } (x \text{ OR } z) = x \text{ OR } (y \text{ OR } z)$
 $(y::int) \text{ XOR } (x \text{ XOR } z) = x \text{ XOR } (y \text{ XOR } z)$
by (*auto simp add: bin-eq-iff bin-nth-ops*)

lemma *bbw-not-dist*:
 $!!y::int. \text{NOT } (x \text{ OR } y) = (\text{NOT } x) \text{ AND } (\text{NOT } y)$
 $!!y::int. \text{NOT } (x \text{ AND } y) = (\text{NOT } x) \text{ OR } (\text{NOT } y)$
by (*auto simp add: bin-eq-iff bin-nth-ops*)

lemma *bbw-oa-dist*:

!!y z::int. (x AND y) OR z =
 (x OR z) AND (y OR z)
by (auto simp add: bin-eq-iff bin-nth-ops)

lemma *bbw-ao-dist*:

!!y z::int. (x OR y) AND z =
 (x AND z) OR (y AND z)
by (auto simp add: bin-eq-iff bin-nth-ops)

5.1.4 Simplification with numerals

Cases for 0 and -1 are already covered by other simp rules.

lemma *bin-rl-eqI*: $\llbracket \text{bin-rest } x = \text{bin-rest } y; \text{bin-last } x = \text{bin-last } y \rrbracket \implies x = y$
by (metis (mono-tags) BIT-eq-iff bin-ex-rl bin-last-BIT bin-rest-BIT)

lemma *bin-rest-neg-numeral-BitM* [simp]:

bin-rest (– numeral (Num.BitM w)) = – numeral w
by (simp only: BIT-bin-simps [symmetric] bin-rest-BIT)

lemma *bin-last-neg-numeral-BitM* [simp]:

bin-last (– numeral (Num.BitM w))
by (simp only: BIT-bin-simps [symmetric] bin-last-BIT)

FIXME: The rule sets below are very large (24 rules for each operator). Is there a simpler way to do this?

lemma *int-and-numerals* [simp]:

numeral (Num.Bit0 x) AND numeral (Num.Bit0 y) = (numeral x AND numeral y) BIT False
 numeral (Num.Bit0 x) AND numeral (Num.Bit1 y) = (numeral x AND numeral y) BIT False
 numeral (Num.Bit1 x) AND numeral (Num.Bit0 y) = (numeral x AND numeral y) BIT False
 numeral (Num.Bit1 x) AND numeral (Num.Bit1 y) = (numeral x AND numeral y) BIT True
 numeral (Num.Bit0 x) AND – numeral (Num.Bit0 y) = (numeral x AND – numeral y) BIT False
 numeral (Num.Bit0 x) AND – numeral (Num.Bit1 y) = (numeral x AND – numeral (y + Num.One)) BIT False
 numeral (Num.Bit1 x) AND – numeral (Num.Bit0 y) = (numeral x AND – numeral y) BIT False
 numeral (Num.Bit1 x) AND – numeral (Num.Bit1 y) = (numeral x AND – numeral (y + Num.One)) BIT True
 – numeral (Num.Bit0 x) AND numeral (Num.Bit0 y) = (– numeral x AND numeral y) BIT False
 – numeral (Num.Bit0 x) AND numeral (Num.Bit1 y) = (– numeral x AND numeral y) BIT False

$\text{numeral (Num.Bit1 } x) \text{ AND numeral (Num.Bit0 } y) = (\text{numeral } (x + \text{Num.One}) \text{ AND numeral } y) \text{ BIT False}$
 $\text{numeral (Num.Bit1 } x) \text{ AND numeral (Num.Bit1 } y) = (\text{numeral } (x + \text{Num.One}) \text{ AND numeral } y) \text{ BIT True}$
 $\text{numeral (Num.Bit0 } x) \text{ AND } \neg \text{numeral (Num.Bit0 } y) = (\neg \text{numeral } x \text{ AND } \neg \text{numeral } y) \text{ BIT False}$
 $\text{numeral (Num.Bit0 } x) \text{ AND } \neg \text{numeral (Num.Bit1 } y) = (\neg \text{numeral } x \text{ AND } \neg \text{numeral } (y + \text{Num.One})) \text{ BIT False}$
 $\text{numeral (Num.Bit1 } x) \text{ AND } \neg \text{numeral (Num.Bit0 } y) = (\neg \text{numeral } (x + \text{Num.One}) \text{ AND } \neg \text{numeral } y) \text{ BIT False}$
 $\text{numeral (Num.Bit1 } x) \text{ AND } \neg \text{numeral (Num.Bit1 } y) = (\neg \text{numeral } (x + \text{Num.One}) \text{ AND } \neg \text{numeral } (y + \text{Num.One})) \text{ BIT True}$
 $(1::\text{int}) \text{ AND numeral (Num.Bit0 } y) = 0$
 $(1::\text{int}) \text{ AND numeral (Num.Bit1 } y) = 1$
 $(1::\text{int}) \text{ AND } \neg \text{numeral (Num.Bit0 } y) = 0$
 $(1::\text{int}) \text{ AND } \neg \text{numeral (Num.Bit1 } y) = 1$
 $\text{numeral (Num.Bit0 } x) \text{ AND } (1::\text{int}) = 0$
 $\text{numeral (Num.Bit1 } x) \text{ AND } (1::\text{int}) = 1$
 $\neg \text{numeral (Num.Bit0 } x) \text{ AND } (1::\text{int}) = 0$
 $\neg \text{numeral (Num.Bit1 } x) \text{ AND } (1::\text{int}) = 1$
by (rule bin-rl-eqI, simp, simp)+

lemma *int-or-numerals* [simp]:

$\text{numeral (Num.Bit0 } x) \text{ OR numeral (Num.Bit0 } y) = (\text{numeral } x \text{ OR numeral } y) \text{ BIT False}$
 $\text{numeral (Num.Bit0 } x) \text{ OR numeral (Num.Bit1 } y) = (\text{numeral } x \text{ OR numeral } y) \text{ BIT True}$
 $\text{numeral (Num.Bit1 } x) \text{ OR numeral (Num.Bit0 } y) = (\text{numeral } x \text{ OR numeral } y) \text{ BIT True}$
 $\text{numeral (Num.Bit1 } x) \text{ OR numeral (Num.Bit1 } y) = (\text{numeral } x \text{ OR numeral } y) \text{ BIT True}$
 $\text{numeral (Num.Bit0 } x) \text{ OR } \neg \text{numeral (Num.Bit0 } y) = (\text{numeral } x \text{ OR } \neg \text{numeral } y) \text{ BIT False}$
 $\text{numeral (Num.Bit0 } x) \text{ OR } \neg \text{numeral (Num.Bit1 } y) = (\text{numeral } x \text{ OR } \neg \text{numeral } (y + \text{Num.One})) \text{ BIT True}$
 $\text{numeral (Num.Bit1 } x) \text{ OR } \neg \text{numeral (Num.Bit0 } y) = (\text{numeral } x \text{ OR } \neg \text{numeral } y) \text{ BIT True}$
 $\text{numeral (Num.Bit1 } x) \text{ OR } \neg \text{numeral (Num.Bit1 } y) = (\text{numeral } x \text{ OR } \neg \text{numeral } (y + \text{Num.One})) \text{ BIT True}$
 $\neg \text{numeral (Num.Bit0 } x) \text{ OR numeral (Num.Bit0 } y) = (\neg \text{numeral } x \text{ OR numeral } y) \text{ BIT False}$
 $\neg \text{numeral (Num.Bit0 } x) \text{ OR numeral (Num.Bit1 } y) = (\neg \text{numeral } x \text{ OR numeral } y) \text{ BIT True}$
 $\neg \text{numeral (Num.Bit1 } x) \text{ OR numeral (Num.Bit0 } y) = (\neg \text{numeral } (x + \text{Num.One}) \text{ OR numeral } y) \text{ BIT True}$
 $\neg \text{numeral (Num.Bit1 } x) \text{ OR numeral (Num.Bit1 } y) = (\neg \text{numeral } (x + \text{Num.One}) \text{ OR numeral } y) \text{ BIT True}$
 $\neg \text{numeral (Num.Bit0 } x) \text{ OR } \neg \text{numeral (Num.Bit0 } y) = (\neg \text{numeral } x \text{ OR } \neg \text{numeral } y) \text{ BIT False}$

– numeral (Num.Bit0 x) OR – numeral (Num.Bit1 y) = (– numeral x OR – numeral (y + Num.One)) BIT True
 – numeral (Num.Bit1 x) OR – numeral (Num.Bit0 y) = (– numeral (x + Num.One) OR – numeral y) BIT True
 – numeral (Num.Bit1 x) OR – numeral (Num.Bit1 y) = (– numeral (x + Num.One) OR – numeral (y + Num.One)) BIT True
 (1::int) OR numeral (Num.Bit0 y) = numeral (Num.Bit1 y)
 (1::int) OR numeral (Num.Bit1 y) = numeral (Num.Bit1 y)
 (1::int) OR – numeral (Num.Bit0 y) = – numeral (Num.BitM y)
 (1::int) OR – numeral (Num.Bit1 y) = – numeral (Num.Bit1 y)
 numeral (Num.Bit0 x) OR (1::int) = numeral (Num.Bit1 x)
 numeral (Num.Bit1 x) OR (1::int) = numeral (Num.Bit1 x)
 – numeral (Num.Bit0 x) OR (1::int) = – numeral (Num.BitM x)
 – numeral (Num.Bit1 x) OR (1::int) = – numeral (Num.Bit1 x)
 by (rule bin-rl-eqI, simp, simp)+

lemma int-xor-numerals [simp]:

numeral (Num.Bit0 x) XOR numeral (Num.Bit0 y) = (numeral x XOR numeral y) BIT False
 numeral (Num.Bit0 x) XOR numeral (Num.Bit1 y) = (numeral x XOR numeral y) BIT True
 numeral (Num.Bit1 x) XOR numeral (Num.Bit0 y) = (numeral x XOR numeral y) BIT True
 numeral (Num.Bit1 x) XOR numeral (Num.Bit1 y) = (numeral x XOR numeral y) BIT False
 numeral (Num.Bit0 x) XOR – numeral (Num.Bit0 y) = (numeral x XOR – numeral y) BIT False
 numeral (Num.Bit0 x) XOR – numeral (Num.Bit1 y) = (numeral x XOR – numeral (y + Num.One)) BIT True
 numeral (Num.Bit1 x) XOR – numeral (Num.Bit0 y) = (numeral x XOR – numeral y) BIT True
 numeral (Num.Bit1 x) XOR – numeral (Num.Bit1 y) = (numeral x XOR – numeral (y + Num.One)) BIT False
 – numeral (Num.Bit0 x) XOR numeral (Num.Bit0 y) = (– numeral x XOR numeral y) BIT False
 – numeral (Num.Bit0 x) XOR numeral (Num.Bit1 y) = (– numeral x XOR numeral y) BIT True
 – numeral (Num.Bit1 x) XOR numeral (Num.Bit0 y) = (– numeral (x + Num.One) XOR numeral y) BIT True
 – numeral (Num.Bit1 x) XOR numeral (Num.Bit1 y) = (– numeral (x + Num.One) XOR numeral y) BIT False
 – numeral (Num.Bit0 x) XOR – numeral (Num.Bit0 y) = (– numeral x XOR – numeral y) BIT False
 – numeral (Num.Bit0 x) XOR – numeral (Num.Bit1 y) = (– numeral x XOR – numeral (y + Num.One)) BIT True
 – numeral (Num.Bit1 x) XOR – numeral (Num.Bit0 y) = (– numeral (x + Num.One) XOR – numeral y) BIT True
 – numeral (Num.Bit1 x) XOR – numeral (Num.Bit1 y) = (– numeral (x + Num.One) XOR – numeral (y + Num.One)) BIT False

```

(1::int) XOR numeral (Num.Bit0 y) = numeral (Num.Bit1 y)
(1::int) XOR numeral (Num.Bit1 y) = numeral (Num.Bit0 y)
(1::int) XOR - numeral (Num.Bit0 y) = - numeral (Num.BitM y)
(1::int) XOR - numeral (Num.Bit1 y) = - numeral (Num.Bit0 (y + Num.One))
numeral (Num.Bit0 x) XOR (1::int) = numeral (Num.Bit1 x)
numeral (Num.Bit1 x) XOR (1::int) = numeral (Num.Bit0 x)
- numeral (Num.Bit0 x) XOR (1::int) = - numeral (Num.BitM x)
- numeral (Num.Bit1 x) XOR (1::int) = - numeral (Num.Bit0 (x + Num.One))
by (rule bin-rl-eqI, simp, simp)+

```

5.1.5 Interactions with arithmetic

```

lemma plus-and-or [rule-format]:
  ALL y::int. (x AND y) + (x OR y) = x + y
  apply (induct x rule: bin-induct)
  apply clarsimp
  apply clarsimp
  apply clarsimp
  apply (case-tac y rule: bin-exhaust)
  apply clarsimp
  apply (unfold Bit-def)
  apply clarsimp
  apply (erule-tac x = x in allE)
  apply simp
done

```

```

lemma le-int-or:
  bin-sign (y::int) = 0 ==> x <= x OR y
  apply (induct y arbitrary: x rule: bin-induct)
  apply clarsimp
  apply clarsimp
  apply (case-tac x rule: bin-exhaust)
  apply (case-tac b)
  apply (case-tac [!] bit)
  apply (auto simp: le-Bits)
done

```

```

lemmas int-and-le =
  xtrans(3) [OF bbw-ao-absorbs (2) [THEN conjunct2, symmetric] le-int-or]

```

```

lemma bin-add-not: x + NOT x = (-1::int)
  apply (induct x rule: bin-induct)
  apply clarsimp
  apply clarsimp
  apply (case-tac bit, auto)
done

```

5.1.6 Truncating results of bit-wise operations

lemma *bin-trunc-ao*:

!! $x\ y.$ $(\text{bintrunc } n\ x)\ \text{AND}\ (\text{bintrunc } n\ y) = \text{bintrunc } n\ (x\ \text{AND}\ y)$
 !! $x\ y.$ $(\text{bintrunc } n\ x)\ \text{OR}\ (\text{bintrunc } n\ y) = \text{bintrunc } n\ (x\ \text{OR}\ y)$
by (*auto simp add: bin-eq-iff bin-nth-ops nth-bintr*)

lemma *bin-trunc-xor*:

!! $x\ y.$ $\text{bintrunc } n\ (\text{bintrunc } n\ x\ \text{XOR}\ \text{bintrunc } n\ y) =$
 $\text{bintrunc } n\ (x\ \text{XOR}\ y)$
by (*auto simp add: bin-eq-iff bin-nth-ops nth-bintr*)

lemma *bin-trunc-not*:

!! $x.$ $\text{bintrunc } n\ (\text{NOT}\ (\text{bintrunc } n\ x)) = \text{bintrunc } n\ (\text{NOT}\ x)$
by (*auto simp add: bin-eq-iff bin-nth-ops nth-bintr*)

lemma *bintr-bintr-i*:

$x = \text{bintrunc } n\ y \implies \text{bintrunc } n\ x = \text{bintrunc } n\ y$
by *auto*

lemmas *bin-trunc-and* = *bin-trunc-ao*(1) [*THEN bintr-bintr-i*]

lemmas *bin-trunc-or* = *bin-trunc-ao*(2) [*THEN bintr-bintr-i*]

5.2 Setting and clearing bits

primrec

bin-sc :: $\text{nat} \Rightarrow \text{bool} \Rightarrow \text{int} \Rightarrow \text{int}$

where

$Z:$ *bin-sc* 0 $b\ w = \text{bin-rest } w\ \text{BIT } b$

| *Suc*: *bin-sc* (*Suc* n) $b\ w = \text{bin-sc } n\ b\ (\text{bin-rest } w)\ \text{BIT } \text{bin-last } w$

lemma *bin-nth-sc* [*simp*]:

$\text{bin-nth}\ (\text{bin-sc } n\ b\ w)\ n \longleftrightarrow b$
by (*induct n arbitrary: w*) *auto*

lemma *bin-sc-sc-same* [*simp*]:

$\text{bin-sc } n\ c\ (\text{bin-sc } n\ b\ w) = \text{bin-sc } n\ c\ w$
by (*induct n arbitrary: w*) *auto*

lemma *bin-sc-sc-diff*:

$m \sim n \implies$
 $\text{bin-sc } m\ c\ (\text{bin-sc } n\ b\ w) = \text{bin-sc } n\ b\ (\text{bin-sc } m\ c\ w)$
apply (*induct n arbitrary: w m*)
apply (*case-tac* [!] *m*)
apply *auto*
done

lemma *bin-nth-sc-gen*:

$\text{bin-nth}\ (\text{bin-sc } n\ b\ w)\ m = (\text{if } m = n \text{ then } b \text{ else } \text{bin-nth } w\ m)$

by (*induct n arbitrary: w m*) (*case-tac* [!] *m, auto*)

lemma *bin-sc-nth* [*simp*]:
 (*bin-sc n (bin-nth w n) w*) = *w*
by (*induct n arbitrary: w*) *auto*

lemma *bin-sign-sc* [*simp*]:
bin-sign (bin-sc n b w) = *bin-sign w*
by (*induct n arbitrary: w*) *auto*

lemma *bin-sc-bintr* [*simp*]:
bintrunc m (bin-sc n x (bintrunc m (w))) = *bintrunc m (bin-sc n x w)*
apply (*induct n arbitrary: w m*)
apply (*case-tac* [!] *w rule: bin-exhaust*)
apply (*case-tac* [!] *m, auto*)
done

lemma *bin-clr-le*:
bin-sc n False w <= *w*
apply (*induct n arbitrary: w*)
apply (*case-tac* [!] *w rule: bin-exhaust*)
apply (*auto simp: le-Bits*)
done

lemma *bin-set-ge*:
bin-sc n True w >= *w*
apply (*induct n arbitrary: w*)
apply (*case-tac* [!] *w rule: bin-exhaust*)
apply (*auto simp: le-Bits*)
done

lemma *bintr-bin-clr-le*:
bintrunc n (bin-sc m False w) <= *bintrunc n w*
apply (*induct n arbitrary: w m*)
apply *simp*
apply (*case-tac w rule: bin-exhaust*)
apply (*case-tac m*)
apply (*auto simp: le-Bits*)
done

lemma *bintr-bin-set-ge*:
bintrunc n (bin-sc m True w) >= *bintrunc n w*
apply (*induct n arbitrary: w m*)
apply *simp*
apply (*case-tac w rule: bin-exhaust*)
apply (*case-tac m*)
apply (*auto simp: le-Bits*)
done

lemma *bin-sc-FP* [simp]: *bin-sc* *n* *False* 0 = 0
 by (induct *n*) auto

lemma *bin-sc-TM* [simp]: *bin-sc* *n* *True* (− 1) = − 1
 by (induct *n*) auto

lemmas *bin-sc-simps* = *bin-sc.Z bin-sc.Suc bin-sc-TM bin-sc-FP*

lemma *bin-sc-minus*:
 0 < *n* ==> *bin-sc* (*Suc* (*n* − 1)) *b w* = *bin-sc* *n b w*
 by auto

lemmas *bin-sc-Suc-minus* =
trans [*OF bin-sc-minus* [*symmetric*] *bin-sc.Suc*]

lemma *bin-sc-numeral* [simp]:
bin-sc (*numeral* *k*) *b w* =
bin-sc (*pred-numeral* *k*) *b* (*bin-rest* *w*) *BIT bin-last w*
 by (simp add: *numeral-eq-Suc*)

5.3 Splitting and concatenation

definition *bin-rcat* :: *nat* ⇒ *int list* ⇒ *int*
where

bin-rcat *n* = *foldl* (λ *u v*. *bin-cat* *u n v*) 0

fun *bin-rsplit-aux* :: *nat* ⇒ *nat* ⇒ *int* ⇒ *int list* ⇒ *int list*
where

bin-rsplit-aux *n m c bs* =
 (if *m* = 0 | *n* = 0 then *bs* else
 let (*a*, *b*) = *bin-split* *n c*
 in *bin-rsplit-aux* *n* (*m* − *n*) *a* (*b* # *bs*))

definition *bin-rsplit* :: *nat* ⇒ *nat* × *int* ⇒ *int list*
where

bin-rsplit *n w* = *bin-rsplit-aux* *n* (*fst w*) (*snd w*) []

fun *bin-rsplittl-aux* :: *nat* ⇒ *nat* ⇒ *int* ⇒ *int list* ⇒ *int list*
where

bin-rsplittl-aux *n m c bs* =
 (if *m* = 0 | *n* = 0 then *bs* else
 let (*a*, *b*) = *bin-split* (*min* *m n*) *c*
 in *bin-rsplittl-aux* *n* (*m* − *n*) *a* (*b* # *bs*))

definition *bin-rsplittl* :: *nat* ⇒ *nat* × *int* ⇒ *int list*
where

bin-rsplittl *n w* = *bin-rsplittl-aux* *n* (*fst w*) (*snd w*) []

declare *bin-rsplit-aux.simps* [simp del]

declare *bin-rsplitl-aux.simps* [*simp del*]

lemma *bin-sign-cat*:

bin-sign (bin-cat x n y) = bin-sign x
by (*induct n arbitrary: y*) *auto*

lemma *bin-cat-Suc-Bit*:

bin-cat w (Suc n) (v BIT b) = bin-cat w n v BIT b
by *auto*

lemma *bin-nth-cat*:

bin-nth (bin-cat x k y) n =
(if n < k then bin-nth y n else bin-nth x (n - k))
apply (*induct k arbitrary: n y*)
apply *clarsimp*
apply (*case-tac n, auto*)
done

lemma *bin-nth-split*:

bin-split n c = (a, b) ==>
(ALL k. bin-nth a k = bin-nth c (n + k)) &
(ALL k. bin-nth b k = (k < n & bin-nth c k))
apply (*induct n arbitrary: b c*)
apply *clarsimp*
apply (*clarsimp simp: Let-def split: prod.split-asm*)
apply (*case-tac k*)
apply *auto*
done

lemma *bin-cat-assoc*:

bin-cat (bin-cat x m y) n z = bin-cat x (m + n) (bin-cat y n z)
by (*induct n arbitrary: z*) *auto*

lemma *bin-cat-assoc-sym*:

bin-cat x m (bin-cat y n z) = bin-cat (bin-cat x (m - n) y) (min m n) z
apply (*induct n arbitrary: z m, clarsimp*)
apply (*case-tac m, auto*)
done

lemma *bin-cat-zero* [*simp*]: *bin-cat 0 n w = bintrunc n w*

by (*induct n arbitrary: w*) *auto*

lemma *bintr-cat1*:

bintrunc (k + n) (bin-cat a n b) = bin-cat (bintrunc k a) n b
by (*induct n arbitrary: b*) *auto*

lemma *bintr-cat*: *bintrunc m (bin-cat a n b) =*

bin-cat (bintrunc (m - n) a) n (bintrunc (min m n) b)
by (*rule bin-eqI*) (*auto simp: bin-nth-cat nth-bintr*)

lemma *bintr-cat-same* [simp]:

$\text{bintrunc } n \ (\text{bin-cat } a \ n \ b) = \text{bintrunc } n \ b$
by (auto simp add : bintr-cat)

lemma *cat-bintr* [simp]:

$\text{bin-cat } a \ n \ (\text{bintrunc } n \ b) = \text{bin-cat } a \ n \ b$
by (induct n arbitrary: b) auto

lemma *split-bintrunc*:

$\text{bin-split } n \ c = (a, b) \implies b = \text{bintrunc } n \ c$
by (induct n arbitrary: b c) (auto simp: Let-def split: prod.split-asm)

lemma *bin-cat-split*:

$\text{bin-split } n \ w = (u, v) \implies w = \text{bin-cat } u \ n \ v$
by (induct n arbitrary: v w) (auto simp: Let-def split: prod.split-asm)

lemma *bin-split-cat*:

$\text{bin-split } n \ (\text{bin-cat } v \ n \ w) = (v, \text{bintrunc } n \ w)$
by (induct n arbitrary: w) auto

lemma *bin-split-zero* [simp]: $\text{bin-split } n \ 0 = (0, 0)$

by (induct n) auto

lemma *bin-split-minus1* [simp]:

$\text{bin-split } n \ (-1) = (-1, \text{bintrunc } n \ (-1))$
by (induct n) auto

lemma *bin-split-trunc*:

$\text{bin-split } (\text{min } m \ n) \ c = (a, b) \implies$
 $\text{bin-split } n \ (\text{bintrunc } m \ c) = (\text{bintrunc } (m - n) \ a, b)$
apply (induct n arbitrary: m b c, clarsimp)
apply (simp add: bin-rest-trunc Let-def split: prod.split-asm)
apply (case-tac m)
apply (auto simp: Let-def split: prod.split-asm)
done

lemma *bin-split-trunc1*:

$\text{bin-split } n \ c = (a, b) \implies$
 $\text{bin-split } n \ (\text{bintrunc } m \ c) = (\text{bintrunc } (m - n) \ a, \text{bintrunc } m \ b)$
apply (induct n arbitrary: m b c, clarsimp)
apply (simp add: bin-rest-trunc Let-def split: prod.split-asm)
apply (case-tac m)
apply (auto simp: Let-def split: prod.split-asm)
done

lemma *bin-cat-num*:

$\text{bin-cat } a \ n \ b = a * 2^n + \text{bintrunc } n \ b$
apply (induct n arbitrary: b, clarsimp)

```

apply (simp add: Bit-def)
done

```

```

lemma bin-split-num:
  bin-split n b = (b div 2 ^ n, b mod 2 ^ n)
apply (induct n arbitrary: b, simp)
apply (simp add: bin-rest-def zdiv-zmult2-eq)
apply (case-tac b rule: bin-exhaust)
apply simp
apply (simp add: Bit-def mod-mult-mult1 p1mod22k)
done

```

5.4 Miscellaneous lemmas

```

lemma nth-2p-bin:
  bin-nth (2 ^ n) m = (m = n)
apply (induct n arbitrary: m)
apply clarsimp
apply safe
apply (case-tac m)
apply (auto simp: Bit-B0-2t [symmetric])
done

```

```

lemma ex-eq-or:
  (EX m. n = Suc m & (m = k | P m)) = (n = Suc k | (EX m. n = Suc m & P m))
by auto

```

```

lemma power-BIT: 2 ^ (Suc n) - 1 = (2 ^ n - 1) BIT True
unfolding Bit-B1
by (induct n) simp-all

```

```

lemma mod-BIT:
  bin BIT bit mod 2 ^ Suc n = (bin mod 2 ^ n) BIT bit
proof -
  have 2 * (bin mod 2 ^ n) + 1 = (2 * bin mod 2 ^ Suc n) + 1
    by (simp add: mod-mult-mult1)
  also have ... = ((2 * bin mod 2 ^ Suc n) + 1) mod 2 ^ Suc n
    by (simp add: ac-simps p1mod22k')
  also have ... = (2 * bin + 1) mod 2 ^ Suc n
    by (simp only: mod-simps)
  finally show ?thesis
    by (auto simp add: Bit-def)
qed

```

```

lemma AND-mod:
  fixes x :: int

```

```

  shows  $x \text{ AND } 2^n - 1 = x \text{ mod } 2^n$ 
proof (induct x arbitrary: n rule: bin-induct)
  case 1
  then show ?case
    by simp
next
  case 2
  then show ?case
    by (simp, simp add: m1mod2k)
next
  case (3 bin bit)
  show ?case
  proof (cases n)
    case 0
    then show ?thesis by simp
  next
    case (Suc m)
    with 3 show ?thesis
      by (simp only: power-BIT mod-BIT int-and-Bits) simp
  qed
qed
end

```

6 Bool lists and integers

```

theory Bool-List-Representation
  imports Bits-Int
begin

```

```

definition map2 :: ('a  $\Rightarrow$  'b  $\Rightarrow$  'c)  $\Rightarrow$  'a list  $\Rightarrow$  'b list  $\Rightarrow$  'c list
  where map2 f as bs = map (case-prod f) (zip as bs)

```

```

lemma map2-Nil [simp, code]: map2 f [] ys = []
  by (auto simp: map2-def)

```

```

lemma map2-Nil2 [simp, code]: map2 f xs [] = []
  by (auto simp: map2-def)

```

```

lemma map2-Cons [simp, code]: map2 f (x # xs) (y # ys) = f x y # map2 f xs
ys
  by (auto simp: map2-def)

```

6.1 Operations on lists of booleans

```

primrec bl-to-bin-aux :: bool list  $\Rightarrow$  int  $\Rightarrow$  int
  where
    Nil: bl-to-bin-aux [] w = w
    | Cons: bl-to-bin-aux (b # bs) w = bl-to-bin-aux bs (w BIT b)

```

definition $bl\text{-}to\text{-}bin :: bool\ list \Rightarrow int$
where $bl\text{-}to\text{-}bin\ bs = bl\text{-}to\text{-}bin\text{-}aux\ bs\ 0$

primrec $bin\text{-}to\text{-}bl\text{-}aux :: nat \Rightarrow int \Rightarrow bool\ list \Rightarrow bool\ list$
where
 $Z: bin\text{-}to\text{-}bl\text{-}aux\ 0\ w\ bl = bl$
 $| Suc: bin\text{-}to\text{-}bl\text{-}aux\ (Suc\ n)\ w\ bl = bin\text{-}to\text{-}bl\text{-}aux\ n\ (bin\text{-}rest\ w)\ ((bin\text{-}last\ w)\ \# bl)$

definition $bin\text{-}to\text{-}bl :: nat \Rightarrow int \Rightarrow bool\ list$
where $bin\text{-}to\text{-}bl\ n\ w = bin\text{-}to\text{-}bl\text{-}aux\ n\ w\ []$

primrec $bl\text{-}of\text{-}nth :: nat \Rightarrow (nat \Rightarrow bool) \Rightarrow bool\ list$
where
 $Suc: bl\text{-}of\text{-}nth\ (Suc\ n)\ f = f\ n\ \# bl\text{-}of\text{-}nth\ n\ f$
 $| Z: bl\text{-}of\text{-}nth\ 0\ f = []$

primrec $takefill :: 'a \Rightarrow nat \Rightarrow 'a\ list \Rightarrow 'a\ list$
where
 $Z: takefill\ fill\ 0\ xs = []$
 $| Suc: takefill\ fill\ (Suc\ n)\ xs =$
 $(case\ xs\ of$
 $\quad [] \Rightarrow fill\ \# takefill\ fill\ n\ xs$
 $\quad | y\ \# ys \Rightarrow y\ \# takefill\ fill\ n\ ys)$

6.2 Arithmetic in terms of bool lists

Arithmetic operations in terms of the reversed bool list, assuming input list(s) the same length, and don't extend them.

primrec $rbl\text{-}succ :: bool\ list \Rightarrow bool\ list$
where
 $Nil: rbl\text{-}succ\ Nil = Nil$
 $| Cons: rbl\text{-}succ\ (x\ \# xs) = (if\ x\ then\ False\ \# rbl\text{-}succ\ xs\ else\ True\ \# xs)$

primrec $rbl\text{-}pred :: bool\ list \Rightarrow bool\ list$
where
 $Nil: rbl\text{-}pred\ Nil = Nil$
 $| Cons: rbl\text{-}pred\ (x\ \# xs) = (if\ x\ then\ False\ \# xs\ else\ True\ \# rbl\text{-}pred\ xs)$

primrec $rbl\text{-}add :: bool\ list \Rightarrow bool\ list \Rightarrow bool\ list$
where — result is length of first arg, second arg may be longer
 $Nil: rbl\text{-}add\ Nil\ x = Nil$
 $| Cons: rbl\text{-}add\ (y\ \# ys)\ x =$
 $(let\ ws = rbl\text{-}add\ ys\ (tl\ x)$
 $\quad in\ (y\ \neq\ hd\ x)\ \# (if\ hd\ x\ \wedge\ y\ then\ rbl\text{-}succ\ ws\ else\ ws))$

primrec $rbl\text{-}mult :: bool\ list \Rightarrow bool\ list \Rightarrow bool\ list$
where — result is length of first arg, second arg may be longer

Nil: $rbl_mult\ Nil\ x = Nil$
 | *Cons*: $rbl_mult\ (y \# ys)\ x =$
 $(let\ ws = False \# rbl_mult\ ys\ x$
 $in\ if\ y\ then\ rbl_add\ ws\ x\ else\ ws)$

lemma *butlast-power*: $(butlast\ \wedge^{\wedge} n)\ bl = take\ (length\ bl - n)\ bl$
by (*induct* n) (*auto simp: butlast-take*)

lemma *bin-to-bl-aux-zero-minus-simp* [*simp*]:
 $0 < n \implies bin_to_bl_aux\ n\ 0\ bl = bin_to_bl_aux\ (n - 1)\ 0\ (False \# bl)$
by (*cases* n) *auto*

lemma *bin-to-bl-aux-minus1-minus-simp* [*simp*]:
 $0 < n \implies bin_to_bl_aux\ n\ (-1)\ bl = bin_to_bl_aux\ (n - 1)\ (-1)\ (True \# bl)$
by (*cases* n) *auto*

lemma *bin-to-bl-aux-one-minus-simp* [*simp*]:
 $0 < n \implies bin_to_bl_aux\ n\ 1\ bl = bin_to_bl_aux\ (n - 1)\ 0\ (True \# bl)$
by (*cases* n) *auto*

lemma *bin-to-bl-aux-Bit-minus-simp* [*simp*]:
 $0 < n \implies bin_to_bl_aux\ n\ (w\ BIT\ b)\ bl = bin_to_bl_aux\ (n - 1)\ w\ (b \# bl)$
by (*cases* n) *auto*

lemma *bin-to-bl-aux-Bit0-minus-simp* [*simp*]:
 $0 < n \implies$
 $bin_to_bl_aux\ n\ (numeral\ (Num.Bit0\ w))\ bl = bin_to_bl_aux\ (n - 1)\ (numeral$
 $w)\ (False \# bl)$
by (*cases* n) *auto*

lemma *bin-to-bl-aux-Bit1-minus-simp* [*simp*]:
 $0 < n \implies$
 $bin_to_bl_aux\ n\ (numeral\ (Num.Bit1\ w))\ bl = bin_to_bl_aux\ (n - 1)\ (numeral$
 $w)\ (True \# bl)$
by (*cases* n) *auto*

Link between bin and bool list.

lemma *bl-to-bin-aux-append*: $bl_to_bin_aux\ (bs @ cs)\ w = bl_to_bin_aux\ cs\ (bl_to_bin_aux\ bs\ w)$
by (*induct* bs *arbitrary: w*) *auto*

lemma *bin-to-bl-aux-append*: $bin_to_bl_aux\ n\ w\ bs @ cs = bin_to_bl_aux\ n\ w\ (bs @ cs)$
by (*induct* n *arbitrary: w bs*) *auto*

lemma *bl-to-bin-append*: $bl_to_bin\ (bs @ cs) = bl_to_bin_aux\ cs\ (bl_to_bin\ bs)$
unfolding *bl-to-bin-def* **by** (*rule* *bl-to-bin-aux-append*)

lemma *bin-to-bl-aux-alt*: $bin_to_bl_aux\ n\ w\ bs = bin_to_bl\ n\ w @ bs$

```

    by (simp add: bin-to-bl-def bin-to-bl-aux-append)

lemma bin-to-bl-0 [simp]: bin-to-bl 0 bs = []
  by (auto simp: bin-to-bl-def)

lemma size-bin-to-bl-aux: size (bin-to-bl-aux n w bs) = n + length bs
  by (induct n arbitrary: w bs) auto

lemma size-bin-to-bl [simp]: size (bin-to-bl n w) = n
  by (simp add: bin-to-bl-def size-bin-to-bl-aux)

lemma bin-bl-bin': bl-to-bin (bin-to-bl-aux n w bs) = bl-to-bin-aux bs (bintrunc n w)
  by (induct n arbitrary: w bs) (auto simp: bl-to-bin-def)

lemma bin-bl-bin [simp]: bl-to-bin (bin-to-bl n w) = bintrunc n w
  by (auto simp: bin-to-bl-def bin-bl-bin')

lemma bl-bin-bl': bin-to-bl (n + length bs) (bl-to-bin-aux bs w) = bin-to-bl-aux n w bs
  apply (induct bs arbitrary: w n)
  apply auto
  apply (simp-all only: add-Suc [symmetric])
  apply (auto simp add: bin-to-bl-def)
  done

lemma bl-bin-bl [simp]: bin-to-bl (length bs) (bl-to-bin bs) = bs
  unfolding bl-to-bin-def
  apply (rule box-equals)
  apply (rule bl-bin-bl')
  prefer 2
  apply (rule bin-to-bl-aux.Z)
  apply simp
  done

lemma bl-to-bin-inj: bl-to-bin bs = bl-to-bin cs  $\implies$  length bs = length cs  $\implies$  bs = cs
  apply (rule-tac box-equals)
  defer
  apply (rule bl-bin-bl)
  apply (rule bl-bin-bl)
  apply simp
  done

lemma bl-to-bin-False [simp]: bl-to-bin (False # bl) = bl-to-bin bl
  by (auto simp: bl-to-bin-def)

lemma bl-to-bin-Nil [simp]: bl-to-bin [] = 0
  by (auto simp: bl-to-bin-def)

```

lemma *bin-to-bl-zero-aux*: $\text{bin-to-bl-aux } n \ 0 \ bl = \text{replicate } n \ \text{False} \ @ \ bl$
by (*induct* n *arbitrary*: bl) (*auto simp*: *replicate-app-Cons-same*)

lemma *bin-to-bl-zero*: $\text{bin-to-bl } n \ 0 = \text{replicate } n \ \text{False}$
by (*simp add*: *bin-to-bl-def bin-to-bl-zero-aux*)

lemma *bin-to-bl-minus1-aux*: $\text{bin-to-bl-aux } n \ (-1) \ bl = \text{replicate } n \ \text{True} \ @ \ bl$
by (*induct* n *arbitrary*: bl) (*auto simp*: *replicate-app-Cons-same*)

lemma *bin-to-bl-minus1*: $\text{bin-to-bl } n \ (-1) = \text{replicate } n \ \text{True}$
by (*simp add*: *bin-to-bl-def bin-to-bl-minus1-aux*)

lemma *bl-to-bin-rep-F*: $\text{bl-to-bin } (\text{replicate } n \ \text{False} \ @ \ bl) = \text{bl-to-bin } bl$
by (*simp add*: *bin-to-bl-zero-aux [symmetric] bin-bl-bin'*) (*simp add*: *bl-to-bin-def*)

lemma *bin-to-bl-trunc* [*simp*]: $n \leq m \implies \text{bin-to-bl } n \ (\text{bintrunc } m \ w) = \text{bin-to-bl } n \ w$
by (*auto intro*: *bl-to-bin-inj*)

lemma *bin-to-bl-aux-bintr*:
 $\text{bin-to-bl-aux } n \ (\text{bintrunc } m \ bin) \ bl =$
 $\text{replicate } (n - m) \ \text{False} \ @ \ \text{bin-to-bl-aux } (\min n \ m) \ bin \ bl$
apply (*induct* n *arbitrary*: $m \ bin \ bl$)
apply *clarsimp*
apply *clarsimp*
apply (*case-tac* m)
apply (*clarsimp simp*: *bin-to-bl-zero-aux*)
apply (*erule thin-rl*)
apply (*induct-tac* n)
apply *auto*
done

lemma *bin-to-bl-bintr*:
 $\text{bin-to-bl } n \ (\text{bintrunc } m \ bin) = \text{replicate } (n - m) \ \text{False} \ @ \ \text{bin-to-bl } (\min n \ m) \ bin$
unfolding *bin-to-bl-def* **by** (*rule bin-to-bl-aux-bintr*)

lemma *bl-to-bin-rep-False*: $\text{bl-to-bin } (\text{replicate } n \ \text{False}) = 0$
by (*induct* n) *auto*

lemma *len-bin-to-bl-aux*: $\text{length } (\text{bin-to-bl-aux } n \ w \ bs) = n + \text{length } bs$
by (*fact size-bin-to-bl-aux*)

lemma *len-bin-to-bl*: $\text{length } (\text{bin-to-bl } n \ w) = n$
by (*fact size-bin-to-bl*)

lemma *sign-bl-bin'*: $\text{bin-sign } (\text{bl-to-bin-aux } bs \ w) = \text{bin-sign } w$
by (*induct* bs *arbitrary*: w) *auto*

lemma *sign-bl-bin*: $\text{bin-sign } (\text{bl-to-bin } bs) = 0$
by (*simp add: bl-to-bin-def sign-bl-bin*)

lemma *bl-sbin-sign-aux*: $\text{hd } (\text{bin-to-bl-aux } (\text{Suc } n) w bs) = (\text{bin-sign } (\text{sbintrunc } n w) = -1)$
apply (*induct n arbitrary: w bs*)
apply *clarsimp*
apply (*cases w rule: bin-exhaust*)
apply *simp*
done

lemma *bl-sbin-sign*: $\text{hd } (\text{bin-to-bl } (\text{Suc } n) w) = (\text{bin-sign } (\text{sbintrunc } n w) = -1)$
unfolding *bin-to-bl-def* **by** (*rule bl-sbin-sign-aux*)

lemma *bin-nth-of-bl-aux*:
 $\text{bin-nth } (\text{bl-to-bin-aux } bl w) n =$
 $(n < \text{size } bl \wedge \text{rev } bl ! n \mid n \geq \text{length } bl \wedge \text{bin-nth } w (n - \text{size } bl))$
apply (*induct bl arbitrary: w*)
apply *clarsimp*
apply *clarsimp*
apply (*cut-tac x=n and y=size bl in linorder-less-linear*)
apply (*erule disjE, simp add: nth-append*)
apply *auto*
done

lemma *bin-nth-of-bl*: $\text{bin-nth } (\text{bl-to-bin } bl) n = (n < \text{length } bl \wedge \text{rev } bl ! n)$
by (*simp add: bl-to-bin-def bin-nth-of-bl-aux*)

lemma *bin-nth-bl*: $n < m \implies \text{bin-nth } w n = \text{nth } (\text{rev } (\text{bin-to-bl } m w)) n$
apply (*induct n arbitrary: m w*)
apply *clarsimp*
apply (*case-tac m, clarsimp*)
apply (*clarsimp simp: bin-to-bl-def*)
apply (*simp add: bin-to-bl-aux-alt*)
apply *clarsimp*
apply (*case-tac m, clarsimp*)
apply (*clarsimp simp: bin-to-bl-def*)
apply (*simp add: bin-to-bl-aux-alt*)
done

lemma *nth-rev*: $n < \text{length } xs \implies \text{rev } xs ! n = xs ! (\text{length } xs - 1 - n)$
apply (*induct xs*)
apply *simp*
apply (*clarsimp simp add: nth-append nth.simps split: nat.split*)
apply (*rule-tac f = $\lambda n. xs ! n$ in arg-cong*)
apply *arith*
done

lemma *nth-rev-alt*: $n < \text{length } ys \implies ys ! n = \text{rev } ys ! (\text{length } ys - \text{Suc } n)$
by (*simp add: nth-rev*)

lemma *nth-bin-to-bl-aux*:
 $n < m + \text{length } bl \implies (\text{bin-to-bl-aux } m \ w \ bl) ! n =$
 $(\text{if } n < m \text{ then } \text{bin-nth } w \ (m - 1 - n) \text{ else } bl ! (n - m))$
apply (*induct m arbitrary: w n bl*)
apply *clarsimp*
apply *clarsimp*
apply (*case-tac w rule: bin-exhaust*)
apply *simp*
done

lemma *nth-bin-to-bl*: $n < m \implies (\text{bin-to-bl } m \ w) ! n = \text{bin-nth } w \ (m - \text{Suc } n)$
by (*simp add: bin-to-bl-def nth-bin-to-bl-aux*)

lemma *bl-to-bin-lt2p-aux*: $\text{bl-to-bin-aux } bs \ w < (w + 1) * (2 ^ \text{length } bs)$
apply (*induct bs arbitrary: w*)
apply *clarsimp*
apply *clarsimp*
apply (*drule meta-spec, erule xtrans(8) [rotated], simp add: Bit-def*)+
done

lemma *bl-to-bin-lt2p-drop*: $\text{bl-to-bin } bs < 2 ^ \text{length } (\text{dropWhile Not } bs)$
proof (*induct bs*)
case *Nil*
then show ?*case* **by** *simp*
next
case (*Cons b bs*) **with** *bl-to-bin-lt2p-aux* [**where** $w=1$]
show ?*case* **unfolding** *bl-to-bin-def* **by** *simp*
qed

lemma *bl-to-bin-lt2p*: $\text{bl-to-bin } bs < 2 ^ \text{length } bs$
by (*metis bin-bl-bin bintr-lt2p bl-bin-bl*)

lemma *bl-to-bin-ge2p-aux*: $\text{bl-to-bin-aux } bs \ w \geq w * (2 ^ \text{length } bs)$
apply (*induct bs arbitrary: w*)
apply *clarsimp*
apply *clarsimp*
apply (*drule meta-spec, erule order-trans [rotated],*
 $\text{simp add: Bit-B0-2t Bit-B1-2t algebra-simps}$) +
apply (*simp add: Bit-def*)
done

lemma *bl-to-bin-ge0*: $\text{bl-to-bin } bs \geq 0$
apply (*unfold bl-to-bin-def*)
apply (*rule xtrans(4)*)
apply (*rule bl-to-bin-ge2p-aux*)
apply *simp*

done

lemma *butlast-rest-bin*: $\text{butlast } (\text{bin-to-bl } n \ w) = \text{bin-to-bl } (n - 1) \ (\text{bin-rest } w)$
apply (*unfold bin-to-bl-def*)
apply (*cases w rule: bin-exhaust*)
apply (*cases n, clarsimp*)
apply *clarsimp*
apply (*auto simp add: bin-to-bl-aux-alt*)
done

lemma *butlast-bin-rest*: $\text{butlast } bl = \text{bin-to-bl } (\text{length } bl - \text{Suc } 0) \ (\text{bin-rest } (\text{bl-to-bin } bl))$
using *butlast-rest-bin* [**where** $w = \text{bl-to-bin } bl$ **and** $n = \text{length } bl$] **by** *simp*

lemma *butlast-rest-bl2bin-aux*:
 $bl \neq [] \implies \text{bl-to-bin-aux } (\text{butlast } bl) \ w = \text{bin-rest } (\text{bl-to-bin-aux } bl \ w)$
by (*induct bl arbitrary: w*) *auto*

lemma *butlast-rest-bl2bin*: $\text{bl-to-bin } (\text{butlast } bl) = \text{bin-rest } (\text{bl-to-bin } bl)$
by (*cases bl*) (*auto simp: bl-to-bin-def butlast-rest-bl2bin-aux*)

lemma *trunc-bl2bin-aux*:
 $\text{bintrunc } m \ (\text{bl-to-bin-aux } bl \ w) =$
 $\text{bl-to-bin-aux } (\text{drop } (\text{length } bl - m) \ bl) \ (\text{bintrunc } (m - \text{length } bl) \ w)$

proof (*induct bl arbitrary: w*)

case *Nil*

show *?case* **by** *simp*

next

case (*Cons b bl*)

show *?case*

proof (*cases m - length bl*)

case *0*

then have $\text{Suc } (\text{length } bl) - m = \text{Suc } (\text{length } bl - m)$ **by** *simp*

with *Cons* **show** *?thesis* **by** *simp*

next

case (*Suc n*)

then have $m - \text{Suc } (\text{length } bl) = n$ **by** *simp*

with *Cons Suc* **show** *?thesis* **by** *simp*

qed

qed

lemma *trunc-bl2bin*: $\text{bintrunc } m \ (\text{bl-to-bin } bl) = \text{bl-to-bin } (\text{drop } (\text{length } bl - m) \ bl)$
by (*simp add: bl-to-bin-def trunc-bl2bin-aux*)

lemma *trunc-bl2bin-len* [*simp*]: $\text{bintrunc } (\text{length } bl) \ (\text{bl-to-bin } bl) = \text{bl-to-bin } bl$
by (*simp add: trunc-bl2bin*)

lemma *bl2bin-drop*: $\text{bl-to-bin } (\text{drop } k \ bl) = \text{bintrunc } (\text{length } bl - k) \ (\text{bl-to-bin } bl)$

```

apply (rule trans)
prefer 2
apply (rule trunc-bl2bin [symmetric])
apply (cases  $k \leq \text{length } bl$ )
apply auto
done

lemma nth-rest-power-bin:  $\text{bin-nth } ((\text{bin-rest } \wedge^k) w) n = \text{bin-nth } w (n + k)$ 
apply (induct k arbitrary: n)
apply clarsimp
apply clarsimp
apply (simp only: bin-nth.Suc [symmetric] add-Suc)
done

lemma take-rest-power-bin:  $m \leq n \implies \text{take } m (\text{bin-to-bl } n w) = \text{bin-to-bl } m$ 
 $((\text{bin-rest } \wedge^{(n - m)}) w)$ 
apply (rule nth-equalityI)
apply simp
apply (clarsimp simp add: nth-bin-to-bl nth-rest-power-bin)
done

lemma hd-butlast:  $\text{size } xs > 1 \implies \text{hd } (\text{butlast } xs) = \text{hd } xs$ 
by (cases xs) auto

lemma last-bin-last':  $\text{size } xs > 0 \implies \text{last } xs \longleftrightarrow \text{bin-last } (\text{bl-to-bin-aux } xs w)$ 
by (induct xs arbitrary: w) auto

lemma last-bin-last:  $\text{size } xs > 0 \implies \text{last } xs \longleftrightarrow \text{bin-last } (\text{bl-to-bin } xs)$ 
unfolding bl-to-bin-def by (erule last-bin-last')

lemma bin-last-last:  $\text{bin-last } w \longleftrightarrow \text{last } (\text{bin-to-bl } (\text{Suc } n) w)$ 
by (simp add: bin-to-bl-def) (auto simp: bin-to-bl-aux-alt)

```

6.3 Links between bit-wise operations and operations on bool lists

```

lemma bl-xor-aux-bin:
   $\text{map2 } (\lambda x y. x \neq y) (\text{bin-to-bl-aux } n v bs) (\text{bin-to-bl-aux } n w cs) =$ 
 $\text{bin-to-bl-aux } n (v \text{ XOR } w) (\text{map2 } (\lambda x y. x \neq y) bs cs)$ 
apply (induct n arbitrary: v w bs cs)
apply simp
apply (case-tac v rule: bin-exhaust)
apply (case-tac w rule: bin-exhaust)
apply clarsimp
apply (case-tac b)
apply auto
done

lemma bl-or-aux-bin:

```

```

map2 (op ∨) (bin-to-bl-aux n v bs) (bin-to-bl-aux n w cs) =
  bin-to-bl-aux n (v OR w) (map2 (op ∨) bs cs)
apply (induct n arbitrary: v w bs cs)
apply simp
apply (case-tac v rule: bin-exhaust)
apply (case-tac w rule: bin-exhaust)
apply clarsimp
done

```

lemma *bl-and-aux-bin*:

```

map2 (op ∧) (bin-to-bl-aux n v bs) (bin-to-bl-aux n w cs) =
  bin-to-bl-aux n (v AND w) (map2 (op ∧) bs cs)
apply (induct n arbitrary: v w bs cs)
apply simp
apply (case-tac v rule: bin-exhaust)
apply (case-tac w rule: bin-exhaust)
apply clarsimp
done

```

lemma *bl-not-aux-bin*: $\text{map Not } (\text{bin-to-bl-aux } n \ w \ cs) = \text{bin-to-bl-aux } n \ (NOT \ w)$
 $(\text{map Not } cs)$

by (induct n arbitrary: w cs) auto

lemma *bl-not-bin*: $\text{map Not } (\text{bin-to-bl } n \ w) = \text{bin-to-bl } n \ (NOT \ w)$

by (simp add: bin-to-bl-def bl-not-aux-bin)

lemma *bl-and-bin*: $\text{map2 } (op \ \wedge) \ (\text{bin-to-bl } n \ v) \ (\text{bin-to-bl } n \ w) = \text{bin-to-bl } n \ (v \ AND \ w)$

by (simp add: bin-to-bl-def bl-and-aux-bin)

lemma *bl-or-bin*: $\text{map2 } (op \ \vee) \ (\text{bin-to-bl } n \ v) \ (\text{bin-to-bl } n \ w) = \text{bin-to-bl } n \ (v \ OR \ w)$

by (simp add: bin-to-bl-def bl-or-aux-bin)

lemma *bl-xor-bin*: $\text{map2 } (\lambda x \ y. \ x \neq y) \ (\text{bin-to-bl } n \ v) \ (\text{bin-to-bl } n \ w) = \text{bin-to-bl } n \ (v \ XOR \ w)$

by (simp only: bin-to-bl-def bl-xor-aux-bin map2-Nil)

lemma *drop-bin2bl-aux*:

```

drop m (bin-to-bl-aux n bin bs) =
  bin-to-bl-aux (n - m) bin (drop (m - n) bs)
apply (induct n arbitrary: m bin bs, clarsimp)
apply clarsimp
apply (case-tac bin rule: bin-exhaust)
apply (case-tac m ≤ n, simp)
apply (case-tac m - n, simp)
apply simp
apply (rule-tac f = λnat. drop nat bs in arg-cong)
apply simp

```

done

lemma *drop-bin2bl*: $\text{drop } m \text{ (bin-to-bl } n \text{ bin)} = \text{bin-to-bl } (n - m) \text{ bin}$
 by (*simp add: bin-to-bl-def drop-bin2bl-aux*)

lemma *take-bin2bl-lem1*: $\text{take } m \text{ (bin-to-bl-aux } m \text{ w bs)} = \text{bin-to-bl } m \text{ w}$
 apply (*induct m arbitrary: w bs*)
 apply *clarsimp*
 apply *clarsimp*
 apply (*simp add: bin-to-bl-aux-alt*)
 apply (*simp add: bin-to-bl-def*)
 apply (*simp add: bin-to-bl-aux-alt*)
 done

lemma *take-bin2bl-lem*: $\text{take } m \text{ (bin-to-bl-aux } (m + n) \text{ w bs)} = \text{take } m \text{ (bin-to-bl } (m + n) \text{ w)}$
 by (*induct n arbitrary: w bs*) (*simp-all (no-asm) add: bin-to-bl-def take-bin2bl-lem1, simp*)

lemma *bin-split-take*: $\text{bin-split } n \text{ c} = (a, b) \implies \text{bin-to-bl } m \text{ a} = \text{take } m \text{ (bin-to-bl } (m + n) \text{ c)}$
 apply (*induct n arbitrary: b c*)
 apply *clarsimp*
 apply (*clarsimp simp: Let-def split: prod.split-asm*)
 apply (*simp add: bin-to-bl-def*)
 apply (*simp add: take-bin2bl-lem*)
 done

lemma *bin-split-take1*:
 $k = m + n \implies \text{bin-split } n \text{ c} = (a, b) \implies \text{bin-to-bl } m \text{ a} = \text{take } m \text{ (bin-to-bl } k \text{ c)}$
 by (*auto elim: bin-split-take*)

lemma *nth-takefill*: $m < n \implies \text{takefill fill } n \text{ l ! } m = (\text{if } m < \text{length } l \text{ then } l ! m \text{ else fill})$
 apply (*induct n arbitrary: m l*)
 apply *clarsimp*
 apply *clarsimp*
 apply (*case-tac m*)
 apply (*simp split: list.split*)
 apply (*simp split: list.split*)
 done

lemma *takefill-alt*: $\text{takefill fill } n \text{ l} = \text{take } n \text{ l @ replicate } (n - \text{length } l) \text{ fill}$
 by (*induct n arbitrary: l*) (*auto split: list.split*)

lemma *takefill-replicate* [*simp*]: $\text{takefill fill } n \text{ (replicate } m \text{ fill)} = \text{replicate } n \text{ fill}$
 by (*simp add: takefill-alt replicate-add [symmetric]*)

lemma *takefill-le'*: $n = m + k \implies \text{takefill } x \text{ m (takefill } x \text{ n l)} = \text{takefill } x \text{ m l}$

by (*induct* *m* *arbitrary*: *l* *n*) (*auto* *split*: *list.split*)

lemma *length-takefill* [*simp*]: *length* (*takefill* *fill* *n* *l*) = *n*
by (*simp* *add*: *takefill-alt*)

lemma *take-takefill'*: $\bigwedge w\ n. \ n = k + m \implies \text{take } k \ (\text{takefill } \text{fill } n \ w) = \text{takefill } \text{fill } k \ w$
by (*induct* *k*) (*auto* *split*: *list.split*)

lemma *drop-takefill*: $\bigwedge w. \ \text{drop } k \ (\text{takefill } \text{fill } (m + k) \ w) = \text{takefill } \text{fill } m \ (\text{drop } k \ w)$
by (*induct* *k*) (*auto* *split*: *list.split*)

lemma *takefill-le* [*simp*]: $m \leq n \implies \text{takefill } x \ m \ (\text{takefill } x \ n \ l) = \text{takefill } x \ m \ l$
by (*auto* *simp*: *le-iff-add* *takefill-le'*)

lemma *take-takefill* [*simp*]: $m \leq n \implies \text{take } m \ (\text{takefill } \text{fill } n \ w) = \text{takefill } \text{fill } m \ w$
by (*auto* *simp*: *le-iff-add* *take-takefill'*)

lemma *takefill-append*: *takefill* *fill* (*m* + *length* *xs*) (*xs* @ *w*) = *xs* @ (*takefill* *fill* *m* *w*)
by (*induct* *xs*) *auto*

lemma *takefill-same'*: $l = \text{length } xs \implies \text{takefill } \text{fill } l \ xs = xs$
by (*induct* *xs* *arbitrary*: *l*) *auto*

lemmas *takefill-same* [*simp*] = *takefill-same'* [*OF refl*]

lemma *takefill-bintrunc*: *takefill* *False* *n* *bl* = *rev* (*bin-to-bl* *n* (*bl-to-bin* (*rev* *bl*)))
apply (*rule* *nth-equalityI*)
apply *simp*
apply (*clarsimp* *simp*: *nth-takefill* *nth-rev* *nth-bin-to-bl* *bin-nth-of-bl*)
done

lemma *bl-bin-bl-rtf*: *bin-to-bl* *n* (*bl-to-bin* *bl*) = *rev* (*takefill* *False* *n* (*rev* *bl*))
by (*simp* *add*: *takefill-bintrunc*)

lemma *bl-bin-bl-rep-drop*:
bin-to-bl *n* (*bl-to-bin* *bl*) =
replicate (*n* - *length* *bl*) *False* @ *drop* (*length* *bl* - *n*) *bl*
by (*simp* *add*: *bl-bin-bl-rtf* *takefill-alt* *rev-take*)

lemma *tf-rev*:
 $n + k = m + \text{length } bl \implies \text{takefill } x \ m \ (\text{rev } (\text{takefill } y \ n \ bl)) =$
 $\text{rev } (\text{takefill } y \ m \ (\text{rev } (\text{takefill } x \ k \ (\text{rev } bl))))$
apply (*rule* *nth-equalityI*)
apply (*auto* *simp* *add*: *nth-takefill* *nth-rev*)
apply (*rule-tac* *f* = $\lambda n. \ bl ! n$ **in** *arg-cong*)
apply *arith*

done

lemma *takefill-minus*: $0 < n \implies \text{takefill fill (Suc (n - 1)) } w = \text{takefill fill } n \ w$
by *auto*

lemmas *takefill-Suc-cases* =
list.cases [THEN takefill.Suc [THEN trans]]

lemmas *takefill-Suc-Nil* = *takefill-Suc-cases* (1)
lemmas *takefill-Suc-Cons* = *takefill-Suc-cases* (2)

lemmas *takefill-minus-simps* = *takefill-Suc-cases [THEN [2]*
takefill-minus [symmetric, THEN trans]]

lemma *takefill-numeral-Nil* [*simp*]:
 $\text{takefill fill (numeral } k) [] = \text{fill } \# \text{ takefill fill (pred-numeral } k) []$
by (*simp add: numeral-eq-Suc*)

lemma *takefill-numeral-Cons* [*simp*]:
 $\text{takefill fill (numeral } k) (x \# xs) = x \# \text{ takefill fill (pred-numeral } k) xs$
by (*simp add: numeral-eq-Suc*)

6.4 Links with function *bl-to-bin*

lemma *bl-to-bin-aux-cat*:
 $\bigwedge nv \ v. \text{bl-to-bin-aux } bs \ (\text{bin-cat } w \ nv \ v) =$
 $\text{bin-cat } w \ (nv + \text{length } bs) \ (\text{bl-to-bin-aux } bs \ v)$
by (*induct bs*) (*simp, simp add: bin-cat-Suc-Bit [symmetric] del: bin-cat.simps*)

lemma *bin-to-bl-aux-cat*:
 $\bigwedge w \ bs. \text{bin-to-bl-aux } (nv + nw) \ (\text{bin-cat } v \ nw \ w) \ bs =$
 $\text{bin-to-bl-aux } nv \ v \ (\text{bin-to-bl-aux } nw \ w \ bs)$
by (*induct nw*) *auto*

lemma *bl-to-bin-aux-alt*: $\text{bl-to-bin-aux } bs \ w = \text{bin-cat } w \ (\text{length } bs) \ (\text{bl-to-bin } bs)$
using *bl-to-bin-aux-cat [where nv = 0 and v = 0]*
by (*simp add: bl-to-bin-def [symmetric]*)

lemma *bin-to-bl-cat*:
 $\text{bin-to-bl } (nv + nw) \ (\text{bin-cat } v \ nw \ w) =$
 $\text{bin-to-bl-aux } nv \ v \ (\text{bin-to-bl } nw \ w)$
by (*simp add: bin-to-bl-def bin-to-bl-aux-cat*)

lemmas *bl-to-bin-aux-app-cat* =
trans [OF bl-to-bin-aux-append bl-to-bin-aux-alt]

lemmas *bin-to-bl-aux-cat-app* =
trans [OF bin-to-bl-aux-cat bin-to-bl-aux-alt]

lemma *bl-to-bin-app-cat*:

bl-to-bin (*bsa* @ *bs*) = *bin-cat* (*bl-to-bin* *bsa*) (*length* *bs*) (*bl-to-bin* *bs*)

by (*simp only*: *bl-to-bin-aux-app-cat bl-to-bin-def*)

lemma *bin-to-bl-cat-app*:

bin-to-bl (*n* + *nw*) (*bin-cat* *w nw wa*) = *bin-to-bl* *n w* @ *bin-to-bl* *nw wa*

by (*simp only*: *bin-to-bl-def bin-to-bl-aux-cat-app*)

bl-to-bin-app-cat-alt and *bl-to-bin-app-cat* are easily interderivable.

lemma *bl-to-bin-app-cat-alt*: *bin-cat* (*bl-to-bin* *cs*) *n w* = *bl-to-bin* (*cs* @ *bin-to-bl* *n w*)

by (*simp add*: *bl-to-bin-app-cat*)

lemma *mask-lem*: (*bl-to-bin* (*True* # *replicate* *n False*)) = *bl-to-bin* (*replicate* *n True*) + 1

apply (*unfold bl-to-bin-def*)

apply (*induct* *n*)

apply *simp*

apply (*simp only*: *Suc-eq-plus1 replicate-add append-Cons [symmetric] bl-to-bin-aux-append*)

apply (*simp add*: *Bit-B0-2t Bit-B1-2t*)

done

6.5 Function *bl-of-nth*

lemma *length-bl-of-nth* [*simp*]: *length* (*bl-of-nth* *n f*) = *n*

by (*induct* *n*) *auto*

lemma *nth-bl-of-nth* [*simp*]: *m* < *n* $\implies$ *rev* (*bl-of-nth* *n f*) ! *m* = *f m*

apply (*induct* *n*)

apply *simp*

apply (*clarsimp simp add*: *nth-append*)

apply (*rule-tac* *f = f in arg-cong*)

apply *simp*

done

lemma *bl-of-nth-inj*: ($\bigwedge k. k < n \implies f k = g k$) $\implies$ *bl-of-nth* *n f* = *bl-of-nth* *n g*

by (*induct* *n*) *auto*

lemma *bl-of-nth-nth-le*: *n* $\leq$ *length* *xs* $\implies$ *bl-of-nth* *n* (*nth* (*rev* *xs*)) = *drop* (*length* *xs* - *n*) *xs*

apply (*induct* *n arbitrary*: *xs*)

apply *clarsimp*

apply *clarsimp*

apply (*rule* *trans* [*OF* - *hd-Cons-tl*])

apply (*frule* *Suc-le-lessD*)

apply (*simp add*: *nth-rev trans* [*OF* *drop-Suc drop-tl*, *symmetric*])

apply (*subst* *hd-drop-conv-nth*)

apply *force*

apply *simp-all*

```

apply (rule-tac f =  $\lambda n. \text{drop } n \text{ xs}$  in arg-cong)
apply simp
done

lemma bl-of-nth-nth [simp]: bl-of-nth (length xs) (op ! (rev xs)) = xs
by (simp add: bl-of-nth-nth-le)

lemma size-rbl-pred: length (rbl-pred bl) = length bl
by (induct bl) auto

lemma size-rbl-succ: length (rbl-succ bl) = length bl
by (induct bl) auto

lemma size-rbl-add: length (rbl-add bl cl) = length bl
by (induct bl arbitrary: cl) (auto simp: Let-def size-rbl-succ)

lemma size-rbl-mult: length (rbl-mult bl cl) = length bl
by (induct bl arbitrary: cl) (auto simp add: Let-def size-rbl-add)

lemmas rbl-sizes [simp] =
  size-rbl-pred size-rbl-succ size-rbl-add size-rbl-mult

lemmas rbl-Nils =
  rbl-pred.Nil rbl-succ.Nil rbl-add.Nil rbl-mult.Nil

lemma rbl-pred: rbl-pred (rev (bin-to-bl n bin)) = rev (bin-to-bl n (bin - 1))
apply (unfold bin-to-bl-def)
apply (induct n arbitrary: bin)
apply simp
apply clarsimp
apply (case-tac bin rule: bin-exhaust)
apply (case-tac b)
apply (clarsimp simp: bin-to-bl-aux-alt)+
done

lemma rbl-succ: rbl-succ (rev (bin-to-bl n bin)) = rev (bin-to-bl n (bin + 1))
apply (unfold bin-to-bl-def)
apply (induct n arbitrary: bin)
apply simp
apply clarsimp
apply (case-tac bin rule: bin-exhaust)
apply (case-tac b)
apply (clarsimp simp: bin-to-bl-aux-alt)+
done

lemma rbl-add:
   $\bigwedge \text{bina binb. rbl-add (rev (bin-to-bl n bina)) (rev (bin-to-bl n binb)) =}$ 
   $\text{rev (bin-to-bl n (bina + binb))}$ 
apply (unfold bin-to-bl-def)

```

```

apply (induct n)
  apply simp
apply clarsimp
apply (case-tac bina rule: bin-exhaust)
apply (case-tac binb rule: bin-exhaust)
apply (case-tac b)
  apply (case-tac [!] ba)
    apply (auto simp: rbl-succ bin-to-bl-aux-alt Let-def ac-simps)
done

```

lemma *rbl-add-app2*: $\text{length } bbl \geq \text{length } bla \implies \text{rbl-add } bla \ (bbl @ bbl) = \text{rbl-add } bla \ bbl$

```

apply (induct bla arbitrary: bbl)
  apply simp
apply clarsimp
apply (case-tac bbl, clarsimp)
apply (clarsimp simp: Let-def)
done

```

lemma *rbl-add-take2*:

$\text{length } bbl \geq \text{length } bla \implies \text{rbl-add } bla \ (\text{take } (\text{length } bla) \ bbl) = \text{rbl-add } bla \ bbl$

```

apply (induct bla arbitrary: bbl)
  apply simp
apply clarsimp
apply (case-tac bbl, clarsimp)
apply (clarsimp simp: Let-def)
done

```

lemma *rbl-add-long*:

```

 $m \geq n \implies \text{rbl-add } (\text{rev } (\text{bin-to-bl } n \ bina)) \ (\text{rev } (\text{bin-to-bl } m \ binb)) =$ 
 $\text{rev } (\text{bin-to-bl } n \ (bina + binb))$ 
apply (rule box-equals [OF - rbl-add-take2 rbl-add])
apply (rule-tac f = rbl-add (rev (bin-to-bl n bina)) in arg-cong)
apply (rule rev-swap [THEN iffD1])
apply (simp add: rev-take drop-bin2bl)
apply simp
done

```

lemma *rbl-mult-app2*: $\text{length } bbl \geq \text{length } bla \implies \text{rbl-mult } bla \ (bbl @ bbl) = \text{rbl-mult } bla \ bbl$

```

apply (induct bla arbitrary: bbl)
  apply simp
apply clarsimp
apply (case-tac bbl, clarsimp)
apply (clarsimp simp: Let-def rbl-add-app2)
done

```

lemma *rbl-mult-take2*:

```

length blb ≥ length bla ⇒ rbl-mult bla (take (length bla) blb) = rbl-mult bla blb
apply (rule trans)
apply (rule rbl-mult-app2 [symmetric])
apply simp
apply (rule-tac f = rbl-mult bla in arg-cong)
apply (rule append-take-drop-id)
done

```

```

lemma rbl-mult-gt1:
  m ≥ length bl ⇒
    rbl-mult bl (rev (bin-to-bl m binb)) =
    rbl-mult bl (rev (bin-to-bl (length bl) binb))
apply (rule trans)
apply (rule rbl-mult-take2 [symmetric])
apply simp-all
apply (rule-tac f = rbl-mult bl in arg-cong)
apply (rule rev-swap [THEN iffD1])
apply (simp add: rev-take drop-bin2bl)
done

```

```

lemma rbl-mult-gt:
  m > n ⇒
    rbl-mult (rev (bin-to-bl n bina)) (rev (bin-to-bl m binb)) =
    rbl-mult (rev (bin-to-bl n bina)) (rev (bin-to-bl n binb))
by (auto intro: trans [OF rbl-mult-gt1])

```

```

lemmas rbl-mult-Suc = lessI [THEN rbl-mult-gt]

```

```

lemma rbbl-Cons: b # rev (bin-to-bl n x) = rev (bin-to-bl (Suc n) (x BIT b))
by (simp add: bin-to-bl-def) (simp add: bin-to-bl-aux-alt)

```

```

lemma rbl-mult:
  rbl-mult (rev (bin-to-bl n bina)) (rev (bin-to-bl n binb)) =
  rev (bin-to-bl n (bina * binb))
apply (induct n arbitrary: bina binb)
apply simp
apply (unfold bin-to-bl-def)
apply clarsimp
apply (case-tac bina rule: bin-exhaust)
apply (case-tac binb rule: bin-exhaust)
apply (case-tac b)
apply (case-tac [!] ba)
apply (auto simp: bin-to-bl-aux-alt Let-def)
apply (auto simp: rbbl-Cons rbl-mult-Suc rbl-add)
done

```

```

lemma rbl-add-split:
  P (rbl-add (y # ys) (x # xs)) =
  (∀ ws. length ws = length ys ⇒ ws = rbl-add ys xs ⇒

```

```

      (y → ((x → P (False # rbl-succ ws)) ∧ (¬ x → P (True # ws)))) ∧
      (¬ y → P (x # ws)))
  by (cases y) (auto simp: Let-def)

```

lemma *rbl-mult-split*:

```

P (rbl-mult (y # ys) xs) =
  (∀ ws. length ws = Suc (length ys) → ws = False # rbl-mult ys xs →
    (y → P (rbl-add ws xs)) ∧ (¬ y → P ws))
  by (auto simp: Let-def)

```

6.6 Repeated splitting or concatenation

lemma *sclem*: $\text{size} (\text{concat} (\text{map} (\text{bin-to-bl } n) \text{ xs})) = \text{length } \text{xs} * n$
 by (induct xs) auto

lemma *bin-cat-foldl-lem*:

```

foldl (λu. bin-cat u n) x xs =
  bin-cat x (size xs * n) (foldl (λu. bin-cat u n) y xs)
  apply (induct xs arbitrary: x)
  apply simp
  apply (simp (no-asm))
  apply (frule asm-rl)
  apply (drule meta-spec)
  apply (erule trans)
  apply (drule-tac x = bin-cat y n a in meta-spec)
  apply (simp add: bin-cat-assoc-sym min.absorb2)
  done

```

lemma *bin-rcat-bl*: $\text{bin-rcat } n \text{ wl} = \text{bl-to-bin} (\text{concat} (\text{map} (\text{bin-to-bl } n) \text{ wl}))$
 apply (unfold bin-rcat-def)
 apply (rule sym)
 apply (induct wl)
 apply (auto simp add: bl-to-bin-append)
 apply (simp add: bl-to-bin-aux-alt sclem)
 apply (simp add: bin-cat-foldl-lem [symmetric])
 done

lemmas *bin-rsplit-aux-simps* = *bin-rsplit-aux.simps* *bin-rsplitl-aux.simps*

lemmas *rsplit-aux-simps* = *bin-rsplit-aux-simps*

lemmas *th-if-simp1* = *if-split* [where $P = \text{op} = l$, THEN *iffD1*, THEN *conjunct1*,
 THEN *mp*] for *l*

lemmas *th-if-simp2* = *if-split* [where $P = \text{op} = l$, THEN *iffD1*, THEN *conjunct2*,
 THEN *mp*] for *l*

lemmas *rsplit-aux-simp1s* = *rsplit-aux-simps* [THEN *th-if-simp1*]

lemmas *rsplit-aux-simp2ls* = *rsplit-aux-simps* [THEN *th-if-simp2*]

lemmas *bin-rsplit-aux-simp2s* [*simp*] = *rsplit-aux-simp2ls* [*unfolded Let-def*]
lemmas *rbscl* = *bin-rsplit-aux-simp2s* (2)

lemmas *rsplit-aux-0-simps* [*simp*] =
rsplit-aux-simp1s [*OF disjI1*] *rsplit-aux-simp1s* [*OF disjI2*]

lemma *bin-rsplit-aux-append*: *bin-rsplit-aux* *n m c* (*bs @ cs*) = *bin-rsplit-aux* *n m c bs @ cs*
apply (*induct* *n m c bs* *rule*: *bin-rsplit-aux.induct*)
apply (*subst* *bin-rsplit-aux.simps*)
apply (*subst* *bin-rsplit-aux.simps*)
apply (*clarsimp* *split*: *prod.split*)
done

lemma *bin-rsplittl-aux-append*: *bin-rsplittl-aux* *n m c* (*bs @ cs*) = *bin-rsplittl-aux* *n m c bs @ cs*
apply (*induct* *n m c bs* *rule*: *bin-rsplittl-aux.induct*)
apply (*subst* *bin-rsplittl-aux.simps*)
apply (*subst* *bin-rsplittl-aux.simps*)
apply (*clarsimp* *split*: *prod.split*)
done

lemmas *rsplit-aux-apps* [**where** *bs* = []] =
bin-rsplit-aux-append *bin-rsplittl-aux-append*

lemmas *rsplit-def-auxs* = *bin-rsplit-def* *bin-rsplittl-def*

lemmas *rsplit-aux-alt*s = *rsplit-aux-apps*
[*unfolded append-Nil* *rsplit-def-auxs* [*symmetric*]]

lemma *bin-split-minus*: $0 < n \implies \text{bin-split } (\text{Suc } (n - 1)) \ w = \text{bin-split } n \ w$
by *auto*

lemmas *bin-split-minus-simp* =
bin-split.Suc [*THEN* [2] *bin-split-minus* [*symmetric*, *THEN* *trans*]]

lemma *bin-split-pred-simp* [*simp*]:
 $(0::\text{nat}) < \text{numeral } \text{bin} \implies$
bin-split (*numeral bin*) *w* =
(*let* (*w1*, *w2*) = *bin-split* (*numeral bin* - 1) (*bin-rest w*)
in (*w1*, *w2 BIT bin-last w*))
by (*simp only*: *bin-split-minus-simp*)

lemma *bin-rsplit-aux-simp-alt*:
bin-rsplit-aux *n m c bs* =
(*if* *m* = 0 $\vee$ *n* = 0 *then* *bs*
else *let* (*a*, *b*) = *bin-split* *n c* in *bin-rsplit* *n* (*m* - *n*, *a*) @ *b* # *bs*)
apply (*simp add*: *bin-rsplit-aux.simps* [*of* *n m c bs*])
apply (*subst* *rsplit-aux-alt*s)

apply (*simp add: bin-rsplit-def*)
done

lemmas *bin-rsplit-simp-alt* =
trans [OF bin-rsplit-def bin-rsplit-aux-simp-alt]

lemmas *bthrs* = *bin-rsplit-simp-alt [THEN [2] trans]*

lemma *bin-rsplit-size-sign'* [*rule-format*]:
 $n > 0 \implies \text{rev } sw = \text{bin-rsplit } n \ (nw, w) \implies \forall v \in \text{set } sw. \text{bintrunc } n \ v = v$
apply (*induct sw arbitrary: nw w*)
apply *clarsimp*
apply *clarsimp*
apply (*drule bthrs*)
apply (*simp (no-asm-use) add: Let-def split: prod.split-asm if-split-asm*)
apply *clarify*
apply (*drule split-bintrunc*)
apply *simp*
done

lemmas *bin-rsplit-size-sign* = *bin-rsplit-size-sign' [OF asm-rl
rev-rev-ident [THEN trans] set-rev [THEN equalityD2 [THEN subsetD]]]*

lemma *bin-nth-rsplit* [*rule-format*] :
 $n > 0 \implies m < n \implies$
 $\forall w \ k \ nw. \text{rev } sw = \text{bin-rsplit } n \ (nw, w) \longrightarrow$
 $k < \text{size } sw \longrightarrow \text{bin-nth } (sw ! k) \ m = \text{bin-nth } w \ (k * n + m)$
apply (*induct sw*)
apply *clarsimp*
apply *clarsimp*
apply (*drule bthrs*)
apply (*simp (no-asm-use) add: Let-def split: prod.split-asm if-split-asm*)
apply *clarify*
apply (*erule allE, erule impE, erule exI*)
apply (*case-tac k*)
apply *clarsimp*
prefer 2
apply *clarsimp*
apply (*erule allE*)
apply (*erule (1) impE*)
apply (*drule bin-nth-split, erule conjE, erule allE, erule trans, simp add:*
ac-simps)
done

lemma *bin-rsplit-all*: $0 < nw \implies nw \leq n \implies \text{bin-rsplit } n \ (nw, w) = [\text{bintrunc } n \ w]$
by (*auto simp: bin-rsplit-def rsplit-aux-simp2ls split: prod.split dest!: split-bintrunc*)

```

lemma bin-rsplit-l [rule-format]:
   $\forall \text{bin. bin-rsplittl } n (m, \text{bin}) = \text{bin-rsplit } n (m, \text{bintrunc } m \text{ bin})$ 
  apply (rule-tac a = m in wf-less-than [THEN wf-induct])
  apply (simp (no-asm) add: bin-rsplittl-def bin-rsplit-def)
  apply (rule allI)
  apply (subst bin-rsplittl-aux.simps)
  apply (subst bin-rsplit-aux.simps)
  apply (clarsimp simp: Let-def split: prod.split)
  apply (drule bin-split-trunc)
  apply (drule sym [THEN trans], assumption)
  apply (subst rsplit-aux-alts(1))
  apply (subst rsplit-aux-alts(2))
  apply clarsimp
  unfolding bin-rsplit-def bin-rsplittl-def
  apply simp
  done

lemma bin-rsplit-rcat [rule-format]:
   $n > 0 \longrightarrow \text{bin-rsplit } n (n * \text{size } ws, \text{bin-rcat } n \text{ ws}) = \text{map } (\text{bintrunc } n) \text{ ws}$ 
  apply (unfold bin-rsplit-def bin-rcat-def)
  apply (rule-tac xs = ws in rev-induct)
  apply clarsimp
  apply clarsimp
  apply (subst rsplit-aux-alts)
  unfolding bin-split-cat
  apply simp
  done

lemma bin-rsplit-aux-len-le [rule-format] :
   $\forall ws \ m. \ n \neq 0 \longrightarrow ws = \text{bin-rsplit-aux } n \text{ nw } w \text{ bs} \longrightarrow$ 
   $\text{length } ws \leq m \longleftrightarrow \text{nw} + \text{length } bs * n \leq m * n$ 
proof –
  have *: R
  if d:  $i \leq j \vee m < j'$ 
  and R1:  $i * k \leq j * k \implies R$ 
  and R2:  $\text{Suc } m * k' \leq j' * k' \implies R$ 
  for  $i \ j \ j' \ k \ k' \ m :: \text{nat}$  and R
  using d
  apply safe
  apply (rule R1, erule mult-le-mono1)
  apply (rule R2, erule Suc-le-eq [THEN iffD2 [THEN mult-le-mono1]])
  done
have **:  $0 < sc \implies sc - n + (n + lb * n) \leq m * n \longleftrightarrow sc + lb * n \leq m * n$ 
for  $sc \ m \ n \ lb :: \text{nat}$ 
  apply safe
  apply arith
  apply (case-tac sc  $\geq n$ )
  apply arith
  apply (insert linorder-le-less-linear [of m lb])

```

```

  apply (erule-tac k=n and k'=n in *)
  apply arith
  apply simp
  done
show ?thesis
  apply (induct n nw w bs rule: bin-rsplit-aux.induct)
  apply (subst bin-rsplit-aux.simps)
  apply (simp add: ** Let-def split: prod.split)
  done
qed

```

lemma *bin-rsplit-len-le*: $n \neq 0 \longrightarrow ws = \text{bin-rsplit } n \ (nw, w) \longrightarrow \text{length } ws \leq m$
 $\longleftrightarrow nw \leq m * n$
by (*auto simp: bin-rsplit-def bin-rsplit-aux-len-le*)

lemma *bin-rsplit-aux-len*:
 $n \neq 0 \implies \text{length } (\text{bin-rsplit-aux } n \ nw \ w \ cs) = (nw + n - 1) \text{ div } n + \text{length } cs$
apply (*induct n nw w cs rule: bin-rsplit-aux.induct*)
apply (*subst bin-rsplit-aux.simps*)
apply (*clarsimp simp: Let-def split: prod.split*)
apply (*erule thin-rl*)
apply (*case-tac m*)
apply *simp*
apply (*case-tac m ≤ n*)
apply (*auto simp add: div-add-self2*)
done

lemma *bin-rsplit-len*: $n \neq 0 \implies \text{length } (\text{bin-rsplit } n \ (nw, w)) = (nw + n - 1) \text{ div } n$
by (*auto simp: bin-rsplit-def bin-rsplit-aux-len*)

lemma *bin-rsplit-aux-len-indep*:
 $n \neq 0 \implies \text{length } bs = \text{length } cs \implies$
 $\text{length } (\text{bin-rsplit-aux } n \ nw \ v \ bs) =$
 $\text{length } (\text{bin-rsplit-aux } n \ nw \ w \ cs)$
proof (*induct n nw w cs arbitrary: v bs rule: bin-rsplit-aux.induct*)
case (*1 n m w cs v bs*)
show ?case
proof (*cases m = 0*)
case *True*
with $\langle \text{length } bs = \text{length } cs \rangle$ **show** ?thesis **by** *simp*
next
case *False*
from *1.hyps* $\langle m \neq 0 \rangle \langle n \neq 0 \rangle$ **have** *hyp*: $\bigwedge v \ bs. \text{length } bs = \text{Suc } (\text{length } cs)$
 $\implies$
 $\text{length } (\text{bin-rsplit-aux } n \ (m - n) \ v \ bs) =$
 $\text{length } (\text{bin-rsplit-aux } n \ (m - n) \ (\text{fst } (\text{bin-split } n \ w)) \ (\text{snd } (\text{bin-split } n \ w) \ \#$
 $cs))$
by *auto*

```

from  $\langle \text{length } bs = \text{length } cs \rangle \langle n \neq 0 \rangle$  show ?thesis
by (auto simp add: bin-rsplit-aux-simp-alt Let-def bin-rsplit-len split: prod.split)
qed
qed

```

```

lemma bin-rsplit-len-indep:
   $n \neq 0 \implies \text{length } (\text{bin-rsplit } n \ (nw, v)) = \text{length } (\text{bin-rsplit } n \ (nw, w))$ 
apply (unfold bin-rsplit-def)
apply (simp (no-asm))
apply (erule bin-rsplit-aux-len-indep)
apply (rule refl)
done

```

Even more bit operations

```

instantiation int :: bitss
begin

```

```

definition [iff]:  $i !! n \longleftrightarrow \text{bin-nth } i \ n$ 

```

```

definition  $\text{lsb } i = i !! 0$  for  $i :: \text{int}$ 

```

```

definition  $\text{set-bit } i \ n \ b = \text{bin-sc } n \ b \ i$ 

```

```

definition
  set-bits  $f =$ 
    (if  $\exists n. \forall n' \geq n. \neg f \ n'$  then
      let  $n = \text{LEAST } n. \forall n' \geq n. \neg f \ n'$ 
      in  $\text{bl-to-bin } (\text{rev } (\text{map } f \ [0..<n]))$ 
    else if  $\exists n. \forall n' \geq n. f \ n'$  then
      let  $n = \text{LEAST } n. \forall n' \geq n. f \ n'$ 
      in  $\text{sbintrunc } n \ (\text{bl-to-bin } (\text{True} \ \# \ \text{rev } (\text{map } f \ [0..<n])))$ 
    else  $0 :: \text{int}$ )

```

```

definition  $\text{shiffl } x \ n = x * 2 ^ n$  for  $x :: \text{int}$ 

```

```

definition  $\text{shiftr } x \ n = x \text{ div } 2 ^ n$  for  $x :: \text{int}$ 

```

```

definition  $\text{msb } x \longleftrightarrow x < 0$  for  $x :: \text{int}$ 

```

```

instance ..

```

```

end

```

```

end

```

7 Type Definition Theorems

```

theory Misc-Typedef
imports Main

```

begin

8 More lemmas about normal type definitions

lemma

tdD1: type-definition Rep Abs A $\implies \forall x. \text{Rep } x \in A$ and
tdD2: type-definition Rep Abs A $\implies \forall x. \text{Abs } (\text{Rep } x) = x$ and
tdD3: type-definition Rep Abs A $\implies \forall y. y \in A \longrightarrow \text{Rep } (\text{Abs } y) = y$
by (*auto simp: type-definition-def*)

lemma *td-nat-int:*

type-definition int nat (Collect (op <= 0))
unfolding *type-definition-def* **by** *auto*

context *type-definition*

begin

declare *Rep* [*iff*] *Rep-inverse* [*simp*] *Rep-inject* [*simp*]

lemma *Abs-eqD*: *Abs x = Abs y $\implies x \in A \implies y \in A \implies x = y$*
by (*simp add: Abs-inject*)

lemma *Abs-inverse'*:

r : A $\implies \text{Abs } r = a \implies \text{Rep } a = r$
by (*safe elim!: Abs-inverse*)

lemma *Rep-comp-inverse*:

Rep $\circ f = g \implies \text{Abs} \circ g = f$
using *Rep-inverse* **by** *auto*

lemma *Rep-eqD* [*elim!*]: *Rep x = Rep y $\implies x = y$*
by *simp*

lemma *Rep-inverse'*: *Rep a = r $\implies \text{Abs } r = a$*
by (*safe intro!: Rep-inverse*)

lemma *comp-Abs-inverse*:

f $\circ \text{Abs} = g \implies g \circ \text{Rep} = f$
using *Rep-inverse* **by** *auto*

lemma *set-Rep*:

A = range Rep

proof (*rule set-eqI*)

fix *x*

show (*x* $\in A$) = (*x* $\in \text{range Rep}$)

by (*auto dest: Abs-inverse [of x, symmetric]*)

qed

lemma *set-Rep-Abs*: *A = range (Rep $\circ$ Abs)*

```

proof (rule set-eqI)
  fix x
  show  $(x \in A) = (x \in \text{range } (Rep \circ Abs))$ 
    by (auto dest: Abs-inverse [of x, symmetric])
qed

lemma Abs-inj-on: inj-on Abs A
  unfolding inj-on-def
  by (auto dest: Abs-inject [THEN iffD1])

lemma image: Abs ‘ A = UNIV
  by (auto intro!: image-eqI)

lemmas td-thm = type-definition-axioms

lemma fns1:
   $Rep \circ fa = fr \circ Rep \mid fa \circ Abs = Abs \circ fr \implies Abs \circ fr \circ Rep = fa$ 
  by (auto dest: Rep-comp-inverse elim: comp-Abs-inverse simp: o-assoc)

lemmas fns1a = disjI1 [THEN fns1]
lemmas fns1b = disjI2 [THEN fns1]

lemma fns4:
   $Rep \circ fa \circ Abs = fr \implies$ 
   $Rep \circ fa = fr \circ Rep \ \& \ fa \circ Abs = Abs \circ fr$ 
  by auto

end

interpretation nat-int: type-definition int nat Collect (op <= 0)
  by (rule td-nat-int)

declare
  nat-int.Rep-cases [cases del]
  nat-int.Abs-cases [cases del]
  nat-int.Rep-induct [induct del]
  nat-int.Abs-induct [induct del]

```

8.1 Extended form of type definition predicate

```

lemma td-conds:
   $norm \circ norm = norm \implies (fr \circ norm = norm \circ fr) =$ 
   $(norm \circ fr \circ norm = fr \circ norm \ \& \ norm \circ fr \circ norm = norm \circ fr)$ 
  apply safe
  apply (simp-all add: comp-assoc)
  apply (simp-all add: o-assoc)
  done

lemma fn-comm-power:

```

```

fa ∘ tr = tr ∘ fr ==> fa ^ n ∘ tr = tr ∘ fr ^ n
apply (rule ext)
apply (induct n)
apply (auto dest: fun-cong)
done

lemmas fn-comm-power' =
  ext [THEN fn-comm-power, THEN fun-cong, unfolded o-def]

locale td-ext = type-definition +
  fixes norm
  assumes eq-norm:  $\bigwedge x. \text{Rep } (\text{Abs } x) = \text{norm } x$ 
begin

lemma Abs-norm [simp]:
  Abs (norm x) = Abs x
  using eq-norm [of x] by (auto elim: Rep-inverse')

lemma td-th:
  g ∘ Abs = f ==> f (Rep x) = g x
  by (drule comp-Abs-inverse [symmetric]) simp

lemma eq-norm': Rep ∘ Abs = norm
  by (auto simp: eq-norm)

lemma norm-Rep [simp]: norm (Rep x) = Rep x
  by (auto simp: eq-norm' intro: td-th)

lemmas td = td-thm

lemma set-iff-norm: w : A  $\longleftrightarrow$  w = norm w
  by (auto simp: set-Rep-Abs eq-norm' eq-norm [symmetric])

lemma inverse-norm:
  (Abs n = w) = (Rep w = norm n)
  apply (rule iffI)
  apply (clarsimp simp add: eq-norm)
  apply (simp add: eq-norm' [symmetric])
  done

lemma norm-eq-iff:
  (norm x = norm y) = (Abs x = Abs y)
  by (simp add: eq-norm' [symmetric])

lemma norm-comps:
  Abs ∘ norm = Abs
  norm ∘ Rep = Rep
  norm ∘ norm = norm

```

by (*auto simp: eq-norm' [symmetric] o-def*)

lemmas *norm-norm [simp] = norm-comps*

lemma *fns5:*

Rep $\circ$ *fa* $\circ$ *Abs* = *fr* ==>
fr $\circ$ *norm* = *fr* & *norm* $\circ$ *fr* = *fr*
by (*fold eq-norm'*) *auto*

lemma *fns2:*

Abs $\circ$ *fr* $\circ$ *Rep* = *fa* ==>
(*norm* $\circ$ *fr* $\circ$ *norm* = *fr* $\circ$ *norm*) = (*Rep* $\circ$ *fa* = *fr* $\circ$ *Rep*)
apply (*fold eq-norm'*)
apply *safe*
prefer 2
apply (*simp add: o-assoc*)
apply (*rule ext*)
apply (*drule-tac x=Rep x in fun-cong*)
apply *auto*
done

lemma *fns3:*

Abs $\circ$ *fr* $\circ$ *Rep* = *fa* ==>
(*norm* $\circ$ *fr* $\circ$ *norm* = *norm* $\circ$ *fr*) = (*fa* $\circ$ *Abs* = *Abs* $\circ$ *fr*)
apply (*fold eq-norm'*)
apply *safe*
prefer 2
apply (*simp add: comp-assoc*)
apply (*rule ext*)
apply (*drule-tac f=a $\circ$ b for a b in fun-cong*)
apply *simp*
done

lemma *fns:*

fr $\circ$ *norm* = *norm* $\circ$ *fr* ==>
(*fa* $\circ$ *Abs* = *Abs* $\circ$ *fr*) = (*Rep* $\circ$ *fa* = *fr* $\circ$ *Rep*)
apply *safe*
apply (*frule fns1b*)
prefer 2
apply (*frule fns1a*)
apply (*rule fns3 [THEN iffD1]*)
prefer 3
apply (*rule fns2 [THEN iffD1]*)
apply (*simp-all add: comp-assoc*)
apply (*simp-all add: o-assoc*)
done

lemma *range-norm:*

```

    range (Rep ◦ Abs) = A
  by (simp add: set-Rep-Abs)

end

lemmas td-ext-def' =
  td-ext-def [unfolded type-definition-def td-ext-axioms-def]

end

```

9 Miscellaneous lemmas, of at least doubtful value

```

theory Word-Miscellaneous
  imports HOL-Library.Bit Misc-Numeric
begin

lemma power-minus-simp:  $0 < n \implies a ^ n = a * a ^ (n - 1)$ 
  by (auto dest: gr0-implies-Suc)

lemma funpow-minus-simp:  $0 < n \implies f ^ n = f \circ f ^ (n - 1)$ 
  by (auto dest: gr0-implies-Suc)

lemma power-numeral:  $a ^ \text{numeral } k = a * a ^ (\text{pred-numeral } k)$ 
  by (simp add: numeral-eq-Suc)

lemma funpow-numeral [simp]:  $f ^ \text{numeral } k = f \circ f ^ (\text{pred-numeral } k)$ 
  by (simp add: numeral-eq-Suc)

lemma replicate-numeral [simp]:  $\text{replicate } (\text{numeral } k) \ x = x \# \text{replicate } (\text{pred-numeral } k) \ x$ 
  by (simp add: numeral-eq-Suc)

lemma rco-alt:  $(f \circ g) ^ n \circ f = f \circ (g \circ f) ^ n$ 
  apply (rule ext)
  apply (induct n)
  apply (simp-all add: o-def)
done

lemma list-exhaust-size-gt0:
  assumes  $y: \bigwedge a \text{ list. } y = a \# \text{list} \implies P$ 
  shows  $0 < \text{length } y \implies P$ 
  apply (cases y)
  apply simp
  apply (rule y)
  apply fastforce
done

lemma list-exhaust-size-eq0:
  assumes  $y: y = [] \implies P$ 

```

```

shows  $\text{length } y = 0 \implies P$ 
apply (cases y)
  apply (rule y)
  apply simp
  apply simp
done

```

```

lemma size-Cons-lem-eq:  $y = xa \# \text{list} \implies \text{size } y = \text{Suc } k \implies \text{size list} = k$ 
  by auto

```

```

lemmas ls-splits = prod.split prod.split-asm if-split-asm

```

```

lemma not-B1-is-B0:  $y \neq 1 \implies y = 0$ 
  for  $y :: \text{bit}$ 
  by (cases y) auto

```

```

lemma B1-ass-B0:
  fixes  $y :: \text{bit}$ 
  assumes  $y: y = 0 \implies y = 1$ 
  shows  $y = 1$ 
  apply (rule classical)
  apply (drule not-B1-is-B0)
  apply (erule y)
done

```

— simplifications for specific word lengths

```

lemmas n2s-ths [THEN eq-reflection] = add-2-eq-Suc add-2-eq-Suc'

```

```

lemmas s2n-ths = n2s-ths [symmetric]

```

```

lemma and-len:  $xs = ys \implies xs = ys \wedge \text{length } xs = \text{length } ys$ 
  by auto

```

```

lemma size-if:  $\text{size } (\text{if } p \text{ then } xs \text{ else } ys) = (\text{if } p \text{ then } \text{size } xs \text{ else } \text{size } ys)$ 
  by auto

```

```

lemma tl-if:  $\text{tl } (\text{if } p \text{ then } xs \text{ else } ys) = (\text{if } p \text{ then } \text{tl } xs \text{ else } \text{tl } ys)$ 
  by auto

```

```

lemma hd-if:  $\text{hd } (\text{if } p \text{ then } xs \text{ else } ys) = (\text{if } p \text{ then } \text{hd } xs \text{ else } \text{hd } ys)$ 
  by auto

```

```

lemma if-Not-x:  $(\text{if } p \text{ then } \neg x \text{ else } x) = (p = (\neg x))$ 
  by auto

```

```

lemma if-x-Not:  $(\text{if } p \text{ then } x \text{ else } \neg x) = (p = x)$ 
  by auto

```

```

lemma if-same-and:  $(\text{If } p \ x \ y \wedge \text{If } p \ u \ v) = (\text{if } p \text{ then } x \wedge u \text{ else } y \wedge v)$ 

```

by *auto*

lemma *if-same-eq*: $(\text{If } p \ x \ y = (\text{If } p \ u \ v)) = (\text{if } p \ \text{then } x = u \ \text{else } y = v)$
by *auto*

lemma *if-same-eq-not*: $(\text{If } p \ x \ y = (\neg \text{If } p \ u \ v)) = (\text{if } p \ \text{then } x = (\neg u) \ \text{else } y = (\neg v))$
by *auto*

lemma *if-Cons*: $(\text{if } p \ \text{then } x \# \ xs \ \text{else } y \# \ ys) = \text{If } p \ x \ y \# \ \text{If } p \ xs \ ys$
by *auto*

lemma *if-single*: $(\text{if } xc \ \text{then } [xab] \ \text{else } [an]) = [\text{if } xc \ \text{then } xab \ \text{else } an]$
by *auto*

lemma *if-bool-simps*:
 $\text{If } p \ \text{True } y = (p \vee y) \wedge \text{If } p \ \text{False } y = (\neg p \wedge y) \wedge$
 $\text{If } p \ y \ \text{True} = (p \longrightarrow y) \wedge \text{If } p \ y \ \text{False} = (p \wedge y)$
by *auto*

lemmas *if-simps* =
if-x-Not if-Not-x if-cancel if-True if-False if-bool-simps

lemmas *seqr* = *eq-reflection* [**where** $x = \text{size } w$] **for** w

lemma *the-elemI*: $y = \{x\} \Longrightarrow \text{the-elem } y = x$
by *simp*

lemma *nonemptyE*: $S \neq \{\} \Longrightarrow (\bigwedge x. x \in S \Longrightarrow R) \Longrightarrow R$
by *auto*

lemma *gt-or-eq-0*: $0 < y \vee 0 = y$
for $y :: \text{nat}$
by *arith*

lemmas *xtr1* = *xtrans*(1)
lemmas *xtr2* = *xtrans*(2)
lemmas *xtr3* = *xtrans*(3)
lemmas *xtr4* = *xtrans*(4)
lemmas *xtr5* = *xtrans*(5)
lemmas *xtr6* = *xtrans*(6)
lemmas *xtr7* = *xtrans*(7)
lemmas *xtr8* = *xtrans*(8)

lemmas *nat-simps* = *diff-add-inverse2* *diff-add-inverse*
lemmas *nat-iffs* = *le-add1* *le-add2*

lemma *sum-imp-diff*: $j = k + i \Longrightarrow j - i = k$

```

for  $k :: nat$ 
by arith

lemmas  $pos\text{-}mod\text{-}sign2 = zless2 \ [THEN\ pos\text{-}mod\text{-}sign \ [where\ b = 2::int]]$ 
lemmas  $pos\text{-}mod\text{-}bound2 = zless2 \ [THEN\ pos\text{-}mod\text{-}bound \ [where\ b = 2::int]]$ 

lemma  $nmod2: n \bmod 2 = 0 \vee n \bmod 2 = 1$ 
for  $n :: int$ 
by arith

lemmas  $eme1p = emep1 \ [simplified\ add.commute]$ 

lemma  $le\text{-}diff\text{-}eq': a \leq c - b \longleftrightarrow b + a \leq c$ 
for  $a\ b\ c :: int$ 
by arith

lemma  $less\text{-}diff\text{-}eq': a < c - b \longleftrightarrow b + a < c$ 
for  $a\ b\ c :: int$ 
by arith

lemma  $diff\text{-}less\text{-}eq': a - b < c \longleftrightarrow a < b + c$ 
for  $a\ b\ c :: int$ 
by arith

lemmas  $m1mod22k = mult\text{-}pos\text{-}pos \ [OF\ zless2\ zless2p,\ THEN\ zmod\text{-}minus1]$ 

lemma  $z1pdiv2: (2 * b + 1) \div 2 = b$ 
for  $b :: int$ 
by arith

lemmas  $zdiv\text{-}le\text{-}dividend = xtr3 \ [OF\ div\text{-}by\text{-}1 \ [symmetric] \ zdiv\text{-}mono2,$ 
 $\quad simplified\ int\text{-}one\text{-}le\text{-}iff\text{-}zero\text{-}less,\ simplified]$ 

lemma  $axbbyy: a + m + m = b + n + n \implies a = 0 \vee a = 1 \implies b = 0 \vee b =$ 
 $1 \implies a = b \wedge m = n$ 
for  $a\ b\ m\ n :: int$ 
by arith

lemma  $axxmod2: (1 + x + x) \bmod 2 = 1 \wedge (0 + x + x) \bmod 2 = 0$ 
for  $x :: int$ 
by arith

lemma  $axxdiv2: (1 + x + x) \div 2 = x \wedge (0 + x + x) \div 2 = x$ 
for  $x :: int$ 
by arith

lemmas  $iszero\text{-}minus =$ 
 $\quad trans \ [THEN\ trans,\ OF\ iszero\text{-}def\ neg\text{-}equal\text{-}0\text{-}iff\text{-}equal\ iszero\text{-}def \ [symmetric]]$ 

```

```

lemmas zadd-diff-inverse =
  trans [OF diff-add-cancel [symmetric] add commute]

lemmas add-diff-cancel2 =
  add commute [THEN diff-eq-eq [THEN iffD2]]

lemmas rdmodes [symmetric] = mod-minus-eq
  mod-diff-left-eq mod-diff-right-eq mod-add-left-eq
  mod-add-right-eq mod-mult-right-eq mod-mult-left-eq

lemma mod-plus-right:  $(a + x) \bmod m = (b + x) \bmod m \longleftrightarrow a \bmod m = b \bmod m$ 
  for  $a\ b\ m\ x :: \text{nat}$ 
  by (induct x) (simp-all add: mod-Suc, arith)

lemma nat-minus-mod:  $(n - n \bmod m) \bmod m = 0$ 
  for  $m\ n :: \text{nat}$ 
  by (induct n) (simp-all add: mod-Suc)

lemmas nat-minus-mod-plus-right =
  trans [OF nat-minus-mod mod-0 [symmetric],
    THEN mod-plus-right [THEN iffD2], simplified]

lemmas push-mods' = mod-add-eq
  mod-mult-eq mod-diff-eq
  mod-minus-eq

lemmas push-mods = push-mods' [THEN eq-reflection]
lemmas pull-mods = push-mods [symmetric] rdmodes [THEN eq-reflection]

lemma nat-mod-eq:  $b < n \implies a \bmod n = b \bmod n \implies a \bmod n = b$ 
  for  $a\ b\ n :: \text{nat}$ 
  by (induct a) auto

lemmas nat-mod-eq' = refl [THEN [2] nat-mod-eq]

lemma nat-mod-lem:  $0 < n \implies b < n \longleftrightarrow b \bmod n = b$ 
  for  $b\ n :: \text{nat}$ 
  apply safe
  apply (erule nat-mod-eq')
  apply (erule subst)
  apply (erule mod-less-divisor)
  done

lemma mod-nat-add:  $x < z \implies y < z \implies (x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$ 
  for  $x\ y\ z :: \text{nat}$ 
  apply (rule nat-mod-eq)
  apply auto

```

```

apply (rule trans)
apply (rule le-mod-geq)
apply simp
apply (rule nat-mod-eq')
apply arith
done

lemma mod-nat-sub:  $x < z \implies (x - y) \bmod z = x - y$ 
  for  $x\ y :: \text{nat}$ 
  by (rule nat-mod-eq') arith

lemma int-mod-eq:  $0 \leq b \implies b < n \implies a \bmod n = b \bmod n \implies a \bmod n = b$ 
  for  $a\ b\ n :: \text{int}$ 
  by (metis mod-pos-pos-trivial)

lemmas int-mod-eq' = mod-pos-pos-trivial

lemma int-mod-le:  $0 \leq a \implies a \bmod n \leq a$ 
  for  $a :: \text{int}$ 
  by (fact Divides.semiring-numeral-div-class.mod-less-eq-dividend)

lemma mod-add-if-z:
   $x < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$ 
   $(x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$ 
  for  $x\ y\ z :: \text{int}$ 
  by (auto intro: int-mod-eq)

lemma mod-sub-if-z:
   $x < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$ 
   $(x - y) \bmod z = (\text{if } y \leq x \text{ then } x - y \text{ else } x - y + z)$ 
  for  $x\ y\ z :: \text{int}$ 
  by (auto intro: int-mod-eq)

lemmas zmde = mult-div-mod-eq [symmetric, THEN diff-eq-eq [THEN iffD2],
symmetric]
lemmas mcl = mult-cancel-left [THEN iffD1, THEN make-pos-rule]

lemma zdiv-mult-self:  $m \neq 0 \implies (a + m * n) \text{ div } m = a \text{ div } m + n$ 
  for  $a\ m\ n :: \text{int}$ 
  apply (rule mcl)
  prefer 2
  apply (erule asm-rl)
  apply (simp add: zmde ring-distrib)
  done

lemma mod-power-lem:  $a > 1 \implies a ^ n \bmod a ^ m = (\text{if } m \leq n \text{ then } 0 \text{ else } a ^$ 
 $n)$ 
  for  $a :: \text{int}$ 

```

```

apply clarsimp
apply safe
apply (simp add: dvd-eq-mod-eq-0 [symmetric])
apply (drule le-iff-add [THEN iffD1])
apply (force simp: power-add)
apply (rule mod-pos-pos-trivial)
apply (simp)
apply (rule power-strict-increasing)
apply auto
done

lemma pl-pl-rels:  $a + b = c + d \implies a \geq c \wedge b \leq d \vee a \leq c \wedge b \geq d$ 
  for  $a\ b\ c\ d :: \text{nat}$ 
  by arith

lemmas pl-pl-rels' = add.commute [THEN [2] trans, THEN pl-pl-rels]

lemma minus-eq:  $m - k = m \longleftrightarrow k = 0 \vee m = 0$ 
  for  $k\ m :: \text{nat}$ 
  by arith

lemma pl-pl-mm:  $a + b = c + d \implies a - c = d - b$ 
  for  $a\ b\ c\ d :: \text{nat}$ 
  by arith

lemmas pl-pl-mm' = add.commute [THEN [2] trans, THEN pl-pl-mm]

lemmas dme = div-mult-mod-eq
lemmas dtle = xtr3 [OF dme [symmetric] le-add1]
lemmas th2 = order-trans [OF order-refl [THEN [2] mult-le-mono] dtle]

lemma td-gal:  $0 < c \implies a \geq b * c \longleftrightarrow a \text{ div } c \geq b$ 
  for  $a\ b\ c :: \text{nat}$ 
  apply safe
  apply (erule (1) xtr4 [OF div-le-mono div-mult-self-is-m])
  apply (erule th2)
  done

lemmas td-gal-lt = td-gal [simplified not-less [symmetric], simplified]

lemma div-mult-le:  $a \text{ div } b * b \leq a$ 
  for  $a\ b :: \text{nat}$ 
  by (fact dtle)

lemmas sdl = split-div-lemma [THEN iffD1, symmetric]

lemma given-quot:  $f > 0 \implies (f * l + (f - 1)) \text{ div } f = l$ 
  for  $f\ l :: \text{nat}$ 
  by (rule sdl, assumption) (simp (no-asm))

```

```

lemma given-quot-alt:  $f > 0 \implies (l * f + f - \text{Suc } 0) \text{ div } f = l$ 
  for  $f\ l :: \text{nat}$ 
  apply (frule given-quot)
  apply (rule trans)
  prefer 2
  apply (erule asm-rl)
  apply (rule-tac f = λn. n div f in arg-cong)
  apply (simp add : ac-simps)
  done

```

```

lemma diff-mod-le:  $a < d \implies b \text{ dvd } d \implies a - a \text{ mod } b \leq d - b$ 
  for  $a\ b\ d :: \text{nat}$ 
  apply (unfold dvd-def)
  apply clarify
  apply (case-tac k)
  apply clarsimp
  apply clarify
  apply (cases b > 0)
  apply (drule mult.commute [THEN xtr1])
  apply (frule (1) td-gal-lt [THEN iffD1])
  apply (clarsimp simp: le-simps)
  apply (rule minus-mod-eq-mult-div [symmetric, THEN [2] xtr4])
  apply (rule mult-mono)
  apply auto
  done

```

```

lemma less-le-mult':  $w * c < b * c \implies 0 \leq c \implies (w + 1) * c \leq b * c$ 
  for  $b\ c\ w :: \text{int}$ 
  apply (rule mult-right-mono)
  apply (rule zless-imp-add1-zle)
  apply (erule (1) mult-right-less-imp-less)
  apply assumption
  done

```

```

lemma less-le-mult:  $w * c < b * c \implies 0 \leq c \implies w * c + c \leq b * c$ 
  for  $b\ c\ w :: \text{int}$ 
  using less-le-mult' [of w c b] by (simp add: algebra-simps)

```

```

lemmas less-le-mult-minus = iffD2 [OF le-diff-eq less-le-mult,
  simplified left-diff-distrib]

```

```

lemma gen-minus:  $0 < n \implies f\ n = f\ (\text{Suc } (n - 1))$ 
  by auto

```

```

lemma mpl-lem:  $j \leq i \implies k < j \implies i - j + k < i$ 
  for  $i\ j\ k :: \text{nat}$ 
  by arith

```

```

lemma nonneg-mod-div:  $0 \leq a \implies 0 \leq b \implies 0 \leq (a \bmod b) \wedge 0 \leq a \operatorname{div} b$ 
  for  $a\ b :: \text{int}$ 
  by (cases  $b = 0$ ) (auto intro: pos-imp-zdiv-nonneg-iff [THEN iffD2])

declare iszero-0 [intro]

lemma min-pm [simp]:  $\min a\ b + (a - b) = a$ 
  for  $a\ b :: \text{nat}$ 
  by arith

lemma min-pm1 [simp]:  $a - b + \min a\ b = a$ 
  for  $a\ b :: \text{nat}$ 
  by arith

lemma rev-min-pm [simp]:  $\min b\ a + (a - b) = a$ 
  for  $a\ b :: \text{nat}$ 
  by arith

lemma rev-min-pm1 [simp]:  $a - b + \min b\ a = a$ 
  for  $a\ b :: \text{nat}$ 
  by arith

lemma min-minus [simp]:  $\min m\ (m - k) = m - k$ 
  for  $m\ k :: \text{nat}$ 
  by arith

lemma min-minus' [simp]:  $\min (m - k)\ m = m - k$ 
  for  $m\ k :: \text{nat}$ 
  by arith

end

```

10 A type of finite bit strings

```

theory Word
imports
  HOL-Library.Type-Length
  HOL-Library.Boolean-Algebra
  Bits-Bit
  Bool-List-Representation
  Misc-Typedef
  Word-Miscellaneous
begin

```

See `Examples/WordExamples.thy` for examples.

10.1 Type definition

```

typedef (overloaded)  $'a\ \text{word} = \{(0 :: \text{int}) ..< 2^{\text{len-of } \text{TYPE}('a :: \text{len0})}\}$ 

```

morphisms *uint Abs-word* **by** *auto*

lemma *uint-nonnegative*: $0 \leq \text{uint } w$
using *word.uint [of w]* **by** *simp*

lemma *uint-bounded*: $\text{uint } w < 2^{\text{len-of TYPE('a)}}$
for $w :: 'a::\text{len0 word}$
using *word.uint [of w]* **by** *simp*

lemma *uint-idem*: $\text{uint } w \bmod 2^{\text{len-of TYPE('a)}} = \text{uint } w$
for $w :: 'a::\text{len0 word}$
using *uint-nonnegative uint-bounded* **by** (*rule mod-pos-pos-trivial*)

lemma *word-uint-eq-iff*: $a = b \iff \text{uint } a = \text{uint } b$
by (*simp add: uint-inject*)

lemma *word-uint-eqI*: $\text{uint } a = \text{uint } b \implies a = b$
by (*simp add: word-uint-eq-iff*)

definition *word-of-int* :: $\text{int} \Rightarrow 'a::\text{len0 word}$
 — representation of words using unsigned or signed bins, only difference in these
 is the type class
where *word-of-int* $k = \text{Abs-word } (k \bmod 2^{\text{len-of TYPE('a)}})$

lemma *uint-word-of-int*: $\text{uint } (\text{word-of-int } k :: 'a::\text{len0 word}) = k \bmod 2^{\text{len-of TYPE('a)}}$
by (*auto simp add: word-of-int-def intro: Abs-word-inverse*)

lemma *word-of-int-uint*: $\text{word-of-int } (\text{uint } w) = w$
by (*simp add: word-of-int-def uint-idem uint-inverse*)

lemma *split-word-all*: $(\bigwedge x :: 'a::\text{len0 word}. \text{PROP } P x) \equiv (\bigwedge x. \text{PROP } P (\text{word-of-int } x))$

proof
fix $x :: 'a \text{ word}$
assume $\bigwedge x. \text{PROP } P (\text{word-of-int } x)$
then have $\text{PROP } P (\text{word-of-int } (\text{uint } x))$.
then show $\text{PROP } P x$ **by** (*simp add: word-of-int-uint*)
qed

10.2 Type conversions and casting

definition *sint* :: $'a::\text{len word} \Rightarrow \text{int}$
 — treats the most-significant-bit as a sign bit
where *sint-uint*: $\text{sint } w = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) (\text{uint } w)$

definition *unat* :: $'a::\text{len0 word} \Rightarrow \text{nat}$
where *unat* $w = \text{nat } (\text{uint } w)$

definition *uints* :: *nat* $\Rightarrow$ *int set*
 — the sets of integers representing the words
where *uints* *n* = *range* (*bintrunc* *n*)

definition *sints* :: *nat* $\Rightarrow$ *int set*
where *sints* *n* = *range* (*sbintrunc* (*n* - 1))

lemma *uints-num*: *uints* *n* = {*i*. $0 \leq i \wedge i < 2^n$ }
by (*simp add: uints-def range-bintrunc*)

lemma *sints-num*: *sints* *n* = {*i*. $-(2^{n-1}) \leq i \wedge i < 2^n$ }
by (*simp add: sints-def range-sbintrunc*)

definition *unats* :: *nat* $\Rightarrow$ *nat set*
where *unats* *n* = {*i*. $i < 2^n$ }

definition *norm-sint* :: *nat* $\Rightarrow$ *int* $\Rightarrow$ *int*
where *norm-sint* *n w* = (*w* + 2^{n-1}) *mod* $2^n - 2^{n-1}$

definition *scast* :: '*a*::*len* word $\Rightarrow$ '*b*::*len* word
 — cast a word to a different length
where *scast* *w* = *word-of-int* (*sint* *w*)

definition *ucast* :: '*a*::*len0* word $\Rightarrow$ '*b*::*len0* word
where *ucast* *w* = *word-of-int* (*uint* *w*)

instantiation *word* :: (*len0*) *size*
begin

definition *word-size*: *size* (*w* :: '*a* word) = *len-of TYPE*('a)

instance ..

end

lemma *word-size-gt-0* [*iff*]: $0 < \text{size } w$
for *w* :: '*a*::*len* word
by (*simp add: word-size*)

lemmas *lens-gt-0* = *word-size-gt-0 len-gt-0*

lemma *lens-not-0* [*iff*]:
fixes *w* :: '*a*::*len* word
shows *size* *w* $\neq 0$
and *len-of TYPE*('a) $\neq 0$
by *auto*

definition *source-size* :: ('a::*len0* word $\Rightarrow$ '*b*) $\Rightarrow$ *nat*
 — whether a cast (or other) function is to a longer or shorter length

where $[code\ del]: source\text{-}size\ c = (let\ arb = undefined; x = c\ arb\ in\ size\ arb)$

definition $target\text{-}size :: ('a \Rightarrow 'b::len0\ word) \Rightarrow nat$
where $[code\ del]: target\text{-}size\ c = size\ (c\ undefined)$

definition $is\text{-}up :: ('a::len0\ word \Rightarrow 'b::len0\ word) \Rightarrow bool$
where $is\text{-}up\ c \longleftrightarrow source\text{-}size\ c \leq target\text{-}size\ c$

definition $is\text{-}down :: ('a::len0\ word \Rightarrow 'b::len0\ word) \Rightarrow bool$
where $is\text{-}down\ c \longleftrightarrow target\text{-}size\ c \leq source\text{-}size\ c$

definition $of\text{-}bl :: bool\ list \Rightarrow 'a::len0\ word$
where $of\text{-}bl\ bl = word\text{-}of\text{-}int\ (bl\text{-}to\text{-}bin\ bl)$

definition $to\text{-}bl :: 'a::len0\ word \Rightarrow bool\ list$
where $to\text{-}bl\ w = bin\text{-}to\text{-}bl\ (len\text{-}of\ TYPE('a))\ (uint\ w)$

definition $word\text{-}reverse :: 'a::len0\ word \Rightarrow 'a\ word$
where $word\text{-}reverse\ w = of\text{-}bl\ (rev\ (to\text{-}bl\ w))$

definition $word\text{-}int\text{-}case :: (int \Rightarrow 'b) \Rightarrow 'a::len0\ word \Rightarrow 'b$
where $word\text{-}int\text{-}case\ f\ w = f\ (uint\ w)$

translations

$case\ x\ of\ XCONST\ of\text{-}int\ y \Rightarrow b \Rightarrow CONST\ word\text{-}int\text{-}case\ (\lambda y. b)\ x$
 $case\ x\ of\ (XCONST\ of\text{-}int :: 'a)\ y \Rightarrow b \rightarrow CONST\ word\text{-}int\text{-}case\ (\lambda y. b)\ x$

10.3 Correspondence relation for theorem transfer

definition $cr\text{-}word :: int \Rightarrow 'a::len0\ word \Rightarrow bool$
where $cr\text{-}word = (\lambda x\ y. word\text{-}of\text{-}int\ x = y)$

lemma *Quotient-word:*

$Quotient\ (\lambda x\ y. bintrunc\ (len\text{-}of\ TYPE('a))\ x = bintrunc\ (len\text{-}of\ TYPE('a))\ y)$
 $word\text{-}of\text{-}int\ uint\ (cr\text{-}word :: - \Rightarrow 'a::len0\ word \Rightarrow bool)$

unfolding *Quotient-alt-def cr-word-def*

by $(simp\ add: no\text{-}bintr\text{-}alt1\ word\text{-}of\text{-}int\text{-}uint)\ (simp\ add: word\text{-}of\text{-}int\text{-}def\ Abs\text{-}word\text{-}inject)$

lemma *reflp-word:*

$reflp\ (\lambda x\ y. bintrunc\ (len\text{-}of\ TYPE('a::len0))\ x = bintrunc\ (len\text{-}of\ TYPE('a))\ y)$

by $(simp\ add: reflp\text{-}def)$

setup-lifting *Quotient-word reflp-word*

TODO: The next lemma could be generated automatically.

lemma *uint-transfer [transfer-rule]:*

$(rel\text{-}fun\ pcr\text{-}word\ op =) (bintrunc\ (len\text{-}of\ TYPE('a)))\ (uint :: 'a::len0\ word \Rightarrow int)$

unfolding *rel-fun-def word.pcr-cr-eq cr-word-def*
by (*simp add: no-bintr-alt1 uint-word-of-int*)

10.4 Basic code generation setup

definition *Word* :: *int* $\Rightarrow$ '*a::len0* *word*
where [*code-post*]: *Word* = *word-of-int*

lemma [*code abstype*]: *Word* (*uint w*) = *w*
by (*simp add: Word-def word-of-int-uint*)

declare *uint-word-of-int* [*code abstract*]

instantiation *word* :: (*len0*) *equal*
begin

definition *equal-word* :: '*a word* $\Rightarrow$ '*a word* $\Rightarrow$ *bool*
where *equal-word* *k l* $\longleftrightarrow$ *HOL.equal* (*uint k*) (*uint l*)

instance
by *standard* (*simp add: equal equal-word-def word-uint-eq-iff*)

end

notation *fcomp* (**infixl** $\circ>$ 60)
notation *scomp* (**infixl** $\circ\rightarrow$ 60)

instantiation *word* :: ({*len0*, *typerep*}) *random*
begin

definition
random-word *i* = *Random.range* *i* $\circ\rightarrow$ ($\lambda k.$ *Pair* (
let j = word-of-int (int-of-integer (integer-of-natural k)) :: '*a word*
in (j, $\lambda::unit.$ Code-Evaluation.term-of j)))

instance ..

end

no-notation *fcomp* (**infixl** $\circ>$ 60)
no-notation *scomp* (**infixl** $\circ\rightarrow$ 60)

10.5 Type-definition locale instantiations

lemmas *uint-0* = *uint-nonnegative*
lemmas *uint-lt* = *uint-bounded*
lemmas *uint-mod-same* = *uint-idem*

lemma *td-ext-uint*:
td-ext (*uint* :: '*a word* $\Rightarrow$ *int*) *word-of-int* (*uints* (*len-of* *TYPE*('a::len0)))

```

    ( $\lambda w :: \text{int}. w \bmod 2 \wedge \text{len-of TYPE}('a)$ )
  apply (unfold td-ext-def')
  apply (simp add: uints-num word-of-int-def bintrunc-mod2p)
  apply (simp add: uint-mod-same uint-0 uint-lt
    word.uint-inverse word.Abs-word-inverse int-mod-lem)
done

```

interpretation word-uint:

```

td-ext
uint :: 'a::len0 word  $\Rightarrow$  int
word-of-int
uints (len-of TYPE('a::len0))
 $\lambda w. w \bmod 2 \wedge \text{len-of TYPE}('a::len0)$ 
by (fact td-ext-uint)

```

lemmas td-uint = word-uint.td-thm

lemmas int-word-uint = word-uint.eq-norm

lemma td-ext-ubin:

```

td-ext (uint :: 'a word  $\Rightarrow$  int) word-of-int (uints (len-of TYPE('a::len0)))
(bintrunc (len-of TYPE('a)))
by (unfold no-bintr-alt1) (fact td-ext-uint)

```

interpretation word-ubin:

```

td-ext
uint :: 'a::len0 word  $\Rightarrow$  int
word-of-int
uints (len-of TYPE('a::len0))
bintrunc (len-of TYPE('a::len0))
by (fact td-ext-ubin)

```

10.6 Arithmetic operations

lift-definition word-succ :: 'a::len0 word $\Rightarrow$ 'a word **is** $\lambda x. x + 1$
 by (auto simp add: bintrunc-mod2p intro: mod-add-cong)

lift-definition word-pred :: 'a::len0 word $\Rightarrow$ 'a word **is** $\lambda x. x - 1$
 by (auto simp add: bintrunc-mod2p intro: mod-diff-cong)

instantiation word :: (len0) {neg-numeral, modulo, comm-monoid-mult, comm-ring}
begin

lift-definition zero-word :: 'a word **is** 0 .

lift-definition one-word :: 'a word **is** 1 .

lift-definition plus-word :: 'a word $\Rightarrow$ 'a word **is** op +
 by (auto simp add: bintrunc-mod2p intro: mod-add-cong)

lift-definition *minus-word* :: 'a word $\Rightarrow$ 'a word $\Rightarrow$ 'a word **is** op −
 by (auto simp add: bintrunc-mod2p intro: mod-diff-cong)

lift-definition *uminus-word* :: 'a word $\Rightarrow$ 'a word **is** uminus
 by (auto simp add: bintrunc-mod2p intro: mod-minus-cong)

lift-definition *times-word* :: 'a word $\Rightarrow$ 'a word $\Rightarrow$ 'a word **is** op *
 by (auto simp add: bintrunc-mod2p intro: mod-mult-cong)

definition *word-div-def*: $a \text{ div } b = \text{word-of-int } (\text{uint } a \text{ div uint } b)$

definition *word-mod-def*: $a \text{ mod } b = \text{word-of-int } (\text{uint } a \text{ mod uint } b)$

instance
 by standard (transfer, simp add: algebra-simps)+

end

Legacy theorems:

lemma *word-arith-wis* [code]:
 shows *word-add-def*: $a + b = \text{word-of-int } (\text{uint } a + \text{uint } b)$
 and *word-sub-wi*: $a - b = \text{word-of-int } (\text{uint } a - \text{uint } b)$
 and *word-mult-def*: $a * b = \text{word-of-int } (\text{uint } a * \text{uint } b)$
 and *word-minus-def*: $- a = \text{word-of-int } (- \text{uint } a)$
 and *word-succ-alt*: $\text{word-succ } a = \text{word-of-int } (\text{uint } a + 1)$
 and *word-pred-alt*: $\text{word-pred } a = \text{word-of-int } (\text{uint } a - 1)$
 and *word-0-wi*: $0 = \text{word-of-int } 0$
 and *word-1-wi*: $1 = \text{word-of-int } 1$
unfolding *plus-word-def minus-word-def times-word-def uminus-word-def*
unfolding *word-succ-def word-pred-def zero-word-def one-word-def*
 by simp-all

lemma *wi-homs*:
 shows *wi-hom-add*: $\text{word-of-int } a + \text{word-of-int } b = \text{word-of-int } (a + b)$
 and *wi-hom-sub*: $\text{word-of-int } a - \text{word-of-int } b = \text{word-of-int } (a - b)$
 and *wi-hom-mult*: $\text{word-of-int } a * \text{word-of-int } b = \text{word-of-int } (a * b)$
 and *wi-hom-neg*: $- \text{word-of-int } a = \text{word-of-int } (- a)$
 and *wi-hom-succ*: $\text{word-succ } (\text{word-of-int } a) = \text{word-of-int } (a + 1)$
 and *wi-hom-pred*: $\text{word-pred } (\text{word-of-int } a) = \text{word-of-int } (a - 1)$
 by (transfer, simp)+

lemmas *wi-hom-syms* = *wi-homs* [symmetric]

lemmas *word-of-int-homs* = *wi-homs word-0-wi word-1-wi*

lemmas *word-of-int-hom-syms* = *word-of-int-homs* [symmetric]

instance *word* :: (len) comm-ring-1
proof

```

have *: 0 < len-of TYPE('a) by (rule len-gt-0)
show (0::'a word) ≠ 1
  by transfer (use * in ⟨auto simp add: gr0-conv-Suc⟩)
qed

```

```

lemma word-of-nat: of-nat n = word-of-int (int n)
  by (induct n) (auto simp add : word-of-int-hom-syms)

```

```

lemma word-of-int: of-int = word-of-int
  apply (rule ext)
  apply (case-tac x rule: int-diff-cases)
  apply (simp add: word-of-nat wi-hom-sub)
  done

```

```

definition udvd :: 'a::len word ⇒ 'a::len word ⇒ bool (infixl udvd 50)
  where a udvd b = (∃ n ≥ 0. uint b = n * uint a)

```

10.7 Ordering

```

instantiation word :: (len0) linorder
begin

```

```

definition word-le-def: a ≤ b ⟷ uint a ≤ uint b

```

```

definition word-less-def: a < b ⟷ uint a < uint b

```

```

instance
  by standard (auto simp: word-less-def word-le-def)

```

```

end

```

```

definition word-sle :: 'a::len word ⇒ 'a word ⇒ bool ((-/ <=s -) [50, 51] 50)
  where a <=s b ⟷ sint a ≤ sint b

```

```

definition word-sless :: 'a::len word ⇒ 'a word ⇒ bool ((-/ <s -) [50, 51] 50)
  where x <s y ⟷ x <=s y ∧ x ≠ y

```

10.8 Bit-wise operations

```

instantiation word :: (len0) bits
begin

```

```

lift-definition bitNOT-word :: 'a word ⇒ 'a word is bitNOT
  by (metis bin-trunc-not)

```

```

lift-definition bitAND-word :: 'a word ⇒ 'a word ⇒ 'a word is bitAND
  by (metis bin-trunc-and)

```

```

lift-definition bitOR-word :: 'a word ⇒ 'a word ⇒ 'a word is bitOR
  by (metis bin-trunc-or)

```

lift-definition *bitXOR-word* :: 'a word $\Rightarrow$ 'a word $\Rightarrow$ 'a word **is** *bitXOR*
by (*metis bin-trunc-xor*)

definition *word-test-bit-def*: *test-bit a* = *bin-nth (uint a)*

definition *word-set-bit-def*: *set-bit a n x* = *word-of-int (bin-sc n x (uint a))*

definition *word-set-bits-def*: (*BITS n. f n*) = *of-bl (bl-of-nth (len-of TYPE('a)) f)*

definition *word-lsb-def*: *lsb a* $\longleftrightarrow$ *bin-last (uint a)*

definition *shiftrl1* :: 'a word $\Rightarrow$ 'a word
where *shiftrl1 w* = *word-of-int (uint w BIT False)*

definition *shiftr1* :: 'a word $\Rightarrow$ 'a word
 — shift right as unsigned or as signed, ie logical or arithmetic
where *shiftr1 w* = *word-of-int (bin-rest (uint w))*

definition *shiftrl-def*: *w << n* = (*shiftrl1* $\hat{\hat{}}$ *n*) *w*

definition *shiftr-def*: *w >> n* = (*shiftr1* $\hat{\hat{}}$ *n*) *w*

instance ..

end

lemma [*code*]:

shows *word-not-def*: *NOT (a::'a::len0 word)* = *word-of-int (NOT (uint a))*
and *word-and-def*: (*a::'a word*) *AND b* = *word-of-int (uint a AND uint b)*
and *word-or-def*: (*a::'a word*) *OR b* = *word-of-int (uint a OR uint b)*
and *word-xor-def*: (*a::'a word*) *XOR b* = *word-of-int (uint a XOR uint b)*
by (*simp-all add: bitNOT-word-def bitAND-word-def bitOR-word-def bitXOR-word-def*)

instantiation *word* :: (*len*) *bitss*

begin

definition *word-msb-def*: *msb a* $\longleftrightarrow$ *bin-sign (sint a) = -1*

instance ..

end

definition *setBit* :: 'a::len0 word $\Rightarrow$ nat $\Rightarrow$ 'a word
where *setBit w n* = *set-bit w n True*

definition *clearBit* :: 'a::len0 word $\Rightarrow$ nat $\Rightarrow$ 'a word
where *clearBit w n* = *set-bit w n False*

10.9 Shift operations

definition *sshiftr1* :: 'a::len word $\Rightarrow$ 'a word
 where *sshiftr1* *w* = word-of-int (bin-rest (sint *w*))

definition *bshiftr1* :: bool $\Rightarrow$ 'a::len word $\Rightarrow$ 'a word
 where *bshiftr1* *b w* = of-bl (b # butlast (to-bl *w*))

definition *sshiftr* :: 'a::len word $\Rightarrow$ nat $\Rightarrow$ 'a word (infixl >>> 55)
 where *w* >>> *n* = (*sshiftr1* ^^ *n*) *w*

definition *mask* :: nat $\Rightarrow$ 'a::len word
 where *mask* *n* = (1 << *n*) - 1

definition *revcast* :: 'a::len0 word $\Rightarrow$ 'b::len0 word
 where *revcast* *w* = of-bl (takefill False (len-of TYPE('b)) (to-bl *w*))

definition *slice1* :: nat $\Rightarrow$ 'a::len0 word $\Rightarrow$ 'b::len0 word
 where *slice1* *n w* = of-bl (takefill False *n* (to-bl *w*))

definition *slice* :: nat $\Rightarrow$ 'a::len0 word $\Rightarrow$ 'b::len0 word
 where *slice* *n w* = *slice1* (size *w* - *n*) *w*

10.10 Rotation

definition *rotater1* :: 'a list $\Rightarrow$ 'a list
 where *rotater1* *ys* =
 (case *ys* of [] $\Rightarrow$ [] | *x* # *xs* $\Rightarrow$ last *ys* # butlast *ys*)

definition *rotater* :: nat $\Rightarrow$ 'a list $\Rightarrow$ 'a list
 where *rotater* *n* = *rotater1* ^^ *n*

definition *word-rotr* :: nat $\Rightarrow$ 'a::len0 word $\Rightarrow$ 'a::len0 word
 where *word-rotr* *n w* = of-bl (*rotater* *n* (to-bl *w*))

definition *word-rotl* :: nat $\Rightarrow$ 'a::len0 word $\Rightarrow$ 'a::len0 word
 where *word-rotl* *n w* = of-bl (rotate *n* (to-bl *w*))

definition *word-roti* :: int $\Rightarrow$ 'a::len0 word $\Rightarrow$ 'a::len0 word
 where *word-roti* *i w* =
 (if *i* $\geq$ 0 then *word-rotr* (nat *i*) *w* else *word-rotl* (nat (- *i*)) *w*)

10.11 Split and cat operations

definition *word-cat* :: 'a::len0 word $\Rightarrow$ 'b::len0 word $\Rightarrow$ 'c::len0 word
 where *word-cat* *a b* = word-of-int (bin-cat (uint *a*) (len-of TYPE('b)) (uint *b*))

definition *word-split* :: 'a::len0 word $\Rightarrow$ 'b::len0 word $\times$ 'c::len0 word
 where *word-split* *a* =
 (case bin-split (len-of TYPE('c)) (uint *a*) of

$$(u, v) \Rightarrow (\text{word-of-int } u, \text{word-of-int } v)$$

definition *word-rcat* :: 'a::len0 word list $\Rightarrow$ 'b::len0 word
where *word-rcat* ws = *word-of-int* (*bin-rcat* (*len-of TYPE*('a)) (*map uint ws*))

definition *word-rsplit* :: 'a::len0 word $\Rightarrow$ 'b::len word list
where *word-rsplit* w = *map word-of-int* (*bin-rsplit* (*len-of TYPE*('b)) (*len-of TYPE*('a), *uint w*))

definition *max-word* :: 'a::len word
 — Largest representable machine integer.
where *max-word* = *word-of-int* ($2 \wedge \text{len-of TYPE}('a) - 1$)

lemmas *of-nth-def* = *word-set-bits-def*

10.12 Theorems about typedefs

lemma *sint-sbintrunc'*: *sint* (*word-of-int* bin :: 'a word) = *sbintrunc* (*len-of TYPE*('a::len) – 1) bin
by (*auto simp: sint-uint word-ubin.eq-norm sbintrunc-bintrunc-lt*)

lemma *uint-sint*: *uint* w = *bintrunc* (*len-of TYPE*('a)) (*sint* w)
for w :: 'a::len word
by (*auto simp: sint-uint bintrunc-sbintrunc-le*)

lemma *bintr-uint*: *len-of TYPE*('a) $\leq n \implies$ *bintrunc* n (*uint* w) = *uint* w
for w :: 'a::len0 word
apply (*subst word-ubin.norm-Rep [symmetric]*)
apply (*simp only: bintrunc-bintrunc-min word-size*)
apply (*simp add: min.absorb2*)
done

lemma *wi-bintr*:
len-of TYPE('a::len0) $\leq n \implies$
word-of-int (*bintrunc* n w) = (*word-of-int* w :: 'a word)
by (*auto simp: word-ubin.norm-eq-iff [symmetric] min.absorb1*)

lemma *td-ext-sbin*:
td-ext (*sint* :: 'a word $\Rightarrow$ int) *word-of-int* (*sints* (*len-of TYPE*('a::len)))
 (*sbintrunc* (*len-of TYPE*('a) – 1))
apply (*unfold td-ext-def' sint-uint*)
apply (*simp add : word-ubin.eq-norm*)
apply (*cases len-of TYPE*('a))
apply (*auto simp add : sints-def*)
apply (*rule sym [THEN trans]*)
apply (*rule word-ubin.Abs-norm*)
apply (*simp only: bintrunc-sbintrunc*)
apply (*drule sym*)
apply *simp*

done

lemma *td-ext-sint*:

td-ext (*sint* :: 'a word $\Rightarrow$ int) *word-of-int* (*sints* (len-of TYPE('a::len)))
 ($\lambda w. (w + 2^{\wedge} (\text{len-of TYPE('a)} - 1)) \bmod 2^{\wedge} \text{len-of TYPE('a)} -$
 $2^{\wedge} (\text{len-of TYPE('a)} - 1)$)
using *td-ext-sbin* [where ?'a = 'a] **by** (*simp add: no-sbintr-alt2*)

interpretation *word-sint*:

td-ext
sint :: 'a::len word $\Rightarrow$ int
word-of-int
sints (len-of TYPE('a::len))
 $\lambda w. (w + 2^{\wedge} (\text{len-of TYPE('a::len)} - 1)) \bmod 2^{\wedge} \text{len-of TYPE('a::len)} -$
 $2^{\wedge} (\text{len-of TYPE('a::len)} - 1)$
by (rule *td-ext-sint*)

interpretation *word-sbin*:

td-ext
sint :: 'a::len word $\Rightarrow$ int
word-of-int
sints (len-of TYPE('a::len))
sbintrunc (len-of TYPE('a::len) - 1)
by (rule *td-ext-sbin*)

lemmas *int-word-sint* = *td-ext-sint* [THEN *td-ext.eq-norm*]

lemmas *td-sint* = *word-sint.td*

lemma *to-bl-def'*: (*to-bl* :: 'a::len0 word $\Rightarrow$ bool list) = *bin-to-bl* (len-of TYPE('a))
 $\circ$ *uint*
by (*auto simp: to-bl-def*)

lemmas *word-reverse-no-def* [*simp*] =
word-reverse-def [of numeral *w*] **for** *w*

lemma *uints-mod*: *uints* *n* = range ($\lambda w. w \bmod 2^{\wedge} n$)
by (fact *uints-def* [unfolded *no-bintr-alt1*])

lemma *word-numeral-alt*: numeral *b* = *word-of-int* (numeral *b*)
by (*induct b, simp-all only: numeral.simps word-of-int-homs*)

declare *word-numeral-alt* [*symmetric, code-abbrev*]

lemma *word-neg-numeral-alt*: $-$ numeral *b* = *word-of-int* ($-$ numeral *b*)
by (*simp only: word-numeral-alt wi-hom-neg*)

declare *word-neg-numeral-alt* [*symmetric, code-abbrev*]

```

lemma word-numeral-transfer [transfer-rule]:
  (rel-fun op = pcr-word) numeral numeral
  (rel-fun op = pcr-word) (– numeral) (– numeral)
  apply (simp-all add: rel-fun-def word.pcr-cr-eq cr-word-def)
  using word-numeral-alt [symmetric] word-neg-numeral-alt [symmetric] by auto

lemma uint-bintrunc [simp]:
  uint (numeral bin :: 'a word) =
    bintrunc (len-of TYPE('a::len0)) (numeral bin)
  unfolding word-numeral-alt by (rule word-ubin.eq-norm)

lemma uint-bintrunc-neg [simp]:
  uint (– numeral bin :: 'a word) = bintrunc (len-of TYPE('a::len0)) (– numeral bin)
  by (simp only: word-neg-numeral-alt word-ubin.eq-norm)

lemma sint-sbintrunc [simp]:
  sint (numeral bin :: 'a word) = sbintrunc (len-of TYPE('a::len) – 1) (numeral bin)
  by (simp only: word-numeral-alt word-sbin.eq-norm)

lemma sint-sbintrunc-neg [simp]:
  sint (– numeral bin :: 'a word) = sbintrunc (len-of TYPE('a::len) – 1) (– numeral bin)
  by (simp only: word-neg-numeral-alt word-sbin.eq-norm)

lemma unat-bintrunc [simp]:
  unat (numeral bin :: 'a::len0 word) = nat (bintrunc (len-of TYPE('a)) (numeral bin))
  by (simp only: unat-def uint-bintrunc)

lemma unat-bintrunc-neg [simp]:
  unat (– numeral bin :: 'a::len0 word) = nat (bintrunc (len-of TYPE('a)) (– numeral bin))
  by (simp only: unat-def uint-bintrunc-neg)

lemma size-0-eq: size w = 0  $\implies$  v = w
  for v w :: 'a::len0 word
  apply (unfold word-size)
  apply (rule word-uint.Rep-eqD)
  apply (rule box-equals)
  defer
  apply (rule word-ubin.norm-Rep)+
  apply simp
  done

lemma uint-ge-0 [iff]: 0  $\leq$  uint x
  for x :: 'a::len0 word

```

```

using word-uint.Rep [of x] by (simp add: uints-num)

lemma uint-lt2p [iff]: uint x < 2 ^ len-of TYPE('a)
  for x :: 'a::len0 word
  using word-uint.Rep [of x] by (simp add: uints-num)

lemma sint-ge: - (2 ^ (len-of TYPE('a) - 1)) ≤ sint x
  for x :: 'a::len word
  using word-sint.Rep [of x] by (simp add: sints-num)

lemma sint-lt: sint x < 2 ^ (len-of TYPE('a) - 1)
  for x :: 'a::len word
  using word-sint.Rep [of x] by (simp add: sints-num)

lemma sign-uint-Pls [simp]: bin-sign (uint x) = 0
  by (simp add: sign-Pls-ge-0)

lemma uint-m2p-neg: uint x - 2 ^ len-of TYPE('a) < 0
  for x :: 'a::len0 word
  by (simp only: diff-less-0-iff-less uint-lt2p)

lemma uint-m2p-not-non-neg: ¬ 0 ≤ uint x - 2 ^ len-of TYPE('a)
  for x :: 'a::len0 word
  by (simp only: not-le uint-m2p-neg)

lemma lt2p-lem: len-of TYPE('a) ≤ n ⇒ uint w < 2 ^ n
  for w :: 'a::len0 word
  by (metis bintr-uint bintrunc-mod2p int-mod-lem zless2p)

lemma uint-le-0-iff [simp]: uint x ≤ 0 ⇔ uint x = 0
  by (fact uint-ge-0 [THEN leD, THEN linorder-antisym-conv1])

lemma uint-nat: uint w = int (unat w)
  by (auto simp: unat-def)

lemma uint-numeral: uint (numeral b :: 'a::len0 word) = numeral b mod 2 ^ len-of
  TYPE('a)
  by (simp only: word-numeral-alt int-word-uint)

lemma uint-neg-numeral: uint (- numeral b :: 'a::len0 word) = - numeral b mod
  2 ^ len-of TYPE('a)
  by (simp only: word-neg-numeral-alt int-word-uint)

lemma unat-numeral: unat (numeral b :: 'a::len0 word) = numeral b mod 2 ^
  len-of TYPE('a)
  apply (unfold unat-def)
  apply (clarsimp simp only: uint-numeral)
  apply (rule nat-mod-distrib [THEN trans])
  apply (rule zero-le-numeral)

```

apply (*simp-all add: nat-power-eq*)
done

lemma *sint-numeral*:
 $\text{sint } (\text{numeral } b :: 'a::\text{len word}) =$
 $(\text{numeral } b +$
 $2^{\wedge} (\text{len-of TYPE('a)} - 1)) \bmod 2^{\wedge} \text{len-of TYPE('a)} -$
 $2^{\wedge} (\text{len-of TYPE('a)} - 1)$
unfolding *word-numeral-alt* **by** (*rule int-word-sint*)

lemma *word-of-int-0* [*simp, code-post*]: $\text{word-of-int } 0 = 0$
unfolding *word-0-wi* ..

lemma *word-of-int-1* [*simp, code-post*]: $\text{word-of-int } 1 = 1$
unfolding *word-1-wi* ..

lemma *word-of-int-neg-1* [*simp*]: $\text{word-of-int } (-1) = -1$
by (*simp add: wi-hom-syms*)

lemma *word-of-int-numeral* [*simp*] : $(\text{word-of-int } (\text{numeral } bin) :: 'a::\text{len0 word})$
 $= \text{numeral } bin$
by (*simp only: word-numeral-alt*)

lemma *word-of-int-neg-numeral* [*simp*]:
 $(\text{word-of-int } (- \text{numeral } bin) :: 'a::\text{len0 word}) = - \text{numeral } bin$
by (*simp only: word-numeral-alt wi-hom-syms*)

lemma *word-int-case-wi*:
 $\text{word-int-case } f (\text{word-of-int } i :: 'b \text{ word}) = f (i \bmod 2^{\wedge} \text{len-of TYPE('b::len0)})$
by (*simp add: word-int-case-def word-uint.eq-norm*)

lemma *word-int-split*:
 $P (\text{word-int-case } f x) =$
 $(\forall i. x = (\text{word-of-int } i :: 'b::\text{len0 word}) \wedge 0 \leq i \wedge i < 2^{\wedge} \text{len-of TYPE('b)})$
 $\longrightarrow P (f i))$
by (*auto simp: word-int-case-def word-uint.eq-norm mod-pos-pos-trivial*)

lemma *word-int-split-asm*:
 $P (\text{word-int-case } f x) =$
 $(\nexists n. x = (\text{word-of-int } n :: 'b::\text{len0 word}) \wedge 0 \leq n \wedge n < 2^{\wedge} \text{len-of TYPE('b::len0)})$
 $\wedge \neg P (f n))$
by (*auto simp: word-int-case-def word-uint.eq-norm mod-pos-pos-trivial*)

lemmas *uint-range'* = *word-uint.Rep* [*unfolded uints-num mem-Collect-eq*]
lemmas *sint-range'* = *word-sint.Rep* [*unfolded One-nat-def sints-num mem-Collect-eq*]

lemma *uint-range-size*: $0 \leq \text{uint } w \wedge \text{uint } w < 2^{\wedge} \text{size } w$
unfolding *word-size* **by** (*rule uint-range'*)

lemma *sint-range-size*: $-(2 \wedge (\text{size } w - \text{Suc } 0)) \leq \text{sint } w \wedge \text{sint } w < 2 \wedge (\text{size } w - \text{Suc } 0)$

unfolding *word-size* **by** (*rule sint-range'*)

lemma *sint-above-size*: $2 \wedge (\text{size } w - 1) \leq x \implies \text{sint } w < x$

for $w :: 'a::\text{len0 word}$

unfolding *word-size* **by** (*rule less-le-trans [OF sint-lt]*)

lemma *sint-below-size*: $x \leq -(2 \wedge (\text{size } w - 1)) \implies x \leq \text{sint } w$

for $w :: 'a::\text{len0 word}$

unfolding *word-size* **by** (*rule order-trans [OF - sint-ge]*)

10.13 Testing bits

lemma *test-bit-eq-iff*: $\text{test-bit } u = \text{test-bit } v \longleftrightarrow u = v$

for $u v :: 'a::\text{len0 word}$

unfolding *word-test-bit-def* **by** (*simp add: bin-nth-eq-iff*)

lemma *test-bit-size [rule-format]*: $w !! n \longrightarrow n < \text{size } w$

for $w :: 'a::\text{len0 word}$

apply (*unfold word-test-bit-def*)

apply (*subst word-ubin.norm-Rep [symmetric]*)

apply (*simp only: nth-bintr word-size*)

apply *fast*

done

lemma *word-eq-iff*: $x = y \longleftrightarrow (\forall n < \text{len-of TYPE}('a). x !! n = y !! n)$

for $x y :: 'a::\text{len0 word}$

unfolding *uint-inject [symmetric] bin-eq-iff word-test-bit-def [symmetric]*

by (*metis test-bit-size [unfolded word-size]*)

lemma *word-eqI*: $(\bigwedge n. n < \text{size } u \longrightarrow u !! n = v !! n) \implies u = v$

for $u :: 'a::\text{len0 word}$

by (*simp add: word-size word-eq-iff*)

lemma *word-eqD*: $u = v \implies u !! x = v !! x$

for $u v :: 'a::\text{len0 word}$

by *simp*

lemma *test-bit-bin'*: $w !! n \longleftrightarrow n < \text{size } w \wedge \text{bin-nth } (\text{uint } w) n$

by (*simp add: word-test-bit-def word-size nth-bintr [symmetric]*)

lemmas *test-bit-bin = test-bit-bin' [unfolded word-size]*

lemma *bin-nth-uint-imp*: $\text{bin-nth } (\text{uint } w) n \implies n < \text{len-of TYPE}('a)$

for $w :: 'a::\text{len0 word}$

apply (*rule nth-bintr [THEN iffD1, THEN conjunct1]*)

apply (*subst word-ubin.norm-Rep*)

apply *assumption*

done

lemma *bin-nth-sint*:

len-of TYPE('a) ≤ n ⇒
bin-nth (sint w) n = bin-nth (sint w) (len-of TYPE('a) - 1)
for *w :: 'a::len word*
apply (*subst word-sbin.norm-Rep [symmetric]*)
apply (*auto simp add: nth-sbintr*)
done

lemma *td-bl*:

type-definition
(to-bl :: 'a::len0 word ⇒ bool list)
of-bl
{bl. length bl = len-of TYPE('a)}
apply (*unfold type-definition-def of-bl-def to-bl-def*)
apply (*simp add: word-ubin.eq-norm*)
apply *safe*
apply (*drule sym*)
apply *simp*
done

interpretation *word-bl*:

type-definition
to-bl :: 'a::len0 word ⇒ bool list
of-bl
{bl. length bl = len-of TYPE('a::len0)}
by (*fact td-bl*)

lemmas *word-bl-Rep' = word-bl.Rep [unfolded mem-Collect-eq, iff]*

lemma *word-size-bl*: *size w = size (to-bl w)*

by (*auto simp: word-size*)

lemma *to-bl-use-of-bl*: *to-bl w = bl ⟷ w = of-bl bl ∧ length bl = length (to-bl w)*

by (*fastforce elim!: word-bl.Abs-inverse [unfolded mem-Collect-eq]*)

lemma *to-bl-word-rev*: *to-bl (word-reverse w) = rev (to-bl w)*

by (*simp add: word-reverse-def word-bl.Abs-inverse*)

lemma *word-rev-rev [simp]* : *word-reverse (word-reverse w) = w*

by (*simp add: word-reverse-def word-bl.Abs-inverse*)

lemma *word-rev-gal*: *word-reverse w = u ⟹ word-reverse u = w*

by (*metis word-rev-rev*)

lemma *word-rev-gal'*: *u = word-reverse w ⟹ w = word-reverse u*

by *simp*

lemma *length-bl-gt-0* [iff]: $0 < \text{length } (\text{to-bl } x)$
for $x :: 'a::\text{len word}$
unfolding *word-bl-Rep'* **by** (rule *len-gt-0*)

lemma *bl-not-Nil* [iff]: $\text{to-bl } x \neq []$
for $x :: 'a::\text{len word}$
by (fact *length-bl-gt-0* [unfolded *length-greater-0-conv*])

lemma *length-bl-neq-0* [iff]: $\text{length } (\text{to-bl } x) \neq 0$
for $x :: 'a::\text{len word}$
by (fact *length-bl-gt-0* [THEN *gr-implies-not0*])

lemma *hd-bl-sign-sint*: $\text{hd } (\text{to-bl } w) = (\text{bin-sign } (\text{sint } w) = -1)$
apply (*unfold to-bl-def sint-uint*)
apply (rule *trans* [*OF* - *bl-sbin-sign*])
apply *simp*
done

lemma *of-bl-drop'*:
 $\text{lend} = \text{length bl} - \text{len-of TYPE}('a::\text{len0}) \implies$
 $\text{of-bl } (\text{drop lend bl}) = (\text{of-bl bl} :: 'a \text{ word})$
by (*auto simp: of-bl-def trunc-bl2bin* [symmetric])

lemma *test-bit-of-bl*:
 $(\text{of-bl bl} :: 'a::\text{len0 word}) !! n = (\text{rev bl} ! n \wedge n < \text{len-of TYPE}('a) \wedge n < \text{length bl})$
by (*auto simp add: of-bl-def word-test-bit-def word-size word-ubin.eq-norm nth-bintr bin-nth-of-bl*)

lemma *no-of-bl*: $(\text{numeral bin} :: 'a::\text{len0 word}) = \text{of-bl } (\text{bin-to-bl } (\text{len-of TYPE}('a)) (\text{numeral bin}))$
by (*simp add: of-bl-def*)

lemma *uint-bl*: $\text{to-bl } w = \text{bin-to-bl } (\text{size } w) (\text{uint } w)$
by (*auto simp: word-size to-bl-def*)

lemma *to-bl-bin*: $\text{bl-to-bin } (\text{to-bl } w) = \text{uint } w$
by (*simp add: uint-bl word-size*)

lemma *to-bl-of-bin*: $\text{to-bl } (\text{word-of-int bin} :: 'a::\text{len0 word}) = \text{bin-to-bl } (\text{len-of TYPE}('a)) \text{ bin}$
by (*auto simp: uint-bl word-ubin.eq-norm word-size*)

lemma *to-bl-numeral* [simp]:
 $\text{to-bl } (\text{numeral bin} :: 'a::\text{len0 word}) =$
 $\text{bin-to-bl } (\text{len-of TYPE}('a)) (\text{numeral bin})$
unfolding *word-numeral-alt* **by** (rule *to-bl-of-bin*)

```

lemma to-bl-neg-numeral [simp]:
  to-bl (− numeral bin::'a::len0 word) =
    bin-to-bl (len-of TYPE('a)) (− numeral bin)
unfolding word-neg-numeral-alt by (rule to-bl-of-bin)

lemma to-bl-to-bin [simp] : bl-to-bin (to-bl w) = uint w
by (simp add: uint-bl word-size)

lemma uint-bl-bin: bl-to-bin (bin-to-bl (len-of TYPE('a)) (uint x)) = uint x
for x :: 'a::len0 word
by (rule trans [OF bin-bl-bin word-ubin.norm-Rep])

lemma uints-unats: uints n = int ‘ unats n
apply (unfold unats-def uints-num)
apply safe
apply (rule-tac image-eqI)
apply (erule-tac nat-0-le [symmetric])
apply auto
apply (erule-tac nat-less-iff [THEN iffD2])
apply (rule-tac [2] zless-nat-eq-int-zless [THEN iffD1])
apply (auto simp: nat-power-eq)
done

lemma unats-uints: unats n = nat ‘ uints n
by (auto simp: uints-unats image-iff)

lemmas bintr-num =
  word-ubin.norm-eq-iff [of numeral a numeral b, symmetric, folded word-numeral-alt]
for a b
lemmas sbintr-num =
  word-sbin.norm-eq-iff [of numeral a numeral b, symmetric, folded word-numeral-alt]
for a b

lemma num-of-bintr':
  bintrunc (len-of TYPE('a::len0)) (numeral a) = (numeral b)  $\implies$ 
    numeral a = (numeral b :: 'a word)
unfolding bintr-num by (erule subst, simp)

lemma num-of-sbintr':
  sbintrunc (len-of TYPE('a::len) − 1) (numeral a) = (numeral b)  $\implies$ 
    numeral a = (numeral b :: 'a word)
unfolding sbintr-num by (erule subst, simp)

lemma num-abs-bintr:
  (numeral x :: 'a word) =
    word-of-int (bintrunc (len-of TYPE('a::len0)) (numeral x))
by (simp only: word-ubin.Abs-norm word-numeral-alt)

```

lemma *num-abs-sbintr*:

(*numeral* *x* :: 'a word) =
 word-of-int (sbintrunc (len-of TYPE('a::len) - 1) (*numeral* *x*))
 by (simp only: word-sbin.Abs-norm word-numeral-alt)

lemma *ucast-id*: *ucast* *w* = *w*

by (auto simp: ucast-def)

lemma *scast-id*: *scast* *w* = *w*

by (auto simp: scast-def)

lemma *ucast-bl*: *ucast* *w* = of-bl (to-bl *w*)

by (auto simp: ucast-def of-bl-def uint-bl word-size)

lemma *nth-ucast*: (*ucast* *w*::'a::len0 word) !! *n* = (*w* !! *n* ∧ *n* < len-of TYPE('a))

by (simp add: ucast-def test-bit-bin word-ubin.eq-norm nth-bintr word-size)
 (fast elim!: bin-nth-uint-imp)

lemma *ucast-bintr* [simp]:

ucast (*numeral* *w* :: 'a::len0 word) =
 word-of-int (bintrunc (len-of TYPE('a)) (*numeral* *w*))
 by (simp add: ucast-def)

lemma *scast-sbintr* [simp]:

scast (*numeral* *w* :: 'a::len word) =
 word-of-int (sbintrunc (len-of TYPE('a) - Suc 0) (*numeral* *w*))
 by (simp add: scast-def)

lemma *source-size*: *source-size* (*c*::'a::len0 word ⇒ -) = len-of TYPE('a)

unfolding *source-size-def* *word-size* *Let-def* ..

lemma *target-size*: *target-size* (*c*::- ⇒ 'b::len0 word) = len-of TYPE('b)

unfolding *target-size-def* *word-size* *Let-def* ..

lemma *is-down*: *is-down* *c* ⇔ len-of TYPE('b) ≤ len-of TYPE('a)

for *c* :: 'a::len0 word ⇒ 'b::len0 word

by (simp only: is-down-def *source-size* *target-size*)

lemma *is-up*: *is-up* *c* ⇔ len-of TYPE('a) ≤ len-of TYPE('b)

for *c* :: 'a::len0 word ⇒ 'b::len0 word

by (simp only: is-up-def *source-size* *target-size*)

lemmas *is-up-down* = *trans* [*OF is-up is-down* [*symmetric*]]

lemma *down-cast-same* [*OF refl*]: *uc* = *ucast* $\implies$ *is-down* *uc* $\implies$ *uc* = *scast*
apply (*unfold is-down*)
apply *safe*
apply (*rule ext*)
apply (*unfold ucast-def scast-def uint-sint*)
apply (*rule word-ubin.norm-eq-iff* [*THEN iffD1*])
apply *simp*
done

lemma *word-rev-tf*:
to-bl (*of-bl* *bl*::'*a*::*len0* *word*) =
rev (*takefill* *False* (*len-of TYPE*('a')) (*rev bl*))
by (*auto simp: of-bl-def uint-bl bl-bin-bl-rtf word-ubin.eq-norm word-size*)

lemma *word-rep-drop*:
to-bl (*of-bl* *bl*::'*a*::*len0* *word*) =
replicate (*len-of TYPE*('a') - *length bl*) *False* @
drop (*length bl* - *len-of TYPE*('a')) *bl*
by (*simp add: word-rev-tf takefill-alt rev-take*)

lemma *to-bl-ucast*:
to-bl (*ucast* (*w*::'*b*::*len0* *word*) ::'*a*::*len0* *word*) =
replicate (*len-of TYPE*('a') - *len-of TYPE*('b')) *False* @
drop (*len-of TYPE*('b') - *len-of TYPE*('a')) (*to-bl w*)
apply (*unfold ucast-bl*)
apply (*rule trans*)
apply (*rule word-rep-drop*)
apply *simp*
done

lemma *ucast-up-app* [*OF refl*]:
uc = *ucast* $\implies$ *source-size* *uc* + *n* = *target-size* *uc* $\implies$
to-bl (*uc w*) = *replicate* *n* *False* @ (*to-bl w*)
by (*auto simp add : source-size target-size to-bl-ucast*)

lemma *ucast-down-drop* [*OF refl*]:
uc = *ucast* $\implies$ *source-size* *uc* = *target-size* *uc* + *n* $\implies$
to-bl (*uc w*) = *drop* *n* (*to-bl w*)
by (*auto simp add : source-size target-size to-bl-ucast*)

lemma *scast-down-drop* [*OF refl*]:
sc = *scast* $\implies$ *source-size* *sc* = *target-size* *sc* + *n* $\implies$
to-bl (*sc w*) = *drop* *n* (*to-bl w*)
apply (*subgoal-tac sc = ucast*)
apply *safe*
apply *simp*
apply (*erule ucast-down-drop*)

```

apply (rule down-cast-same [symmetric])
apply (simp add : source-size target-size is-down)
done

```

```

lemma sint-up-scast [OF refl]: sc = scast  $\implies$  is-up sc  $\implies$  sint (sc w) = sint w
apply (unfold is-up)
apply safe
apply (simp add: scast-def word-sbin.eq-norm)
apply (rule box-equals)
  prefer 3
  apply (rule word-sbin.norm-Rep)
  apply (rule sbintrunc-sbintrunc-l)
defer
  apply (subst word-sbin.norm-Rep)
  apply (rule refl)
apply simp
done

```

```

lemma uint-up-ucast [OF refl]: uc = ucast  $\implies$  is-up uc  $\implies$  uint (uc w) = uint
w
apply (unfold is-up)
apply safe
apply (rule bin-eqI)
apply (fold word-test-bit-def)
apply (auto simp add: nth-ucast)
apply (auto simp add: test-bit-bin)
done

```

```

lemma ucast-up-ucast [OF refl]: uc = ucast  $\implies$  is-up uc  $\implies$  ucast (uc w) = ucast
w
apply (simp (no-asm) add: ucast-def)
apply (clarsimp simp add: uint-up-ucast)
done

```

```

lemma scast-up-scast [OF refl]: sc = scast  $\implies$  is-up sc  $\implies$  scast (sc w) = scast
w
apply (simp (no-asm) add: scast-def)
apply (clarsimp simp add: sint-up-scast)
done

```

```

lemma ucast-of-bl-up [OF refl]: w = of-bl bl  $\implies$  size bl  $\leq$  size w  $\implies$  ucast w =
of-bl bl
  by (auto simp add : nth-ucast word-size test-bit-of-bl intro!: word-eqI)

```

```

lemmas ucast-up-ucast-id = trans [OF ucast-up-ucast ucast-id]
lemmas scast-up-scast-id = trans [OF scast-up-scast scast-id]

```

```

lemmas isduu = is-up-down [where c = ucast, THEN iffD2]
lemmas isdus = is-up-down [where c = scast, THEN iffD2]

```

lemmas *ucast-down-ucast-id* = *isduu* [*THEN* *ucast-up-ucast-id*]
lemmas *scast-down-scast-id* = *isdus* [*THEN* *ucast-up-ucast-id*]

lemma *up-ucast-surj*:
is-up (*ucast* :: 'b::len0 word $\Rightarrow$ 'a::len0 word) $\implies$
surj (*ucast* :: 'a word $\Rightarrow$ 'b word)
by (*rule surjI*) (*erule ucast-up-ucast-id*)

lemma *up-scast-surj*:
is-up (*scast* :: 'b::len word $\Rightarrow$ 'a::len word) $\implies$
surj (*scast* :: 'a word $\Rightarrow$ 'b word)
by (*rule surjI*) (*erule scast-up-scast-id*)

lemma *down-scast-inj*:
is-down (*scast* :: 'b::len word $\Rightarrow$ 'a::len word) $\implies$
inj-on (*ucast* :: 'a word $\Rightarrow$ 'b word) *A*
by (*rule inj-on-inverseI*, *erule scast-down-scast-id*)

lemma *down-ucast-inj*:
is-down (*ucast* :: 'b::len0 word $\Rightarrow$ 'a::len0 word) $\implies$
inj-on (*ucast* :: 'a word $\Rightarrow$ 'b word) *A*
by (*rule inj-on-inverseI*) (*erule ucast-down-ucast-id*)

lemma *of-bl-append-same*: *of-bl* (*X* @ *to-bl* *w*) = *w*
by (*rule word-bl.Rep-eqD*) (*simp add: word-rep-drop*)

lemma *ucast-down-wi* [*OF refl*]: *uc* = *ucast* $\implies$ *is-down* *uc* $\implies$ *uc* (*word-of-int* *x*) = *word-of-int* *x*
apply (*unfold is-down*)
apply (*clarsimp simp add: ucast-def word-ubin.eq-norm*)
apply (*rule word-ubin.norm-eq-iff* [*THEN iffD1*])
apply (*erule bintrunc-bintrunc-ge*)
done

lemma *ucast-down-no* [*OF refl*]: *uc* = *ucast* $\implies$ *is-down* *uc* $\implies$ *uc* (*numeral* *bin*) = *numeral* *bin*
unfolding *word-numeral-alt* **by** *clarify* (*rule ucast-down-wi*)

lemma *ucast-down-bl* [*OF refl*]: *uc* = *ucast* $\implies$ *is-down* *uc* $\implies$ *uc* (*of-bl* *bl*) = *of-bl* *bl*
unfolding *of-bl-def* **by** *clarify* (*erule ucast-down-wi*)

lemmas *slice-def'* = *slice-def* [*unfolded word-size*]
lemmas *test-bit-def'* = *word-test-bit-def* [*THEN fun-cong*]

lemmas *word-log-defs* = *word-and-def* *word-or-def* *word-xor-def* *word-not-def*

10.14 Word Arithmetic

lemma *word-less-alt*: $a < b \longleftrightarrow \text{uint } a < \text{uint } b$
by (*fact word-less-def*)

lemma *signed-linorder*: *class.linorder word-sle word-sless*
by *standard* (*auto simp: word-sle-def word-sless-def*)

interpretation *signed*: *linorder word-sle word-sless*
by (*rule signed-linorder*)

lemma *udvdI*: $0 \leq n \implies \text{uint } b = n * \text{uint } a \implies a \text{ udvd } b$
by (*auto simp: udvd-def*)

lemmas *word-div-no* [*simp*] = *word-div-def* [*of numeral a numeral b*] **for** *a b*
lemmas *word-mod-no* [*simp*] = *word-mod-def* [*of numeral a numeral b*] **for** *a b*
lemmas *word-less-no* [*simp*] = *word-less-def* [*of numeral a numeral b*] **for** *a b*
lemmas *word-le-no* [*simp*] = *word-le-def* [*of numeral a numeral b*] **for** *a b*
lemmas *word-sless-no* [*simp*] = *word-sless-def* [*of numeral a numeral b*] **for** *a b*
lemmas *word-sle-no* [*simp*] = *word-sle-def* [*of numeral a numeral b*] **for** *a b*

lemma *word-m1-wi*: $-1 = \text{word-of-int } (-1)$
by (*simp add: word-neg-numeral-alt* [*of Num.One*])

lemma *word-0-bl* [*simp*]: *of-bl* [] = 0
by (*simp add: of-bl-def*)

lemma *word-1-bl*: *of-bl* [True] = 1
by (*simp add: of-bl-def bl-to-bin-def*)

lemma *uint-eq-0* [*simp*]: *uint* 0 = 0
unfolding *word-0-wi word-ubin.eq-norm* **by** *simp*

lemma *of-bl-0* [*simp*]: *of-bl* (*replicate n False*) = 0
by (*simp add: of-bl-def bl-to-bin-rep-False*)

lemma *to-bl-0* [*simp*]: *to-bl* (0::'a::len0 word) = *replicate* (*len-of TYPE('a)*) *False*
by (*simp add: uint-bl word-size bin-to-bl-zero*)

lemma *uint-0-iff*: $\text{uint } x = 0 \longleftrightarrow x = 0$
by (*simp add: word-uint-eq-iff*)

lemma *unat-0-iff*: $\text{unat } x = 0 \longleftrightarrow x = 0$
by (*auto simp: unat-def nat-eq-iff uint-0-iff*)

lemma *unat-0* [*simp*]: *unat* 0 = 0
by (*auto simp: unat-def*)

lemma *size-0-same'*: $\text{size } w = 0 \implies w = v$
for $v \text{ word} :: 'a::\text{len0}$

```

apply (unfold word-size)
apply (rule box-equals)
  defer
    apply (rule word-uint.Rep-inverse)+
  apply (rule word-ubin.norm-eq-iff [THEN iffD1])
apply simp
done

lemmas size-0-same = size-0-same' [unfolded word-size]

lemmas unat-eq-0 = unat-0-iff
lemmas unat-eq-zero = unat-0-iff

lemma unat-gt-0:  $0 < \text{unat } x \longleftrightarrow x \neq 0$ 
  by (auto simp: unat-0-iff [symmetric])

lemma ucast-0 [simp]:  $\text{ucast } 0 = 0$ 
  by (simp add: ucast-def)

lemma sint-0 [simp]:  $\text{sint } 0 = 0$ 
  by (simp add: sint-uint)

lemma scast-0 [simp]:  $\text{scast } 0 = 0$ 
  by (simp add: scast-def)

lemma sint-n1 [simp]:  $\text{sint } (-1) = -1$ 
  by (simp only: word-m1-wi word-sbin.eq-norm) simp

lemma scast-n1 [simp]:  $\text{scast } (-1) = -1$ 
  by (simp add: scast-def)

lemma uint-1 [simp]:  $\text{uint } (1::'a::\text{len word}) = 1$ 
  by (simp only: word-1-wi word-ubin.eq-norm) (simp add: bintrunc-minus-simps(4))

lemma unat-1 [simp]:  $\text{unat } (1::'a::\text{len word}) = 1$ 
  by (simp add: unat-def)

lemma ucast-1 [simp]:  $\text{ucast } (1::'a::\text{len word}) = 1$ 
  by (simp add: ucast-def)

```

10.15 Transferring goals from words to ints

```

lemma word-ths:
  shows word-succ-p1:  $\text{word-succ } a = a + 1$ 
    and word-pred-m1:  $\text{word-pred } a = a - 1$ 
    and word-pred-succ:  $\text{word-pred } (\text{word-succ } a) = a$ 
    and word-succ-pred:  $\text{word-succ } (\text{word-pred } a) = a$ 
    and word-mult-succ:  $\text{word-succ } a * b = b + a * b$ 
  by (transfer, simp add: algebra-simps)+

```

lemma *uint-cong*: $x = y \implies \text{uint } x = \text{uint } y$
by *simp*

lemma *uint-word-ariths*:
fixes $a\ b :: 'a::\text{len0 word}$
shows $\text{uint } (a + b) = (\text{uint } a + \text{uint } b) \bmod 2^{\text{len-of TYPE('a::len0)}}$
and $\text{uint } (a - b) = (\text{uint } a - \text{uint } b) \bmod 2^{\text{len-of TYPE('a)}}$
and $\text{uint } (a * b) = \text{uint } a * \text{uint } b \bmod 2^{\text{len-of TYPE('a)}}$
and $\text{uint } (-a) = -\text{uint } a \bmod 2^{\text{len-of TYPE('a)}}$
and $\text{uint } (\text{word-succ } a) = (\text{uint } a + 1) \bmod 2^{\text{len-of TYPE('a)}}$
and $\text{uint } (\text{word-pred } a) = (\text{uint } a - 1) \bmod 2^{\text{len-of TYPE('a)}}$
and $\text{uint } (0 :: 'a \text{ word}) = 0 \bmod 2^{\text{len-of TYPE('a)}}$
and $\text{uint } (1 :: 'a \text{ word}) = 1 \bmod 2^{\text{len-of TYPE('a)}}$
by (*simp-all add: word-arith-wis [THEN trans [OF uint-cong int-word-uint]]*)

lemma *uint-word-arith-bintrunc*:
fixes $a\ b :: 'a::\text{len0 word}$
shows $\text{uint } (a + b) = \text{bintrunc } (\text{len-of TYPE('a)}) (\text{uint } a + \text{uint } b)$
and $\text{uint } (a - b) = \text{bintrunc } (\text{len-of TYPE('a)}) (\text{uint } a - \text{uint } b)$
and $\text{uint } (a * b) = \text{bintrunc } (\text{len-of TYPE('a)}) (\text{uint } a * \text{uint } b)$
and $\text{uint } (-a) = \text{bintrunc } (\text{len-of TYPE('a)}) (-\text{uint } a)$
and $\text{uint } (\text{word-succ } a) = \text{bintrunc } (\text{len-of TYPE('a)}) (\text{uint } a + 1)$
and $\text{uint } (\text{word-pred } a) = \text{bintrunc } (\text{len-of TYPE('a)}) (\text{uint } a - 1)$
and $\text{uint } (0 :: 'a \text{ word}) = \text{bintrunc } (\text{len-of TYPE('a)}) 0$
and $\text{uint } (1 :: 'a \text{ word}) = \text{bintrunc } (\text{len-of TYPE('a)}) 1$
by (*simp-all add: uint-word-ariths bintrunc-mod2p*)

lemma *sint-word-ariths*:
fixes $a\ b :: 'a::\text{len word}$
shows $\text{sint } (a + b) = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) (\text{sint } a + \text{sint } b)$
and $\text{sint } (a - b) = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) (\text{sint } a - \text{sint } b)$
and $\text{sint } (a * b) = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) (\text{sint } a * \text{sint } b)$
and $\text{sint } (-a) = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) (-\text{sint } a)$
and $\text{sint } (\text{word-succ } a) = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) (\text{sint } a + 1)$
and $\text{sint } (\text{word-pred } a) = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) (\text{sint } a - 1)$
and $\text{sint } (0 :: 'a \text{ word}) = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) 0$
and $\text{sint } (1 :: 'a \text{ word}) = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) 1$
apply (*simp-all only: word-sbin.inverse-norm [symmetric]*)
apply (*simp-all add: wi-hom-syms*)
apply *transfer* **apply** *simp*
apply *transfer* **apply** *simp*
done

lemmas *uint-div-alt* = *word-div-def* [*THEN trans [OF uint-cong int-word-uint]*]
lemmas *uint-mod-alt* = *word-mod-def* [*THEN trans [OF uint-cong int-word-uint]*]

lemma *word-pred-0-n1*: $\text{word-pred } 0 = \text{word-of-int } (-1)$
unfolding *word-pred-m1* **by** *simp*

lemma *succ-pred-no* [simp]:

word-succ (numeral *w*) = numeral *w* + 1
word-pred (numeral *w*) = numeral *w* - 1
word-succ (- numeral *w*) = - numeral *w* + 1
word-pred (- numeral *w*) = - numeral *w* - 1
by (simp-all add: *word-succ-p1 word-pred-m1*)

lemma *word-sp-01* [simp]:

word-succ (- 1) = 0 ∧ *word-succ* 0 = 1 ∧ *word-pred* 0 = - 1 ∧ *word-pred* 1 = 0
by (simp-all add: *word-succ-p1 word-pred-m1*)

lemma *word-of-int-Ex*: ∃ *y*. *x* = *word-of-int* *y*

by (rule-tac *x=uint x in exI*) simp

10.16 Order on fixed-length words

lemma *word-zero-le* [simp]: 0 ≤ *y*

for *y* :: 'a::len0 word
unfolding *word-le-def* **by** auto

lemma *word-m1-ge* [simp]: *word-pred* 0 ≥ *y*

by (simp only: *word-le-def word-pred-0-n1 word-uint.eq-norm m1mod2k*) auto

lemma *word-n1-ge* [simp]: *y* ≤ -1

for *y* :: 'a::len0 word
by (simp only: *word-le-def word-m1-wi word-uint.eq-norm m1mod2k*) auto

lemmas *word-not-simps* [simp] =

word-zero-le [THEN *leD*] *word-m1-ge* [THEN *leD*] *word-n1-ge* [THEN *leD*]

lemma *word-gt-0*: 0 < *y* ⟷ 0 ≠ *y*

for *y* :: 'a::len0 word
by (simp add: *less-le*)

lemmas *word-gt-0-no* [simp] = *word-gt-0* [of numeral *y*] **for** *y*

lemma *word-sless-alt*: *a* < *s* *b* ⟷ *sint* *a* < *sint* *b*

by (auto simp add: *word-sle-def word-sless-def less-le*)

lemma *word-le-nat-alt*: *a* ≤ *b* ⟷ *unat* *a* ≤ *unat* *b*

unfolding *unat-def word-le-def*
by (rule *nat-le-eq-zle* [symmetric]) simp

lemma *word-less-nat-alt*: *a* < *b* ⟷ *unat* *a* < *unat* *b*

unfolding *unat-def word-less-alt*
by (rule *nat-less-eq-zless* [symmetric]) simp

lemma *wi-less*:

(*word-of-int* $n < (\text{word-of-int } m :: 'a::\text{len0 word})$) =
 ($n \bmod 2 \wedge \text{len-of TYPE('a)} < m \bmod 2 \wedge \text{len-of TYPE('a)}$)
unfolding *word-less-alt* **by** (*simp add: word-uint.eq-norm*)

lemma *wi-le*:

(*word-of-int* $n \leq (\text{word-of-int } m :: 'a::\text{len0 word})$) =
 ($n \bmod 2 \wedge \text{len-of TYPE('a)} \leq m \bmod 2 \wedge \text{len-of TYPE('a)}$)
unfolding *word-le-def* **by** (*simp add: word-uint.eq-norm*)

lemma *udvd-nat-alt*: $a \text{ udvd } b \longleftrightarrow (\exists n \geq 0. \text{unat } b = n * \text{unat } a)$

apply (*unfold udvd-def*)
apply *safe*
apply (*simp add: unat-def nat-mult-distrib*)
apply (*simp add: uint-nat*)
apply (*rule exI*)
apply *safe*
prefer 2
apply (*erule notE*)
apply (*rule refl*)
apply *force*
done

lemma *udvd-iff-dvd*: $x \text{ udvd } y \longleftrightarrow \text{unat } x \text{ dvd unat } y$

unfolding *dvd-def udvd-nat-alt* **by** *force*

lemmas *unat-mono* = *word-less-nat-alt* [*THEN iffD1*]

lemma *unat-minus-one*:

assumes $w \neq 0$
shows $\text{unat } (w - 1) = \text{unat } w - 1$
proof –
have $0 \leq \text{uint } w$ **by** (*fact uint-nonnegative*)
moreover from *assms* **have** $0 \neq \text{uint } w$
by (*simp add: uint-0-iff*)
ultimately have $1 \leq \text{uint } w$
by *arith*
from *uint-lt2p* [*of w*] **have** $\text{uint } w - 1 < 2 \wedge \text{len-of TYPE('a)}$
by *arith*
with $\langle 1 \leq \text{uint } w \rangle$ **have** $(\text{uint } w - 1) \bmod 2 \wedge \text{len-of TYPE('a)} = \text{uint } w - 1$
by (*auto intro: mod-pos-pos-trivial*)
with $\langle 1 \leq \text{uint } w \rangle$ **have** $\text{nat } ((\text{uint } w - 1) \bmod 2 \wedge \text{len-of TYPE('a)}) = \text{nat } (\text{uint } w) - 1$
by *auto*
then show *?thesis*
by (*simp only: unat-def int-word-uint word-arith-wis mod-diff-right-eq*)
qed

lemma *measure-unat*: $p \neq 0 \implies \text{unat } (p - 1) < \text{unat } p$
by (*simp add: unat-minus-one*) (*simp add: unat-0-iff [symmetric]*)

lemmas *uint-add-ge0* [*simp*] = *add-nonneg-nonneg* [*OF uint-ge-0 uint-ge-0*]
lemmas *uint-mult-ge0* [*simp*] = *mult-nonneg-nonneg* [*OF uint-ge-0 uint-ge-0*]

lemma *uint-sub-lt2p* [*simp*]: $\text{uint } x - \text{uint } y < 2^{\text{len-of TYPE('a)}}$
for $x :: 'a::\text{len0 word}$ **and** $y :: 'b::\text{len0 word}$
using *uint-ge-0* [*of y*] *uint-lt2p* [*of x*] **by** *arith*

10.17 Conditions for the addition (etc) of two words to overflow

lemma *uint-add-lem*:
 $(\text{uint } x + \text{uint } y < 2^{\text{len-of TYPE('a)}} =$
 $(\text{uint } (x + y) = \text{uint } x + \text{uint } y)$
for $x y :: 'a::\text{len0 word}$
by (*unfold uint-word-ariths*) (*auto intro!: trans [OF - int-mod-lem]*)

lemma *uint-mult-lem*:
 $(\text{uint } x * \text{uint } y < 2^{\text{len-of TYPE('a)}} =$
 $(\text{uint } (x * y) = \text{uint } x * \text{uint } y)$
for $x y :: 'a::\text{len0 word}$
by (*unfold uint-word-ariths*) (*auto intro!: trans [OF - int-mod-lem]*)

lemma *uint-sub-lem*: $\text{uint } x \geq \text{uint } y \iff \text{uint } (x - y) = \text{uint } x - \text{uint } y$
by (*auto simp: uint-word-ariths intro!: trans [OF - int-mod-lem]*)

lemma *uint-add-le*: $\text{uint } (x + y) \leq \text{uint } x + \text{uint } y$
unfolding *uint-word-ariths* **by** (*metis uint-add-ge0 zmod-le-nonneg-dividend*)

lemma *uint-sub-ge*: $\text{uint } (x - y) \geq \text{uint } x - \text{uint } y$
unfolding *uint-word-ariths* **by** (*metis int-mod-ge uint-sub-lt2p zless2p*)

lemma *mod-add-if-z*:
 $x < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$
for $x y z :: \text{int}$
by (*auto intro: int-mod-eq*)

lemma *uint-plus-if'*:
 $\text{uint } (a + b) =$
 $(\text{if } \text{uint } a + \text{uint } b < 2^{\text{len-of TYPE('a)}} \text{ then } \text{uint } a + \text{uint } b$
 $\text{else } \text{uint } a + \text{uint } b - 2^{\text{len-of TYPE('a)}})$
for $a b :: 'a::\text{len0 word}$
using *mod-add-if-z* [*of uint a - uint b*] **by** (*simp add: uint-word-ariths*)

lemma *mod-sub-if-z*:
 $x < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$

```

  (x - y) mod z = (if y ≤ x then x - y else x - y + z)
for x y z :: int
by (auto intro: int-mod-eq)

```

```

lemma uint-sub-if':
  uint (a - b) =
    (if uint b ≤ uint a then uint a - uint b
     else uint a - uint b + 2 ^ len-of TYPE('a))
for a b :: 'a::len0 word
using mod-sub-if-z [of uint a - uint b] by (simp add: uint-word-ariths)

```

10.18 Definition of *uint-arith*

```

lemma word-of-int-inverse:
  word-of-int r = a ⟹ 0 ≤ r ⟹ r < 2 ^ len-of TYPE('a) ⟹ uint a = r
for a :: 'a::len0 word
apply (erule word-uint.Abs-inverse' [rotated])
apply (simp add: uints-num)
done

```

```

lemma uint-split:
  P (uint x) = (∀ i. word-of-int i = x ∧ 0 ≤ i ∧ i < 2 ^ len-of TYPE('a) ⟶ P i)
for x :: 'a::len0 word
apply (fold word-int-case-def)
apply (auto dest!: word-of-int-inverse simp: int-word-uint mod-pos-pos-trivial
        split: word-int-split)
done

```

```

lemma uint-split-asm:
  P (uint x) = (∃ i. word-of-int i = x ∧ 0 ≤ i ∧ i < 2 ^ len-of TYPE('a) ∧ ¬ P i)
for x :: 'a::len0 word
by (auto dest!: word-of-int-inverse
        simp: int-word-uint mod-pos-pos-trivial
        split: uint-split)

```

```

lemmas uint-splits = uint-split uint-split-asm

```

```

lemmas uint-arith-simps =
  word-le-def word-less-alt
  word-uint.Rep-inject [symmetric]
  uint-sub-if' uint-plus-if'

```

```

lemma power-False-cong: False ⟹ a ^ b = c ^ d
by auto

```

```

ML ⟨
  fun uint-arith-simpset ctxt =

```

```

    ctxt addsimps @{thms uint-arith-simps}
    delsimps @{thms word-uint.Rep-inject}
    |> fold Splitter.add-split @{thms if-split-asm}
    |> fold Simplifier.add-cong @{thms power-False-cong}

fun uint-arith-tacs ctxt =
  let
    fun arith-tac' n t =
      Arith-Data.arith-tac ctxt n t
      handle Cooper.COOPER - => Seq.empty;
  in
    [ clarify-tac ctxt 1,
      full-simp-tac (uint-arith-simpset ctxt) 1,
      ALLGOALS (full-simp-tac
        (put-simpset HOL-ss ctxt
          |> fold Splitter.add-split @{thms uint-splits}
          |> fold Simplifier.add-cong @{thms power-False-cong}))),
      rewrite-goals-tac ctxt @{thms word-size},
      ALLGOALS (fn n => REPEAT (resolve-tac ctxt [allI, impI] n) THEN
        REPEAT (eresolve-tac ctxt [conjE] n) THEN
        REPEAT (dresolve-tac ctxt @{thms word-of-int-inverse} n
          THEN assume-tac ctxt n
          THEN assume-tac ctxt n)),
      TRYALL arith-tac' ]
  end

fun uint-arith-tac ctxt = SELECT-GOAL (EVERY (uint-arith-tacs ctxt))
)

method-setup uint-arith =
  ⟨Scan.succeed (SIMPLE-METHOD' o uint-arith-tac)⟩
  solving word arithmetic via integers and arith

```

10.19 More on overflows and monotonicity

lemma *no-plus-overflow-uint-size*: $x \leq x + y \longleftrightarrow \text{uint } x + \text{uint } y < 2^{\text{size } x}$
 for $x \ y :: 'a::\text{len0 word}$
 unfolding *word-size* by *uint-arith*

lemmas *no-olen-add* = *no-plus-overflow-uint-size* [unfolded *word-size*]

lemma *no-ulen-sub*: $x \geq x - y \longleftrightarrow \text{uint } y \leq \text{uint } x$
 for $x \ y :: 'a::\text{len0 word}$
 by *uint-arith*

lemma *no-olen-add'*: $x \leq y + x \longleftrightarrow \text{uint } y + \text{uint } x < 2^{\text{len-of TYPE('a)}}$
 for $x \ y :: 'a::\text{len0 word}$
 by (*simp add: ac-simps no-olen-add*)

lemmas *olen-add-equiv* = *trans* [*OF no-olen-add no-olen-add'* [*symmetric*]]

lemmas *uint-plus-simple-iff* = *trans* [*OF no-olen-add uint-add-lem*]

lemmas *uint-plus-simple* = *uint-plus-simple-iff* [*THEN iffD1*]

lemmas *uint-minus-simple-iff* = *trans* [*OF no-ulen-sub uint-sub-lem*]

lemmas *uint-minus-simple-alt* = *uint-sub-lem* [*folded word-le-def*]

lemmas *word-sub-le-iff* = *no-ulen-sub* [*folded word-le-def*]

lemmas *word-sub-le* = *word-sub-le-iff* [*THEN iffD2*]

lemma *word-less-sub1*: $x \neq 0 \implies 1 < x \longleftrightarrow 0 < x - 1$

for $x :: 'a::len$ *word*

by *uint-arith*

lemma *word-le-sub1*: $x \neq 0 \implies 1 \leq x \longleftrightarrow 0 \leq x - 1$

for $x :: 'a::len$ *word*

by *uint-arith*

lemma *sub-wrap-lt*: $x < x - z \longleftrightarrow x < z$

for $x z :: 'a::len0$ *word*

by *uint-arith*

lemma *sub-wrap*: $x \leq x - z \longleftrightarrow z = 0 \vee x < z$

for $x z :: 'a::len0$ *word*

by *uint-arith*

lemma *plus-minus-not-NULL-ab*: $x \leq ab - c \implies c \leq ab \implies c \neq 0 \implies x + c \neq 0$

for $x ab c :: 'a::len0$ *word*

by *uint-arith*

lemma *plus-minus-no-overflow-ab*: $x \leq ab - c \implies c \leq ab \implies x \leq x + c$

for $x ab c :: 'a::len0$ *word*

by *uint-arith*

lemma *le-minus'*: $a + c \leq b \implies a \leq a + c \implies c \leq b - a$

for $a b c :: 'a::len0$ *word*

by *uint-arith*

lemma *le-plus'*: $a \leq b \implies c \leq b - a \implies a + c \leq b$

for $a b c :: 'a::len0$ *word*

by *uint-arith*

lemmas *le-plus* = *le-plus'* [*rotated*]

lemmas *le-minus* = *leD* [*THEN thin-rl*, *THEN le-minus'*]

lemma *word-plus-mono-right*: $y \leq z \implies x \leq x + z \implies x + y \leq x + z$

for $x y z :: 'a::len0$ *word*

by *uint-arith*

lemma *word-less-minus-cancel*: $y - x < z - x \implies x \leq z \implies y < z$
for $x\ y\ z :: 'a::len0\ word$
by *uint-arith*

lemma *word-less-minus-mono-left*: $y < z \implies x \leq y \implies y - x < z - x$
for $x\ y\ z :: 'a::len0\ word$
by *uint-arith*

lemma *word-less-minus-mono*: $a < c \implies d < b \implies a - b < a \implies c - d < c$
 $\implies a - b < c - d$
for $a\ b\ c\ d :: 'a::len\ word$
by *uint-arith*

lemma *word-le-minus-cancel*: $y - x \leq z - x \implies x \leq z \implies y \leq z$
for $x\ y\ z :: 'a::len0\ word$
by *uint-arith*

lemma *word-le-minus-mono-left*: $y \leq z \implies x \leq y \implies y - x \leq z - x$
for $x\ y\ z :: 'a::len0\ word$
by *uint-arith*

lemma *word-le-minus-mono*:
 $a \leq c \implies d \leq b \implies a - b \leq a \implies c - d \leq c \implies a - b \leq c - d$
for $a\ b\ c\ d :: 'a::len\ word$
by *uint-arith*

lemma *plus-le-left-cancel-wrap*: $x + y' < x \implies x + y < x \implies x + y' < x + y$
 $\longleftrightarrow y' < y$
for $x\ y\ y' :: 'a::len0\ word$
by *uint-arith*

lemma *plus-le-left-cancel-nowrap*: $x \leq x + y' \implies x \leq x + y \implies x + y' < x + y$
 $y \longleftrightarrow y' < y$
for $x\ y\ y' :: 'a::len0\ word$
by *uint-arith*

lemma *word-plus-mono-right2*: $a \leq a + b \implies c \leq b \implies a \leq a + c$
for $a\ b\ c :: 'a::len0\ word$
by *uint-arith*

lemma *word-less-add-right*: $x < y - z \implies z \leq y \implies x + z < y$
for $x\ y\ z :: 'a::len0\ word$
by *uint-arith*

lemma *word-less-sub-right*: $x < y + z \implies y \leq x \implies x - y < z$
for $x\ y\ z :: 'a::len0\ word$
by *uint-arith*

lemma *word-le-plus-either*: $x \leq y \vee x \leq z \implies y \leq y + z \implies x \leq y + z$
for $x\ y\ z :: 'a::len0\ word$
by *uint-arith*

lemma *word-less-nowrapI*: $x < z - k \implies k \leq z \implies 0 < k \implies x < x + k$
for $x\ z\ k :: 'a::len0\ word$
by *uint-arith*

lemma *inc-le*: $i < m \implies i + 1 \leq m$
for $i\ m :: 'a::len\ word$
by *uint-arith*

lemma *inc-i*: $1 \leq i \implies i < m \implies 1 \leq i + 1 \wedge i + 1 \leq m$
for $i\ m :: 'a::len\ word$
by *uint-arith*

lemma *udvd-incr-lem*:
 $up < uq \implies up = ua + n * uint\ K \implies$
 $uq = ua + n' * uint\ K \implies up + uint\ K \leq uq$
apply *clarsimp*
apply (*drule less-le-mult*)
apply *safe*
done

lemma *udvd-incr'*:
 $p < q \implies uint\ p = ua + n * uint\ K \implies$
 $uint\ q = ua + n' * uint\ K \implies p + K \leq q$
apply (*unfold word-less-alt word-le-def*)
apply (*drule* (2) *udvd-incr-lem*)
apply (*erule uint-add-le [THEN order-trans]*)
done

lemma *udvd-decr'*:
 $p < q \implies uint\ p = ua + n * uint\ K \implies$
 $uint\ q = ua + n' * uint\ K \implies p \leq q - K$
apply (*unfold word-less-alt word-le-def*)
apply (*drule* (2) *udvd-incr-lem*)
apply (*drule le-diff-eq [THEN iffD2]*)
apply (*erule order-trans*)
apply (*rule uint-sub-ge*)
done

lemmas *udvd-incr-lem0* = *udvd-incr-lem* [**where** *ua=0, unfolded add-0-left*]
lemmas *udvd-incr0* = *udvd-incr'* [**where** *ua=0, unfolded add-0-left*]
lemmas *udvd-decr0* = *udvd-decr'* [**where** *ua=0, unfolded add-0-left*]

lemma *udvd-minus-le'*: $xy < k \implies z\ udvd\ xy \implies z\ udvd\ k \implies xy \leq k - z$
apply (*unfold udvd-def*)
apply *clarify*

apply (*erule* (2) *udvd-decr0*)
done

lemma *udvd-incr2-K*:

$p < a + s \implies a \leq a + s \implies K \text{ udvd } s \implies K \text{ udvd } p - a \implies a \leq p \implies$
 $0 < K \implies p \leq p + K \wedge p + K \leq a + s$
supply [[*simplproc del: linordered-ring-less-cancel-factor*]]
apply (*unfold udvd-def*)
apply *clarify*
apply (*simp add: uint-arith-simps split: if-split-asm*)
prefer 2
apply (*insert uint-range' [of s]*)[1]
apply *arith*
apply (*drule add.commute [THEN xtr1]*)
apply (*simp add: diff-less-eq [symmetric]*)
apply (*drule less-le-mult*)
apply *arith*
apply *simp*
done

lemma *word-succ-rbl*: $\text{to-bl } w = \text{bl} \implies \text{to-bl } (\text{word-succ } w) = \text{rev } (\text{rbl-succ } (\text{rev } \text{bl}))$

apply (*unfold word-succ-def*)
apply *clarify*
apply (*simp add: to-bl-of-bin*)
apply (*simp add: to-bl-def rbl-succ*)
done

lemma *word-pred-rbl*: $\text{to-bl } w = \text{bl} \implies \text{to-bl } (\text{word-pred } w) = \text{rev } (\text{rbl-pred } (\text{rev } \text{bl}))$

apply (*unfold word-pred-def*)
apply *clarify*
apply (*simp add: to-bl-of-bin*)
apply (*simp add: to-bl-def rbl-pred*)
done

lemma *word-add-rbl*:

$\text{to-bl } v = \text{vbl} \implies \text{to-bl } w = \text{wbl} \implies$
 $\text{to-bl } (v + w) = \text{rev } (\text{rbl-add } (\text{rev } \text{vbl}) (\text{rev } \text{wbl}))$
apply (*unfold word-add-def*)
apply *clarify*
apply (*simp add: to-bl-of-bin*)
apply (*simp add: to-bl-def rbl-add*)
done

lemma *word-mult-rbl*:

$\text{to-bl } v = \text{vbl} \implies \text{to-bl } w = \text{wbl} \implies$
 $\text{to-bl } (v * w) = \text{rev } (\text{rbl-mult } (\text{rev } \text{vbl}) (\text{rev } \text{wbl}))$

```

apply (unfold word-mult-def)
apply clarify
apply (simp add: to-bl-of-bin)
apply (simp add: to-bl-def rbl-mult)
done

```

lemma *rtb-rbl-ariths*:

```

  rev (to-bl w) = ys  $\implies$  rev (to-bl (word-succ w)) = rbl-succ ys
  rev (to-bl w) = ys  $\implies$  rev (to-bl (word-pred w)) = rbl-pred ys
  rev (to-bl v) = ys  $\implies$  rev (to-bl w) = xs  $\implies$  rev (to-bl (v * w)) = rbl-mult ys
  xs
  rev (to-bl v) = ys  $\implies$  rev (to-bl w) = xs  $\implies$  rev (to-bl (v + w)) = rbl-add ys xs
  by (auto simp: rev-swap [symmetric] word-succ-rbl word-pred-rbl word-mult-rbl
word-add-rbl)

```

10.20 Arithmetic type class instantiations

lemmas *word-le-0-iff* [simp] =
word-zero-le [THEN leD, THEN linorder-antisym-conv1]

lemma *word-of-int-nat*: $0 \leq x \implies \text{word-of-int } x = \text{of-nat } (\text{nat } x)$
by (simp add: word-of-int)

lemma *iszero-word-no* [simp]:
iszero (numeral bin :: 'a::len word) =
iszero (bintrunc (len-of TYPE('a)) (numeral bin))
using word-ubin.norm-eq-iff [where 'a='a, of numeral bin 0]
by (simp add: iszero-def [symmetric])

Use *iszero* to simplify equalities between word numerals.

lemmas *word-eq-numeral-iff-iszero* [simp] =
eq-numeral-iff-iszero [where 'a='a::len word]

10.21 Word and nat

lemma *td-ext-unat* [OF refl]:
 $n = \text{len-of TYPE('a::len)} \implies$
 $\text{td-ext } (\text{unat} :: 'a \text{ word} \Rightarrow \text{nat}) \text{ of-nat } (\text{unats } n) (\lambda i. i \bmod 2 \wedge n)$
apply (unfold td-ext-def' unat-def word-of-nat unats-uints)
apply (auto intro!: imageI simp add : word-of-int-hom-syms)
apply (erule word-uint.Abs-inverse [THEN arg-cong])
apply (simp add: int-word-uint nat-mod-distrib nat-power-eq)
done

lemmas *unat-of-nat* = *td-ext-unat* [THEN td-ext.eq-norm]

interpretation *word-unat*:
td-ext

```

    unat::'a::len word  $\Rightarrow$  nat
  of-nat
  unats (len-of TYPE('a::len))
   $\lambda i. i \bmod 2 \wedge \text{len-of TYPE('a::len)}$ 
  by (rule td-ext-unat)

```

lemmas *td-unat* = *word-unat.td-thm*

lemmas *unat-lt2p* [iff] = *word-unat.Rep* [unfolded *unats-def* *mem-Collect-eq*]

lemma *unat-le*: $y \leq \text{unat } z \implies y \in \text{unats } (\text{len-of TYPE('a)})$
 for $z :: 'a::\text{len word}$
 apply (unfold *unats-def*)
 apply *clarsimp*
 apply (rule *xtrans*, rule *unat-lt2p*, *assumption*)
 done

lemma *word-nchotomy*: $\forall w :: 'a::\text{len word}. \exists n. w = \text{of-nat } n \wedge n < 2 \wedge \text{len-of TYPE('a)}$
 apply (rule *allI*)
 apply (rule *word-unat.Abs-cases*)
 apply (unfold *unats-def*)
 apply *auto*
 done

lemma *of-nat-eq*: $\text{of-nat } n = w \longleftrightarrow (\exists q. n = \text{unat } w + q * 2 \wedge \text{len-of TYPE('a)})$
 for $w :: 'a::\text{len word}$
 apply (rule *trans*)
 apply (rule *word-unat.inverse-norm*)
 apply (rule *iffI*)
 apply (rule *mod-eqD*)
 apply *simp*
 apply *clarsimp*
 done

lemma *of-nat-eq-size*: $\text{of-nat } n = w \longleftrightarrow (\exists q. n = \text{unat } w + q * 2 \wedge \text{size } w)$
 unfolding *word-size* by (rule *of-nat-eq*)

lemma *of-nat-0*: $\text{of-nat } m = (0::'a::\text{len word}) \longleftrightarrow (\exists q. m = q * 2 \wedge \text{len-of TYPE('a)})$
 by (*simp add: of-nat-eq*)

lemma *of-nat-2p* [*simp*]: $\text{of-nat } (2 \wedge \text{len-of TYPE('a)}) = (0::'a::\text{len word})$
 by (*fact mult-1* [*symmetric*, *THEN iffD2* [*OF of-nat-0 exI*]])

lemma *of-nat-gt-0*: $\text{of-nat } k \neq 0 \implies 0 < k$
 by (*cases k*) *auto*

lemma *of-nat-neq-0*: $0 < k \implies k < 2 \wedge \text{len-of TYPE('a::len)} \implies \text{of-nat } k \neq (0$

```

:: 'a word)
  by (auto simp add : of-nat-0)

lemma Abs-fnat-hom-add: of-nat a + of-nat b = of-nat (a + b)
  by simp

lemma Abs-fnat-hom-mult: of-nat a * of-nat b = (of-nat (a * b) :: 'a::len word)
  by (simp add: word-of-nat wi-hom-mult)

lemma Abs-fnat-hom-Suc: word-succ (of-nat a) = of-nat (Suc a)
  by (simp add: word-of-nat wi-hom-succ ac-simps)

lemma Abs-fnat-hom-0: (0::'a::len word) = of-nat 0
  by simp

lemma Abs-fnat-hom-1: (1::'a::len word) = of-nat (Suc 0)
  by simp

lemmas Abs-fnat-homs =
  Abs-fnat-hom-add Abs-fnat-hom-mult Abs-fnat-hom-Suc
  Abs-fnat-hom-0 Abs-fnat-hom-1

lemma word-arith-nat-add: a + b = of-nat (unat a + unat b)
  by simp

lemma word-arith-nat-mult: a * b = of-nat (unat a * unat b)
  by simp

lemma word-arith-nat-Suc: word-succ a = of-nat (Suc (unat a))
  by (subst Abs-fnat-hom-Suc [symmetric]) simp

lemma word-arith-nat-div: a div b = of-nat (unat a div unat b)
  by (simp add: word-div-def word-of-nat zdiv-int uint-nat)

lemma word-arith-nat-mod: a mod b = of-nat (unat a mod unat b)
  by (simp add: word-mod-def word-of-nat zmod-int uint-nat)

lemmas word-arith-nat-defs =
  word-arith-nat-add word-arith-nat-mult
  word-arith-nat-Suc Abs-fnat-hom-0
  Abs-fnat-hom-1 word-arith-nat-div
  word-arith-nat-mod

lemma unat-cong: x = y  $\implies$  unat x = unat y
  by simp

lemmas unat-word-ariths = word-arith-nat-defs
  [THEN trans [OF unat-cong unat-of-nat]]

```

lemmas *word-sub-less-iff* = *word-sub-le-iff*
 [unfolding linorder-not-less [symmetric] Not-eq-iff]

lemma *unat-add-lem*:
 $\text{unat } x + \text{unat } y < 2^{\text{len-of TYPE('a)}} \longleftrightarrow \text{unat } (x + y) = \text{unat } x + \text{unat } y$
for $x \ y :: 'a::\text{len word}$
by (*auto simp: unat-word-ariths intro!: trans [OF - nat-mod-lem]*)

lemma *unat-mult-lem*:
 $\text{unat } x * \text{unat } y < 2^{\text{len-of TYPE('a)}} \longleftrightarrow \text{unat } (x * y) = \text{unat } x * \text{unat } y$
for $x \ y :: 'a::\text{len word}$
by (*auto simp: unat-word-ariths intro!: trans [OF - nat-mod-lem]*)

lemmas *unat-plus-if'* =
trans [OF unat-word-ariths(1) mod-nat-add, simplified]

lemma *le-no-overflow*: $x \leq b \implies a \leq a + b \implies x \leq a + b$
for $a \ b \ x :: 'a::\text{len0 word}$
apply (*erule order-trans*)
apply (*erule olen-add-eqv [THEN iffD1]*)
done

lemmas *un-ui-le* =
trans [OF word-le-nat-alt [symmetric] word-le-def]

lemma *unat-sub-if-size*:
 $\text{unat } (x - y) =$
 (*if* $\text{unat } y \leq \text{unat } x$
 then $\text{unat } x - \text{unat } y$
 else $\text{unat } x + 2^{\text{size } x} - \text{unat } y$)
apply (*unfold word-size*)
apply (*simp add: un-ui-le*)
apply (*auto simp add: unat-def uint-sub-if'*)
apply (*rule nat-diff-distrib*)
prefer 3
apply (*simp add: algebra-simps*)
apply (*rule nat-diff-distrib [THEN trans]*)
prefer 3
apply (*subst nat-add-distrib*)
prefer 3
apply (*simp add: nat-power-eq*)
apply *auto*
apply *uint-arith*
done

lemmas *unat-sub-if'* = *unat-sub-if-size* [unfolding word-size]

lemma *unat-div*: $\text{unat } (x \text{ div } y) = \text{unat } x \text{ div } \text{unat } y$
for $x \ y :: 'a::\text{len word}$

```

apply (simp add : unat-word-ariths)
apply (rule unat-lt2p [THEN xtr7, THEN nat-mod-eq])
apply (rule div-le-dividend)
done

```

```

lemma unat-mod: unat (x mod y) = unat x mod unat y
  for x y :: 'a::len word
  apply (clarsimp simp add : unat-word-ariths)
  apply (cases unat y)
  prefer 2
  apply (rule unat-lt2p [THEN xtr7, THEN nat-mod-eq])
  apply (rule mod-le-divisor)
  apply auto
done

```

```

lemma uint-div: uint (x div y) = uint x div uint y
  for x y :: 'a::len word
  by (simp add: uint-nat unat-div zdiv-int)

```

```

lemma uint-mod: uint (x mod y) = uint x mod uint y
  for x y :: 'a::len word
  by (simp add: uint-nat unat-mod zmod-int)

```

10.22 Definition of unat-arith tactic

```

lemma unat-split: P (unat x)  $\longleftrightarrow$  ( $\forall n. \text{of-nat } n = x \wedge n < 2^{\text{len-of TYPE('a)}}$ 
 $\longrightarrow P n$ )
  for x :: 'a::len word
  by (auto simp: unat-of-nat)

```

```

lemma unat-split-asm: P (unat x)  $\longleftrightarrow$  ( $\nexists n. \text{of-nat } n = x \wedge n < 2^{\text{len-of TYPE('a)}}$ 
 $\wedge \neg P n$ )
  for x :: 'a::len word
  by (auto simp: unat-of-nat)

```

```

lemmas of-nat-inverse =
  word-unat.Abs-inverse' [rotated, unfolded unats-def, simplified]

```

```

lemmas unat-splits = unat-split unat-split-asm

```

```

lemmas unat-arith-simps =
  word-le-nat-alt word-less-nat-alt
  word-unat.Rep-inject [symmetric]
  unat-sub-if' unat-plus-if' unat-div unat-mod

```

```

ML <
fun unat-arith-simpset ctxt =
  ctxt addsimps @ {thms unat-arith-simps}

```

```

    delsimps @{thms word-unat.Rep-inject}
    |> fold Splitter.add-split @{thms if-split-asm}
    |> fold Simplifier.add-cong @{thms power-False-cong}

fun unat-arith-tacs ctxt =
  let
    fun arith-tac' n t =
      Arith-Data.arith-tac ctxt n t
      handle Cooper.COOPER - => Seq.empty;
  in
    [ clarify-tac ctxt 1,
      full-simp-tac (unat-arith-simpset ctxt) 1,
      ALLGOALS (full-simp-tac
        (put-simpset HOL-ss ctxt
          |> fold Splitter.add-split @{thms unat-splits}
          |> fold Simplifier.add-cong @{thms power-False-cong}))),
      rewrite-goals-tac ctxt @{thms word-size},
      ALLGOALS (fn n => REPEAT (resolve-tac ctxt [allI, impI] n) THEN
        REPEAT (eresolve-tac ctxt [conjE] n) THEN
        REPEAT (dresolve-tac ctxt @{thms of-nat-inverse} n) THEN
        assume-tac ctxt n)),
      TRYALL arith-tac' ]
  end

fun unat-arith-tac ctxt = SELECT-GOAL (EVERY (unat-arith-tacs ctxt))
)

method-setup unat-arith =
  ⟨Scan.succeed (SIMPLE-METHOD' o unat-arith-tac)⟩
  solving word arithmetic via natural numbers and arith

lemma no-plus-overflow-unat-size:  $x \leq x + y \longleftrightarrow \text{unat } x + \text{unat } y < 2^{\text{size } x}$ 
  for  $x \ y :: 'a::\text{len word}$ 
  unfolding word-size by unat-arith

lemmas no-olen-add-nat =
  no-plus-overflow-unat-size [unfolded word-size]

lemmas unat-plus-simple =
  trans [OF no-olen-add-nat unat-add-lem]

lemma word-div-mult:  $0 < y \implies \text{unat } x * \text{unat } y < 2^{\text{len-of TYPE('a)}} \implies x * y \text{ div } y = x$ 
  for  $x \ y :: 'a::\text{len word}$ 
  apply unat-arith
  apply clarsimp
  apply (subst unat-mult-lem [THEN iffD1])
  apply auto
  done

```

```

lemma div-lt':  $i \leq k \text{ div } x \implies \text{unat } i * \text{unat } x < 2^{\text{len-of TYPE('a)}}$ 
  for  $i \ k \ x :: 'a::\text{len word}$ 
  apply unat-arith
  apply clarsimp
  apply (drule mult-le-mono1)
  apply (erule order-le-less-trans)
  apply (rule xtr7 [OF unat-lt2p div-mult-le])
  done

```

```

lemmas div-lt'' = order-less-imp-le [THEN div-lt']

```

```

lemma div-lt-mult:  $i < k \text{ div } x \implies 0 < x \implies i * x < k$ 
  for  $i \ k \ x :: 'a::\text{len word}$ 
  apply (frule div-lt'' [THEN unat-mult-lem [THEN iffD1]])
  apply (simp add: unat-arith-simps)
  apply (drule (1) mult-less-mono1)
  apply (erule order-less-le-trans)
  apply (rule div-mult-le)
  done

```

```

lemma div-le-mult:  $i \leq k \text{ div } x \implies 0 < x \implies i * x \leq k$ 
  for  $i \ k \ x :: 'a::\text{len word}$ 
  apply (frule div-lt' [THEN unat-mult-lem [THEN iffD1]])
  apply (simp add: unat-arith-simps)
  apply (drule mult-le-mono1)
  apply (erule order-trans)
  apply (rule div-mult-le)
  done

```

```

lemma div-lt-uint':  $i \leq k \text{ div } x \implies \text{uint } i * \text{uint } x < 2^{\text{len-of TYPE('a)}}$ 
  for  $i \ k \ x :: 'a::\text{len word}$ 
  apply (unfold uint-nat)
  apply (drule div-lt')
  apply (metis of-nat-less-iff of-nat-mult of-nat-numeral of-nat-power)
  done

```

```

lemmas div-lt-uint'' = order-less-imp-le [THEN div-lt-uint']

```

```

lemma word-le-exists':  $x \leq y \implies \exists z. y = x + z \wedge \text{uint } x + \text{uint } z < 2^{\text{len-of TYPE('a)}}$ 
  for  $x \ y \ z :: 'a::\text{len0 word}$ 
  apply (rule exI)
  apply (rule conjI)
  apply (rule zadd-diff-inverse)
  apply uint-arith
  done

```

```

lemmas plus-minus-not-NULL = order-less-imp-le [THEN plus-minus-not-NULL-ab]

```

```

lemmas plus-minus-no-overflow =
  order-less-imp-le [THEN plus-minus-no-overflow-ab]

lemmas mcs = word-less-minus-cancel word-less-minus-mono-left
  word-le-minus-cancel word-le-minus-mono-left

lemmas word-l-diffs = mcs [where  $y = w + x$ , unfolded add-diff-cancel] for  $w\ x$ 
lemmas word-diff-ls = mcs [where  $z = w + x$ , unfolded add-diff-cancel] for  $w\ x$ 
lemmas word-plus-mcs = word-diff-ls [where  $y = v + x$ , unfolded add-diff-cancel]
for  $v\ x$ 

lemmas le-unat-uo = unat-le [THEN word-unat.Abs-inverse]

lemmas thd = refl [THEN [2] split-div-lemma [THEN iffD2], THEN conjunct1]

lemmas uno-simps [THEN le-unat-uo] = mod-le-divisor div-le-dividend dtl

lemma word-mod-div-equality:  $(n \text{ div } b) * b + (n \text{ mod } b) = n$ 
  for  $n\ b :: 'a::len\ word$ 
  apply (unfold word-less-nat-alt word-arith-nat-defs)
  apply (cut-tac y=unat b in gt-or-eq-0)
  apply (erule disjE)
  apply (simp only: div-mult-mod-eq uno-simps Word.word-unat.Rep-inverse)
  apply simp
  done

lemma word-div-mult-le:  $a \text{ div } b * b \leq a$ 
  for  $a\ b :: 'a::len\ word$ 
  apply (unfold word-le-nat-alt word-arith-nat-defs)
  apply (cut-tac y=unat b in gt-or-eq-0)
  apply (erule disjE)
  apply (simp only: div-mult-le uno-simps Word.word-unat.Rep-inverse)
  apply simp
  done

lemma word-mod-less-divisor:  $0 < n \implies m \text{ mod } n < n$ 
  for  $m\ n :: 'a::len\ word$ 
  apply (simp only: word-less-nat-alt word-arith-nat-defs)
  apply (auto simp: uno-simps)
  done

lemma word-of-int-power-hom:  $\text{word-of-int } a ^ n = (\text{word-of-int } (a ^ n) :: 'a::len\ word)$ 
  by (induct n) (simp-all add: wi-hom-mult [symmetric])

lemma word-arith-power-alt:  $a ^ n = (\text{word-of-int } (\text{uint } a ^ n) :: 'a::len\ word)$ 
  by (simp add : word-of-int-power-hom [symmetric])

```

```

lemma of-bl-length-less:
  length x = k  $\implies$  k < len-of TYPE('a)  $\implies$  (of-bl x :: 'a::len word) < 2 ^ k
apply (unfold of-bl-def word-less-alt word-numeral-alt)
apply safe
apply (simp (no-asm) add: word-of-int-power-hom word-uint.eq-norm
  del: word-of-int-numeral)
apply (simp add: mod-pos-pos-trivial)
apply (subst mod-pos-pos-trivial)
  apply (rule bl-to-bin-ge0)
  apply (rule order-less-trans)
  apply (rule bl-to-bin-lt2p)
apply simp
apply (rule bl-to-bin-lt2p)
done

```

10.23 Cardinality, finiteness of set of words

```

instance word :: (len0) finite
  by standard (simp add: type-definition.univ [OF type-definition-word])

lemma card-word: CARD('a::len0 word) = 2 ^ len-of TYPE('a)
  by (simp add: type-definition.card [OF type-definition-word] nat-power-eq)

lemma card-word-size: card (UNIV :: 'a word set) = (2 ^ size x)
  for x :: 'a::len0 word
  unfolding word-size by (rule card-word)

```

10.24 Bitwise Operations on Words

```

lemmas bin-log-bintrs = bin-trunc-not bin-trunc-xor bin-trunc-and bin-trunc-or

```

```

lemmas wils1 = bin-log-bintrs [THEN word-ubin.norm-eq-iff [THEN iffD1],
  folded word-ubin.eq-norm, THEN eq-reflection]

```

```

lemmas word-log-binary-defs =
  word-and-def word-or-def word-xor-def

```

```

lemma word-wi-log-defs:
  NOT word-of-int a = word-of-int (NOT a)
  word-of-int a AND word-of-int b = word-of-int (a AND b)
  word-of-int a OR word-of-int b = word-of-int (a OR b)
  word-of-int a XOR word-of-int b = word-of-int (a XOR b)
  by (transfer, rule refl)+

```

```

lemma word-no-log-defs [simp]:

```

$NOT (numeral\ a) = word-of-int\ (NOT\ (numeral\ a))$
 $NOT\ (-\ numeral\ a) = word-of-int\ (NOT\ (-\ numeral\ a))$
 $numeral\ a\ AND\ numeral\ b = word-of-int\ (numeral\ a\ AND\ numeral\ b)$
 $numeral\ a\ AND\ -\ numeral\ b = word-of-int\ (numeral\ a\ AND\ -\ numeral\ b)$
 $-\ numeral\ a\ AND\ numeral\ b = word-of-int\ (-\ numeral\ a\ AND\ numeral\ b)$
 $-\ numeral\ a\ AND\ -\ numeral\ b = word-of-int\ (-\ numeral\ a\ AND\ -\ numeral\ b)$
 $numeral\ a\ OR\ numeral\ b = word-of-int\ (numeral\ a\ OR\ numeral\ b)$
 $numeral\ a\ OR\ -\ numeral\ b = word-of-int\ (numeral\ a\ OR\ -\ numeral\ b)$
 $-\ numeral\ a\ OR\ numeral\ b = word-of-int\ (-\ numeral\ a\ OR\ numeral\ b)$
 $-\ numeral\ a\ OR\ -\ numeral\ b = word-of-int\ (-\ numeral\ a\ OR\ -\ numeral\ b)$
 $numeral\ a\ XOR\ numeral\ b = word-of-int\ (numeral\ a\ XOR\ numeral\ b)$
 $numeral\ a\ XOR\ -\ numeral\ b = word-of-int\ (numeral\ a\ XOR\ -\ numeral\ b)$
 $-\ numeral\ a\ XOR\ numeral\ b = word-of-int\ (-\ numeral\ a\ XOR\ numeral\ b)$
 $-\ numeral\ a\ XOR\ -\ numeral\ b = word-of-int\ (-\ numeral\ a\ XOR\ -\ numeral\ b)$
by (transfer, rule refl)+

Special cases for when one of the arguments equals 1.

lemma *word-bitwise-1-simps* [simp]:

$NOT\ (1::'a::len0\ word) = -2$
 $1\ AND\ numeral\ b = word-of-int\ (1\ AND\ numeral\ b)$
 $1\ AND\ -\ numeral\ b = word-of-int\ (1\ AND\ -\ numeral\ b)$
 $numeral\ a\ AND\ 1 = word-of-int\ (numeral\ a\ AND\ 1)$
 $-\ numeral\ a\ AND\ 1 = word-of-int\ (-\ numeral\ a\ AND\ 1)$
 $1\ OR\ numeral\ b = word-of-int\ (1\ OR\ numeral\ b)$
 $1\ OR\ -\ numeral\ b = word-of-int\ (1\ OR\ -\ numeral\ b)$
 $numeral\ a\ OR\ 1 = word-of-int\ (numeral\ a\ OR\ 1)$
 $-\ numeral\ a\ OR\ 1 = word-of-int\ (-\ numeral\ a\ OR\ 1)$
 $1\ XOR\ numeral\ b = word-of-int\ (1\ XOR\ numeral\ b)$
 $1\ XOR\ -\ numeral\ b = word-of-int\ (1\ XOR\ -\ numeral\ b)$
 $numeral\ a\ XOR\ 1 = word-of-int\ (numeral\ a\ XOR\ 1)$
 $-\ numeral\ a\ XOR\ 1 = word-of-int\ (-\ numeral\ a\ XOR\ 1)$
by (transfer, simp)+

Special cases for when one of the arguments equals -1.

lemma *word-bitwise-m1-simps* [simp]:

$NOT\ (-1::'a::len0\ word) = 0$
 $(-1::'a::len0\ word)\ AND\ x = x$
 $x\ AND\ (-1::'a::len0\ word) = x$
 $(-1::'a::len0\ word)\ OR\ x = -1$
 $x\ OR\ (-1::'a::len0\ word) = -1$
 $(-1::'a::len0\ word)\ XOR\ x = NOT\ x$
 $x\ XOR\ (-1::'a::len0\ word) = NOT\ x$
by (transfer, simp)+

lemma *uint-or*: $uint\ (x\ OR\ y) = uint\ x\ OR\ uint\ y$

by (transfer, simp add: bin-trunc-ao)

lemma *uint-and*: $uint\ (x\ AND\ y) = uint\ x\ AND\ uint\ y$

by (transfer, simp add: bin-trunc-ao)

lemma *test-bit-wi* [*simp*]:
 (word-of-int $x :: 'a::len0$ word) !! $n \longleftrightarrow n < len\text{-of } TYPE('a) \wedge bin\text{-nth } x n$
by (*simp add: word-test-bit-def word-ubin.eq-norm nth-bintr*)

lemma *word-test-bit-transfer* [*transfer-rule*]:
 (rel-fun pcr-word (rel-fun op = op =))
 ($\lambda x n. n < len\text{-of } TYPE('a) \wedge bin\text{-nth } x n$) (test-bit :: $'a::len0$ word $\Rightarrow$ -)
unfolding rel-fun-def word.pcr-cr-eq cr-word-def **by** *simp*

lemma *word-ops-nth-size*:
 $n < size\ x \implies$
 ($x\ OR\ y$) !! $n = (x\ !!\ n \mid y\ !!\ n) \wedge$
 ($x\ AND\ y$) !! $n = (x\ !!\ n \wedge y\ !!\ n) \wedge$
 ($x\ XOR\ y$) !! $n = (x\ !!\ n \neq y\ !!\ n) \wedge$
 ($NOT\ x$) !! $n = (\neg x\ !!\ n)$
for $x :: 'a::len0$ word
unfolding word-size **by** transfer (*simp add: bin-nth-ops*)

lemma *word-ao-nth*:
 ($x\ OR\ y$) !! $n = (x\ !!\ n \mid y\ !!\ n) \wedge$
 ($x\ AND\ y$) !! $n = (x\ !!\ n \wedge y\ !!\ n)$
for $x :: 'a::len0$ word
by transfer (*auto simp add: bin-nth-ops*)

lemma *test-bit-numeral* [*simp*]:
 (numeral $w :: 'a::len0$ word) !! $n \longleftrightarrow$
 $n < len\text{-of } TYPE('a) \wedge bin\text{-nth } (numeral\ w)\ n$
by transfer (*rule refl*)

lemma *test-bit-neg-numeral* [*simp*]:
 ($\neg$ numeral $w :: 'a::len0$ word) !! $n \longleftrightarrow$
 $n < len\text{-of } TYPE('a) \wedge bin\text{-nth } (\neg\ numeral\ w)\ n$
by transfer (*rule refl*)

lemma *test-bit-1* [*simp*]: ($1 :: 'a::len0$ word) !! $n \longleftrightarrow n = 0$
by transfer *auto*

lemma *nth-0* [*simp*]: $\neg (0 :: 'a::len0$ word) !! n
by transfer *simp*

lemma *nth-minus1* [*simp*]: ($-1 :: 'a::len0$ word) !! $n \longleftrightarrow n < len\text{-of } TYPE('a)$
by transfer *simp*

lemmas *bwsimps* =
 wi-hom-add
 word-wi-log-defs

lemma *word-bw-assocs*:

$(x \text{ AND } y) \text{ AND } z = x \text{ AND } y \text{ AND } z$
 $(x \text{ OR } y) \text{ OR } z = x \text{ OR } y \text{ OR } z$
 $(x \text{ XOR } y) \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$
for $x :: 'a::\text{len0 word}$
by (*auto simp: word-eq-iff word-ops-nth-size [unfolded word-size]*)

lemma *word-bw-comms*:

$x \text{ AND } y = y \text{ AND } x$
 $x \text{ OR } y = y \text{ OR } x$
 $x \text{ XOR } y = y \text{ XOR } x$
for $x :: 'a::\text{len0 word}$
by (*auto simp: word-eq-iff word-ops-nth-size [unfolded word-size]*)

lemma *word-bw-lcs*:

$y \text{ AND } x \text{ AND } z = x \text{ AND } y \text{ AND } z$
 $y \text{ OR } x \text{ OR } z = x \text{ OR } y \text{ OR } z$
 $y \text{ XOR } x \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$
for $x :: 'a::\text{len0 word}$
by (*auto simp: word-eq-iff word-ops-nth-size [unfolded word-size]*)

lemma *word-log-esimps [simp]*:

$x \text{ AND } 0 = 0$
 $x \text{ AND } -1 = x$
 $x \text{ OR } 0 = x$
 $x \text{ OR } -1 = -1$
 $x \text{ XOR } 0 = x$
 $x \text{ XOR } -1 = \text{NOT } x$
 $0 \text{ AND } x = 0$
 $-1 \text{ AND } x = x$
 $0 \text{ OR } x = x$
 $-1 \text{ OR } x = -1$
 $0 \text{ XOR } x = x$
 $-1 \text{ XOR } x = \text{NOT } x$
for $x :: 'a::\text{len0 word}$
by (*auto simp: word-eq-iff word-ops-nth-size [unfolded word-size]*)

lemma *word-not-dist*:

$\text{NOT } (x \text{ OR } y) = \text{NOT } x \text{ AND } \text{NOT } y$
 $\text{NOT } (x \text{ AND } y) = \text{NOT } x \text{ OR } \text{NOT } y$
for $x :: 'a::\text{len0 word}$
by (*auto simp: word-eq-iff word-ops-nth-size [unfolded word-size]*)

lemma *word-bw-same*:

$x \text{ AND } x = x$
 $x \text{ OR } x = x$
 $x \text{ XOR } x = 0$
for $x :: 'a::\text{len0 word}$

by (*auto simp: word-eq-iff word-ops-nth-size [unfolded word-size]*)

lemma *word-ao-absorbs* [*simp*]:

$x \text{ AND } (y \text{ OR } x) = x$

$x \text{ OR } y \text{ AND } x = x$

$x \text{ AND } (x \text{ OR } y) = x$

$y \text{ AND } x \text{ OR } x = x$

$(y \text{ OR } x) \text{ AND } x = x$

$x \text{ OR } x \text{ AND } y = x$

$(x \text{ OR } y) \text{ AND } x = x$

$x \text{ AND } y \text{ OR } x = x$

for $x :: 'a::\text{len0 word}$

by (*auto simp: word-eq-iff word-ops-nth-size [unfolded word-size]*)

lemma *word-not-not* [*simp*]: $\text{NOT NOT } x = x$

for $x :: 'a::\text{len0 word}$

by (*auto simp: word-eq-iff word-ops-nth-size [unfolded word-size]*)

lemma *word-ao-dist*: $(x \text{ OR } y) \text{ AND } z = x \text{ AND } z \text{ OR } y \text{ AND } z$

for $x :: 'a::\text{len0 word}$

by (*auto simp: word-eq-iff word-ops-nth-size [unfolded word-size]*)

lemma *word-oa-dist*: $x \text{ AND } y \text{ OR } z = (x \text{ OR } z) \text{ AND } (y \text{ OR } z)$

for $x :: 'a::\text{len0 word}$

by (*auto simp: word-eq-iff word-ops-nth-size [unfolded word-size]*)

lemma *word-add-not* [*simp*]: $x + \text{NOT } x = -1$

for $x :: 'a::\text{len0 word}$

by *transfer (simp add: bin-add-not)*

lemma *word-plus-and-or* [*simp*]: $(x \text{ AND } y) + (x \text{ OR } y) = x + y$

for $x :: 'a::\text{len0 word}$

by *transfer (simp add: plus-and-or)*

lemma *leoa*: $w = x \text{ OR } y \implies y = w \text{ AND } y$

for $x :: 'a::\text{len0 word}$

by *auto*

lemma *leao*: $w' = x' \text{ AND } y' \implies x' = x' \text{ OR } w'$

for $x' :: 'a::\text{len0 word}$

by *auto*

lemma *word-ao-equiv*: $w = w \text{ OR } w' \iff w' = w \text{ AND } w'$

for $w w' :: 'a::\text{len0 word}$

by (*auto intro: leoa leao*)

lemma *le-word-or2*: $x \leq x \text{ OR } y$

for $x y :: 'a::\text{len0 word}$

by (*auto simp: word-le-def uint-or intro: le-int-or*)

```

lemmas le-word-or1 = xtr3 [OF word-bw-comms (2) le-word-or2]
lemmas word-and-le1 = xtr3 [OF word-ao-absorbs (4) [symmetric] le-word-or2]
lemmas word-and-le2 = xtr3 [OF word-ao-absorbs (8) [symmetric] le-word-or2]

lemma bl-word-not: to-bl (NOT w) = map Not (to-bl w)
  unfolding to-bl-def word-log-defs bl-not-bin
  by (simp add: word-ubin.eq-norm)

lemma bl-word-xor: to-bl (v XOR w) = map2 op ≠ (to-bl v) (to-bl w)
  unfolding to-bl-def word-log-defs bl-xor-bin
  by (simp add: word-ubin.eq-norm)

lemma bl-word-or: to-bl (v OR w) = map2 op ∨ (to-bl v) (to-bl w)
  unfolding to-bl-def word-log-defs bl-or-bin
  by (simp add: word-ubin.eq-norm)

lemma bl-word-and: to-bl (v AND w) = map2 op ∧ (to-bl v) (to-bl w)
  unfolding to-bl-def word-log-defs bl-and-bin
  by (simp add: word-ubin.eq-norm)

lemma word-lsb-alt: lsb w = test-bit w 0
  for w :: 'a::len0 word
  by (auto simp: word-test-bit-def word-lsb-def)

lemma word-lsb-1-0 [simp]: lsb (1::'a::len word) ∧ ¬ lsb (0::'b::len0 word)
  unfolding word-lsb-def uint-eq-0 uint-1 by simp

lemma word-lsb-last: lsb w = last (to-bl w)
  for w :: 'a::len word
  apply (unfold word-lsb-def uint-bl bin-to-bl-def)
  apply (rule-tac bin=uint w in bin-exhaust)
  apply (cases size w)
  apply auto
  apply (auto simp add: bin-to-bl-aux-alt)
  done

lemma word-lsb-int: lsb w ⟷ uint w mod 2 = 1
  by (auto simp: word-lsb-def bin-last-def)

lemma word-msb-sint: msb w ⟷ sint w < 0
  by (simp only: word-msb-def sign-Min-lt-0)

lemma msb-word-of-int: msb (word-of-int x::'a::len word) = bin-nth x (len-of
  TYPE('a) - 1)
  by (simp add: word-msb-def word-sbin.eq-norm bin-sign-lem)

lemma word-msb-numeral [simp]:
  msb (numeral w::'a::len word) = bin-nth (numeral w) (len-of TYPE('a) - 1)

```

```

unfolding word-numeral-alt by (rule msb-word-of-int)

lemma word-msb-neg-numeral [simp]:
  msb (– numeral w :: 'a::len word) = bin-nth (– numeral w) (len-of TYPE('a) –
  1)
unfolding word-neg-numeral-alt by (rule msb-word-of-int)

lemma word-msb-0 [simp]: ¬ msb (0 :: 'a::len word)
by (simp add: word-msb-def)

lemma word-msb-1 [simp]: msb (1 :: 'a::len word) ⟷ len-of TYPE('a) = 1
unfolding word-1-wi msb-word-of-int eq-iff [where 'a=nat]
by (simp add: Suc-le-eq)

lemma word-msb-nth: msb w = bin-nth (uint w) (len-of TYPE('a) – 1)
for w :: 'a::len word
by (simp add: word-msb-def sint-uint bin-sign-lem)

lemma word-msb-alt: msb w = hd (to-bl w)
for w :: 'a::len word
apply (unfold word-msb-nth uint-bl)
apply (subst hd-conv-nth)
apply (rule length-greater-0-conv [THEN iffD1])
apply simp
apply (simp add : nth-bin-to-bl word-size)
done

lemma word-set-nth [simp]: set-bit w n (test-bit w n) = w
for w :: 'a::len0 word
by (auto simp: word-test-bit-def word-set-bit-def)

lemma bin-nth-uint': bin-nth (uint w) n ⟷ rev (bin-to-bl (size w) (uint w)) ! n
  ∧ n < size w
apply (unfold word-size)
apply (safe elim!: bin-nth-uint-imp)
apply (frule bin-nth-uint-imp)
apply (fast dest!: bin-nth-bl)+
done

lemmas bin-nth-uint = bin-nth-uint' [unfolded word-size]

lemma test-bit-bl: w !! n ⟷ rev (to-bl w) ! n ∧ n < size w
unfolding to-bl-def word-test-bit-def word-size by (rule bin-nth-uint)

lemma to-bl-nth: n < size w ⟹ to-bl w ! n = w !! (size w – Suc n)
apply (unfold test-bit-bl)
apply clarsimp
apply (rule trans)
apply (rule nth-rev-alt)

```

```

  apply (auto simp add: word-size)
done

lemma test-bit-set: (set-bit w n x) !! n  $\longleftrightarrow$  n < size w  $\wedge$  x
  for w :: 'a::len0 word
  by (auto simp: word-size word-test-bit-def word-set-bit-def word-ubin.eq-norm
nth-bintr)

lemma test-bit-set-gen:
  test-bit (set-bit w n x) m = (if m = n then n < size w  $\wedge$  x else test-bit w m)
  for w :: 'a::len0 word
  apply (unfold word-size word-test-bit-def word-set-bit-def)
  apply (clarsimp simp add: word-ubin.eq-norm nth-bintr bin-nth-sc-gen)
  apply (auto elim!: test-bit-size [unfolded word-size]
    simp add: word-test-bit-def [symmetric])
done

lemma of-bl-rep-False: of-bl (replicate n False @ bs) = of-bl bs
  by (auto simp: of-bl-def bl-to-bin-rep-F)

lemma msb-nth: msb w = w !! (len-of TYPE('a) - 1)
  for w :: 'a::len word
  by (simp add: word-msb-nth word-test-bit-def)

lemmas msb0 = len-gt-0 [THEN diff-Suc-less, THEN word-ops-nth-size [unfolded
word-size]]
lemmas msb1 = msb0 [where i = 0]
lemmas word-ops-msb = msb1 [unfolded msb-nth [symmetric, unfolded One-nat-def]]

lemmas lsb0 = len-gt-0 [THEN word-ops-nth-size [unfolded word-size]]
lemmas word-ops-lsb = lsb0 [unfolded word-lsb-alt]

lemma td-ext-nth [OF refl refl refl, unfolded word-size]:
  n = size w  $\implies$  ofn = set-bits  $\implies$  [w, ofn g] = l  $\implies$ 
  td-ext test-bit ofn {f.  $\forall i. f i \longrightarrow i < n$ } ( $\lambda h i. h i \wedge i < n$ )
  for w :: 'a::len0 word
  apply (unfold word-size td-ext-def)
  apply safe
  apply (rule-tac [3] ext)
  apply (rule-tac [4] ext)
  apply (unfold word-size of-nth-def test-bit-bl)
  apply safe
  defer
  apply (clarsimp simp: word-bl.Abs-inverse)+
  apply (rule word-bl.Rep-inverse')
  apply (rule sym [THEN trans])
  apply (rule bl-of-nth-nth)
  apply simp
  apply (rule bl-of-nth-inj)

```

apply (*clarsimp simp add : test-bit-bl word-size*)
done

interpretation *test-bit*:

td-ext
 $op \ !! :: 'a::len0 \ word \Rightarrow nat \Rightarrow bool$
set-bits
 $\{f. \forall i. f \ i \longrightarrow i < len\text{-of} \ TYPE('a::len0)\}$
 $(\lambda h \ i. h \ i \wedge i < len\text{-of} \ TYPE('a::len0))$
by (*rule td-ext-nth*)

lemmas *td-nth = test-bit.td-thm*

lemma *word-set-set-same* [*simp*]: *set-bit (set-bit w n x) n y = set-bit w n y*
for *w :: 'a::len0 word*
by (*rule word-eqI*) (*simp add : test-bit-set-gen word-size*)

lemma *word-set-set-diff*:

fixes *w :: 'a::len0 word*
assumes $m \neq n$
shows *set-bit (set-bit w m x) n y = set-bit (set-bit w n y) m x*
by (*rule word-eqI*) (*auto simp: test-bit-set-gen word-size assms*)

lemma *nth-sint*:

fixes *w :: 'a::len word*
defines $l \equiv len\text{-of} \ TYPE('a)$
shows $bin\text{-nth} \ (sint \ w) \ n = (if \ n < l - 1 \ then \ w \ !! \ n \ else \ w \ !! \ (l - 1))$
unfolding *sint-uint l-def*
by (*auto simp: nth-sbintr word-test-bit-def [symmetric]*)

lemma *word-lsb-numeral* [*simp*]:

$lsb \ (numeral \ bin :: 'a::len \ word) \longleftrightarrow bin\text{-last} \ (numeral \ bin)$
unfolding *word-lsb-alt test-bit-numeral* **by** *simp*

lemma *word-lsb-neg-numeral* [*simp*]:

$lsb \ (- \ numeral \ bin :: 'a::len \ word) \longleftrightarrow bin\text{-last} \ (- \ numeral \ bin)$
by (*simp add: word-lsb-alt*)

lemma *set-bit-word-of-int*: *set-bit (word-of-int x) n b = word-of-int (bin-sc n b x)*

unfolding *word-set-bit-def*
by (*rule word-eqI*)(*simp add: word-size bin-nth-sc-gen word-ubin.eq-norm nth-bintr*)

lemma *word-set-numeral* [*simp*]:

$set\text{-bit} \ (numeral \ bin :: 'a::len0 \ word) \ n \ b =$
 $word\text{-of-int} \ (bin\text{-sc} \ n \ b \ (numeral \ bin))$
unfolding *word-numeral-alt* **by** (*rule set-bit-word-of-int*)

lemma *word-set-neg-numeral* [*simp*]:

$set\text{-bit} \ (- \ numeral \ bin :: 'a::len0 \ word) \ n \ b =$

```

    word-of-int (bin-sc n b (– numeral bin))
  unfolding word-neg-numeral-alt by (rule set-bit-word-of-int)

lemma word-set-bit-0 [simp]: set-bit 0 n b = word-of-int (bin-sc n b 0)
  unfolding word-0-wi by (rule set-bit-word-of-int)

lemma word-set-bit-1 [simp]: set-bit 1 n b = word-of-int (bin-sc n b 1)
  unfolding word-1-wi by (rule set-bit-word-of-int)

lemma setBit-no [simp]: setBit (numeral bin) n = word-of-int (bin-sc n True
  (numeral bin))
  by (simp add: setBit-def)

lemma clearBit-no [simp]:
  clearBit (numeral bin) n = word-of-int (bin-sc n False (numeral bin))
  by (simp add: clearBit-def)

lemma to-bl-n1: to-bl (–1::'a::len0 word) = replicate (len-of TYPE('a)) True
  apply (rule word-bl.Abs-inverse')
  apply simp
  apply (rule word-eqI)
  apply (clarsimp simp add: word-size)
  apply (auto simp add: word-bl.Abs-inverse test-bit-bl word-size)
  done

lemma word-msb-n1 [simp]: msb (–1::'a::len word)
  unfolding word-msb-alt to-bl-n1 by simp

lemma word-set-nth-iff: set-bit w n b = w ⟷ w !! n = b ∨ n ≥ size w
  for w :: 'a::len0 word
  apply (rule iffI)
  apply (rule disjCI)
  apply (drule word-eqD)
  apply (erule sym [THEN trans])
  apply (simp add: test-bit-set)
  apply (erule disjE)
  apply clarsimp
  apply (rule word-eqI)
  apply (clarsimp simp add: test-bit-set-gen)
  apply (drule test-bit-size)
  apply force
  done

lemma test-bit-2p: (word-of-int (2 ^ n)::'a::len word) !! m ⟷ m = n ∧ m <
  len-of TYPE('a)
  by (auto simp: word-test-bit-def word-ubin.eq-norm nth-bintr nth-2p-bin)

lemma nth-w2p: ((2::'a::len word) ^ n) !! m ⟷ m = n ∧ m < len-of TYPE('a::len)
  by (simp add: test-bit-2p [symmetric] word-of-int [symmetric])

```

```

lemma uint-2p: (0::'a::len word) < 2 ^ n  $\implies$  uint (2 ^ n::'a::len word) = 2 ^ n
  apply (unfold word-arith-power-alt)
  apply (case-tac len-of TYPE('a))
  apply clarsimp
  apply (case-tac nat)
  apply clarsimp
  apply (case-tac n)
  apply clarsimp
  apply clarsimp
  apply (drule word-gt-0 [THEN iffD1])
  apply (safe intro!: word-eqI)
  apply (auto simp add: nth-2p-bin)
  apply (erule notE)
  apply (simp (no-asm-use) add: uint-word-of-int word-size)
  apply (subst mod-pos-pos-trivial)
  apply simp
  apply (rule power-strict-increasing)
  apply simp-all
done

```

```

lemma word-of-int-2p: (word-of-int (2 ^ n) :: 'a::len word) = 2 ^ n
  by (induct n) (simp-all add: wi-hom-syms)

```

```

lemma bang-is-le: x !! m  $\implies$  2 ^ m  $\leq$  x
  for x :: 'a::len word
  apply (rule xtr3)
  apply (rule-tac [2] y = x in le-word-or2)
  apply (rule word-eqI)
  apply (auto simp add: word-ao-nth nth-w2p word-size)
done

```

```

lemma word-clr-le: w  $\geq$  set-bit w n False
  for w :: 'a::len0 word
  apply (unfold word-set-bit-def word-le-def word-ubin.eq-norm)
  apply (rule order-trans)
  apply (rule bintr-bin-clr-le)
  apply simp
done

```

```

lemma word-set-ge: w  $\leq$  set-bit w n True
  for w :: 'a::len word
  apply (unfold word-set-bit-def word-le-def word-ubin.eq-norm)
  apply (rule order-trans [OF - bintr-bin-set-ge])
  apply simp
done

```

10.25 Shifting, Rotating, and Splitting Words

lemma *shiffl1-wi* [*simp*]: *shiffl1* (word-of-int *w*) = word-of-int (*w* BIT False)
unfolding *shiffl1-def*
apply (*simp* add: word-ubin.norm-eq-iff [*symmetric*] word-ubin.eq-norm)
apply (*subst* refl [*THEN* bintrunc-BIT-I, *symmetric*])
apply (*subst* bintrunc-bintrunc-min)
apply *simp*
done

lemma *shiffl1-numeral* [*simp*]: *shiffl1* (numeral *w*) = numeral (Num.Bit0 *w*)
unfolding word-numeral-alt *shiffl1-wi* **by** *simp*

lemma *shiffl1-neg-numeral* [*simp*]: *shiffl1* (– numeral *w*) = – numeral (Num.Bit0 *w*)
unfolding word-neg-numeral-alt *shiffl1-wi* **by** *simp*

lemma *shiffl1-0* [*simp*]: *shiffl1* 0 = 0
by (*simp* add: *shiffl1-def*)

lemma *shiffl1-def-u*: *shiffl1* *w* = word-of-int (uint *w* BIT False)
by (*simp* only: *shiffl1-def*)

lemma *shiffl1-def-s*: *shiffl1* *w* = word-of-int (sint *w* BIT False)
by (*simp* add: *shiffl1-def* Bit-B0 wi-hom-syms)

lemma *shiftr1-0* [*simp*]: *shiftr1* 0 = 0
by (*simp* add: *shiftr1-def*)

lemma *sshiftr1-0* [*simp*]: *sshiftr1* 0 = 0
by (*simp* add: *sshiftr1-def*)

lemma *sshiftr1-n1* [*simp*]: *sshiftr1* (– 1) = – 1
by (*simp* add: *sshiftr1-def*)

lemma *shiffl-0* [*simp*]: (0::'a::len0 word) << *n* = 0
by (*induct* *n*) (*auto simp: shiffl-def*)

lemma *shiftr-0* [*simp*]: (0::'a::len0 word) >> *n* = 0
by (*induct* *n*) (*auto simp: shiftr-def*)

lemma *sshiftr-0* [*simp*]: 0 >>> *n* = 0
by (*induct* *n*) (*auto simp: sshiftr-def*)

lemma *sshiftr-n1* [*simp*]: –1 >>> *n* = –1
by (*induct* *n*) (*auto simp: sshiftr-def*)

lemma *nth-shiffl1*: *shiffl1* *w* !! *n* $\longleftrightarrow$ *n* < size *w* $\wedge$ *n* > 0 $\wedge$ *w* !! (*n* – 1)
apply (*unfold shiffl1-def word-test-bit-def*)
apply (*simp* add: *nth-bintr* word-ubin.eq-norm word-size)

```

apply (cases n)
apply auto
done

lemma nth-shiftl': (w << m) !! n  $\longleftrightarrow$  n < size w  $\wedge$  n >= m  $\wedge$  w !! (n - m)
for w :: 'a::len0 word
apply (unfold shiftl-def)
apply (induct m arbitrary: n)
apply (force elim!: test-bit-size)
apply (clarsimp simp add : nth-shiftl1 word-size)
apply arith
done

lemmas nth-shiftl = nth-shiftl' [unfolded word-size]

lemma nth-shiftr1: shiftr1 w !! n = w !! Suc n
apply (unfold shiftr1-def word-test-bit-def)
apply (simp add: nth-bintr word-ubin.eq-norm)
apply safe
apply (drule bin-nth.Suc [THEN iffD2, THEN bin-nth-uint-imp])
apply simp
done

lemma nth-shiftr: (w >> m) !! n = w !! (n + m)
for w :: 'a::len0 word
apply (unfold shiftr-def)
apply (induct m arbitrary: n)
apply (auto simp add: nth-shiftr1)
done

lemma uint-shiftr1: uint (shiftr1 w) = bin-rest (uint w)
apply (unfold shiftr1-def word-ubin.eq-norm bin-rest-trunc-i)
apply (subst bintr-uint [symmetric, OF order-refl])
apply (simp only : bintrunc-bintrunc-l)
apply simp
done

lemma nth-sshiftr1: sshiftr1 w !! n = (if n = size w - 1 then w !! n else w !! Suc n)
apply (unfold sshiftr1-def word-test-bit-def)
apply (simp add: nth-bintr word-ubin.eq-norm bin-nth.Suc [symmetric] word-size
del: bin-nth.simps)
apply (simp add: nth-bintr uint-sint del : bin-nth.simps)
apply (auto simp add: bin-nth-sint)
done

lemma nth-sshiftr [rule-format] :

```

```

 $\forall n. \text{sshiftr } w \text{ !! } n =$ 
  ( $n < \text{size } w \wedge (\text{if } n + m \geq \text{size } w \text{ then } w \text{ !! } (\text{size } w - 1) \text{ else } w \text{ !! } (n + m))$ )
apply (unfold sshiftr-def)
apply (induct-tac m)
  apply (simp add: test-bit-bl)
apply (clarsimp simp add: nth-sshiftr1 word-size)
apply safe
  apply arith
  apply arith
apply (erule thin-rl)
apply (case-tac n)
  apply safe
  apply simp
  apply simp
apply (erule thin-rl)
apply (case-tac n)
  apply safe
  apply simp
  apply simp
apply arith+
done

lemma shiftr1-div-2:  $\text{uint } (\text{shiftr1 } w) = \text{uint } w \text{ div } 2$ 
apply (unfold shiftr1-def bin-rest-def)
apply (rule word-uint.Abs-inverse)
apply (simp add: uints-num pos-imp-zdiv-nonneg-iff)
apply (rule xtr7)
prefer 2
apply (rule zdiv-le-dividend)
apply auto
done

lemma sshiftr1-div-2:  $\text{sint } (\text{sshiftr1 } w) = \text{sint } w \text{ div } 2$ 
apply (unfold sshiftr1-def bin-rest-def [symmetric])
apply (simp add: word-sbin.eq-norm)
apply (rule trans)
defer
apply (subst word-sbin.norm-Rep [symmetric])
apply (rule refl)
apply (subst word-sbin.norm-Rep [symmetric])
apply (unfold One-nat-def)
apply (rule sbintrunc-rest)
done

lemma shiftr-div-2n:  $\text{uint } (\text{shiftr } w \text{ } n) = \text{uint } w \text{ div } 2 \wedge n$ 
apply (unfold shiftr-def)
apply (induct n)
  apply simp
apply (simp add: shiftr1-div-2 mult.commute zdiv-zmult2-eq [symmetric])

```

done

```
lemma sshiftr-div-2n: sint (sshiftr w n) = sint w div 2 ^ n
  apply (unfold sshiftr-def)
  apply (induct n)
  apply simp
  apply (simp add: sshiftr1-div-2 mult.commute zdiv-zmult2-eq [symmetric])
done
```

10.25.1 shift functions in terms of lists of bools

```
lemmas bshiftr1-numeral [simp] =
  bshiftr1-def [where w=numeral w, unfolded to-bl-numeral] for w
```

```
lemma bshiftr1-bl: to-bl (bshiftr1 b w) = b # butlast (to-bl w)
  unfolding bshiftr1-def by (rule word-bl.Abs-inverse) simp
```

```
lemma shiftrl1-of-bl: shiftrl1 (of-bl bl) = of-bl (bl @ [False])
  by (simp add: of-bl-def bl-to-bin-append)
```

```
lemma shiftrl1-bl: shiftrl1 w = of-bl (to-bl w @ [False])
  for w :: 'a::len0 word
```

```
proof -
  have shiftrl1 w = shiftrl1 (of-bl (to-bl w))
    by simp
  also have ... = of-bl (to-bl w @ [False])
    by (rule shiftrl1-of-bl)
  finally show ?thesis .
qed
```

```
lemma bl-shiftrl1: to-bl (shiftrl1 w) = tl (to-bl w) @ [False]
  for w :: 'a::len word
  by (simp add: shiftrl1-bl word-rep-drop drop-Suc drop-Cons') (fast intro!: Suc-leI)
```

```
lemma bl-shiftrl1': to-bl (shiftrl1 w) = tl (to-bl w @ [False])
  by (simp add: shiftrl1-bl word-rep-drop drop-Suc del: drop-append)
```

```
lemma shiftr1-bl: shiftr1 w = of-bl (butlast (to-bl w))
  apply (unfold shiftr1-def uint-bl of-bl-def)
  apply (simp add: butlast-rest-bin word-size)
  apply (simp add: bin-rest-trunc [symmetric, unfolded One-nat-def])
done
```

```
lemma bl-shiftr1: to-bl (shiftr1 w) = False # butlast (to-bl w)
  for w :: 'a::len word
  by (simp add: shiftr1-bl word-rep-drop len-gt-0 [THEN Suc-leI])
```

```

lemma bl-shiftr1': to-bl (shiftr1 w) = butlast (False # to-bl w)
  apply (rule word-bl.Abs-inverse')
  apply (simp del: butlast.simps)
  apply (simp add: shiftr1-bl of-bl-def)
  done

lemma shifl1-rev: shifl1 w = word-reverse (shiftr1 (word-reverse w))
  apply (unfold word-reverse-def)
  apply (rule word-bl.Rep-inverse' [symmetric])
  apply (simp add: bl-shifl1' bl-shiftr1' word-bl.Abs-inverse)
  apply (cases to-bl w)
  apply auto
  done

lemma shifl-rev: shifl w n = word-reverse (shiftr (word-reverse w) n)
  by (induct n) (auto simp add: shifl-def shiftr-def shifl1-rev)

lemma rev-shifl: word-reverse w << n = word-reverse (w >> n)
  by (simp add: shifl-rev)

lemma shiftr-rev: w >> n = word-reverse (word-reverse w << n)
  by (simp add: rev-shifl)

lemma rev-shiftr: word-reverse w >> n = word-reverse (w << n)
  by (simp add: shiftr-rev)

lemma bl-sshiftr1: to-bl (sshiftr1 w) = hd (to-bl w) # butlast (to-bl w)
  for w :: 'a::len word
  apply (unfold sshiftr1-def uint-bl word-size)
  apply (simp add: butlast-rest-bin word-ubin.eq-norm)
  apply (simp add: sint-uint)
  apply (rule nth-equalityI)
  apply clarsimp
  apply clarsimp
  apply (case-tac i)
  apply (simp-all add: hd-conv-nth length-0-conv [symmetric]
    nth-bin-to-bl bin-nth.Suc [symmetric] nth-sbintr
    del: bin-nth.Suc)
  apply force
  apply (rule impI)
  apply (rule-tac f = bin-nth (uint w) in arg-cong)
  apply simp
  done

lemma drop-shiftr: drop n (to-bl (w >> n)) = take (size w - n) (to-bl w)
  for w :: 'a::len word
  apply (unfold shiftr-def)
  apply (induct n)
  prefer 2

```

```

apply (simp add: drop-Suc bl-shiftr1 butlast-drop [symmetric])
apply (rule butlast-take [THEN trans])
apply (auto simp: word-size)
done

lemma drop-sshiftr: drop n (to-bl (w >>> n)) = take (size w - n) (to-bl w)
for w :: 'a::len word
apply (unfold sshiftr-def)
apply (induct n)
prefer 2
apply (simp add: drop-Suc bl-sshiftr1 butlast-drop [symmetric])
apply (rule butlast-take [THEN trans])
apply (auto simp: word-size)
done

lemma take-shiftr: n ≤ size w ⇒ take n (to-bl (w >> n)) = replicate n False
apply (unfold shiftr-def)
apply (induct n)
prefer 2
apply (simp add: bl-shiftr1' length-0-conv [symmetric] word-size)
apply (rule take-butlast [THEN trans])
apply (auto simp: word-size)
done

lemma take-sshiftr' [rule-format] :
  n ≤ size w ⟶ hd (to-bl (w >>> n)) = hd (to-bl w) ∧
  take n (to-bl (w >>> n)) = replicate n (hd (to-bl w))
for w :: 'a::len word
apply (unfold sshiftr-def)
apply (induct n)
prefer 2
apply (simp add: bl-sshiftr1)
apply (rule impI)
apply (rule take-butlast [THEN trans])
apply (auto simp: word-size)
done

lemmas hd-sshiftr = take-sshiftr' [THEN conjunct1]
lemmas take-sshiftr = take-sshiftr' [THEN conjunct2]

lemma atd-lem: take n xs = t ⟹ drop n xs = d ⟹ xs = t @ d
by (auto intro: append-take-drop-id [symmetric])

lemmas bl-shiftr = atd-lem [OF take-shiftr drop-shiftr]
lemmas bl-sshiftr = atd-lem [OF take-sshiftr drop-sshiftr]

lemma shiftl-of-bl: of-bl bl << n = of-bl (bl @ replicate n False)
by (induct n) (auto simp: shiftl-def shiftl1-of-bl replicate-app-Cons-same)

```

```

lemma shiftrl-bl:  $w << n = \text{of-bl } (to\text{-bl } w @ \text{replicate } n \text{ False})$ 
  for  $w :: 'a::len0 \text{ word}$ 
proof –
  have  $w << n = \text{of-bl } (to\text{-bl } w) << n$ 
    by simp
  also have  $\dots = \text{of-bl } (to\text{-bl } w @ \text{replicate } n \text{ False})$ 
    by (rule shiftrl-of-bl)
  finally show ?thesis .
qed

```

```

lemmas shiftrl-numeral [simp] = shiftrl-def [where  $w = \text{numeral } w$ ] for  $w$ 

```

```

lemma bl-shiftrl:  $to\text{-bl } (w << n) = \text{drop } n (to\text{-bl } w) @ \text{replicate } (\min (\text{size } w) n) \text{ False}$ 
  by (simp add: shiftrl-bl word-rep-drop word-size)

```

```

lemma shiftrl-zero-size:  $\text{size } x \leq n \implies x << n = 0$ 
  for  $x :: 'a::len0 \text{ word}$ 
  apply (unfold word-size)
  apply (rule word-eqI)
  apply (clarsimp simp add: shiftrl-bl word-size test-bit-of-bl nth-append)
  done

```

```

lemma shiftrl1-2t:  $\text{shiftrl1 } w = 2 * w$ 
  for  $w :: 'a::len \text{ word}$ 
  by (simp add: shiftrl1-def Bit-def wi-hom-mult [symmetric])

```

```

lemma shiftrl1-p:  $\text{shiftrl1 } w = w + w$ 
  for  $w :: 'a::len \text{ word}$ 
  by (simp add: shiftrl1-2t)

```

```

lemma shiftrl-t2n:  $\text{shiftrl } w \ n = 2 ^ n * w$ 
  for  $w :: 'a::len \text{ word}$ 
  by (induct n) (auto simp: shiftrl-def shiftrl1-2t)

```

```

lemma shiftr1-bintr [simp]:
  (shiftr1 (numeral  $w$ ) ::  $'a::len0 \text{ word}$ ) =
    word-of-int (bin-rest (bintrunc (len-of TYPE('a)) (numeral  $w$ )))
  unfolding shiftr1-def word-numeral-alt by (simp add: word-ubin.eq-norm)

```

```

lemma sshiftr1-sbintr [simp]:
  (sshiftr1 (numeral  $w$ ) ::  $'a::len \text{ word}$ ) =
    word-of-int (bin-rest (sbintrunc (len-of TYPE('a) - 1) (numeral  $w$ )))
  unfolding sshiftr1-def word-numeral-alt by (simp add: word-sbin.eq-norm)

```

```

lemma shiftr-no [simp]:

```

```

(numeral w::'a::len0 word) >> n = word-of-int
  ((bin-rest ^ n) (bintrunc (len-of TYPE('a)) (numeral w)))
by (rule word-eqI) (auto simp: nth-shiftr nth-rest-power-bin nth-bintr word-size)

```

lemma *sshiftr-no* [simp]:

```

(numeral w::'a::len word) >>> n = word-of-int
  ((bin-rest ^ n) (sbintrunc (len-of TYPE('a) - 1) (numeral w)))
apply (rule word-eqI)
apply (auto simp: nth-sshiftr nth-rest-power-bin nth-sbintr word-size)
apply (subgoal-tac na + n = len-of TYPE('a) - Suc 0, simp, simp)+
done

```

lemma *shiftr1-bl-of*:

```

length bl ≤ len-of TYPE('a) ⇒
  shiftr1 (of-bl bl::'a::len0 word) = of-bl (butlast bl)
by (clarsimp simp: shiftr1-def of-bl-def butlast-rest-bl2bin word-ubin.eq-norm trunc-bl2bin)

```

lemma *shiftr-bl-of*:

```

length bl ≤ len-of TYPE('a) ⇒
  (of-bl bl::'a::len0 word) >> n = of-bl (take (length bl - n) bl)
apply (unfold shiftr-def)
apply (induct n)
apply clarsimp
apply clarsimp
apply (subst shiftr1-bl-of)
apply simp
apply (simp add: butlast-take)
done

```

lemma *shiftr-bl*: $x \gg n \equiv \text{of-bl } (\text{take } (\text{len-of TYPE('a)} - n) (\text{to-bl } x))$
 for $x :: 'a::\text{len0 word}$
 using *shiftr-bl-of* [where $'a = 'a$, of *to-bl* x] by *simp*

lemma *msb-shift*: $\text{msb } w \longleftrightarrow w \gg (\text{len-of TYPE('a)} - 1) \neq 0$
 for $w :: 'a::\text{len word}$
 apply (unfold shiftr-bl word-msb-alt)
 apply (simp add: word-size Suc-le-eq take-Suc)
 apply (cases hd (to-bl w))
 apply (auto simp: word-1-bl of-bl-rep-False [where $n=1$ and $bs=[]$, simplified])
done

lemma *zip-replicate*: $n \geq \text{length } ys \implies \text{zip } (\text{replicate } n x) ys = \text{map } (\lambda y. (x, y)) ys$
 apply (induct ys arbitrary: n)
 apply simp-all
 apply (case-tac n)
 apply simp-all
done

```

lemma align-lem-or [rule-format] :
   $\forall x m. \text{length } x = n + m \longrightarrow \text{length } y = n + m \longrightarrow$ 
   $\text{drop } m \ x = \text{replicate } n \ \text{False} \longrightarrow \text{take } m \ y = \text{replicate } m \ \text{False} \longrightarrow$ 
   $\text{map2 } op \mid x \ y = \text{take } m \ x \ @ \ \text{drop } m \ y$ 
  apply (induct y)
  apply force
  apply clarsimp
  apply (case-tac x)
  apply force
  apply (case-tac m)
  apply auto
  apply (drule-tac t=length xs for xs in sym)
  apply (auto simp: map2-def zip-replicate o-def)
done

```

```

lemma align-lem-and [rule-format] :
   $\forall x m. \text{length } x = n + m \longrightarrow \text{length } y = n + m \longrightarrow$ 
   $\text{drop } m \ x = \text{replicate } n \ \text{False} \longrightarrow \text{take } m \ y = \text{replicate } m \ \text{False} \longrightarrow$ 
   $\text{map2 } op \wedge x \ y = \text{replicate } (n + m) \ \text{False}$ 
  apply (induct y)
  apply force
  apply clarsimp
  apply (case-tac x)
  apply force
  apply (case-tac m)
  apply auto
  apply (drule-tac t=length xs for xs in sym)
  apply (auto simp: map2-def zip-replicate o-def map-replicate-const)
done

```

```

lemma aligned-bl-add-size [OF refl]:
   $\text{size } x - n = m \implies n \leq \text{size } x \implies \text{drop } m \ (\text{to-bl } x) = \text{replicate } n \ \text{False} \implies$ 
   $\text{take } m \ (\text{to-bl } y) = \text{replicate } m \ \text{False} \implies$ 
   $\text{to-bl } (x + y) = \text{take } m \ (\text{to-bl } x) \ @ \ \text{drop } m \ (\text{to-bl } y)$ 
  apply (subgoal-tac x AND y = 0)
  prefer 2
  apply (rule word-bl.Rep-eqD)
  apply (simp add: bl-word-and)
  apply (rule align-lem-and [THEN trans])
  apply (simp-all add: word-size)[5]
  apply simp
  apply (subst word-plus-and-or [symmetric])
  apply (simp add : bl-word-or)
  apply (rule align-lem-or)
  apply (simp-all add: word-size)
done

```

10.25.2 Mask

lemma *nth-mask* [*OF refl, simp*]: $m = \text{mask } n \implies \text{test-bit } m \ i \longleftrightarrow i < n \wedge i < \text{size } m$

apply (*unfold mask-def test-bit-bl*)
apply (*simp only: word-1-bl [symmetric] shiftl-of-bl*)
apply (*clarsimp simp add: word-size*)
apply (*simp only: of-bl-def mask-lem word-of-int-hom-syms add-diff-cancel2*)
apply (*fold of-bl-def*)
apply (*simp add: word-1-bl*)
apply (*rule test-bit-of-bl [THEN trans, unfolded test-bit-bl word-size]*)
apply *auto*
done

lemma *mask-bl*: $\text{mask } n = \text{of-bl } (\text{replicate } n \ \text{True})$
by (*auto simp add : test-bit-of-bl word-size intro: word-eqI*)

lemma *mask-bin*: $\text{mask } n = \text{word-of-int } (\text{bintrunc } n \ (-1))$
by (*auto simp add: nth-bintr word-size intro: word-eqI*)

lemma *and-mask-bintr*: $w \ \text{AND} \ \text{mask } n = \text{word-of-int } (\text{bintrunc } n \ (\text{uint } w))$
apply (*rule word-eqI*)
apply (*simp add: nth-bintr word-size word-ops-nth-size*)
apply (*auto simp add: test-bit-bin*)
done

lemma *and-mask-wi*: $\text{word-of-int } i \ \text{AND} \ \text{mask } n = \text{word-of-int } (\text{bintrunc } n \ i)$
by (*auto simp add: nth-bintr word-size word-ops-nth-size word-eq-iff*)

lemma *and-mask-wi'*:
 $\text{word-of-int } i \ \text{AND} \ \text{mask } n = (\text{word-of-int } (\text{bintrunc } (\min \ \text{LENGTH}('a) \ n) \ i) :: 'a::\text{len word})$
by (*auto simp add: nth-bintr word-size word-ops-nth-size word-eq-iff*)

lemma *and-mask-no*: $\text{numeral } i \ \text{AND} \ \text{mask } n = \text{word-of-int } (\text{bintrunc } n \ (\text{numeral } i))$
unfolding *word-numeral-alt* **by** (*rule and-mask-wi*)

lemma *bl-and-mask'*:
 $\text{to-bl } (w \ \text{AND} \ \text{mask } n :: 'a::\text{len word}) =$
 $\text{replicate } (\text{len-of } \text{TYPE}('a) - n) \ \text{False} \ @$
 $\text{drop } (\text{len-of } \text{TYPE}('a) - n) \ (\text{to-bl } w)$
apply (*rule nth-equalityI*)
apply *simp*
apply (*clarsimp simp add: to-bl-nth word-size*)
apply (*simp add: word-size word-ops-nth-size*)
apply (*auto simp add: word-size test-bit-bl nth-append nth-rev*)
done

lemma *and-mask-mod-2p*: $w \ \text{AND} \ \text{mask } n = \text{word-of-int } (\text{uint } w \ \text{mod } 2^n)$

```

by (simp only: and-mask-bintr bintrunc-mod2p)

lemma and-mask-lt-2p: uint (w AND mask n) < 2 ^ n
  apply (simp add: and-mask-bintr word-ubin.eq-norm)
  apply (simp add: bintrunc-mod2p)
  apply (rule xtr8)
  prefer 2
  apply (rule pos-mod-bound)
  apply auto
done

lemma eq-mod-iff: 0 < n  $\implies$  b = b mod n  $\longleftrightarrow$  0  $\leq$  b  $\wedge$  b < n
  for b n :: int
  by (simp add: int-mod-lem eq-sym-conv)

lemma mask-eq-iff: w AND mask n = w  $\longleftrightarrow$  uint w < 2 ^ n
  apply (simp add: and-mask-bintr)
  apply (simp add: word-ubin.inverse-norm)
  apply (simp add: eq-mod-iff bintrunc-mod2p min-def)
  apply (fast intro!: lt2p-lem)
done

lemma and-mask-dvd: 2 ^ n dvd uint w  $\longleftrightarrow$  w AND mask n = 0
  apply (simp add: dvd-eq-mod-eq-0 and-mask-mod-2p)
  apply (simp add: word-uint.norm-eq-iff [symmetric] word-of-int-homs del: word-of-int-0)
  apply (subst word-uint.norm-Rep [symmetric])
  apply (simp only: bintrunc-bintrunc-min bintrunc-mod2p [symmetric] min-def)
  apply auto
done

lemma and-mask-dvd-nat: 2 ^ n dvd unat w  $\longleftrightarrow$  w AND mask n = 0
  apply (unfold unat-def)
  apply (rule trans [OF - and-mask-dvd])
  apply (unfold dvd-def)
  apply auto
  apply (drule uint-ge-0 [THEN nat-int.Abs-inverse' [simplified], symmetric])
  apply simp
  apply (simp add: nat-mult-distrib nat-power-eq)
done

lemma word-2p-lem: n < size w  $\implies$  w < 2 ^ n = (uint w < 2 ^ n)
  for w :: 'a::len word
  apply (unfold word-size word-less-alt word-numeral-alt)
  apply (auto simp add: word-of-int-power-hom word-uint.eq-norm mod-pos-pos-trivial
    simp del: word-of-int-numeral)
done

lemma less-mask-eq: x < 2 ^ n  $\implies$  x AND mask n = x
  for x :: 'a::len word

```

```

apply (unfold word-less-alt word-numeral-alt)
apply (clarsimp simp add: and-mask-mod-2p word-of-int-power-hom word-uint.eq-norm
      simp del: word-of-int-numeral)
apply (drule xtr8 [rotated])
apply (rule int-mod-le)
apply (auto simp add : mod-pos-pos-trivial)
done

```

```

lemmas mask-eq-iff-w2p = trans [OF mask-eq-iff word-2p-lem [symmetric]]

```

```

lemmas and-mask-less' = iffD2 [OF word-2p-lem and-mask-lt-2p, simplified word-size]

```

```

lemma and-mask-less-size:  $n < \text{size } x \implies x \text{ AND mask } n < 2^n$ 
unfolding word-size by (erule and-mask-less')

```

```

lemma word-mod-2p-is-mask [OF refl]:  $c = 2^n \implies c > 0 \implies x \bmod c = x$ 
AND mask n
for c x :: 'a::len word
by (auto simp: word-mod-def uint-2p and-mask-mod-2p)

```

```

lemma mask-egs:

```

```

  (a AND mask n) + b AND mask n = a + b AND mask n
  a + (b AND mask n) AND mask n = a + b AND mask n
  (a AND mask n) - b AND mask n = a - b AND mask n
  a - (b AND mask n) AND mask n = a - b AND mask n
  a * (b AND mask n) AND mask n = a * b AND mask n
  (b AND mask n) * a AND mask n = b * a AND mask n
  (a AND mask n) + (b AND mask n) AND mask n = a + b AND mask n
  (a AND mask n) - (b AND mask n) AND mask n = a - b AND mask n
  (a AND mask n) * (b AND mask n) AND mask n = a * b AND mask n
  - (a AND mask n) AND mask n = - a AND mask n
  word-succ (a AND mask n) AND mask n = word-succ a AND mask n
  word-pred (a AND mask n) AND mask n = word-pred a AND mask n
using word-of-int-Ex [where x=a] word-of-int-Ex [where x=b]
by (auto simp: and-mask-wi' word-of-int-homs word.abs-eq-iff bintrunc-mod2p
  mod-simps)

```

```

lemma mask-power-eq:  $(x \text{ AND mask } n)^k \text{ AND mask } n = x^k \text{ AND mask } n$ 
using word-of-int-Ex [where x=x]
by (auto simp: and-mask-wi' word-of-int-power-hom word.abs-eq-iff bintrunc-mod2p
  mod-simps)

```

10.25.3 Recast

```

lemmas revcast-def' = revcast-def [simplified]
lemmas revcast-def'' = revcast-def' [simplified word-size]
lemmas revcast-no-def [simp] = revcast-def' [where w=numeral w, unfolded word-size]
for w

```

lemma *to-bl-revcast*:

```

  to-bl (revcast w :: 'a::len0 word) = takefill False (len-of TYPE('a)) (to-bl w)
  apply (unfold revcast-def' word-size)
  apply (rule word-bl.Abs-inverse)
  apply simp
done

```

lemma *revcast-rev-ucast* [OF refl refl refl]:

```

  cs = [rc, uc]  $\implies$  rc = revcast (word-reverse w)  $\implies$  uc = ucast w  $\implies$ 
    rc = word-reverse uc
  apply (unfold ucast-def revcast-def' Let-def word-reverse-def)
  apply (auto simp: to-bl-of-bin takefill-bintrunc)
  apply (simp add: word-bl.Abs-inverse word-size)
done

```

lemma *revcast-ucast*: *revcast w = word-reverse (ucast (word-reverse w))*
 using *revcast-rev-ucast* [of word-reverse w] **by** *simp*

lemma *ucast-revcast*: *ucast w = word-reverse (revcast (word-reverse w))*
by (fact *revcast-rev-ucast* [THEN word-rev-gal])

lemma *ucast-rev-revcast*: *ucast (word-reverse w) = word-reverse (revcast w)*
by (fact *revcast-ucast* [THEN word-rev-gal])

linking revcast and cast via shift

lemmas *wsst-TYs* = *source-size target-size word-size*

lemma *revcast-down-uu* [OF refl]:

```

  rc = revcast  $\implies$  source-size rc = target-size rc + n  $\implies$  rc w = ucast (w >> n)
  for w :: 'a::len word
  apply (simp add: revcast-def')
  apply (rule word-bl.Rep-inverse')
  apply (rule trans, rule ucast-down-drop)
  prefer 2
  apply (rule trans, rule drop-shiftr)
  apply (auto simp: takefill-alt wsst-TYs)
done

```

lemma *revcast-down-us* [OF refl]:

```

  rc = revcast  $\implies$  source-size rc = target-size rc + n  $\implies$  rc w = ucast (w >>>
n)
  for w :: 'a::len word
  apply (simp add: revcast-def')
  apply (rule word-bl.Rep-inverse')
  apply (rule trans, rule ucast-down-drop)
  prefer 2
  apply (rule trans, rule drop-sshiftr)
  apply (auto simp: takefill-alt wsst-TYs)
done

```

lemma *revcast-down-su* [*OF refl*]:

```

rc = revcast  $\implies$  source-size rc = target-size rc + n  $\implies$  rc w = scast (w >> n)
for w :: 'a::len word
  apply (simp add: revcast-def')
  apply (rule word-bl.Rep-inverse')
  apply (rule trans, rule scast-down-drop)
  prefer 2
  apply (rule trans, rule drop-shiftr)
  apply (auto simp: takefill-alt wsst-TYs)
done

```

lemma *revcast-down-ss* [*OF refl*]:

```

rc = revcast  $\implies$  source-size rc = target-size rc + n  $\implies$  rc w = scast (w >>>
n)
for w :: 'a::len word
  apply (simp add: revcast-def')
  apply (rule word-bl.Rep-inverse')
  apply (rule trans, rule scast-down-drop)
  prefer 2
  apply (rule trans, rule drop-sshiftr)
  apply (auto simp: takefill-alt wsst-TYs)
done

```

lemma *cast-down-rev*:

```

uc = ucast  $\implies$  source-size uc = target-size uc + n  $\implies$  uc w = revcast (w <<
n)
for w :: 'a::len word
  apply (unfold shiftl-rev)
  apply clarify
  apply (simp add: revcast-rev-ucast)
  apply (rule word-rev-gal')
  apply (rule trans [OF - revcast-rev-ucast])
  apply (rule revcast-down-uu [symmetric])
  apply (auto simp add: wsst-TYs)
done

```

lemma *revcast-up* [*OF refl*]:

```

rc = revcast  $\implies$  source-size rc + n = target-size rc  $\implies$ 
  rc w = (ucast w :: 'a::len word) << n
  apply (simp add: revcast-def')
  apply (rule word-bl.Rep-inverse')
  apply (simp add: takefill-alt)
  apply (rule bl-shiftl [THEN trans])
  apply (subst ucast-up-app)
  apply (auto simp add: wsst-TYs)
done

```

```

lemmas rc1 = revcast-up [THEN
  revcast-rev-ucast [symmetric, THEN trans, THEN word-rev-gal, symmetric]]
lemmas rc2 = revcast-down-uu [THEN
  revcast-rev-ucast [symmetric, THEN trans, THEN word-rev-gal, symmetric]]

lemmas ucast-up =
  rc1 [simplified rev-shiftr [symmetric] revcast-ucast [symmetric]]
lemmas ucast-down =
  rc2 [simplified rev-shiftr revcast-ucast [symmetric]]

```

10.25.4 Slices

```

lemma slice1-no-bin [simp]:
  slice1 n (numeral w :: 'b word) = of-bl (takefill False n (bin-to-bl (len-of TYPE('b::len0))
    (numeral w)))
  by (simp add: slice1-def)

lemma slice-no-bin [simp]:
  slice n (numeral w :: 'b word) = of-bl (takefill False (len-of TYPE('b::len0) -
    n)
    (bin-to-bl (len-of TYPE('b::len0)) (numeral w)))
  by (simp add: slice-def word-size)

lemma slice1-0 [simp] : slice1 n 0 = 0
  unfolding slice1-def by simp

lemma slice-0 [simp] : slice n 0 = 0
  unfolding slice-def by auto

lemma slice-take': slice n w = of-bl (take (size w - n) (to-bl w))
  unfolding slice-def' slice1-def
  by (simp add : takefill-alt word-size)

lemmas slice-take = slice-take' [unfolded word-size]

```

— shiftr to a word of the same size is just slice, slice is just shiftr then ucast

```

lemmas shiftr-slice = trans [OF shiftr-bl [THEN meta-eq-to-obj-eq] slice-take [symmetric]]

```

```

lemma slice-shiftr: slice n w = ucast (w >> n)
  apply (unfold slice-take shiftr-bl)
  apply (rule ucast-of-bl-up [symmetric])
  apply (simp add: word-size)
done

```

```

lemma nth-slice: (slice n w :: 'a::len0 word) !! m = (w !! (m + n) ∧ m < len-of
  TYPE('a))
  by (simp add: slice-shiftr nth-ucast nth-shiftr)

```

```

lemma slice1-down-alt':

```

$sl = \text{slice1 } n \ w \implies fs = \text{size } sl \implies fs + k = n \implies$
 $\text{to-bl } sl = \text{takefill False } fs \ (\text{drop } k \ (\text{to-bl } w))$
by (*auto simp: slice1-def word-size of-bl-def uint-bl*
 $\text{word-ubin.eq-norm bl-bin-bl-rep-drop drop-takefill}$)

lemma *slice1-up-alt'*:

$sl = \text{slice1 } n \ w \implies fs = \text{size } sl \implies fs = n + k \implies$
 $\text{to-bl } sl = \text{takefill False } fs \ (\text{replicate } k \ \text{False} \ @ \ (\text{to-bl } w))$
apply (*unfold slice1-def word-size of-bl-def uint-bl*)
apply (*clarsimp simp: word-ubin.eq-norm bl-bin-bl-rep-drop takefill-append [symmetric]*)
apply (*rule-tac f = $\lambda k. \text{takefill False } (\text{len-of TYPE('a)})$*
 $(\text{replicate } k \ \text{False} \ @ \ \text{bin-to-bl } (\text{len-of TYPE('b)}) \ (\text{uint } w))$ **in** *arg-cong*)
apply *arith*
done

lemmas *sd1 = slice1-down-alt' [OF refl refl, unfolded word-size]*

lemmas *su1 = slice1-up-alt' [OF refl refl, unfolded word-size]*

lemmas *slice1-down-alt = le-add-diff-inverse [THEN sd1]*

lemmas *slice1-up-alt =*

le-add-diff-inverse [symmetric, THEN su1]

le-add-diff-inverse2 [symmetric, THEN su1]

lemma *ucast-slice1*: $\text{ucast } w = \text{slice1 } (\text{size } w) \ w$

by (*simp add: slice1-def ucast-bl takefill-same' word-size*)

lemma *ucast-slice*: $\text{ucast } w = \text{slice } 0 \ w$

by (*simp add: slice-def ucast-slice1*)

lemma *slice-id*: $\text{slice } 0 \ t = t$

by (*simp only: ucast-slice [symmetric] ucast-id*)

lemma *revcast-slice1* [*OF refl*]: $rc = \text{revcast } w \implies \text{slice1 } (\text{size } rc) \ w = rc$

by (*simp add: slice1-def revcast-def' word-size*)

lemma *slice1-tf-tf'*:

$\text{to-bl } (\text{slice1 } n \ w :: 'a::\text{len0 word}) =$
 $\text{rev } (\text{takefill False } (\text{len-of TYPE('a)}) \ (\text{rev } (\text{takefill False } n \ (\text{to-bl } w))))$
unfolding *slice1-def* **by** (*rule word-rev-tf*)

lemmas *slice1-tf-tf' = slice1-tf-tf' [THEN word-bl.Rep-inverse', symmetric]*

lemma *rev-slice1*:

$n + k = \text{len-of TYPE('a)} + \text{len-of TYPE('b)} \implies$

$\text{slice1 } n \ (\text{word-reverse } w :: 'b::\text{len0 word}) =$

$\text{word-reverse } (\text{slice1 } k \ w :: 'a::\text{len0 word})$

apply (*unfold word-reverse-def slice1-tf-tf'*)

apply (*rule word-bl.Rep-inverse'*)

apply (*rule rev-swap [THEN iffD1]*)

apply (*rule trans [symmetric]*)

```

apply (rule tf-rev)
apply (simp add: word-bl.Abs-inverse)
apply (simp add: word-bl.Abs-inverse)
done

```

```

lemma rev-slice:
   $n + k + \text{len-of } \text{TYPE}('a::\text{len0}) = \text{len-of } \text{TYPE}('b::\text{len0}) \implies$ 
   $\text{slice } n \text{ (word-reverse (} w::'b \text{ word))} = \text{word-reverse (slice } k \text{ } w :: 'a \text{ word)}$ 
apply (unfold slice-def word-size)
apply (rule rev-slice1)
apply arith
done

```

```

lemmas sym-notr =
  not-iff [THEN iffD2, THEN not-sym, THEN not-iff [THEN iffD1]]

```

— problem posed by TPHOLs referee: criterion for overflow of addition of signed integers

```

lemma soft-test:
   $(\text{sint } x + \text{sint } y = \text{sint } (x + y)) =$ 
   $(((((x + y) \text{ XOR } x) \text{ AND } ((x + y) \text{ XOR } y)) >> (\text{size } x - 1) = 0)$ 
for  $x \ y :: 'a::\text{len word}$ 
apply (unfold word-size)
apply (cases len-of TYPE('a), simp)
apply (subst msb-shift [THEN sym-notr])
apply (simp add: word-ops-msb)
apply (simp add: word-msb-sint)
apply safe
  apply simp-all
apply (unfold sint-word-ariths)
apply (unfold word-sbin.set-iff-norm [symmetric] sints-num)
apply safe
  apply (insert sint-range' [where  $x=x$ ])
  apply (insert sint-range' [where  $x=y$ ])
  defer
  apply (simp (no-asm), arith)
  apply (simp (no-asm), arith)
  defer
  defer
  apply (simp (no-asm), arith)
  apply (simp (no-asm), arith)
apply (rule notI [THEN notnotD],
  drule leI not-le-imp-less,
  drule sbintrunc-inc sbintrunc-dec,
  simp)+
done

```

10.26 Split and cat

lemmas *word-split-bin'* = *word-split-def*

lemmas *word-cat-bin'* = *word-cat-def*

lemma *word-rsplit-no*:

(*word-rsplit* (numeral *bin* :: 'b::len0 word) :: 'a word list) =
 map *word-of-int* (*bin-rsplit* (len-of TYPE('a::len))
 (len-of TYPE('b), *bintrunc* (len-of TYPE('b)) (numeral *bin*)))
 by (*simp add*: *word-rsplit-def word-ubin.eq-norm*)

lemmas *word-rsplit-no-cl* [*simp*] = *word-rsplit-no*
 [*unfolded bin-rsplittl-def bin-rsplit-l* [*symmetric*]]

lemma *test-bit-cat*:

wc = *word-cat* *a b* $\implies$ *wc* !! *n* = (*n* < size *wc* $\wedge$
 (if *n* < size *b* then *b* !! *n* else *a* !! (*n* - size *b*)))
apply (*auto simp*: *word-cat-bin'* *test-bit-bin* *word-ubin.eq-norm* *nth-bintr* *bin-nth-cat*
word-size)
apply (*erule bin-nth-uint-imp*)
done

lemma *word-cat-bl*: *word-cat* *a b* = *of-bl* (*to-bl* *a* @ *to-bl* *b*)
 by (*simp add*: *of-bl-def to-bl-def word-cat-bin'* *bl-to-bin-app-cat*)

lemma *of-bl-append*:

(*of-bl* (*xs* @ *ys*) :: 'a::len word) = *of-bl* *xs* * $2^{\text{length } ys}$ + *of-bl* *ys*
apply (*simp add*: *of-bl-def bl-to-bin-app-cat bin-cat-num*)
apply (*simp add*: *word-of-int-power-hom* [*symmetric*] *word-of-int-hom-syms*)
done

lemma *of-bl-False* [*simp*]: *of-bl* (*False* # *xs*) = *of-bl* *xs*
 by (*rule word-eqI*) (*auto simp*: *test-bit-of-bl* *nth-append*)

lemma *of-bl-True* [*simp*]: (*of-bl* (*True* # *xs*) :: 'a::len word) = $2^{\text{length } xs}$ + *of-bl*
xs
 by (*subst of-bl-append* [where *xs*=*True*, *simplified*]) (*simp add*: *word-1-bl*)

lemma *of-bl-Cons*: *of-bl* (*x* # *xs*) = *of-bool* *x* * $2^{\text{length } xs}$ + *of-bl* *xs*
 by (*cases x*) *simp-all*

lemma *split-uint-lem*: *bin-split* *n* (*uint* *w*) = (*a*, *b*) $\implies$

a = *bintrunc* (len-of TYPE('a) - *n*) *a* $\wedge$ *b* = *bintrunc* (len-of TYPE('a)) *b*
 for *w* :: 'a::len0 word
apply (*frule word-ubin.norm-Rep* [*THEN* *ssubst*])
apply (*drule bin-split-trunc1*)
apply (*drule sym* [*THEN* *trans*])
apply *assumption*
apply *safe*
done

```

lemma word-split-bl':
   $std = size\ c - size\ b \implies (word-split\ c = (a, b)) \implies$ 
   $(a = of-bl\ (take\ std\ (to-bl\ c)) \wedge b = of-bl\ (drop\ std\ (to-bl\ c)))$ 
  apply (unfold word-split-bin')
  apply safe
  defer
  apply (clarsimp split: prod.splits)
  apply hypsubst-thin
  apply (drule word-ubin.norm-Rep [THEN ssubst])
  apply (drule split-bintrunc)
  apply (simp add: of-bl-def bl2bin-drop word-size
    word-ubin.norm-eq-iff [symmetric] min-def del: word-ubin.norm-Rep)
  apply (clarsimp split: prod.splits)
  apply (frule split-wint-lem [THEN conjunct1])
  apply (unfold word-size)
  apply (cases len-of TYPE('a) ≥ len-of TYPE('b))
  defer
  apply simp
  apply (simp add : of-bl-def to-bl-def)
  apply (subst bin-split-take1 [symmetric])
  prefer 2
  apply assumption
  apply simp
  apply (erule thin-rl)
  apply (erule arg-cong [THEN trans])
  apply (simp add : word-ubin.norm-eq-iff [symmetric])
  done

```

```

lemma word-split-bl:  $std = size\ c - size\ b \implies$ 
   $(a = of-bl\ (take\ std\ (to-bl\ c)) \wedge b = of-bl\ (drop\ std\ (to-bl\ c))) \longleftrightarrow$ 
   $word-split\ c = (a, b)$ 
  apply (rule iffI)
  defer
  apply (erule (1) word-split-bl')
  apply (case-tac word-split c)
  apply (auto simp add: word-size)
  apply (frule word-split-bl' [rotated])
  apply (auto simp add: word-size)
  done

```

```

lemma word-split-bl-eq:
   $(word-split\ c :: ('c::len0\ word \times 'd::len0\ word)) =$ 
   $(of-bl\ (take\ (len-of\ TYPE('a)::len) - len-of\ TYPE('d)::len0))\ (to-bl\ c),$ 
   $of-bl\ (drop\ (len-of\ TYPE('a) - len-of\ TYPE('d))\ (to-bl\ c)))$ 
  for  $c :: 'a::len\ word$ 
  apply (rule word-split-bl [THEN iffD1])
  apply (unfold word-size)
  apply (rule refl conjI)+

```

done

— keep quantifiers for use in simplification

lemma *test-bit-split'*:

word-split $c = (a, b) \longrightarrow$

$(\forall n\ m.$

$b !! n = (n < \text{size } b \wedge c !! n) \wedge$

$a !! m = (m < \text{size } a \wedge c !! (m + \text{size } b)))$

apply (*unfold word-split-bin' test-bit-bin*)

apply (*clarify*)

apply (*clarsimp simp: word-ubin.eq-norm nth-bintr word-size split: prod.splits*)

apply (*drule bin-nth-split*)

apply *safe*

apply (*simp-all add: add.commute*)

apply (*erule bin-nth-uint-imp*)

done

lemma *test-bit-split*:

word-split $c = (a, b) \implies$

$(\forall n::\text{nat}. b !! n \longleftrightarrow n < \text{size } b \wedge c !! n) \wedge$

$(\forall m::\text{nat}. a !! m \longleftrightarrow m < \text{size } a \wedge c !! (m + \text{size } b))$

by (*simp add: test-bit-split'*)

lemma *test-bit-split-eq*:

word-split $c = (a, b) \longleftrightarrow$

$((\forall n::\text{nat}. b !! n = (n < \text{size } b \wedge c !! n)) \wedge$

$(\forall m::\text{nat}. a !! m = (m < \text{size } a \wedge c !! (m + \text{size } b))))$

apply (*rule-tac iffI*)

apply (*rule-tac conjI*)

apply (*erule test-bit-split [THEN conjunct1]*)

apply (*erule test-bit-split [THEN conjunct2]*)

apply (*case-tac word-split c*)

apply (*frule test-bit-split*)

apply (*erule trans*)

apply (*fastforce intro!: word-eqI simp add: word-size*)

done

— this odd result is analogous to *ucast-id*, result to the length given by the result type

lemma *word-cat-id*: *word-cat* $a\ b = b$

by (*simp add: word-cat-bin' word-ubin.inverse-norm*)

— limited hom result

lemma *word-cat-hom*:

len-of TYPE('a::len0) ≤ len-of TYPE('b::len0) + len-of TYPE('c::len0) $\implies$

(word-cat (word-of-int w :: 'b word) (b :: 'c word) :: 'a word) =

word-of-int (bin-cat w (size b) (uint b))

by (*auto simp: word-cat-def word-size word-ubin.norm-eq-iff [symmetric]*)

word-ubin.eq-norm bintr-cat min.absorb1)

lemma *word-cat-split-alt: size w ≤ size u + size v ⇒ word-split w = (u, v) ⇒ word-cat u v = w*
 apply (rule word-eqI)
 apply (drule test-bit-split)
 apply (clarsimp simp add: test-bit-cat word-size)
 apply safe
 apply arith
 done

lemmas *word-cat-split-size = sym [THEN [2] word-cat-split-alt [symmetric]]*

10.26.1 Split and slice

lemma *split-slices: word-split w = (u, v) ⇒ u = slice (size v) w ∧ v = slice 0 w*
 apply (drule test-bit-split)
 apply (rule conjI)
 apply (rule word-eqI, clarsimp simp: nth-slice word-size)+
 done

lemma *slice-cat1 [OF refl]:*
wc = word-cat a b ⇒ size wc ≥ size a + size b ⇒ slice (size b) wc = a
 apply safe
 apply (rule word-eqI)
 apply (simp add: nth-slice test-bit-cat word-size)
 done

lemmas *slice-cat2 = trans [OF slice-id word-cat-id]*

lemma *cat-slices:*
a = slice n c ⇒ b = slice 0 c ⇒ n = size b ⇒ size a + size b ≥ size c ⇒ word-cat a b = c
 apply safe
 apply (rule word-eqI)
 apply (simp add: nth-slice test-bit-cat word-size)
 apply safe
 apply arith
 done

lemma *word-split-cat-alt:*
w = word-cat u v ⇒ size u + size v ≤ size w ⇒ word-split w = (u, v)
 apply (case-tac word-split w)
 apply (rule trans, assumption)
 apply (drule test-bit-split)
 apply safe
 apply (rule word-eqI, clarsimp simp: test-bit-cat word-size)+
 done

lemmas *word-cat-bl-no-bin* [*simp*] =
word-cat-bl [**where** *a=numeral a* **and** *b=numeral b*, *unfolded to-bl-numeral*]
for *a b*

lemmas *word-split-bl-no-bin* [*simp*] =
word-split-bl-eq [**where** *c=numeral c*, *unfolded to-bl-numeral*] **for** *c*

This odd result arises from the fact that the statement of the result implies that the decoded words are of the same type, and therefore of the same length, as the original word.

lemma *word-rsplit-same*: *word-rsplit w* = [*w*]
by (*simp add*: *word-rsplit-def bin-rsplit-all*)

lemma *word-rsplit-empty-iff-size*: *word-rsplit w* = [] $\longleftrightarrow$ *size w* = 0
by (*simp add*: *word-rsplit-def bin-rsplit-def word-size bin-rsplit-aux-simp-alt Let-def split*: *prod.split*)

lemma *test-bit-rsplit*:
sw = *word-rsplit w* $\implies$ *m* < *size (hd sw)* $\implies$
k < *length sw* $\implies$ (*rev sw* ! *k*) !! *m* = *w* !! (*k* * *size (hd sw)* + *m*)
for *sw* :: '*a*::len *word list*
apply (*unfold word-rsplit-def word-test-bit-def*)
apply (*rule trans*)
apply (*rule-tac f* = $\lambda x. \text{bin-nth } x \text{ } m$ **in** *arg-cong*)
apply (*rule nth-map* [*symmetric*])
apply *simp*
apply (*rule bin-nth-rsplit*)
apply *simp-all*
apply (*simp add* : *word-size rev-map*)
apply (*rule trans*)
defer
apply (*rule map-ident* [*THEN fun-cong*])
apply (*rule refl* [*THEN map-cong*])
apply (*simp add* : *word-ubin.eq-norm*)
apply (*erule bin-rsplit-size-sign* [*OF len-gt-0 refl*])
done

lemma *word-rcat-bl*: *word-rcat wl* = *of-bl (concat (map to-bl wl))*
by (*auto simp*: *word-rcat-def to-bl-def' of-bl-def bin-rcat-bl*)

lemma *size-rcat-lem'*: *size (concat (map to-bl wl))* = *length wl* * *size (hd wl)*
by (*induct wl*) (*auto simp*: *word-size*)

lemmas *size-rcat-lem* = *size-rcat-lem'* [*unfolded word-size*]

lemmas *td-gal-lt-len* = *len-gt-0* [*THEN td-gal-lt*]

lemma *nth-rcat-lem*:
n < *length (wl::'a word list)* * *len-of TYPE('a::len)* $\implies$

```

    rev (concat (map to-bl wl)) ! n =
    rev (to-bl (rev wl ! (n div len-of TYPE('a)))) ! (n mod len-of TYPE('a))
  apply (induct wl)
  apply clarsimp
  apply (clarsimp simp add : nth-append size-rcat-lem)
  apply (simp (no-asm-use) only: mult-Suc [symmetric]
    td-gal-lt-len less-Suc-eq-le minus-div-mult-eq-mod [symmetric])
  apply clarsimp
  done

```

lemma *test-bit-rcat*:

```

  sw = size (hd wl)  $\implies$  rc = word-rcat wl  $\implies$  rc !! n =
  (n < size rc  $\wedge$  n div sw < size wl  $\wedge$  (rev wl) ! (n div sw) !! (n mod sw))
  for wl :: 'a::len word list
  apply (unfold word-rcat-bl word-size)
  apply (clarsimp simp add: test-bit-of-bl size-rcat-lem word-size td-gal-lt-len)
  apply safe
  apply (auto simp: test-bit-bl word-size td-gal-lt-len [THEN iffD2, THEN nth-rcat-lem])
  done

```

lemma *foldl-eq-foldr*: $\text{foldl } op + x \text{ } xs = \text{foldr } op + (x \# xs) \text{ } 0$

```

  for x :: 'a::comm-monoid-add
  by (induct xs arbitrary: x) (auto simp: add.assoc)

```

lemmas *test-bit-cong* = *arg-cong* [where $f = \text{test-bit}$, THEN *fun-cong*]

lemmas *test-bit-rsplit-alt* =

```

  trans [OF nth-rev-alt [THEN test-bit-cong]
    test-bit-rsplit [OF refl asm-rl diff-Suc-less]]

```

— lazy way of expressing that u and v , and su and sv , have same types

lemma *word-rsplit-len-indep* [OF refl refl refl refl]:

```

  [u,v] = p  $\implies$  [su,sv] = q  $\implies$  word-rsplit u = su  $\implies$ 
  word-rsplit v = sv  $\implies$  length su = length sv
  by (auto simp: word-rsplit-def bin-rsplit-len-indep)

```

lemma *length-word-rsplit-size*:

```

  n = len-of TYPE('a::len)  $\implies$ 
  length (word-rsplit w :: 'a word list)  $\leq m \iff$  size w  $\leq m * n$ 
  by (auto simp: word-rsplit-def word-size bin-rsplit-len-le)

```

lemmas *length-word-rsplit-lt-size* =

```

  length-word-rsplit-size [unfolded Not-eq-iff linorder-not-less [symmetric]]

```

lemma *length-word-rsplit-exp-size*:

```

  n = len-of TYPE('a::len)  $\implies$ 
  length (word-rsplit w :: 'a word list) = (size w + n - 1) div n
  by (auto simp: word-rsplit-def word-size bin-rsplit-len)

```

lemma *length-word-rsplit-even-size*:

$n = \text{len-of TYPE}('a::\text{len}) \implies \text{size } w = m * n \implies$
 $\text{length } (\text{word-rsplit } w :: 'a \text{ word list}) = m$
by (*auto simp: length-word-rsplit-exp-size given-quot-alt*)

lemmas *length-word-rsplit-exp-size' = refl [THEN length-word-rsplit-exp-size]*

lemmas *tdle = iffD2 [OF split-div-lemma refl, THEN conjunct1]*

lemmas *dtle = xtr4 [OF tdle mult.commute]*

lemma *word-rcat-rsplit*: $\text{word-rcat } (\text{word-rsplit } w) = w$

apply (*rule word-eqI*)

apply (*clarsimp simp: test-bit-rcat word-size*)

apply (*subst refl [THEN test-bit-rsplit]*)

apply (*simp-all add: word-size*

refl [THEN length-word-rsplit-size [simplified not-less [symmetric], simplified]])

apply *safe*

apply (*erule xtr7, rule len-gt-0 [THEN dtle]*) +

done

lemma *size-word-rsplit-rcat-size*:

$\text{word-rcat } ws = \text{frcw} \implies \text{size frcw} = \text{length } ws * \text{len-of TYPE}('a)$

$\implies \text{length } (\text{word-rsplit } \text{frcw} :: 'a \text{ word list}) = \text{length } ws$

for $ws :: 'a::\text{len} \text{ word list}$ **and** $\text{frcw} :: 'b::\text{len0} \text{ word}$

apply (*clarsimp simp: word-size length-word-rsplit-exp-size'*)

apply (*fast intro: given-quot-alt*)

done

lemma *msreus*:

$0 < n \implies (k * n + m) \text{ div } n = m \text{ div } n + k$

$(k * n + m) \text{ mod } n = m \text{ mod } n$

for $n :: \text{nat}$

by (*auto simp: add.commute*)

lemma *word-rsplit-rcat-size [OF refl]*:

$\text{word-rcat } ws = \text{frcw} \implies$

$\text{size frcw} = \text{length } ws * \text{len-of TYPE}('a) \implies \text{word-rsplit } \text{frcw} = ws$

for $ws :: 'a::\text{len} \text{ word list}$

apply (*frule size-word-rsplit-rcat-size, assumption*)

apply (*clarsimp simp add: word-size*)

apply (*rule nth-equalityI, assumption*)

apply *clarsimp*

apply (*rule word-eqI [rule-format]*)

apply (*rule trans*)

apply (*rule test-bit-rsplit-alt*)

apply (*clarsimp simp: word-size*) +

apply (*rule trans*)

apply (*rule test-bit-rcat [OF refl refl]*)

```

apply (simp add: word-size)
apply (subst nth-rev)
apply arith
apply (simp add: le0 [THEN [2] xtr7, THEN diff-Suc-less])
apply safe
apply (simp add: diff-mult-distrib)
apply (rule mpl-lem)
apply (cases size ws)
apply simp-all
done

```

10.27 Rotation

lemmas rotater-0' [simp] = rotater-def [where $n = 0$, simplified]

lemmas word-rot-defs = word-roti-def word-rotr-def word-rotl-def

lemma rotate-eq-mod: $m \bmod \text{length } xs = n \bmod \text{length } xs \implies \text{rotate } m \text{ } xs = \text{rotate } n \text{ } xs$

```

apply (rule box-equals)
defer
apply (rule rotate-conv-mod [symmetric])+
apply simp
done

```

lemmas rotate-eqs =

```

trans [OF rotate0 [THEN fun-cong] id-apply]
rotate-rotate [symmetric]
rotate-id
rotate-conv-mod
rotate-eq-mod

```

10.27.1 Rotation of list to right

lemma rotate1-rl': $\text{rotater1 } (l @ [a]) = a \# l$

```

by (cases l) (auto simp: rotater1-def)

```

lemma rotate1-rl [simp]: $\text{rotater1 } (\text{rotate1 } l) = l$

```

apply (unfold rotater1-def)
apply (cases l)
apply (case-tac [2] list)
apply auto
done

```

lemma rotate1-lr [simp]: $\text{rotate1 } (\text{rotater1 } l) = l$

```

by (cases l) (auto simp: rotater1-def)

```

lemma rotater1-rev': $\text{rotater1 } (\text{rev } xs) = \text{rev } (\text{rotate1 } xs)$

```

by (cases xs) (simp add: rotater1-def, simp add: rotate1-rl')

```

lemma *rotater-rev'*: $\text{rotater } n \text{ (rev } xs) = \text{rev (rotate } n \text{ } xs)$
by (*induct* *n*) (*auto simp: rotater-def intro: rotater1-rev'*)

lemma *rotater-rev*: $\text{rotater } n \text{ } ys = \text{rev (rotate } n \text{ (rev } ys))$
using *rotater-rev'* [**where** $xs = \text{rev } ys$] **by** *simp*

lemma *rotater-drop-take*:
 $\text{rotater } n \text{ } xs =$
 $\text{drop (length } xs - n \text{ mod length } xs) \text{ } xs \text{ @}$
 $\text{take (length } xs - n \text{ mod length } xs) \text{ } xs$
by (*auto simp: rotater-rev rotate-drop-take rev-take rev-drop*)

lemma *rotater-Suc* [*simp*]: $\text{rotater (Suc } n) \text{ } xs = \text{rotater1 (rotater } n \text{ } xs)$
unfolding *rotater-def* **by** *auto*

lemma *rotate-inv-plus* [*rule-format*] :
 $\forall k. k = m + n \longrightarrow \text{rotater } k \text{ (rotate } n \text{ } xs) = \text{rotater } m \text{ } xs \wedge$
 $\text{rotate } k \text{ (rotater } n \text{ } xs) = \text{rotate } m \text{ } xs \wedge$
 $\text{rotater } n \text{ (rotate } k \text{ } xs) = \text{rotate } m \text{ } xs \wedge$
 $\text{rotate } n \text{ (rotater } k \text{ } xs) = \text{rotater } m \text{ } xs$
by (*induct* *n*) (*auto simp: rotater-def rotate-def intro: funpow-swap1 [THEN trans]*)

lemmas *rotate-inv-rel* = *le-add-diff-inverse2* [*symmetric*, *THEN rotate-inv-plus*]

lemmas *rotate-inv-eq* = *order-refl* [*THEN rotate-inv-rel*, *simplified*]

lemmas *rotate-lr* [*simp*] = *rotate-inv-eq* [*THEN conjunct1*]

lemmas *rotate-rl* [*simp*] = *rotate-inv-eq* [*THEN conjunct2*, *THEN conjunct1*]

lemma *rotate-gal*: $\text{rotater } n \text{ } xs = ys \longleftrightarrow \text{rotate } n \text{ } ys = xs$
by *auto*

lemma *rotate-gal'*: $ys = \text{rotater } n \text{ } xs \longleftrightarrow xs = \text{rotate } n \text{ } ys$
by *auto*

lemma *length-rotater* [*simp*]: $\text{length (rotater } n \text{ } xs) = \text{length } xs$
by (*simp add : rotater-rev*)

lemma *restrict-to-left*: $x = y \implies x = z \longleftrightarrow y = z$
by *simp*

lemmas *rrs0* = *rotate-eqs* [*THEN restrict-to-left*,
simplified rotate-gal [*symmetric*] *rotate-gal'* [*symmetric*]]
lemmas *rrs1* = *rrs0* [*THEN refl* [*THEN rev-iffD1*]]
lemmas *rotater-eqs* = *rrs1* [*simplified length-rotater*]
lemmas *rotater-0* = *rotater-eqs* (1)
lemmas *rotater-add* = *rotater-eqs* (2)

10.27.2 map, map2, commuting with rotate(r)

lemma *butlast-map*: $xs \neq [] \implies \text{butlast } (\text{map } f \text{ } xs) = \text{map } f \text{ } (\text{butlast } xs)$
 by (induct xs) auto

lemma *rotater1-map*: $\text{rotater1 } (\text{map } f \text{ } xs) = \text{map } f \text{ } (\text{rotater1 } xs)$
 by (cases xs) (auto simp: rotater1-def last-map butlast-map)

lemma *rotater-map*: $\text{rotater } n \text{ } (\text{map } f \text{ } xs) = \text{map } f \text{ } (\text{rotater } n \text{ } xs)$
 by (induct n) (auto simp: rotater-def rotater1-map)

lemma *but-last-zip* [rule-format] :
 $\forall ys. \text{length } xs = \text{length } ys \longrightarrow xs \neq [] \longrightarrow$
 $\text{last } (\text{zip } xs \text{ } ys) = (\text{last } xs, \text{last } ys) \wedge$
 $\text{butlast } (\text{zip } xs \text{ } ys) = \text{zip } (\text{butlast } xs) \text{ } (\text{butlast } ys)$
 apply (induct xs)
 apply auto
 apply ((case-tac ys, auto simp: neq-Nil-conv)[1])+
 done

lemma *but-last-map2* [rule-format] :
 $\forall ys. \text{length } xs = \text{length } ys \longrightarrow xs \neq [] \longrightarrow$
 $\text{last } (\text{map2 } f \text{ } xs \text{ } ys) = f \text{ } (\text{last } xs) \text{ } (\text{last } ys) \wedge$
 $\text{butlast } (\text{map2 } f \text{ } xs \text{ } ys) = \text{map2 } f \text{ } (\text{butlast } xs) \text{ } (\text{butlast } ys)$
 apply (induct xs)
 apply auto
 apply (unfold map2-def)
 apply ((case-tac ys, auto simp: neq-Nil-conv)[1])+
 done

lemma *rotater1-zip*:
 $\text{length } xs = \text{length } ys \implies$
 $\text{rotater1 } (\text{zip } xs \text{ } ys) = \text{zip } (\text{rotater1 } xs) \text{ } (\text{rotater1 } ys)$
 apply (unfold rotater1-def)
 apply (cases xs)
 apply auto
 apply ((case-tac ys, auto simp: neq-Nil-conv but-last-zip)[1])+
 done

lemma *rotater1-map2*:
 $\text{length } xs = \text{length } ys \implies$
 $\text{rotater1 } (\text{map2 } f \text{ } xs \text{ } ys) = \text{map2 } f \text{ } (\text{rotater1 } xs) \text{ } (\text{rotater1 } ys)$
 by (simp add: map2-def rotater1-map rotater1-zip)

lemmas *lrth* =
 box-equals [OF asm-rl length-rotater [symmetric]
 length-rotater [symmetric],
 THEN rotater1-map2]

lemma *rotater-map2*:

$length\ xs = length\ ys \implies$
 $rotater\ n\ (map2\ f\ xs\ ys) = map2\ f\ (rotater\ n\ xs)\ (rotater\ n\ ys)$
by (induct n) (auto intro!: lpth)

lemma rotate1-map2:
 $length\ xs = length\ ys \implies$
 $rotate1\ (map2\ f\ xs\ ys) = map2\ f\ (rotate1\ xs)\ (rotate1\ ys)$
by (cases xs; cases ys) (auto simp: map2-def)

lemmas lth = box-equals [OF asm-rl length-rotate [symmetric]
 length-rotate [symmetric], THEN rotate1-map2]

lemma rotate-map2:
 $length\ xs = length\ ys \implies$
 $rotate\ n\ (map2\ f\ xs\ ys) = map2\ f\ (rotate\ n\ xs)\ (rotate\ n\ ys)$
by (induct n) (auto intro!: lth)

— corresponding equalities for word rotation

lemma to-bl-rotl: $to-bl\ (word-rotl\ n\ w) = rotate\ n\ (to-bl\ w)$
by (simp add: word-bl.Abs-inverse' word-rotl-def)

lemmas blrs0 = rotate-egs [THEN to-bl-rotl [THEN trans]]

lemmas word-rotl-egs =
 blrs0 [simplified word-bl-Rep' word-bl.Rep-inject to-bl-rotl [symmetric]]

lemma to-bl-rotr: $to-bl\ (word-rotr\ n\ w) = rotater\ n\ (to-bl\ w)$
by (simp add: word-bl.Abs-inverse' word-rotr-def)

lemmas brrs0 = rotater-egs [THEN to-bl-rotr [THEN trans]]

lemmas word-rotr-egs =
 brrs0 [simplified word-bl-Rep' word-bl.Rep-inject to-bl-rotr [symmetric]]

declare word-rotr-egs (1) [simp]
declare word-rotl-egs (1) [simp]

lemma word-rot-rl [simp]: $word-rotl\ k\ (word-rotr\ k\ v) = v$
and word-rot-lr [simp]: $word-rotr\ k\ (word-rotl\ k\ v) = v$
by (auto simp add: to-bl-rotr to-bl-rotl word-bl.Rep-inject [symmetric])

lemma word-rot-gal: $word-rotr\ n\ v = w \iff word-rotl\ n\ w = v$
and word-rot-gal': $w = word-rotr\ n\ v \iff v = word-rotl\ n\ w$
by (auto simp: to-bl-rotr to-bl-rotl word-bl.Rep-inject [symmetric] dest: sym)

lemma word-rotr-rev: $word-rotr\ n\ w = word-reverse\ (word-rotl\ n\ (word-reverse\ w))$

by (*simp only: word-bl.Rep-inject [symmetric] to-bl-word-rev to-bl-rotr to-bl-rotl rotater-rev*)

lemma *word-roti-0* [*simp*]: *word-roti 0 w = w*
by (*auto simp: word-rot-defs*)

lemmas *abl-cong = arg-cong [where f = of-bl]*

lemma *word-roti-add*: *word-roti (m + n) w = word-roti m (word-roti n w)*

proof –

have *rotater-eq-lem*: $\bigwedge m\ n\ xs.\ m = n \implies \text{rotater } m\ xs = \text{rotater } n\ xs$
by *auto*

have *rotate-eq-lem*: $\bigwedge m\ n\ xs.\ m = n \implies \text{rotate } m\ xs = \text{rotate } n\ xs$
by *auto*

note *rpts [symmetric]* =
rotate-inv-plus [THEN conjunct1]
rotate-inv-plus [THEN conjunct2, THEN conjunct1]
rotate-inv-plus [THEN conjunct2, THEN conjunct2, THEN conjunct1]
rotate-inv-plus [THEN conjunct2, THEN conjunct2, THEN conjunct2]

note *rrp = trans [symmetric, OF rotate-rotate rotate-eq-lem]*

note *rrrp = trans [symmetric, OF rotater-add [symmetric] rotater-eq-lem]*

show *?thesis*

apply (*unfold word-rot-defs*)
apply (*simp only: split: if-split*)
apply (*safe intro!: abl-cong*)
apply (*simp-all only: to-bl-rotl [THEN word-bl.Rep-inverse']*
to-bl-rotl
to-bl-rotr [THEN word-bl.Rep-inverse']
to-bl-rotr)
apply (*rule rrp rrrp rpts,*
simp add: nat-add-distrib [symmetric]
nat-diff-distrib [symmetric])+

done

qed

lemma *word-roti-conv-mod'*: *word-roti n w = word-roti (n mod int (size w)) w*

apply (*unfold word-rot-defs*)
apply (*cut-tac y=size w in gt-or-eq-0*)
apply (*erule disjE*)
apply *simp-all*
apply (*safe intro!: abl-cong*)
apply (*rule rotater-egs*)
apply (*simp add: word-size nat-mod-distrib*)
apply (*simp add: rotater-add [symmetric] rotate-gal [symmetric]*)
apply (*rule rotater-egs*)

```

apply (simp add: word-size nat-mod-distrib)
apply (rule of-nat-eq-0-iff [THEN iffD1])
apply (auto simp add: not-le mod-eq-0-iff-dvd zdvd-int nat-add-distrib [symmetric])
using mod-mod-trivial mod-eq-dvd-iff
apply blast
done

```

```

lemmas word-roti-conv-mod = word-roti-conv-mod' [unfolded word-size]

```

10.27.3 “Word rotation commutes with bit-wise operations

```

locale word-rotate
begin

```

```

lemmas word-rot-defs' = to-bl-rotl to-bl-rotr

```

```

lemmas blwl-syms [symmetric] = bl-word-not bl-word-and bl-word-or bl-word-xor

```

```

lemmas lbl-lbl = trans [OF word-bl-Rep' word-bl-Rep' [symmetric]]

```

```

lemmas ths-map2 [OF lbl-lbl] = rotate-map2 rotater-map2

```

```

lemmas ths-map [where xs = to-bl v] = rotate-map rotater-map for v

```

```

lemmas th1s [simplified word-rot-defs' [symmetric]] = ths-map2 ths-map

```

```

lemma word-rot-logs:

```

```

  word-rotl n (NOT v) = NOT word-rotl n v
  word-rotr n (NOT v) = NOT word-rotr n v
  word-rotl n (x AND y) = word-rotl n x AND word-rotl n y
  word-rotr n (x AND y) = word-rotr n x AND word-rotr n y
  word-rotl n (x OR y) = word-rotl n x OR word-rotl n y
  word-rotr n (x OR y) = word-rotr n x OR word-rotr n y
  word-rotl n (x XOR y) = word-rotl n x XOR word-rotl n y
  word-rotr n (x XOR y) = word-rotr n x XOR word-rotr n y
by (rule word-bl.Rep-eqD,
    rule word-rot-defs' [THEN trans],
    simp only: blwl-syms [symmetric],
    rule th1s [THEN trans],
    rule refl)+

```

```

end

```

```

lemmas word-rot-logs = word-rotate.word-rot-logs

```

```

lemmas bl-word-rotl-dt = trans [OF to-bl-rotl rotate-drop-take,
  simplified word-bl-Rep']

```

```

lemmas bl-word-rotr-dt = trans [OF to-bl-rotr rotater-drop-take,
  simplified word-bl-Rep']

```

```

lemma bl-word-roti-dt':
   $n = \text{nat } ((- i) \bmod \text{int } (\text{size } (w :: 'a::\text{len word}))) \implies$ 
   $\text{to-bl } (\text{word-roti } i \ w) = \text{drop } n \ (\text{to-bl } w) \ @ \ \text{take } n \ (\text{to-bl } w)$ 
  apply (unfold word-roti-def)
  apply (simp add: bl-word-rotl-dt bl-word-rotr-dt word-size)
  apply safe
  apply (simp add: zmod-zminus1-eq-if)
  apply safe
  apply (simp add: nat-mult-distrib)
  apply (simp add: nat-diff-distrib [OF pos-mod-sign pos-mod-conj
    [THEN conjunct2, THEN order-less-imp-le]]
    nat-mod-distrib)
  apply (simp add: nat-mod-distrib)
  done

lemmas bl-word-roti-dt = bl-word-roti-dt' [unfolded word-size]

lemmas word-rotl-dt = bl-word-rotl-dt [THEN word-bl.Rep-inverse' [symmetric]]
lemmas word-rotr-dt = bl-word-rotr-dt [THEN word-bl.Rep-inverse' [symmetric]]
lemmas word-roti-dt = bl-word-roti-dt [THEN word-bl.Rep-inverse' [symmetric]]

lemma word-rotx-0 [simp] : word-rotr i 0 = 0  $\wedge$  word-rotl i 0 = 0
  by (simp add: word-rotr-dt word-rotl-dt replicate-add [symmetric])

lemma word-roti-0' [simp] : word-roti n 0 = 0
  by (auto simp: word-roti-def)

lemmas word-rotr-dt-no-bin' [simp] =
  word-rotr-dt [where w=numeral w, unfolded to-bl-numeral] for w

lemmas word-rotl-dt-no-bin' [simp] =
  word-rotl-dt [where w=numeral w, unfolded to-bl-numeral] for w

declare word-roti-def [simp]

```

10.28 Maximum machine word

```

lemma word-int-cases:
  fixes  $x :: 'a::\text{len0 word}$ 
  obtains  $n$  where  $x = \text{word-of-int } n$  and  $0 \leq n$  and  $n < 2^{\text{len-of TYPE('a)}}$ 
  by (cases x rule: word-uint.Abs-cases) (simp add: uints-num)

lemma word-nat-cases [cases type: word]:
  fixes  $x :: 'a::\text{len word}$ 
  obtains  $n$  where  $x = \text{of-nat } n$  and  $n < 2^{\text{len-of TYPE('a)}}$ 
  by (cases x rule: word-unat.Abs-cases) (simp add: unats-def)

```

lemma *max-word-eq*: $(\text{max-word}::'a::\text{len word}) = 2^{\text{len-of TYPE}('a)} - 1$
by (*simp add: max-word-def word-of-int-hom-syms word-of-int-2p*)

lemma *max-word-max* [*simp,intro!*]: $n \leq \text{max-word}$
by (*cases n rule: word-int-cases*)
(simp add: max-word-def word-le-def int-word-uint mod-pos-pos-trivial del: minus-mod-self1)

lemma *word-of-int-2p-len*: $\text{word-of-int } (2^{\text{len-of TYPE}('a)}) = (0::'a::\text{len0 word})$
by (*subst word-uint.Abs-norm [symmetric]*) *simp*

lemma *word-pow-0*: $(2::'a::\text{len word})^{\text{len-of TYPE}('a)} = 0$

proof –

have $\text{word-of-int } (2^{\text{len-of TYPE}('a)}) = (0::'a \text{ word})$

by (*rule word-of-int-2p-len*)

then show *?thesis* **by** (*simp add: word-of-int-2p*)

qed

lemma *max-word-wrap*: $x + 1 = 0 \implies x = \text{max-word}$

apply (*simp add: max-word-eq*)

apply *uint-arith*

apply (*auto simp: word-pow-0*)

done

lemma *max-word-minus*: $\text{max-word} = (-1::'a::\text{len word})$

proof –

have $-1 + 1 = (0::'a \text{ word})$

by *simp*

then show *?thesis*

by (*rule max-word-wrap [symmetric]*)

qed

lemma *max-word-bl* [*simp*]: $\text{to-bl } (\text{max-word}::'a::\text{len word}) = \text{replicate } (\text{len-of TYPE}('a))$

True

by (*subst max-word-minus to-bl-n1*) *+* *simp*

lemma *max-test-bit* [*simp*]: $(\text{max-word}::'a::\text{len word}) !! n \longleftrightarrow n < \text{len-of TYPE}('a)$

by (*auto simp: test-bit-bl word-size*)

lemma *word-and-max* [*simp*]: $x \text{ AND } \text{max-word} = x$

by (*rule word-eqI*) (*simp add: word-ops-nth-size word-size*)

lemma *word-or-max* [*simp*]: $x \text{ OR } \text{max-word} = \text{max-word}$

by (*rule word-eqI*) (*simp add: word-ops-nth-size word-size*)

lemma *word-ao-dist2*: $x \text{ AND } (y \text{ OR } z) = x \text{ AND } y \text{ OR } x \text{ AND } z$

for $x \ y \ z :: 'a::\text{len0 word}$

by (*rule word-eqI*) (*auto simp add: word-ops-nth-size word-size*)

lemma *word-oa-dist2*: $x \text{ OR } y \text{ AND } z = (x \text{ OR } y) \text{ AND } (x \text{ OR } z)$
for $x \ y \ z :: 'a::\text{len0 word}$
by (rule *word-eqI*) (auto simp add: *word-ops-nth-size word-size*)

lemma *word-and-not* [simp]: $x \text{ AND NOT } x = 0$
for $x :: 'a::\text{len0 word}$
by (rule *word-eqI*) (auto simp add: *word-ops-nth-size word-size*)

lemma *word-or-not* [simp]: $x \text{ OR NOT } x = \text{max-word}$
by (rule *word-eqI*) (auto simp add: *word-ops-nth-size word-size*)

lemma *word-boolean*: *boolean* (*op AND*) (*op OR*) *bitNOT 0 max-word*
apply (rule *boolean.intro*)
apply (rule *word-bw-assocs*)
apply (rule *word-bw-assocs*)
apply (rule *word-bw-comms*)
apply (rule *word-bw-comms*)
apply (rule *word-ao-dist2*)
apply (rule *word-oa-dist2*)
apply (rule *word-and-max*)
apply (rule *word-log-esimps*)
apply (rule *word-and-not*)
apply (rule *word-or-not*)
done

interpretation *word-bool-alg*: *boolean op AND op OR bitNOT 0 max-word*
by (rule *word-boolean*)

lemma *word-xor-and-or*: $x \text{ XOR } y = x \text{ AND NOT } y \text{ OR NOT } x \text{ AND } y$
for $x \ y :: 'a::\text{len0 word}$
by (rule *word-eqI*) (auto simp add: *word-ops-nth-size word-size*)

interpretation *word-bool-alg*: *boolean-xor op AND op OR bitNOT 0 max-word op XOR*
apply (rule *boolean-xor.intro*)
apply (rule *word-boolean*)
apply (rule *boolean-xor-axioms.intro*)
apply (rule *word-xor-and-or*)
done

lemma *shiftr-x-0* [iff]: $x \gg 0 = x$
for $x :: 'a::\text{len0 word}$
by (simp add: *shiftr-bl*)

lemma *shiftr-x-0* [simp]: $x \ll 0 = x$
for $x :: 'a::\text{len word}$
by (simp add: *shiftr-t2n*)

lemma *shiftr-1* [*simp*]: $(1::'a::\text{len word}) << n = 2^n$
by (*simp add: shiftr-t2n*)

lemma *uint-lt-0* [*simp*]: $\text{uint } x < 0 = \text{False}$
by (*simp add: linorder-not-less*)

lemma *shiftr-1* [*simp*]: $\text{shiftr1 } (1::'a::\text{len word}) = 0$
unfolding *shiftr1-def* **by** *simp*

lemma *shiftr-1* [*simp*]: $(1::'a::\text{len word}) >> n = (\text{if } n = 0 \text{ then } 1 \text{ else } 0)$
by (*induct n*) (*auto simp: shiftr-def*)

lemma *word-less-1* [*simp*]: $x < 1 \iff x = 0$
for $x :: 'a::\text{len word}$
by (*simp add: word-less-nat-alt unat-0-iff*)

lemma *to-bl-mask*:
 $\text{to-bl } (\text{mask } n :: 'a::\text{len word}) =$
 $\text{replicate } (\text{len-of TYPE('a)} - n) \text{ False } @$
 $\text{replicate } (\min (\text{len-of TYPE('a)}) n) \text{ True}$
by (*simp add: mask-bl word-rep-drop min-def*)

lemma *map-replicate-True*:
 $n = \text{length } xs \implies$
 $\text{map } (\lambda(x,y). x \wedge y) (\text{zip } xs (\text{replicate } n \text{ True})) = xs$
by (*induct xs arbitrary: n*) *auto*

lemma *map-replicate-False*:
 $n = \text{length } xs \implies \text{map } (\lambda(x,y). x \wedge y)$
 $(\text{zip } xs (\text{replicate } n \text{ False})) = \text{replicate } n \text{ False}$
by (*induct xs arbitrary: n*) *auto*

lemma *bl-and-mask*:
fixes $w :: 'a::\text{len word}$
and $n :: \text{nat}$
defines $n' \equiv \text{len-of TYPE('a)} - n$
shows $\text{to-bl } (w \text{ AND } \text{mask } n) = \text{replicate } n' \text{ False } @ \text{drop } n' (\text{to-bl } w)$
proof –
note [*simp*] = *map-replicate-True map-replicate-False*
have $\text{to-bl } (w \text{ AND } \text{mask } n) = \text{map2 } \text{op } \wedge (\text{to-bl } w) (\text{to-bl } (\text{mask } n::'a::\text{len word}))$
by (*simp add: bl-word-and*)
also have $\text{to-bl } w = \text{take } n' (\text{to-bl } w) @ \text{drop } n' (\text{to-bl } w)$
by *simp*
also have $\text{map2 } \text{op } \wedge \dots (\text{to-bl } (\text{mask } n::'a::\text{len word})) =$
 $\text{replicate } n' \text{ False } @ \text{drop } n' (\text{to-bl } w)$
unfolding *to-bl-mask n'-def map2-def* **by** (*subst zip-append*) *auto*
finally show *?thesis* .
qed

lemma *drop-rev-takefill*:

length xs ≤ n ⇒
drop (n - length xs) (rev (takefill False n (rev xs))) = xs
by (*simp add: takefill-alt rev-take*)

lemma *map-nth-0* [*simp*]: *map (op !! (0::'a::len0 word)) xs = replicate (length xs) False*

by (*induct xs*) *auto*

lemma *uint-plus-if-size*:

uint (x + y) =
(if uint x + uint y < 2^{size x}
then uint x + uint y
else uint x + uint y - 2^{size x})
by (*simp add: word-arith-wis int-word-uint mod-add-if-z word-size*)

lemma *unat-plus-if-size*:

unat (x + y) =
(if unat x + unat y < 2^{size x}
then unat x + unat y
else unat x + unat y - 2^{size x})
for *x y :: 'a::len word*
apply (*subst word-arith-nat-defs*)
apply (*subst unat-of-nat*)
apply (*simp add: mod-nat-add word-size*)
done

lemma *word-neq-0-conv*: *w ≠ 0 ⟷ 0 < w*

for *w :: 'a::len word*
by (*simp add: word-gt-0*)

lemma *max-lt*: *unat (max a b div c) = unat (max a b) div unat c*

for *c :: 'a::len word*
by (*fact unat-div*)

lemma *uint-sub-if-size*:

uint (x - y) =
(if uint y ≤ uint x
then uint x - uint y
else uint x - uint y + 2^{size x})
by (*simp add: word-arith-wis int-word-uint mod-sub-if-z word-size*)

lemma *unat-sub*: *b ≤ a ⇒ unat (a - b) = unat a - unat b*

by (*simp add: unat-def uint-sub-if-size word-le-def nat-diff-distrib*)

lemmas *word-less-sub1-numberof* [*simp*] = *word-less-sub1* [*of numeral w*] **for** *w*

lemmas *word-le-sub1-numberof* [*simp*] = *word-le-sub1* [*of numeral w*] **for** *w*

lemma *word-of-int-minus*: *word-of-int (2^{len-of TYPE('a)} - i) = (word-of-int*

$(-i)::'a::\text{len word}$

proof –

have *: $2^{\text{len-of TYPE}('a)} - i = -i + 2^{\text{len-of TYPE}('a)}$

by *simp*

show ?thesis

apply (*subst* *)

apply (*subst word-uint.Abs-norm [symmetric], subst mod-add-self2*)

apply *simp*

done

qed

lemmas *word-of-int-inj* =

word-uint.Abs-inject [unfolded uints-num, simplified]

lemma *word-le-less-eq*: $x \leq y \longleftrightarrow x = y \vee x < y$

for $x\ y :: 'z::\text{len word}$

by (*auto simp add: order-class.le-less*)

lemma *mod-plus-cong*:

fixes $b\ b' :: \text{int}$

assumes $1: b = b'$

and $2: x \bmod b' = x' \bmod b'$

and $3: y \bmod b' = y' \bmod b'$

and $4: x' + y' = z'$

shows $(x + y) \bmod b = z' \bmod b'$

proof –

from $1\ 2[\text{symmetric}]\ 3[\text{symmetric}]\ \text{have}$ $(x + y) \bmod b = (x' \bmod b' + y' \bmod b') \bmod b'$

by (*simp add: mod-add-eq*)

also have $\dots = (x' + y') \bmod b'$

by (*simp add: mod-add-eq*)

finally show ?thesis

by (*simp add: 4*)

qed

lemma *mod-minus-cong*:

fixes $b\ b' :: \text{int}$

assumes $b = b'$

and $x \bmod b' = x' \bmod b'$

and $y \bmod b' = y' \bmod b'$

and $x' - y' = z'$

shows $(x - y) \bmod b = z' \bmod b'$

using *assms [symmetric]* **by** (*auto intro: mod-diff-cong*)

lemma *word-induct-less*: $P\ 0 \implies (\bigwedge n. n < m \implies P\ n \implies P\ (1 + n)) \implies P\ m$

for $P :: 'a::\text{len word} \Rightarrow \text{bool}$

apply (*cases m*)

apply *atomize*

apply (*erule rev-mp*)**+**

```

apply (rule-tac x=m in spec)
apply (induct-tac n)
apply simp
apply clarsimp
apply (erule impE)
apply clarsimp
apply (erule-tac x=n in allE)
apply (erule impE)
apply (simp add: unat-arith-simps)
apply (clarsimp simp: unat-of-nat)
apply simp
apply (erule-tac x=of-nat na in allE)
apply (erule impE)
apply (simp add: unat-arith-simps)
apply (clarsimp simp: unat-of-nat)
apply simp
done

```

```

lemma word-induct:  $P\ 0 \implies (\bigwedge n. P\ n \implies P\ (1 + n)) \implies P\ m$ 
for  $P :: 'a::len\ word \Rightarrow bool$ 
by (erule word-induct-less) simp

```

```

lemma word-induct2 [induct type]:  $P\ 0 \implies (\bigwedge n. 1 + n \neq 0 \implies P\ n \implies P\ (1 + n)) \implies P\ n$ 
for  $P :: 'b::len\ word \Rightarrow bool$ 
apply (rule word-induct)
apply simp
apply (case-tac 1 + n = 0)
apply auto
done

```

10.29 Recursion combinator for words

```

definition word-rec ::  $'a \Rightarrow ('b::len\ word \Rightarrow 'a \Rightarrow 'a) \Rightarrow 'b\ word \Rightarrow 'a$ 
where word-rec forZero forSuc n = rec-nat forZero (forSuc  $\circ$  of-nat) (unat n)

```

```

lemma word-rec-0: word-rec z s 0 = z
by (simp add: word-rec-def)

```

```

lemma word-rec-Suc:  $1 + n \neq 0 \implies word-rec\ z\ s\ (1 + n) = s\ n\ (word-rec\ z\ s\ n)$ 
for  $n :: 'a::len\ word$ 
apply (simp add: word-rec-def unat-word-ariths)
apply (subst nat-mod-eq')
apply (metis Suc-eq-plus1-left Suc-lessI of-nat-2p unat-1 unat-lt2p word-arith-nat-add)
apply simp
done

```

```

lemma word-rec-Pred:  $n \neq 0 \implies word-rec\ z\ s\ n = s\ (n - 1)\ (word-rec\ z\ s\ (n - 1))$ 

```

```

apply (rule subst[where  $t=n$  and  $s=1 + (n - 1)$ ])
apply simp
apply (subst word-rec-Suc)
apply simp
apply simp
done

lemma word-rec-in:  $f \text{ (word-rec } z \text{ (}\lambda\cdot. f\text{) } n) = \text{word-rec (} f \text{ } z \text{) (}\lambda\cdot. f\text{) } n$ 
  by (induct  $n$ ) (simp-all add: word-rec-0 word-rec-Suc)

lemma word-rec-in2:  $f \text{ } n \text{ (word-rec } z \text{ } f \text{ } n) = \text{word-rec (} f \text{ } 0 \text{ } z \text{) (} f \circ op + 1 \text{) } n$ 
  by (induct  $n$ ) (simp-all add: word-rec-0 word-rec-Suc)

lemma word-rec-twice:
   $m \leq n \implies \text{word-rec } z \text{ } f \text{ } n = \text{word-rec (word-rec } z \text{ } f \text{ (} n - m \text{)) (} f \circ op + (n - m) \text{) } m$ 
  apply (erule rev-mp)
  apply (rule-tac  $x=z$  in spec)
  apply (rule-tac  $x=f$  in spec)
  apply (induct  $n$ )
  apply (simp add: word-rec-0)
  apply clarsimp
  apply (rule-tac  $t=1 + n - m$  and  $s=1 + (n - m)$  in subst)
  apply simp
  apply (case-tac  $1 + (n - m) = 0$ )
  apply (simp add: word-rec-0)
  apply (rule-tac  $f = \text{word-rec } a \text{ } b$  for  $a \text{ } b$  in arg-cong)
  apply (rule-tac  $t=m$  and  $s=m + (1 + (n - m))$  in subst)
  apply simp
  apply (simp (no-asm-use))
  apply (simp add: word-rec-Suc word-rec-in2)
  apply (erule impE)
  apply uint-arith
  apply (drule-tac  $x=x \circ op + 1$  in spec)
  apply (drule-tac  $x=x \text{ } 0 \text{ } xa$  in spec)
  apply simp
  apply (rule-tac  $t=\lambda a. x \text{ (} 1 + (n - m + a) \text{) and } s=\lambda a. x \text{ (} 1 + (n - m) + a \text{) in$  subst)
  apply (clarsimp simp add: fun-eq-iff)
  apply (rule-tac  $t=(1 + (n - m + xb))$  and  $s=1 + (n - m) + xb$  in subst)
  apply simp
  apply (rule refl)
  apply (rule refl)
done

lemma word-rec-id:  $\text{word-rec } z \text{ (}\lambda\cdot. id\text{) } n = z$ 
  by (induct  $n$ ) (auto simp add: word-rec-0 word-rec-Suc)

lemma word-rec-id-eq:  $\forall m < n. f \text{ } m = id \implies \text{word-rec } z \text{ } f \text{ } n = z$ 

```

```

apply (erule rev-mp)
apply (induct n)
apply (auto simp add: word-rec-0 word-rec-Suc)
apply (drule spec, erule mp)
apply uint-arith
apply (drule-tac x=n in spec, erule impE)
apply uint-arith
apply simp
done

lemma word-rec-max:
 $\forall m \geq n. m \neq -1 \longrightarrow f\ m = id \implies word-rec\ z\ f\ (-1) = word-rec\ z\ f\ n$ 
apply (subst word-rec-twice[where n=-1 and m=-1 - n])
apply simp
apply simp
apply (rule word-rec-id-eq)
apply clarsimp
apply (drule spec, rule mp, erule mp)
apply (rule word-plus-mono-right2[OF - order-less-imp-le])
prefer 2
apply assumption
apply simp
apply (erule contrapos-pn)
apply simp
apply (drule arg-cong[where f= $\lambda x. x - n$ ])
apply simp
done

lemma unatSuc:  $1 + n \neq 0 \implies unat\ (1 + n) = Suc\ (unat\ n)$ 
for n :: 'a::len word
by unat-arith

declare bin-to-bl-def [simp]

ML-file Tools/word-lib.ML
ML-file Tools/smt-word.ML

hide-const (open) Word

end

theory WordBitwise
imports Word
begin

Helper constants used in defining addition

definition xor3 :: bool  $\Rightarrow$  bool  $\Rightarrow$  bool  $\Rightarrow$  bool
where xor3 a b c = (a = (b = c))

```

definition *carry* :: *bool* $\Rightarrow$ *bool* $\Rightarrow$ *bool* $\Rightarrow$ *bool*
where *carry* *a* *b* *c* = ((*a* $\wedge$ (*b* $\vee$ *c*)) $\vee$ (*b* $\wedge$ *c*))

lemma *carry-simps*:

carry *True* *a* *b* = (*a* $\vee$ *b*)
carry *a* *True* *b* = (*a* $\vee$ *b*)
carry *a* *b* *True* = (*a* $\vee$ *b*)
carry *False* *a* *b* = (*a* $\wedge$ *b*)
carry *a* *False* *b* = (*a* $\wedge$ *b*)
carry *a* *b* *False* = (*a* $\wedge$ *b*)
by (*auto simp add: carry-def*)

lemma *xor3-simps*:

xor3 *True* *a* *b* = (*a* = *b*)
xor3 *a* *True* *b* = (*a* = *b*)
xor3 *a* *b* *True* = (*a* = *b*)
xor3 *False* *a* *b* = (*a* $\neq$ *b*)
xor3 *a* *False* *b* = (*a* $\neq$ *b*)
xor3 *a* *b* *False* = (*a* $\neq$ *b*)
by (*simp-all add: xor3-def*)

Breaking up word equalities into equalities on their bit lists. Equalities are generated and manipulated in the reverse order to *to-bl*.

lemma *word-eq-rbl-eq*: *x* = *y* $\longleftrightarrow$ *rev* (*to-bl* *x*) = *rev* (*to-bl* *y*)
by *simp*

lemma *rbl-word-or*: *rev* (*to-bl* (*x* *OR* *y*)) = *map2* *op* $\vee$ (*rev* (*to-bl* *x*)) (*rev* (*to-bl* *y*))
by (*simp add: map2-def zip-rev bl-word-or rev-map*)

lemma *rbl-word-and*: *rev* (*to-bl* (*x* *AND* *y*)) = *map2* *op* $\wedge$ (*rev* (*to-bl* *x*)) (*rev* (*to-bl* *y*))
by (*simp add: map2-def zip-rev bl-word-and rev-map*)

lemma *rbl-word-xor*: *rev* (*to-bl* (*x* *XOR* *y*)) = *map2* *op* $\neq$ (*rev* (*to-bl* *x*)) (*rev* (*to-bl* *y*))
by (*simp add: map2-def zip-rev bl-word-xor rev-map*)

lemma *rbl-word-not*: *rev* (*to-bl* (*NOT* *x*)) = *map* *Not* (*rev* (*to-bl* *x*))
by (*simp add: bl-word-not rev-map*)

lemma *bl-word-sub*: *to-bl* (*x* - *y*) = *to-bl* (*x* + (- *y*))
by *simp*

lemma *rbl-word-1*: *rev* (*to-bl* (*1* :: '*a*::len0 word)) = *takefill* *False* (*len-of* *TYPE*('a'))
[*True*]
apply (*rule-tac* *s=rev* (*to-bl* (*word-succ* (*0* :: '*a* word)))) **in** *trans*
apply *simp*

```

apply (simp only: rtb-rbl-ariths(1)[OF refl])
apply simp
apply (case-tac len-of TYPE('a))
  apply simp
apply (simp add: takefill-alt)
done

```

lemma *rbl-word-if*: $\text{rev } (\text{to-bl } (\text{if } P \text{ then } x \text{ else } y)) = \text{map2 } (\text{If } P) (\text{rev } (\text{to-bl } x)) (\text{rev } (\text{to-bl } y))$
by (simp add: map2-def split-def)

lemma *rbl-add-carry-Cons*:
 $(\text{if } \text{car} \text{ then } \text{rbl-succ} \text{ else } \text{id}) (\text{rbl-add } (x \# xs) (y \# ys)) =$
 $\text{xor3 } x \ y \ \text{car} \# (\text{if } \text{carry } x \ y \ \text{car} \text{ then } \text{rbl-succ} \text{ else } \text{id}) (\text{rbl-add } xs \ ys)$
by (simp add: carry-def xor3-def)

lemma *rbl-add-suc-carry-fold*:
 $\text{length } xs = \text{length } ys \implies$
 $\forall \text{car. } (\text{if } \text{car} \text{ then } \text{rbl-succ} \text{ else } \text{id}) (\text{rbl-add } xs \ ys) =$
 $(\text{foldr } (\lambda(x, y) \ \text{res} \ \text{car. } \text{xor3 } x \ y \ \text{car} \# \text{res } (\text{carry } x \ y \ \text{car})) (\text{zip } xs \ ys) (\lambda-. []))$
car
apply (erule list-induct2)
apply simp
apply (simp only: rbl-add-carry-Cons)
apply simp
done

lemma *to-bl-plus-carry*:
 $\text{to-bl } (x + y) =$
 $\text{rev } (\text{foldr } (\lambda(x, y) \ \text{res} \ \text{car. } \text{xor3 } x \ y \ \text{car} \# \text{res } (\text{carry } x \ y \ \text{car}))$
 $(\text{rev } (\text{zip } (\text{to-bl } x) (\text{to-bl } y))) (\lambda-. []) \ \text{False})$
using *rbl-add-suc-carry-fold* **where** $xs = \text{rev } (\text{to-bl } x)$ **and** $ys = \text{rev } (\text{to-bl } y)$
apply (simp add: word-add-rbl[OF refl refl])
apply (drule-tac $x = \text{False}$ **in** *spec*)
apply (simp add: zip-rev)
done

definition *rbl-plus cin xs ys* =
 $\text{foldr } (\lambda(x, y) \ \text{res} \ \text{car. } \text{xor3 } x \ y \ \text{car} \# \text{res } (\text{carry } x \ y \ \text{car})) (\text{zip } xs \ ys) (\lambda-. []) \ \text{cin}$

lemma *rbl-plus-simps*:
 $\text{rbl-plus } \text{cin } (x \# xs) (y \# ys) = \text{xor3 } x \ y \ \text{cin} \# \text{rbl-plus } (\text{carry } x \ y \ \text{cin}) \ xs \ ys$
 $\text{rbl-plus } \text{cin } [] \ ys = []$
 $\text{rbl-plus } \text{cin } xs \ [] = []$
by (simp-all add: rbl-plus-def)

lemma *rbl-word-plus*: $\text{rev } (\text{to-bl } (x + y)) = \text{rbl-plus } \text{False} (\text{rev } (\text{to-bl } x)) (\text{rev } (\text{to-bl } y))$
by (simp add: rbl-plus-def to-bl-plus-carry zip-rev)

definition $rbl-succ2\ b\ xs = (if\ b\ then\ rbl-succ\ xs\ else\ xs)$

lemma $rbl-succ2-simps$:

$rbl-succ2\ b\ [] = []$

$rbl-succ2\ b\ (x \# xs) = (b \neq x) \# rbl-succ2\ (x \wedge b)\ xs$

by ($simp\ add: rbl-succ2-def$)

lemma $twos-complement$: $-x = word-succ\ (NOT\ x)$

using $arg-cong[OF\ word-add-not[where\ x=x],\ where\ f=\lambda a.\ a - x + 1]$

by ($simp\ add: word-succ-p1\ word-sp-01[unfolded\ word-succ-p1]\ del: word-add-not$)

lemma $rbl-word-neg$: $rev\ (to-bl\ (-x)) = rbl-succ2\ True\ (map\ Not\ (rev\ (to-bl\ x)))$

by ($simp\ add: twos-complement\ word-succ-rbl[OF\ refl]\ bl-word-not\ rev-map\ rbl-succ2-def$)

lemma $rbl-word-cat$:

$rev\ (to-bl\ (word-cat\ x\ y :: 'a::len0\ word)) =$

$takefill\ False\ (len-of\ TYPE('a))\ (rev\ (to-bl\ y) @ rev\ (to-bl\ x))$

by ($simp\ add: word-cat-bl\ word-rev-tf$)

lemma $rbl-word-slice$:

$rev\ (to-bl\ (slice\ n\ w :: 'a::len0\ word)) =$

$takefill\ False\ (len-of\ TYPE('a))\ (drop\ n\ (rev\ (to-bl\ w)))$

apply ($simp\ add: slice-take\ word-rev-tf\ rev-take$)

apply ($cases\ n < len-of\ TYPE('b),\ simp-all$)

done

lemma $rbl-word-ucast$:

$rev\ (to-bl\ (ucast\ x :: 'a::len0\ word)) = takefill\ False\ (len-of\ TYPE('a))\ (rev\ (to-bl\ x))$

apply ($simp\ add: to-bl-ucast\ takefill-alt$)

apply ($simp\ add: rev-drop$)

apply ($cases\ len-of\ TYPE('a) < len-of\ TYPE('b)$)

apply $simp-all$

done

lemma $rbl-shiffl$:

$rev\ (to-bl\ (w << n)) = takefill\ False\ (size\ w)\ (replicate\ n\ False @ rev\ (to-bl\ w))$

by ($simp\ add: bl-shiffl\ takefill-alt\ word-size\ rev-drop$)

lemma $rbl-shiftr$:

$rev\ (to-bl\ (w >> n)) = takefill\ False\ (size\ w)\ (drop\ n\ (rev\ (to-bl\ w)))$

by ($simp\ add: shiftr-slice\ rbl-word-slice\ word-size$)

definition $drop-nonempty\ v\ n\ xs = (if\ n < length\ xs\ then\ drop\ n\ xs\ else\ [last\ (v \# xs)])$

lemma $drop-nonempty-simps$:

$drop-nonempty\ v\ (Suc\ n)\ (x \# xs) = drop-nonempty\ x\ n\ xs$

```

drop-nonempty v 0 (x # xs) = (x # xs)
drop-nonempty v n [] = [v]
by (simp-all add: drop-nonempty-def)

```

definition *takefill-last* $x\ n\ xs = \text{takefill}\ (\text{last}\ (x\ \# \ xs))\ n\ xs$

lemma *takefill-last-simps*:

```

takefill-last z (Suc n) (x # xs) = x # takefill-last x n xs
takefill-last z 0 xs = []
takefill-last z n [] = replicate n z
by (simp-all add: takefill-last-def) (simp-all add: takefill-alt)

```

lemma *rbl-sshiftr*:

```

rev (to-bl (w >>> n)) = takefill-last False (size w) (drop-nonempty False n (rev
(to-bl w)))
apply (cases n < size w)
  apply (simp add: bl-sshiftr takefill-last-def word-size
    takefill-alt rev-take last-rev
    drop-nonempty-def)
  apply (subgoal-tac (w >>> n) = of-bl (replicate (size w) (msb w)))
  apply (simp add: word-size takefill-last-def takefill-alt
    last-rev word-msb-alt word-rev-tf
    drop-nonempty-def take-Cons')
  apply (case-tac len-of TYPE('a), simp-all)
  apply (rule word-eqI)
  apply (simp add: nth-sshiftr word-size test-bit-of-bl
    msb-nth)
done

```

lemma *nth-word-of-int*:

```

(word-of-int x :: 'a::len0 word) !! n = (n < len-of TYPE('a) ∧ bin-nth x n)
apply (simp add: test-bit-bl word-size to-bl-of-bin)
apply (subst conj-cong[OF refl], erule bin-nth-bl)
apply auto
done

```

lemma *nth-scast*:

```

(scast (x :: 'a::len word) :: 'b::len word) !! n =
  (n < len-of TYPE('b) ∧
   (if n < len-of TYPE('a) - 1 then x !! n
    else x !! (len-of TYPE('a) - 1)))
by (simp add: scast-def nth-sint)

```

lemma *rbl-word-scast*:

```

rev (to-bl (scast x :: 'a::len word)) = takefill-last False (len-of TYPE('a)) (rev
(to-bl x))
apply (rule nth-equalityI)
  apply (simp add: word-size takefill-last-def)
  apply (clarsimp simp: nth-scast takefill-last-def)

```

```

    nth-takefill word-size nth-rev to-bl-nth)
  apply (cases len-of TYPE('b))
  apply simp
  apply (clarsimp simp: less-Suc-eq-le linorder-not-less
    last-rev word-msb-alt[symmetric]
    msb-nth)
done

```

definition *rbl-mul* :: *bool list* $\Rightarrow$ *bool list* $\Rightarrow$ *bool list*

where *rbl-mul* *xs ys* = foldr ($\lambda x sm. \text{rbl-plus False (map (op } \wedge x) \text{ ys) (False \# sm)}$) *xs* []

lemma *rbl-mul-simps*:

```

  rbl-mul (x # xs) ys = rbl-plus False (map (op } \wedge x) ys) (False # rbl-mul xs ys)
  rbl-mul [] ys = []
  by (simp-all add: rbl-mul-def)

```

lemma *takefill-le2*: *length xs* $\leq n \implies \text{takefill } x \text{ } m \text{ (takefill } x \text{ } n \text{ } xs) = \text{takefill } x \text{ } m \text{ } xs$

by (*simp add: takefill-alt replicate-add[symmetric]*)

lemma *take-rbl-plus*: $\forall n \text{ } b. \text{take } n \text{ (rbl-plus } b \text{ } xs \text{ } ys) = \text{rbl-plus } b \text{ (take } n \text{ } xs) \text{ (take } n \text{ } ys)$

```

  apply (simp add: rbl-plus-def take-zip[symmetric])
  apply (rule-tac list=zip xs ys in list.induct)
  apply simp
  apply (clarsimp simp: split-def)
  apply (case-tac n, simp-all)
done

```

lemma *word-rbl-mul-induct*:

```

  length xs } \leq size y } \implies
  rbl-mul xs (rev (to-bl y)) = take (length xs) (rev (to-bl (of-bl (rev xs) * y)))
  for y :: 'a::len word

```

proof (*induct xs*)

case *Nil*

show ?*case* **by** (*simp add: rbl-mul-simps*)

next

case (*Cons z zs*)

have *rbl-word-plus'*: *to-bl* (*x* + *y*) = *rev* (*rbl-plus False* (*rev* (*to-bl* *x*)) (*rev* (*to-bl* *y*)))

for *x y* :: 'a word

by (*simp add: rbl-word-plus[symmetric]*)

have *mult-bit*: *to-bl* (*of-bl* [*z*] * *y*) = *map* (*op } \wedge z*) (*to-bl* *y*)

by (*cases z*) (*simp cong: map-cong, simp add: map-replicate-const cong: map-cong*)

have *shifftl*: *of-bl* *xs* * 2 * *y* = (*of-bl* *xs* * *y*) << 1 **for** *xs*

```

by (simp add: shiftl-t2n)

have zip-take-triv:  $\bigwedge xs\ ys\ n.\ n = \text{length}\ ys \implies \text{zip}\ (\text{take}\ n\ xs)\ ys = \text{zip}\ xs\ ys$ 
by (rule nth-equalityI) simp-all

from Cons show ?case
  apply (simp add: trans [OF of-bl-append add.commute]
    rbl-mul-simps rbl-word-plus' distrib-right mult-bit shiftl rbl-shiftl)
  apply (simp add: takefill-alt word-size rev-map take-rbl-plus min-def)
  apply (simp add: rbl-plus-def zip-take-triv)
  done
qed

lemma rbl-word-mul:  $\text{rev}\ (\text{to-bl}\ (x * y)) = \text{rbl-mul}\ (\text{rev}\ (\text{to-bl}\ x))\ (\text{rev}\ (\text{to-bl}\ y))$ 
for  $x :: 'a::\text{len}\ \text{word}$ 
  using word-rbl-mul-induct[where  $xs = \text{rev}\ (\text{to-bl}\ x)$  and  $y = y$ ] by (simp add:
word-size)

Breaking up inequalities into bitlist properties.

definition
  rev-bl-order  $F\ xs\ ys =$ 
    ( $\text{length}\ xs = \text{length}\ ys \wedge$ 
      ( $(xs = ys \wedge F)$ 
         $\vee (\exists n < \text{length}\ xs.\ \text{drop}\ (\text{Suc}\ n)\ xs = \text{drop}\ (\text{Suc}\ n)\ ys$ 
           $\wedge \neg xs\ !\ n \wedge ys\ !\ n)))$ )

lemma rev-bl-order-simps:
  rev-bl-order  $F\ []\ [] = F$ 
  rev-bl-order  $F\ (x \# xs)\ (y \# ys) = \text{rev-bl-order}\ ((y \wedge \neg x) \vee ((y \vee \neg x) \wedge F))$ 
   $xs\ ys$ 
  apply (simp-all add: rev-bl-order-def)
  apply (rule conj-cong[OF refl])
  apply (cases  $xs = ys$ )
  apply (simp add: nth-Cons')
  apply blast
  apply (simp add: nth-Cons')
  apply safe
  apply (rule-tac  $x = n - 1$  in  $exI$ )
  apply simp
  apply (rule-tac  $x = \text{Suc}\ n$  in  $exI$ )
  apply simp
  done

lemma rev-bl-order-rev-simp:
   $\text{length}\ xs = \text{length}\ ys \implies$ 
  rev-bl-order  $F\ (xs\ @\ [x])\ (ys\ @\ [y]) = ((y \wedge \neg x) \vee ((y \vee \neg x) \wedge \text{rev-bl-order}\ F$ 
   $xs\ ys))$ 
  by (induct arbitrary:  $F$  rule: list-induct2) (auto simp: rev-bl-order-simps)

```

lemma *rev-bl-order-bl-to-bin*:

```
length xs = length ys  $\implies$ 
  rev-bl-order True xs ys = (bl-to-bin (rev xs)  $\leq$  bl-to-bin (rev ys))  $\wedge$ 
  rev-bl-order False xs ys = (bl-to-bin (rev xs)  $<$  bl-to-bin (rev ys))
apply (induct xs ys rule: list-induct2)
apply (simp-all add: rev-bl-order-simps bl-to-bin-app-cat)
apply (auto simp add: bl-to-bin-def Bit-B0 Bit-B1 add1-zle-eq Bit-def)
done
```

lemma *word-le-rbl*: $x \leq y \iff \text{rev-bl-order True (rev (to-bl x)) (rev (to-bl y))}$
for $x\ y :: 'a::\text{len0 word}$
by (simp add: rev-bl-order-bl-to-bin word-le-def)

lemma *word-less-rbl*: $x < y \iff \text{rev-bl-order False (rev (to-bl x)) (rev (to-bl y))}$
for $x\ y :: 'a::\text{len0 word}$
by (simp add: word-less-alt rev-bl-order-bl-to-bin)

lemma *word-sint-msb-eq*: $\text{sint } x = \text{uint } x - (\text{if msb } x \text{ then } 2^{\text{size } x} \text{ else } 0)$
apply (cases msb x)
apply (rule word-sint.Abs-eqD[**where** 'a='a], simp-all)
apply (simp add: word-size wi-hom-syms word-of-int-2p-len)
apply (simp add: sints-num word-size)
apply (rule conjI)
apply (simp add: le-diff-eq')
apply (rule order-trans[**where** $y=2^{\text{len-of TYPE('a)} - 1}$])
apply (simp add: power-Suc[symmetric])
apply (simp add: linorder-not-less[symmetric] mask-eq-iff[symmetric])
apply (rule notI, drule word-eqD[**where** $x=\text{size } x - 1$])
apply (simp add: msb-nth word-ops-nth-size word-size)
apply (simp add: order-less-le-trans[**where** $y=0$])
apply (rule word-uint.Abs-eqD[**where** 'a='a], simp-all)
apply (simp add: linorder-not-less uints-num word-msb-sint)
apply (rule order-less-le-trans[OF sint-lt])
apply simp
done

lemma *word-sle-msb-le*: $x \leq_s y \iff (\text{msb } y \implies \text{msb } x) \wedge ((\text{msb } x \wedge \neg \text{msb } y) \vee x \leq y)$
apply (simp add: word-sle-def word-sint-msb-eq word-size word-le-def)
apply safe
apply (rule order-trans[OF - uint-ge-0])
apply (simp add: order-less-imp-le)
apply (erule notE[OF leD])
apply (rule order-less-le-trans[OF - uint-ge-0])
apply simp
done

lemma *word-sless-msb-less*: $x <_s y \iff (\text{msb } y \implies \text{msb } x) \wedge ((\text{msb } x \wedge \neg \text{msb } y) \vee x < y)$

by (*auto simp add: word-sless-def word-sle-msb-le*)

definition *map-last* f $xs = (if\ xs = []\ then\ []\ else\ butlast\ xs\ @\ [f\ (last\ xs)])$

lemma *map-last-simps*:

map-last $f\ [] = []$

map-last $f\ [x] = [f\ x]$

map-last $f\ (x\ \# \ y\ \# \ zs) = x\ \# \ map-last\ f\ (y\ \# \ zs)$

by (*simp-all add: map-last-def*)

lemma *word-sle-rbl*:

$x \leq s\ y \longleftrightarrow rev-bl-order\ True\ (map-last\ Not\ (rev\ (to-bl\ x)))\ (map-last\ Not\ (rev\ (to-bl\ y)))$

using *word-msb-alt*[**where** $w=x$] *word-msb-alt*[**where** $w=y$]

apply (*simp add: word-sle-msb-le word-le-rbl*)

apply (*subgoal-tac length (to-bl x) = length (to-bl y)*)

apply (*cases to-bl x, simp*)

apply (*cases to-bl y, simp*)

apply (*clarsimp simp: map-last-def rev-bl-order-rev-simp*)

apply *auto*

done

lemma *word-sless-rbl*:

$x < s\ y \longleftrightarrow rev-bl-order\ False\ (map-last\ Not\ (rev\ (to-bl\ x)))\ (map-last\ Not\ (rev\ (to-bl\ y)))$

using *word-msb-alt*[**where** $w=x$] *word-msb-alt*[**where** $w=y$]

apply (*simp add: word-sless-msb-less word-less-rbl*)

apply (*subgoal-tac length (to-bl x) = length (to-bl y)*)

apply (*cases to-bl x, simp*)

apply (*cases to-bl y, simp*)

apply (*clarsimp simp: map-last-def rev-bl-order-rev-simp*)

apply *auto*

done

Lemmas for unpacking *rev (to-bl n)* for numerals n and also for irreducible values and expressions.

lemma *rev-bin-to-bl-simps*:

rev (bin-to-bl 0 x) = []

rev (bin-to-bl (Suc n) (numeral (num.Bit0 nm))) = False # rev (bin-to-bl n (numeral nm))

rev (bin-to-bl (Suc n) (numeral (num.Bit1 nm))) = True # rev (bin-to-bl n (numeral nm))

rev (bin-to-bl (Suc n) (numeral (num.One))) = True # replicate n False

rev (bin-to-bl (Suc n) (¬ numeral (num.Bit0 nm))) = False # rev (bin-to-bl n (¬ numeral nm))

rev (bin-to-bl (Suc n) (¬ numeral (num.Bit1 nm))) =

True # rev (bin-to-bl n (¬ numeral (nm + num.One)))

rev (bin-to-bl (Suc n) (¬ numeral (num.One))) = True # replicate n True

rev (bin-to-bl (Suc n) (¬ numeral (num.Bit0 nm + num.One))) =

```

    True # rev (bin-to-bl n (- numeral (nm + num.One)))
  rev (bin-to-bl (Suc n) (- numeral (num.Bit1 nm + num.One))) =
    False # rev (bin-to-bl n (- numeral (nm + num.One)))
  rev (bin-to-bl (Suc n) (- numeral (num.One + num.One))) =
    False # rev (bin-to-bl n (- numeral num.One))
apply simp-all
apply (simp-all only: bin-to-bl-aux-alt)
apply (simp-all)
apply (simp-all add: bin-to-bl-zero-aux bin-to-bl-minus1-aux)
done

```

```

lemma to-bl-upt: to-bl x = rev (map (op !! x) [0 ..< size x])
apply (rule nth-equalityI)
apply (simp add: word-size)
apply (auto simp: to-bl-nth word-size nth-rev)
done

```

```

lemma rev-to-bl-upt: rev (to-bl x) = map (op !! x) [0 ..< size x]
by (simp add: to-bl-upt)

```

```

lemma upt-eq-list-intros:
  j ≤ i ⇒ [i ..< j] = []
  i = x ⇒ x < j ⇒ [x + 1 ..< j] = xs ⇒ [i ..< j] = (x # xs)
by (simp-all add: upt-eq-Cons-conv)

```

10.30 Tactic definition

```

ML ⟨
  structure Word-Bitwise-Tac =
  struct

```

```

    val word-ss = simpset-of @{theory-context Word};

```

```

    fun mk-nat-clist ns = List.foldr
      (uncurry (Thm.mk-binop @ {cterm Cons :: nat => -}))
      @ {cterm [] :: nat list} ns;

```

```

    fun upt-conv ctxt ct =
      case Thm.term-of ct of
        (@ {const upt} $ n $ m) =>
          let
            val (i, j) = apply2 (snd o HOLogic.dest-number) (n, m);
            val ns = map (Numeral.mk-cnumber @ {ctyp nat}) (i upto (j - 1))
              |> mk-nat-clist;
            val prop = Thm.mk-binop @ {cterm op = :: nat list => -} ct ns
              |> Thm.apply @ {cterm Trueprop};
          in
            try (fn () =>
              Goal.prove-internal ctxt [] prop

```

```

      (K (REPEAT-DETERM (resolve-tac ctxt @{thms upt-eq-list-intros} 1
        ORELSE simp-tac (put-simpset word-ss ctxt) 1))) |> mk-meta-eq) ()
    end
  | - => NONE;

val expand-upt-simproc =
  Simplifier.make-simproc @{context} expand-upt
  {lhss = [@{term upt x y}], proc = K upt-conv};

fun word-len-simproc-fn ctxt ct =
  case Thm.term-of ct of
    Const (@{const-name len-of}, -) $ t => (let
      val T = fastype-of t |> dest-Type |> snd |> the-single
      val n = Numeral.mk-cnumber @{ctyp nat} (Word-Lib.dest-binT T);
      val prop = Thm.mk-binop @{cterm op = :: nat => -} ct n
        |> Thm.apply @{cterm Trueprop};
      in Goal.prove-internal ctxt [] prop (K (simp-tac (put-simpset word-ss ctxt)
1))
        |> mk-meta-eq |> SOME end
      handle TERM - => NONE | TYPE - => NONE)
    | - => NONE;

val word-len-simproc =
  Simplifier.make-simproc @{context} word-len
  {lhss = [@{term len-of x}], proc = K word-len-simproc-fn};

(* convert 5 or nat 5 to Suc 4 when n-sucs = 1, Suc (Suc 4) when n-sucs = 2,
   or just 5 (discarding nat) when n-sucs = 0 *)

fun nat-get-Suc-simproc-fn n-sucs ctxt ct =
  let
    val (f $ arg) = Thm.term-of ct;
    val n = (case arg of @{term nat} $ n => n | n => n)
      |> HOLogic.dest-number |> snd;
    val (i, j) = if n > n-sucs then (n-sucs, n - n-sucs)
      else (n, 0);
    val arg' = List.foldr (op $) (HOLogic.mk-number @{typ nat} j)
      (replicate i @{term Suc});
    val - = if arg = arg' then raise TERM (, []) else ();
    fun propfn g = HOLogic.mk-eq (g arg, g arg')
      |> HOLogic.mk-Trueprop |> Thm.cterm-of ctxt;
    val eq1 = Goal.prove-internal ctxt [] (propfn I)
      (K (simp-tac (put-simpset word-ss ctxt) 1));
    in Goal.prove-internal ctxt [] (propfn (curry (op $) f))
      (K (simp-tac (put-simpset HOL-ss ctxt addsimps [eq1] 1))
        |> mk-meta-eq |> SOME end
      handle TERM - => NONE;
  end

fun nat-get-Suc-simproc n-sucs ts =

```

```

Simplifier.make-simproc @{context} nat-get-Suc
{lhs = map (fn t => t $ @{term n :: nat}) ts,
 proc = K (nat-get-Suc-simproc-fn n-sucs)};

val no-split-ss =
  simpset-of (put-simpset HOL-ss @{context}
    |> Splitter.del-split @{thm if-split});

val expand-word-eq-sss =
  (simpset-of (put-simpset HOL-basic-ss @{context} addsimps
    @{thms word-eq-rbl-eq word-le-rbl word-less-rbl word-sle-rbl word-sless-rbl}),
  map simpset-of [
    put-simpset no-split-ss @{context} addsimps
      @{thms rbl-word-plus rbl-word-and rbl-word-or rbl-word-not
        rbl-word-neg bl-word-sub rbl-word-xor
        rbl-word-cat rbl-word-slice rbl-word-scast
        rbl-word-ucast rbl-shiftr rbl-shiftr rbl-sshiftr
        rbl-word-if},
    put-simpset no-split-ss @{context} addsimps
      @{thms to-bl-numeral to-bl-neg-numeral to-bl-0 rbl-word-1},
    put-simpset no-split-ss @{context} addsimps
      @{thms rev-rev-ident rev-replicate rev-map to-bl-upt word-size}
      addsimprocs [word-len-simproc],
    put-simpset no-split-ss @{context} addsimps
      @{thms list.simps split-conv replicate.simps list.map
        zip-Cons-Cons zip-Nil drop-Suc-Cons drop-0 drop-Nil
        foldr.simps map2-Cons map2-Nil takefill-Suc-Cons
        takefill-Suc-Nil takefill.Z rbl-succ2-simps
        rbl-plus-simps rev-bin-to-bl-simps append.simps
        takefill-last-simps drop-nonempty-simps
        rev-bl-order-simps}
      addsimprocs [expand-upt-simproc,
        nat-get-Suc-simproc 4
        [@{term replicate}, @{term takefill x},
        @{term drop}, @{term bin-to-bl},
        @{term takefill-last x},
        @{term drop-nonempty x}]],
    put-simpset no-split-ss @{context} addsimps @{thms xor3-simps carry-simps
if-bool-simps}
  ])

fun tac ctxt =
  let
    val (ss, sss) = expand-word-eq-sss;
  in
    foldr1 (op THEN-ALL-NEW)
      ((CHANGED o safe-full-simp-tac (put-simpset ss ctxt)) ::
        map (fn ss => safe-full-simp-tac (put-simpset ss ctxt)) sss)
  end;

```

end

)

method-setup *word-bitwise* =

$\langle \text{Scan.succeed } (\text{fn } \text{ctxt} \Rightarrow \text{Method.SIMPLE-METHOD } (\text{Word-Bitwise-Tac.tac } \text{ctxt } 1)) \rangle$

decomposer for word equalities and inequalities into bit propositions

end

theory *Bit-Comparison*

imports *Bits-Int*

begin

lemma *AND-lower* [*simp*]:

fixes $x :: \text{int}$ **and** $y :: \text{int}$

assumes $0 \leq x$

shows $0 \leq x \text{ AND } y$

using *assms*

proof (*induct x arbitrary: y rule: bin-induct*)

case (3 bin bit)

show ?*case*

proof (*cases y rule: bin-exhaust*)

case (1 bin' bit')

from 3 **have** $0 \leq \text{bin}$

by (*cases bit*) (*simp-all add: Bit-def*)

then have $0 \leq \text{bin AND bin'}$ **by** (*rule 3*)

with 1 show ?*thesis*

by *simp* (*simp add: Bit-def*)

qed

next

case 2

then show ?*case* **by** (*simp only: Min-def*)

qed *simp*

lemma *OR-lower* [*simp*]:

fixes $x :: \text{int}$ **and** $y :: \text{int}$

assumes $0 \leq x$ $0 \leq y$

shows $0 \leq x \text{ OR } y$

using *assms*

proof (*induct x arbitrary: y rule: bin-induct*)

case (3 bin bit)

show ?*case*

proof (*cases y rule: bin-exhaust*)

case (1 bin' bit')

from 3 **have** $0 \leq \text{bin}$

by (*cases bit*) (*simp-all add: Bit-def*)

```

    moreover from 1 3 have  $0 \leq bin'$ 
      by (cases bit') (simp-all add: Bit-def)
    ultimately have  $0 \leq bin$  OR  $bin'$  by (rule 3)
    with 1 show ?thesis
      by simp (simp add: Bit-def)
  qed
qed simp-all

lemma XOR-lower [simp]:
  fixes  $x :: int$  and  $y :: int$ 
  assumes  $0 \leq x$   $0 \leq y$ 
  shows  $0 \leq x XOR y$ 
  using assms
proof (induct x arbitrary: y rule: bin-induct)
  case (3 bin bit)
  show ?case
  proof (cases y rule: bin-exhaust)
    case (1 bin' bit')
    from 3 have  $0 \leq bin$ 
      by (cases bit) (simp-all add: Bit-def)
    moreover from 1 3 have  $0 \leq bin'$ 
      by (cases bit') (simp-all add: Bit-def)
    ultimately have  $0 \leq bin XOR bin'$  by (rule 3)
    with 1 show ?thesis
      by simp (simp add: Bit-def)
  qed
next
  case 2
  then show ?case by (simp only: Min-def)
qed simp

```

```

lemma AND-upper1 [simp]:
  fixes  $x :: int$  and  $y :: int$ 
  assumes  $0 \leq x$ 
  shows  $x AND y \leq x$ 
  using assms
proof (induct x arbitrary: y rule: bin-induct)
  case (3 bin bit)
  show ?case
  proof (cases y rule: bin-exhaust)
    case (1 bin' bit')
    from 3 have  $0 \leq bin$ 
      by (cases bit) (simp-all add: Bit-def)
    then have  $bin AND bin' \leq bin$  by (rule 3)
    with 1 show ?thesis
      by simp (simp add: Bit-def)
  qed
next
  case 2

```

then show ?case by (simp only: Min-def)
qed simp

lemmas AND-upper1' [simp] = order-trans [OF AND-upper1]
lemmas AND-upper1'' [simp] = order-le-less-trans [OF AND-upper1]

lemma AND-upper2 [simp]:
fixes x :: int and y :: int
assumes $0 \leq y$
shows $x \text{ AND } y \leq y$
using assms
proof (induct y arbitrary: x rule: bin-induct)
case (3 bin bit)
show ?case
proof (cases x rule: bin-exhaust)
case (1 bin' bit')
from 3 have $0 \leq \text{bin}$
by (cases bit) (simp-all add: Bit-def)
then have $\text{bin}' \text{ AND } \text{bin} \leq \text{bin}$ by (rule 3)
with 1 show ?thesis
by simp (simp add: Bit-def)
qed
next
case 2
then show ?case by (simp only: Min-def)
qed simp

lemmas AND-upper2' [simp] = order-trans [OF AND-upper2]
lemmas AND-upper2'' [simp] = order-le-less-trans [OF AND-upper2]

lemma OR-upper:
fixes x :: int and y :: int
assumes $0 \leq x$ and $x < 2^n$ and $y < 2^n$
shows $x \text{ OR } y < 2^n$
using assms
proof (induct x arbitrary: y n rule: bin-induct)
case (3 bin bit)
show ?case
proof (cases y rule: bin-exhaust)
case (1 bin' bit')
show ?thesis
proof (cases n)
case 0
with 3 have $\text{bin BIT bit} = 0$ by simp
then have $\text{bin} = 0$ and $\neg \text{bit}$
by (auto simp add: Bit-def split: if-splits) arith
then show ?thesis using 0 1 (y < 2^n)
by simp
next

```

    case (Suc m)
    from 3 have 0 ≤ bin
      by (cases bit) (simp-all add: Bit-def)
    moreover from 3 Suc have bin < 2 ^ m
      by (cases bit) (simp-all add: Bit-def)
    moreover from 1 3 Suc have bin' < 2 ^ m
      by (cases bit') (simp-all add: Bit-def)
    ultimately have bin OR bin' < 2 ^ m by (rule 3)
    with 1 Suc show ?thesis
      by simp (simp add: Bit-def)
  qed
qed
qed simp-all

```

lemma XOR-upper:

```

  fixes x :: int and y :: int
  assumes 0 ≤ x x < 2 ^ n y < 2 ^ n
  shows x XOR y < 2 ^ n
  using assms
proof (induct x arbitrary: y n rule: bin-induct)
  case (3 bin bit)
  show ?case
  proof (cases y rule: bin-exhaust)
    case (1 bin' bit')
    show ?thesis
    proof (cases n)
      case 0
      with 3 have bin BIT bit = 0 by simp
      then have bin = 0 and ¬ bit
        by (auto simp add: Bit-def split: if-splits) arith
      then show ?thesis using 0 1 ⟨y < 2 ^ n⟩
        by simp
    next
      case (Suc m)
      from 3 have 0 ≤ bin
        by (cases bit) (simp-all add: Bit-def)
      moreover from 3 Suc have bin < 2 ^ m
        by (cases bit) (simp-all add: Bit-def)
      moreover from 1 3 Suc have bin' < 2 ^ m
        by (cases bit') (simp-all add: Bit-def)
      ultimately have bin XOR bin' < 2 ^ m by (rule 3)
      with 1 Suc show ?thesis
        by simp (simp add: Bit-def)
    qed
  qed
next
case 2
  then show ?case by (simp only: Min-def)
qed simp

```

end

References

- [1] Jeremy Dawson. Isabelle theories for machine words. In Michael Goldsmith and Bill Roscoe, editors, *Seventh International Workshop on Automated Verification of Critical Systems (AVOCS'07)*, Electronic Notes in Theoretical Computer Science, page 15, Oxford, September 2007. Elsevier. to appear.