

Machine Words in Isabelle/HOL

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Abstract

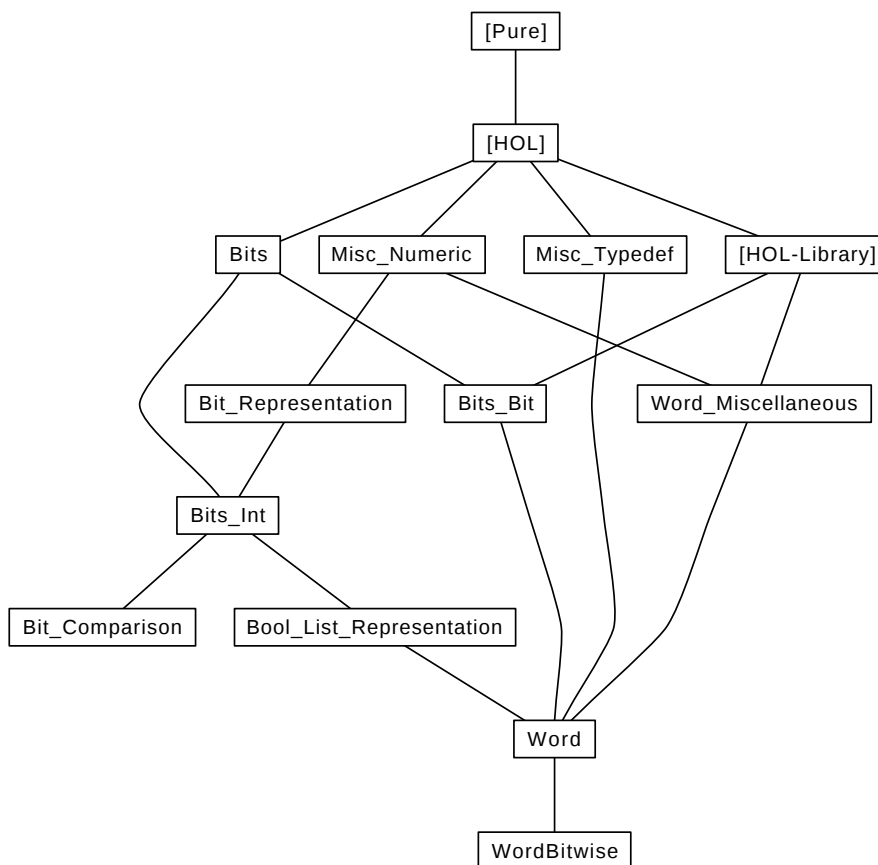
A formalisation of generic, fixed size machine words in Isabelle/HOL.
An earlier version of this formalisation is described in [1].

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1 Syntactic classes for bitwise operations

```

theory Bits
  imports Main
begin

class bit =
  fixes bitNOT :: 'a  $\Rightarrow$  'a (NOT - [70] 71)
    and bitAND :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr AND 64)
    and bitOR :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr OR 59)
    and bitXOR :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr XOR 59)

```

We want the bitwise operations to bind slightly weaker than $+$ and $-$, but $\sim\sim$ to bind slightly stronger than $*$.

Testing and shifting operations.

```

class bits = bit +
  fixes test-bit :: 'a  $\Rightarrow$  nat  $\Rightarrow$  bool (infixl !! 100)
    and lsb :: 'a  $\Rightarrow$  bool
    and set-bit :: 'a  $\Rightarrow$  nat  $\Rightarrow$  bool  $\Rightarrow$  'a
    and set-bits :: (nat  $\Rightarrow$  bool)  $\Rightarrow$  'a (binder BITS 10)
    and shiffl :: 'a  $\Rightarrow$  nat  $\Rightarrow$  'a (infixl << 55)
    and shiftr :: 'a  $\Rightarrow$  nat  $\Rightarrow$  'a (infixl >> 55)

class bitss = bits +
  fixes msb :: 'a  $\Rightarrow$  bool

end

```

2 Bit operations in $\mathcal{Z}_\epsilon$

```

theory Bits-Bit
imports Bits HOL-Library.Bit
begin

```

```

instantiation bit :: bit
begin

```

```

primrec bitNOT-bit
  where
    NOT 0 = (1::bit)
  | NOT 1 = (0::bit)

```

```

primrec bitAND-bit
  where
    0 AND y = (0::bit)
  | 1 AND y = (y::bit)

```

```

primrec bitOR-bit

```

```

where
  0 OR y = (y::bit)
| 1 OR y = (1::bit)

primrec bitXOR-bit
where
  0 XOR y = (y::bit)
| 1 XOR y = (NOT y :: bit)

instance ⟨proof⟩

end

lemmas bit-simps =
  bitNOT-bit.simps bitAND-bit.simps bitOR-bit.simps bitXOR-bit.simps

lemma bit-extra-simps [simp]:
  x AND 0 = 0
  x AND 1 = x
  x OR 1 = 1
  x OR 0 = x
  x XOR 1 = NOT x
  x XOR 0 = x
for x :: bit
⟨proof⟩

lemma bit-ops-comm:
  x AND y = y AND x
  x OR y = y OR x
  x XOR y = y XOR x
for x :: bit
⟨proof⟩

lemma bit-ops-same [simp]:
  x AND x = x
  x OR x = x
  x XOR x = 0
for x :: bit
⟨proof⟩

lemma bit-not-not [simp]: NOT (NOT x) = x
for x :: bit
⟨proof⟩

lemma bit-or-def: b OR c = NOT (NOT b AND NOT c)
for b c :: bit
⟨proof⟩

lemma bit-xor-def: b XOR c = (b AND NOT c) OR (NOT b AND c)

```

```

for  $b\ c :: \text{bit}$ 
   $\langle \text{proof} \rangle$ 

lemma bit-NOT-eq-1-iff [simp]:  $\text{NOT } b = 1 \longleftrightarrow b = 0$ 
  for  $b :: \text{bit}$ 
   $\langle \text{proof} \rangle$ 

lemma bit-AND-eq-1-iff [simp]:  $a \text{ AND } b = 1 \longleftrightarrow a = 1 \wedge b = 1$ 
  for  $a\ b :: \text{bit}$ 
   $\langle \text{proof} \rangle$ 

end

```

3 Useful Numerical Lemmas

```

theory Misc-Numeric
  imports Main
begin

lemma one-mod-exp-eq-one [simp]:  $1 \bmod (2 * 2 ^ n) = (1 :: \text{int})$ 
   $\langle \text{proof} \rangle$ 

lemma mod-2-neq-1-eq-eq-0:  $k \bmod 2 \neq 1 \longleftrightarrow k \bmod 2 = 0$ 
  for  $k :: \text{int}$ 
   $\langle \text{proof} \rangle$ 

lemma z1pmod2:  $(2 * b + 1) \bmod 2 = (1 :: \text{int})$ 
  for  $b :: \text{int}$ 
   $\langle \text{proof} \rangle$ 

lemma diff-le-eq':  $a - b \leq c \longleftrightarrow a \leq b + c$ 
  for  $a\ b\ c :: \text{int}$ 
   $\langle \text{proof} \rangle$ 

lemma emep1:  $\text{even } n \implies \text{even } d \implies 0 \leq d \implies (n + 1) \bmod d = (n \bmod d) + 1$ 
  for  $n\ d :: \text{int}$ 
   $\langle \text{proof} \rangle$ 

lemma int-mod-ge:  $a < n \implies 0 < n \implies a \leq a \bmod n$ 
  for  $a\ n :: \text{int}$ 
   $\langle \text{proof} \rangle$ 

lemma int-mod-ge':  $b < 0 \implies 0 < n \implies b + n \leq b \bmod n$ 
  for  $b\ n :: \text{int}$ 
   $\langle \text{proof} \rangle$ 

lemma int-mod-le':  $0 \leq b - n \implies b \bmod n \leq b - n$ 
  for  $b\ n :: \text{int}$ 

```

$\langle proof \rangle$

lemma *zless2*: $0 < (2 :: int)$
 $\langle proof \rangle$

lemma *zless2p*: $0 < (2^n :: int)$
 $\langle proof \rangle$

lemma *zle2p*: $0 \leq (2^n :: int)$
 $\langle proof \rangle$

lemma *m1mod2k*: $-1 \bmod 2^n = (2^n - 1 :: int)$
 $\langle proof \rangle$

lemma *p1mod2k'*: $(1 + 2 * b) \bmod (2 * 2^n) = 1 + 2 * (b \bmod 2^n)$
for $b :: int$
 $\langle proof \rangle$

lemma *p1mod2k*: $(2 * b + 1) \bmod (2 * 2^n) = 2 * (b \bmod 2^n) + 1$
for $b :: int$
 $\langle proof \rangle$

lemma *int-mod-lem*: $0 < n \implies 0 \leq b \wedge b < n \iff b \bmod n = b$
for $b \ n :: int$
 $\langle proof \rangle$

end

4 Integers as implicit bit strings

theory *Bit-Representation*
imports *Misc-Numeric*
begin

4.1 Constructors and destructors for binary integers

definition *Bit* :: $int \Rightarrow bool \Rightarrow int$ (**infixl** *BIT* 90)
where $k \text{ BIT } b = (if\ b\ then\ 1\ else\ 0) + k + k$

lemma *Bit-B0*: $k \text{ BIT } False = k + k$
 $\langle proof \rangle$

lemma *Bit-B1*: $k \text{ BIT } True = k + k + 1$
 $\langle proof \rangle$

lemma *Bit-B0-2t*: $k \text{ BIT } False = 2 * k$
 $\langle proof \rangle$

lemma *Bit-B1-2t*: $k \text{ BIT } True = 2 * k + 1$

$\langle proof \rangle$

definition $bin\text{-}last :: int \Rightarrow bool$
where $bin\text{-}last\ w \longleftrightarrow w \bmod 2 = 1$

lemma $bin\text{-}last\text{-}odd$: $bin\text{-}last = odd$
 $\langle proof \rangle$

definition $bin\text{-}rest :: int \Rightarrow int$
where $bin\text{-}rest\ w = w \div 2$

lemma $bin\text{-}rl\text{-}simp$ $[simp]$: $bin\text{-}rest\ w\ BIT\ bin\text{-}last\ w = w$
 $\langle proof \rangle$

lemma $bin\text{-}rest\text{-}BIT$ $[simp]$: $bin\text{-}rest\ (x\ BIT\ b) = x$
 $\langle proof \rangle$

lemma $bin\text{-}last\text{-}BIT$ $[simp]$: $bin\text{-}last\ (x\ BIT\ b) = b$
 $\langle proof \rangle$

lemma $BIT\text{-}eq\text{-}iff$ $[iff]$: $u\ BIT\ b = v\ BIT\ c \longleftrightarrow u = v \wedge b = c$
 $\langle proof \rangle$

lemma $BIT\text{-}bin\text{-}simps$ $[simp]$:
 $numeral\ k\ BIT\ False = numeral\ (Num.Bit0\ k)$
 $numeral\ k\ BIT\ True = numeral\ (Num.Bit1\ k)$
 $(- numeral\ k)\ BIT\ False = - numeral\ (Num.Bit0\ k)$
 $(- numeral\ k)\ BIT\ True = - numeral\ (Num.BitM\ k)$
 $\langle proof \rangle$

lemma $BIT\text{-}special\text{-}simps$ $[simp]$:
shows $0\ BIT\ False = 0$
and $0\ BIT\ True = 1$
and $1\ BIT\ False = 2$
and $1\ BIT\ True = 3$
and $(- 1)\ BIT\ False = - 2$
and $(- 1)\ BIT\ True = - 1$
 $\langle proof \rangle$

lemma $Bit\text{-}eq\text{-}0\text{-}iff$: $w\ BIT\ b = 0 \longleftrightarrow w = 0 \wedge \neg b$
 $\langle proof \rangle$

lemma $Bit\text{-}eq\text{-}m1\text{-}iff$: $w\ BIT\ b = -1 \longleftrightarrow w = -1 \wedge b$
 $\langle proof \rangle$

lemma $BitM\text{-}inc$: $Num.BitM\ (Num.inc\ w) = Num.Bit1\ w$
 $\langle proof \rangle$

lemma $expand\text{-}BIT$:

$\text{numeral } (\text{Num.Bit0 } w) = \text{numeral } w \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit1 } w) = \text{numeral } w \text{ BIT True}$
 $\neg \text{numeral } (\text{Num.Bit0 } w) = (\neg \text{numeral } w) \text{ BIT False}$
 $\neg \text{numeral } (\text{Num.Bit1 } w) = (\neg \text{numeral } (w + \text{Num.One})) \text{ BIT True}$
 $\langle \text{proof} \rangle$

lemma *bin-last-numeral-simps* [simp]:

$\neg \text{bin-last } 0$
 $\text{bin-last } 1$
 $\text{bin-last } (-1)$
 bin-last Numeral1
 $\neg \text{bin-last } (\text{numeral } (\text{Num.Bit0 } w))$
 $\text{bin-last } (\text{numeral } (\text{Num.Bit1 } w))$
 $\neg \text{bin-last } (\neg \text{numeral } (\text{Num.Bit0 } w))$
 $\text{bin-last } (\neg \text{numeral } (\text{Num.Bit1 } w))$
 $\langle \text{proof} \rangle$

lemma *bin-rest-numeral-simps* [simp]:

$\text{bin-rest } 0 = 0$
 $\text{bin-rest } 1 = 0$
 $\text{bin-rest } (-1) = -1$
 $\text{bin-rest Numeral1} = 0$
 $\text{bin-rest } (\text{numeral } (\text{Num.Bit0 } w)) = \text{numeral } w$
 $\text{bin-rest } (\text{numeral } (\text{Num.Bit1 } w)) = \text{numeral } w$
 $\text{bin-rest } (\neg \text{numeral } (\text{Num.Bit0 } w)) = \neg \text{numeral } w$
 $\text{bin-rest } (\neg \text{numeral } (\text{Num.Bit1 } w)) = \neg \text{numeral } (w + \text{Num.One})$
 $\langle \text{proof} \rangle$

lemma *less-Bits*: $v \text{ BIT } b < w \text{ BIT } c \longleftrightarrow v < w \vee v \leq w \wedge \neg b \wedge c$
 $\langle \text{proof} \rangle$

lemma *le-Bits*: $v \text{ BIT } b \leq w \text{ BIT } c \longleftrightarrow v < w \vee v \leq w \wedge (\neg b \vee c)$
 $\langle \text{proof} \rangle$

lemma *pred-BIT-simps* [simp]:

$x \text{ BIT False} - 1 = (x - 1) \text{ BIT True}$
 $x \text{ BIT True} - 1 = x \text{ BIT False}$
 $\langle \text{proof} \rangle$

lemma *succ-BIT-simps* [simp]:

$x \text{ BIT False} + 1 = x \text{ BIT True}$
 $x \text{ BIT True} + 1 = (x + 1) \text{ BIT False}$
 $\langle \text{proof} \rangle$

lemma *add-BIT-simps* [simp]:

$x \text{ BIT False} + y \text{ BIT False} = (x + y) \text{ BIT False}$
 $x \text{ BIT False} + y \text{ BIT True} = (x + y) \text{ BIT True}$
 $x \text{ BIT True} + y \text{ BIT False} = (x + y) \text{ BIT True}$
 $x \text{ BIT True} + y \text{ BIT True} = (x + y + 1) \text{ BIT False}$

$\langle \text{proof} \rangle$

lemma *mult-BIT-simps* [simp]:

$x \text{ BIT False} * y = (x * y) \text{ BIT False}$
 $x * y \text{ BIT False} = (x * y) \text{ BIT False}$
 $x \text{ BIT True} * y = (x * y) \text{ BIT False} + y$
 $\langle \text{proof} \rangle$

lemma *B-mod-2'*: $X = 2 \implies (w \text{ BIT True}) \bmod X = 1 \wedge (w \text{ BIT False}) \bmod X = 0$
 $\langle \text{proof} \rangle$

lemma *bin-ex-rl*: $\exists w \ b. \ w \text{ BIT } b = \text{bin}$
 $\langle \text{proof} \rangle$

lemma *bin-exhaust*:

assumes *that*: $\bigwedge x \ b. \ \text{bin} = x \text{ BIT } b \implies Q$
shows Q
 $\langle \text{proof} \rangle$

primrec *bin-nth*

where

$Z: \text{bin-nth } w \ 0 \longleftrightarrow \text{bin-last } w$
 $| \text{Suc}: \text{bin-nth } w \ (\text{Suc } n) \longleftrightarrow \text{bin-nth } (\text{bin-rest } w) \ n$

lemma *bin-abs-lem*: $\text{bin} = (w \text{ BIT } b) \implies \text{bin} \neq -1 \longrightarrow \text{bin} \neq 0 \longrightarrow \text{nat } |w| < \text{nat } |\text{bin}|$
 $\langle \text{proof} \rangle$

lemma *bin-induct*:

assumes *PPls*: $P \ 0$
and *PMin*: $P \ (-1)$
and *PBit*: $\bigwedge \text{bin } \text{bit}. \ P \ \text{bin} \implies P \ (\text{bin BIT bit})$
shows $P \ \text{bin}$
 $\langle \text{proof} \rangle$

lemma *Bit-div2* [simp]: $(w \text{ BIT } b) \text{ div } 2 = w$
 $\langle \text{proof} \rangle$

lemma *bin-nth-eq-iff*: $\text{bin-nth } x = \text{bin-nth } y \longleftrightarrow x = y$
 $\langle \text{proof} \rangle$

lemmas *bin-eqI* = *ext* [THEN *bin-nth-eq-iff*] [THEN *iffD1*]

lemma *bin-eq-iff*: $x = y \longleftrightarrow (\forall n. \ \text{bin-nth } x \ n = \text{bin-nth } y \ n)$
 $\langle \text{proof} \rangle$

lemma *bin-nth-zero* [simp]: $\neg \text{bin-nth } 0 \ n$
 $\langle \text{proof} \rangle$

lemma *bin-nth-1* [simp]: $\text{bin-nth } 1 \ n \longleftrightarrow n = 0$
 ⟨proof⟩

lemma *bin-nth-minus1* [simp]: $\text{bin-nth } (-1) \ n$
 ⟨proof⟩

lemma *bin-nth-0-BIT*: $\text{bin-nth } (w \text{ BIT } b) \ 0 \longleftrightarrow b$
 ⟨proof⟩

lemma *bin-nth-Suc-BIT*: $\text{bin-nth } (w \text{ BIT } b) \ (\text{Suc } n) = \text{bin-nth } w \ n$
 ⟨proof⟩

lemma *bin-nth-minus* [simp]: $0 < n \implies \text{bin-nth } (w \text{ BIT } b) \ n = \text{bin-nth } w \ (n - 1)$
 ⟨proof⟩

lemma *bin-nth-numeral*: $\text{bin-rest } x = y \implies \text{bin-nth } x \ (\text{numeral } n) = \text{bin-nth } y \ (\text{pred-numeral } n)$
 ⟨proof⟩

lemmas *bin-nth-numeral-simps* [simp] =
bin-nth-numeral [OF *bin-rest-numeral-simps*(2)]
bin-nth-numeral [OF *bin-rest-numeral-simps*(5)]
bin-nth-numeral [OF *bin-rest-numeral-simps*(6)]
bin-nth-numeral [OF *bin-rest-numeral-simps*(7)]
bin-nth-numeral [OF *bin-rest-numeral-simps*(8)]

lemmas *bin-nth-simps* =
bin-nth.Z *bin-nth.Suc* *bin-nth-zero* *bin-nth-minus1*
bin-nth-numeral-simps

4.2 Truncating binary integers

definition *bin-sign* :: $\text{int} \Rightarrow \text{int}$
 where *bin-sign* $k = (\text{if } k \geq 0 \text{ then } 0 \text{ else } -1)$

lemma *bin-sign-simps* [simp]:
bin-sign $0 = 0$
bin-sign $1 = 0$
bin-sign $(-1) = -1$
bin-sign $(\text{numeral } k) = 0$
bin-sign $(-\text{numeral } k) = -1$
bin-sign $(w \text{ BIT } b) = \text{bin-sign } w$
 ⟨proof⟩

lemma *bin-sign-rest* [simp]: $\text{bin-sign } (\text{bin-rest } w) = \text{bin-sign } w$
 ⟨proof⟩

primrec *bintrunc* :: *nat* $\Rightarrow$ *int* $\Rightarrow$ *int*

where

Z : *bintrunc* 0 *bin* = 0

| *Suc* : *bintrunc* (*Suc* *n*) *bin* = *bintrunc* *n* (*bin-rest* *bin*) *BIT* (*bin-last* *bin*)

primrec *sbintrunc* :: *nat* $\Rightarrow$ *int* $\Rightarrow$ *int*

where

Z : *sbintrunc* 0 *bin* = (if *bin-last* *bin* then -1 else 0)

| *Suc* : *sbintrunc* (*Suc* *n*) *bin* = *sbintrunc* *n* (*bin-rest* *bin*) *BIT* (*bin-last* *bin*)

lemma *sign-bintr*: *bin-sign* (*bintrunc* *n* *w*) = 0

<proof>

lemma *bintrunc-mod2p*: *bintrunc* *n* *w* = *w* mod $2^{\wedge} n$

<proof>

lemma *sbintrunc-mod2p*: *sbintrunc* *n* *w* = (*w* + $2^{\wedge} n$) mod $2^{\wedge} (\text{Suc } n) - 2^{\wedge} n$

<proof>

4.3 Simplifications for (s)bintrunc

lemma *bintrunc-n-0* [*simp*]: *bintrunc* *n* 0 = 0

<proof>

lemma *sbintrunc-n-0* [*simp*]: *sbintrunc* *n* 0 = 0

<proof>

lemma *sbintrunc-n-minus1* [*simp*]: *sbintrunc* *n* (- 1) = -1

<proof>

lemma *bintrunc-Suc-numeral*:

bintrunc (*Suc* *n*) 1 = 1

bintrunc (*Suc* *n*) (- 1) = *bintrunc* *n* (- 1) *BIT* *True*

bintrunc (*Suc* *n*) (*numeral* (*Num.Bit0* *w*)) = *bintrunc* *n* (*numeral* *w*) *BIT* *False*

bintrunc (*Suc* *n*) (*numeral* (*Num.Bit1* *w*)) = *bintrunc* *n* (*numeral* *w*) *BIT* *True*

bintrunc (*Suc* *n*) (- *numeral* (*Num.Bit0* *w*)) = *bintrunc* *n* (- *numeral* *w*) *BIT* *False*

bintrunc (*Suc* *n*) (- *numeral* (*Num.Bit1* *w*)) = *bintrunc* *n* (- *numeral* (*w* + *Num.One*)) *BIT* *True*

<proof>

lemma *sbintrunc-0-numeral* [*simp*]:

sbintrunc 0 1 = -1

sbintrunc 0 (*numeral* (*Num.Bit0* *w*)) = 0

sbintrunc 0 (*numeral* (*Num.Bit1* *w*)) = -1

sbintrunc 0 (- *numeral* (*Num.Bit0* *w*)) = 0

sbintrunc 0 (- *numeral* (*Num.Bit1* *w*)) = -1

<proof>

lemma *sbintrunc-Suc-numeral*:

$sbintrunc (Suc\ n)\ 1 = 1$
 $sbintrunc (Suc\ n)\ (numeral\ (Num.Bit0\ w)) = sbintrunc\ n\ (numeral\ w)\ BIT\ False$
 $sbintrunc (Suc\ n)\ (numeral\ (Num.Bit1\ w)) = sbintrunc\ n\ (numeral\ w)\ BIT\ True$
 $sbintrunc (Suc\ n)\ (-\ numeral\ (Num.Bit0\ w)) = sbintrunc\ n\ (-\ numeral\ w)\ BIT\ False$
 $sbintrunc (Suc\ n)\ (-\ numeral\ (Num.Bit1\ w)) = sbintrunc\ n\ (-\ numeral\ (w + Num.One))\ BIT\ True$
 $\langle proof \rangle$

lemma *bin-sign-lem*: $(bin_sign\ (sbintrunc\ n\ bin) = -1) = bin_nth\ bin\ n$
 $\langle proof \rangle$

lemma *nth-bintr*: $bin_nth\ (bintrunc\ m\ w)\ n \longleftrightarrow n < m \wedge bin_nth\ w\ n$
 $\langle proof \rangle$

lemma *nth-sbintr*: $bin_nth\ (sbintrunc\ m\ w)\ n = (if\ n < m\ then\ bin_nth\ w\ n\ else\ bin_nth\ w\ m)$
 $\langle proof \rangle$

lemma *bin-nth-Bit*: $bin_nth\ (w\ BIT\ b)\ n \longleftrightarrow n = 0 \wedge b \vee (\exists\ m.\ n = Suc\ m \wedge bin_nth\ w\ m)$
 $\langle proof \rangle$

lemma *bin-nth-Bit0*:
 $bin_nth\ (numeral\ (Num.Bit0\ w))\ n \longleftrightarrow$
 $(\exists\ m.\ n = Suc\ m \wedge bin_nth\ (numeral\ w)\ m)$
 $\langle proof \rangle$

lemma *bin-nth-Bit1*:
 $bin_nth\ (numeral\ (Num.Bit1\ w))\ n \longleftrightarrow$
 $n = 0 \vee (\exists\ m.\ n = Suc\ m \wedge bin_nth\ (numeral\ w)\ m)$
 $\langle proof \rangle$

lemma *bintrunc-bintrunc-l*: $n \leq m \implies bintrunc\ m\ (bintrunc\ n\ w) = bintrunc\ n\ w$
 $\langle proof \rangle$

lemma *sbintrunc-sbintrunc-l*: $n \leq m \implies sbintrunc\ m\ (sbintrunc\ n\ w) = sbintrunc\ n\ w$
 $\langle proof \rangle$

lemma *bintrunc-bintrunc-ge*: $n \leq m \implies bintrunc\ n\ (bintrunc\ m\ w) = bintrunc\ n\ w$
 $\langle proof \rangle$

lemma *bintrunc-bintrunc-min [simp]*: $bintrunc\ m\ (bintrunc\ n\ w) = bintrunc\ (min\ m\ n)\ w$
 $\langle proof \rangle$

lemma *sbintrunc-sbintrunc-min* [simp]: *sbintrunc m (sbintrunc n w) = sbintrunc (min m n) w*
 ⟨proof⟩

lemmas *bintrunc-Pls* =
bintrunc.Suc [where *bin=0*, simplified *bin-last-numeral-simps bin-rest-numeral-simps*]

lemmas *bintrunc-Min* [simp] =
bintrunc.Suc [where *bin=-1*, simplified *bin-last-numeral-simps bin-rest-numeral-simps*]

lemmas *bintrunc-BIT* [simp] =
bintrunc.Suc [where *bin=w BIT b*, simplified *bin-last-BIT bin-rest-BIT*] for *w b*

lemmas *bintrunc-Sucs* = *bintrunc-Pls bintrunc-Min bintrunc-BIT*
bintrunc-Suc-numeral

lemmas *sbintrunc-Suc-Pls* =
sbintrunc.Suc [where *bin=0*, simplified *bin-last-numeral-simps bin-rest-numeral-simps*]

lemmas *sbintrunc-Suc-Min* =
sbintrunc.Suc [where *bin=-1*, simplified *bin-last-numeral-simps bin-rest-numeral-simps*]

lemmas *sbintrunc-Suc-BIT* [simp] =
sbintrunc.Suc [where *bin=w BIT b*, simplified *bin-last-BIT bin-rest-BIT*] for *w b*

lemmas *sbintrunc-Sucs* = *sbintrunc-Suc-Pls sbintrunc-Suc-Min sbintrunc-Suc-BIT*
sbintrunc-Suc-numeral

lemmas *sbintrunc-Pls* =
sbintrunc.Z [where *bin=0*, simplified *bin-last-numeral-simps bin-rest-numeral-simps*]

lemmas *sbintrunc-Min* =
sbintrunc.Z [where *bin=-1*, simplified *bin-last-numeral-simps bin-rest-numeral-simps*]

lemmas *sbintrunc-0-BIT-B0* [simp] =
sbintrunc.Z [where *bin=w BIT False*, simplified *bin-last-numeral-simps bin-rest-numeral-simps*]
 for *w*

lemmas *sbintrunc-0-BIT-B1* [simp] =
sbintrunc.Z [where *bin=w BIT True*, simplified *bin-last-BIT bin-rest-numeral-simps*]
 for *w*

lemmas *sbintrunc-0-simps* =
sbintrunc-Pls sbintrunc-Min sbintrunc-0-BIT-B0 sbintrunc-0-BIT-B1

lemmas *bintrunc-simps* = *bintrunc.Z bintrunc-Sucs*

lemmas *sbintrunc-simps* = *sbintrunc-0-simps sbintrunc-Sucs*

lemma *bintrunc-minus*: $0 < n \implies \text{bintrunc } (\text{Suc } (n - 1)) \ w = \text{bintrunc } n \ w$
 $\langle \text{proof} \rangle$

lemma *sbintrunc-minus*: $0 < n \implies \text{sbintrunc } (\text{Suc } (n - 1)) \ w = \text{sbintrunc } n \ w$
 $\langle \text{proof} \rangle$

lemmas *bintrunc-minus-simps* =
 $\text{bintrunc-Sucs } [\text{THEN } [2] \ \text{bintrunc-minus } [\text{symmetric}, \ \text{THEN } \text{trans}]]$
lemmas *sbintrunc-minus-simps* =
 $\text{sbintrunc-Sucs } [\text{THEN } [2] \ \text{sbintrunc-minus } [\text{symmetric}, \ \text{THEN } \text{trans}]]$

lemmas *thobini1* = $\text{arg-cong } [\text{where } f = \lambda w. \ w \ \text{BIT } b] \ \text{for } b$

lemmas *bintrunc-BIT-I* = $\text{trans } [\text{OF } \text{bintrunc-BIT } \text{thobini1}]$
lemmas *bintrunc-Min-I* = $\text{trans } [\text{OF } \text{bintrunc-Min } \text{thobini1}]$

lemmas *bmsts* = $\text{bintrunc-minus-simps}(1-3) \ [\text{THEN } \text{thobini1} \ [\text{THEN } [2] \ \text{trans}]]$
lemmas *bintrunc-Pls-minus-I* = $\text{bmsts}(1)$
lemmas *bintrunc-Min-minus-I* = $\text{bmsts}(2)$
lemmas *bintrunc-BIT-minus-I* = $\text{bmsts}(3)$

lemma *bintrunc-Suc-lem*: $\text{bintrunc } (\text{Suc } n) \ x = y \implies m = \text{Suc } n \implies \text{bintrunc } m \ x = y$
 $\langle \text{proof} \rangle$

lemmas *bintrunc-Suc-Ialts* =
 $\text{bintrunc-Min-I } [\text{THEN } \text{bintrunc-Suc-lem}]$
 $\text{bintrunc-BIT-I } [\text{THEN } \text{bintrunc-Suc-lem}]$

lemmas *sbintrunc-BIT-I* = $\text{trans } [\text{OF } \text{sbintrunc-Suc-BIT } \text{thobini1}]$

lemmas *sbintrunc-Suc-Is* =
 $\text{sbintrunc-Sucs}(1-3) \ [\text{THEN } \text{thobini1} \ [\text{THEN } [2] \ \text{trans}]]$

lemmas *sbintrunc-Suc-minus-Is* =
 $\text{sbintrunc-minus-simps}(1-3) \ [\text{THEN } \text{thobini1} \ [\text{THEN } [2] \ \text{trans}]]$

lemma *sbintrunc-Suc-lem*: $\text{sbintrunc } (\text{Suc } n) \ x = y \implies m = \text{Suc } n \implies \text{sbintrunc } m \ x = y$
 $\langle \text{proof} \rangle$

lemmas *sbintrunc-Suc-Ialts* =
 $\text{sbintrunc-Suc-Is } [\text{THEN } \text{sbintrunc-Suc-lem}]$

lemma *sbintrunc-bintrunc-lt*: $m > n \implies \text{sbintrunc } n \ (\text{bintrunc } m \ w) = \text{sbintrunc } n \ w$
 $\langle \text{proof} \rangle$

lemma *bintrunc-sbintrunc-le*: $m \leq \text{Suc } n \implies \text{bintrunc } m (\text{sbintrunc } n \ w) = \text{bintrunc } m \ w$

<proof>

lemmas *bintrunc-sbintrunc* [simp] = order-refl [THEN *bintrunc-sbintrunc-le*]

lemmas *sbintrunc-bintrunc* [simp] = lessI [THEN *sbintrunc-bintrunc-lt*]

lemmas *bintrunc-bintrunc* [simp] = order-refl [THEN *bintrunc-bintrunc-l*]

lemmas *sbintrunc-sbintrunc* [simp] = order-refl [THEN *sbintrunc-sbintrunc-l*]

lemma *bintrunc-sbintrunc'* [simp]: $0 < n \implies \text{bintrunc } n (\text{sbintrunc } (n - 1) \ w) = \text{bintrunc } n \ w$

<proof>

lemma *sbintrunc-bintrunc'* [simp]: $0 < n \implies \text{sbintrunc } (n - 1) (\text{bintrunc } n \ w) = \text{sbintrunc } (n - 1) \ w$

<proof>

lemma *bin-sbin-eq-iff*: $\text{bintrunc } (\text{Suc } n) \ x = \text{bintrunc } (\text{Suc } n) \ y \longleftrightarrow \text{sbintrunc } n \ x = \text{sbintrunc } n \ y$

<proof>

lemma *bin-sbin-eq-iff'*:

$0 < n \implies \text{bintrunc } n \ x = \text{bintrunc } n \ y \longleftrightarrow \text{sbintrunc } (n - 1) \ x = \text{sbintrunc } (n - 1) \ y$

<proof>

lemmas *bintrunc-sbintruncS0* [simp] = *bintrunc-sbintrunc'* [unfolded *One-nat-def*]

lemmas *sbintrunc-bintruncS0* [simp] = *sbintrunc-bintrunc'* [unfolded *One-nat-def*]

lemmas *bintrunc-bintrunc-l'* = *le-add1* [THEN *bintrunc-bintrunc-l*]

lemmas *sbintrunc-sbintrunc-l'* = *le-add1* [THEN *sbintrunc-sbintrunc-l*]

lemmas *nat-non0-gr* =

trans [*OF iszero-def* [THEN *Not-eq-iff* [THEN *iffD2*]] *refl*]

lemma *bintrunc-numeral*:

$\text{bintrunc } (\text{numeral } k) \ x = \text{bintrunc } (\text{pred-numeral } k) (\text{bin-rest } x) \ \text{BIT } \text{bin-last } x$

<proof>

lemma *sbintrunc-numeral*:

$\text{sbintrunc } (\text{numeral } k) \ x = \text{sbintrunc } (\text{pred-numeral } k) (\text{bin-rest } x) \ \text{BIT } \text{bin-last } x$

<proof>

lemma *bintrunc-numeral-simps* [simp]:

$\text{bintrunc } (\text{numeral } k) (\text{numeral } (\text{Num.Bit0 } w)) = \text{bintrunc } (\text{pred-numeral } k) (\text{numeral } w) \ \text{BIT } \text{False}$

$\text{bintrunc } (\text{numeral } k) (\text{numeral } (\text{Num.Bit1 } w)) = \text{bintrunc } (\text{pred-numeral } k)$
 $(\text{numeral } w) \text{ BIT True}$
 $\text{bintrunc } (\text{numeral } k) (- \text{numeral } (\text{Num.Bit0 } w)) = \text{bintrunc } (\text{pred-numeral } k)$
 $(- \text{numeral } w) \text{ BIT False}$
 $\text{bintrunc } (\text{numeral } k) (- \text{numeral } (\text{Num.Bit1 } w)) =$
 $\text{bintrunc } (\text{pred-numeral } k) (- \text{numeral } (w + \text{Num.One})) \text{ BIT True}$
 $\text{bintrunc } (\text{numeral } k) 1 = 1$
 $\langle \text{proof} \rangle$

lemma *sbintrunc-numeral-simps* [simp]:
 $\text{sbintrunc } (\text{numeral } k) (\text{numeral } (\text{Num.Bit0 } w)) = \text{sbintrunc } (\text{pred-numeral } k)$
 $(\text{numeral } w) \text{ BIT False}$
 $\text{sbintrunc } (\text{numeral } k) (\text{numeral } (\text{Num.Bit1 } w)) = \text{sbintrunc } (\text{pred-numeral } k)$
 $(\text{numeral } w) \text{ BIT True}$
 $\text{sbintrunc } (\text{numeral } k) (- \text{numeral } (\text{Num.Bit0 } w)) =$
 $\text{sbintrunc } (\text{pred-numeral } k) (- \text{numeral } w) \text{ BIT False}$
 $\text{sbintrunc } (\text{numeral } k) (- \text{numeral } (\text{Num.Bit1 } w)) =$
 $\text{sbintrunc } (\text{pred-numeral } k) (- \text{numeral } (w + \text{Num.One})) \text{ BIT True}$
 $\text{sbintrunc } (\text{numeral } k) 1 = 1$
 $\langle \text{proof} \rangle$

lemma *no-bintr-alt1*: $\text{bintrunc } n = (\lambda w. w \bmod 2^n :: \text{int})$
 $\langle \text{proof} \rangle$

lemma *range-bintrunc*: $\text{range } (\text{bintrunc } n) = \{i. 0 \leq i \wedge i < 2^n\}$
 $\langle \text{proof} \rangle$

lemma *no-sbintr-alt2*: $\text{sbintrunc } n = (\lambda w. (w + 2^n) \bmod 2^{Suc\ n} - 2^n :: \text{int})$
 $\langle \text{proof} \rangle$

lemma *range-sbintrunc*: $\text{range } (\text{sbintrunc } n) = \{i. -(2^n) \leq i \wedge i < 2^n\}$
 $\langle \text{proof} \rangle$

lemma *sb-inc-lem*: $a + 2^k < 0 \implies a + 2^k + 2^{(Suc\ k)} \leq (a + 2^k) \bmod 2^{(Suc\ k)}$
for $a :: \text{int}$
 $\langle \text{proof} \rangle$

lemma *sb-inc-lem'*: $a < -(2^k) \implies a + 2^k + 2^{(Suc\ k)} \leq (a + 2^k) \bmod 2^{(Suc\ k)}$
for $a :: \text{int}$
 $\langle \text{proof} \rangle$

lemma *sbintrunc-inc*: $x < -(2^n) \implies x + 2^n \leq \text{sbintrunc } n\ x$
 $\langle \text{proof} \rangle$

lemma *sb-dec-lem*: $0 \leq -(2^k) + a \implies (a + 2^k) \bmod (2 * 2^k) \leq -(2^k) + a$

for $a :: \text{int}$
 $\langle \text{proof} \rangle$

lemma *sb-dec-lem'*: $2^k \leq a \implies (a + 2^k) \bmod (2 * 2^k) \leq - (2^k) + a$
for $a :: \text{int}$
 $\langle \text{proof} \rangle$

lemma *sbintrunc-dec*: $x \geq (2^n) \implies x - 2^n (\text{Suc } n) \geq \text{sbintrunc } n \ x$
 $\langle \text{proof} \rangle$

lemmas *m2pths* = *pos-mod-sign pos-mod-bound* [*OF zless2p*]

lemma *bintr-ge0*: $0 \leq \text{bintrunc } n \ w$
 $\langle \text{proof} \rangle$

lemma *bintr-lt2p*: $\text{bintrunc } n \ w < 2^n$
 $\langle \text{proof} \rangle$

lemma *bintr-Min*: $\text{bintrunc } n \ (-1) = 2^n - 1$
 $\langle \text{proof} \rangle$

lemma *sbintr-ge*: $-(2^n) \leq \text{sbintrunc } n \ w$
 $\langle \text{proof} \rangle$

lemma *sbintr-lt*: $\text{sbintrunc } n \ w < 2^n$
 $\langle \text{proof} \rangle$

lemma *sign-Pls-ge-0*: $\text{bin-sign } \text{bin} = 0 \longleftrightarrow \text{bin} \geq 0$
for $\text{bin} :: \text{int}$
 $\langle \text{proof} \rangle$

lemma *sign-Min-lt-0*: $\text{bin-sign } \text{bin} = -1 \longleftrightarrow \text{bin} < 0$
for $\text{bin} :: \text{int}$
 $\langle \text{proof} \rangle$

lemma *bin-rest-trunc*: $\text{bin-rest } (\text{bintrunc } n \ \text{bin}) = \text{bintrunc } (n - 1) \ (\text{bin-rest } \text{bin})$
 $\langle \text{proof} \rangle$

lemma *bin-rest-power-trunc*:
 $(\text{bin-rest } ^k) \ (\text{bintrunc } n \ \text{bin}) = \text{bintrunc } (n - k) \ ((\text{bin-rest } ^k) \ \text{bin})$
 $\langle \text{proof} \rangle$

lemma *bin-rest-trunc-i*: $\text{bintrunc } n \ (\text{bin-rest } \text{bin}) = \text{bin-rest } (\text{bintrunc } (\text{Suc } n) \ \text{bin})$
 $\langle \text{proof} \rangle$

lemma *bin-rest-strunc*: $\text{bin-rest } (\text{sbintrunc } (\text{Suc } n) \ \text{bin}) = \text{sbintrunc } n \ (\text{bin-rest } \text{bin})$
 $\langle \text{proof} \rangle$

lemma *bintrunc-rest* [simp]: $\text{bintrunc } n \ (\text{bin-rest } (\text{bintrunc } n \ \text{bin})) = \text{bin-rest } (\text{bintrunc } n \ \text{bin})$
 ⟨proof⟩

lemma *sbintrunc-rest* [simp]: $\text{sbintrunc } n \ (\text{bin-rest } (\text{sbintrunc } n \ \text{bin})) = \text{bin-rest } (\text{sbintrunc } n \ \text{bin})$
 ⟨proof⟩

lemma *bintrunc-rest'*: $\text{bintrunc } n \circ \text{bin-rest} \circ \text{bintrunc } n = \text{bin-rest} \circ \text{bintrunc } n$
 ⟨proof⟩

lemma *sbintrunc-rest'*: $\text{sbintrunc } n \circ \text{bin-rest} \circ \text{sbintrunc } n = \text{bin-rest} \circ \text{sbintrunc } n$
 ⟨proof⟩

lemma *rco-lem*: $f \circ g \circ f = g \circ f \implies f \circ (g \circ f) \ \wedge \wedge \ n = g \ \wedge \wedge \ n \circ f$
 ⟨proof⟩

lemmas *rco-bintr* = *bintrunc-rest'*
 [THEN *rco-lem* [THEN *fun-cong*], unfolded *o-def*]

lemmas *rco-sbintr* = *sbintrunc-rest'*
 [THEN *rco-lem* [THEN *fun-cong*], unfolded *o-def*]

4.4 Splitting and concatenation

primrec *bin-split* :: $\text{nat} \Rightarrow \text{int} \Rightarrow \text{int} \times \text{int}$

where

Z : $\text{bin-split } 0 \ w = (w, 0)$
 $| \text{Suc}$: $\text{bin-split } (\text{Suc } n) \ w =$
 $(\text{let } (w1, w2) = \text{bin-split } n \ (\text{bin-rest } w)$
 $\text{in } (w1, w2 \ \text{BIT } \text{bin-last } w))$

lemma [code]:
 $\text{bin-split } (\text{Suc } n) \ w = (\text{let } (w1, w2) = \text{bin-split } n \ (\text{bin-rest } w) \ \text{in } (w1, w2 \ \text{BIT } \text{bin-last } w))$
 $\text{bin-split } 0 \ w = (w, 0)$
 ⟨proof⟩

primrec *bin-cat* :: $\text{int} \Rightarrow \text{nat} \Rightarrow \text{int} \Rightarrow \text{int}$

where

Z : $\text{bin-cat } w \ 0 \ v = w$
 $| \text{Suc}$: $\text{bin-cat } w \ (\text{Suc } n) \ v = \text{bin-cat } w \ n \ (\text{bin-rest } v) \ \text{BIT } \text{bin-last } v$

end

5 Bitwise Operations on Binary Integers

theory *Bits-Int*

```

imports Bits Bit-Representation
begin

```

5.1 Logical operations

bit-wise logical operations on the int type

```

instantiation int :: bit
begin

```

```

definition int-not-def: bitNOT = ( $\lambda x::int. - x - 1$ )

```

```

function bitAND-int
  where bitAND-int x y =
    (if x = 0 then 0 else if x = -1 then y
     else (bin-rest x AND bin-rest y) BIT (bin-last x  $\wedge$  bin-last y))
  <proof>

```

```

termination
  <proof>

```

```

declare bitAND-int.simps [simp del]

```

```

definition int-or-def:
  bitOR = ( $\lambda x y::int. NOT (NOT x AND NOT y)$ )

```

```

definition int-xor-def:
  bitXOR = ( $\lambda x y::int. (x AND NOT y) OR (NOT x AND y)$ )

```

```

instance <proof>

```

```

end

```

5.1.1 Basic simplification rules

```

lemma int-not-BIT [simp]:
  NOT (w BIT b) = (NOT w) BIT ( $\neg$  b)
  <proof>

```

```

lemma int-not-simps [simp]:
  NOT (0::int) = -1
  NOT (1::int) = -2
  NOT (- 1::int) = 0
  NOT (numeral w::int) = - numeral (w + Num.One)
  NOT (- numeral (Num.Bit0 w)::int) = numeral (Num.BitM w)
  NOT (- numeral (Num.Bit1 w)::int) = numeral (Num.Bit0 w)
  <proof>

```

```

lemma int-not-not [simp]: NOT (NOT (x::int)) = x
  <proof>

```

lemma *int-and-0* [simp]: $(0::int) \text{ AND } x = 0$
 $\langle \text{proof} \rangle$

lemma *int-and-m1* [simp]: $(-1::int) \text{ AND } x = x$
 $\langle \text{proof} \rangle$

lemma *int-and-Bits* [simp]:
 $(x \text{ BIT } b) \text{ AND } (y \text{ BIT } c) = (x \text{ AND } y) \text{ BIT } (b \wedge c)$
 $\langle \text{proof} \rangle$

lemma *int-or-zero* [simp]: $(0::int) \text{ OR } x = x$
 $\langle \text{proof} \rangle$

lemma *int-or-minus1* [simp]: $(-1::int) \text{ OR } x = -1$
 $\langle \text{proof} \rangle$

lemma *int-or-Bits* [simp]:
 $(x \text{ BIT } b) \text{ OR } (y \text{ BIT } c) = (x \text{ OR } y) \text{ BIT } (b \vee c)$
 $\langle \text{proof} \rangle$

lemma *int-xor-zero* [simp]: $(0::int) \text{ XOR } x = x$
 $\langle \text{proof} \rangle$

lemma *int-xor-Bits* [simp]:
 $(x \text{ BIT } b) \text{ XOR } (y \text{ BIT } c) = (x \text{ XOR } y) \text{ BIT } ((b \vee c) \wedge \neg (b \wedge c))$
 $\langle \text{proof} \rangle$

5.1.2 Binary destructors

lemma *bin-rest-NOT* [simp]: $\text{bin-rest } (\text{NOT } x) = \text{NOT } (\text{bin-rest } x)$
 $\langle \text{proof} \rangle$

lemma *bin-last-NOT* [simp]: $\text{bin-last } (\text{NOT } x) \longleftrightarrow \neg \text{bin-last } x$
 $\langle \text{proof} \rangle$

lemma *bin-rest-AND* [simp]: $\text{bin-rest } (x \text{ AND } y) = \text{bin-rest } x \text{ AND } \text{bin-rest } y$
 $\langle \text{proof} \rangle$

lemma *bin-last-AND* [simp]: $\text{bin-last } (x \text{ AND } y) \longleftrightarrow \text{bin-last } x \wedge \text{bin-last } y$
 $\langle \text{proof} \rangle$

lemma *bin-rest-OR* [simp]: $\text{bin-rest } (x \text{ OR } y) = \text{bin-rest } x \text{ OR } \text{bin-rest } y$
 $\langle \text{proof} \rangle$

lemma *bin-last-OR* [simp]: $\text{bin-last } (x \text{ OR } y) \longleftrightarrow \text{bin-last } x \vee \text{bin-last } y$
 $\langle \text{proof} \rangle$

lemma *bin-rest-XOR* [simp]: $\text{bin-rest } (x \text{ XOR } y) = \text{bin-rest } x \text{ XOR } \text{bin-rest } y$

$\langle \text{proof} \rangle$

lemma *bin-last-XOR* [simp]: $\text{bin-last } (x \text{ XOR } y) \longleftrightarrow (\text{bin-last } x \vee \text{bin-last } y) \wedge \neg (\text{bin-last } x \wedge \text{bin-last } y)$
 $\langle \text{proof} \rangle$

lemma *bin-nth-ops*:

$!!x \ y. \text{bin-nth } (x \text{ AND } y) \ n = (\text{bin-nth } x \ n \ \& \ \text{bin-nth } y \ n)$
 $!!x \ y. \text{bin-nth } (x \text{ OR } y) \ n = (\text{bin-nth } x \ n \ | \ \text{bin-nth } y \ n)$
 $!!x \ y. \text{bin-nth } (x \text{ XOR } y) \ n = (\text{bin-nth } x \ n \ \sim = \ \text{bin-nth } y \ n)$
 $!!x. \text{bin-nth } (\text{NOT } x) \ n = (\sim \ \text{bin-nth } x \ n)$
 $\langle \text{proof} \rangle$

5.1.3 Derived properties

lemma *int-xor-minus1* [simp]: $(-1::\text{int}) \text{ XOR } x = \text{NOT } x$
 $\langle \text{proof} \rangle$

lemma *int-xor-extra-simps* [simp]:
 $w \text{ XOR } (0::\text{int}) = w$
 $w \text{ XOR } (-1::\text{int}) = \text{NOT } w$
 $\langle \text{proof} \rangle$

lemma *int-or-extra-simps* [simp]:
 $w \text{ OR } (0::\text{int}) = w$
 $w \text{ OR } (-1::\text{int}) = -1$
 $\langle \text{proof} \rangle$

lemma *int-and-extra-simps* [simp]:
 $w \text{ AND } (0::\text{int}) = 0$
 $w \text{ AND } (-1::\text{int}) = w$
 $\langle \text{proof} \rangle$

lemma *bin-ops-comm*:

shows
int-and-comm: $!!y::\text{int}. x \text{ AND } y = y \text{ AND } x$ **and**
int-or-comm: $!!y::\text{int}. x \text{ OR } y = y \text{ OR } x$ **and**
int-xor-comm: $!!y::\text{int}. x \text{ XOR } y = y \text{ XOR } x$
 $\langle \text{proof} \rangle$

lemma *bin-ops-same* [simp]:
 $(x::\text{int}) \text{ AND } x = x$
 $(x::\text{int}) \text{ OR } x = x$
 $(x::\text{int}) \text{ XOR } x = 0$
 $\langle \text{proof} \rangle$

lemmas *bin-log-esimps* =
int-and-extra-simps int-or-extra-simps int-xor-extra-simps

int-and-0 int-and-m1 int-or-zero int-or-minus1 int-xor-zero int-xor-minus1

lemma *bbw-ao-absorb*:

$!!y::int. x \text{ AND } (y \text{ OR } x) = x \ \& \ x \text{ OR } (y \text{ AND } x) = x$
 $\langle \text{proof} \rangle$

lemma *bbw-ao-absorbs-other*:

$x \text{ AND } (x \text{ OR } y) = x \wedge (y \text{ AND } x) \text{ OR } x = (x::int)$
 $(y \text{ OR } x) \text{ AND } x = x \wedge x \text{ OR } (x \text{ AND } y) = (x::int)$
 $(x \text{ OR } y) \text{ AND } x = x \wedge (x \text{ AND } y) \text{ OR } x = (x::int)$
 $\langle \text{proof} \rangle$

lemmas *bbw-ao-absorbs* [simp] = *bbw-ao-absorb bbw-ao-absorbs-other*

lemma *int-xor-not*:

$!!y::int. (\text{NOT } x) \text{ XOR } y = \text{NOT } (x \text{ XOR } y) \ \&$
 $x \text{ XOR } (\text{NOT } y) = \text{NOT } (x \text{ XOR } y)$
 $\langle \text{proof} \rangle$

lemma *int-and-assoc*:

$(x \text{ AND } y) \text{ AND } (z::int) = x \text{ AND } (y \text{ AND } z)$
 $\langle \text{proof} \rangle$

lemma *int-or-assoc*:

$(x \text{ OR } y) \text{ OR } (z::int) = x \text{ OR } (y \text{ OR } z)$
 $\langle \text{proof} \rangle$

lemma *int-xor-assoc*:

$(x \text{ XOR } y) \text{ XOR } (z::int) = x \text{ XOR } (y \text{ XOR } z)$
 $\langle \text{proof} \rangle$

lemmas *bbw-assocs* = *int-and-assoc int-or-assoc int-xor-assoc*

lemma *bbw-lcs* [simp]:

$(y::int) \text{ AND } (x \text{ AND } z) = x \text{ AND } (y \text{ AND } z)$
 $(y::int) \text{ OR } (x \text{ OR } z) = x \text{ OR } (y \text{ OR } z)$
 $(y::int) \text{ XOR } (x \text{ XOR } z) = x \text{ XOR } (y \text{ XOR } z)$
 $\langle \text{proof} \rangle$

lemma *bbw-not-dist*:

$!!y::int. \text{NOT } (x \text{ OR } y) = (\text{NOT } x) \text{ AND } (\text{NOT } y)$
 $!!y::int. \text{NOT } (x \text{ AND } y) = (\text{NOT } x) \text{ OR } (\text{NOT } y)$
 $\langle \text{proof} \rangle$

lemma *bbw-oa-dist*:

$!!y \ z::int. (x \text{ AND } y) \text{ OR } z =$

$(x \text{ OR } z) \text{ AND } (y \text{ OR } z)$
 $\langle \text{proof} \rangle$

lemma *bbw-ao-dist*:

$!!y \text{ z}::\text{int}. (x \text{ OR } y) \text{ AND } z =$
 $(x \text{ AND } z) \text{ OR } (y \text{ AND } z)$
 $\langle \text{proof} \rangle$

5.1.4 Simplification with numerals

Cases for 0 and -1 are already covered by other simp rules.

lemma *bin-rl-eqI*: $\llbracket \text{bin-rest } x = \text{bin-rest } y; \text{bin-last } x = \text{bin-last } y \rrbracket \implies x = y$
 $\langle \text{proof} \rangle$

lemma *bin-rest-neg-numeral-BitM* [simp]:

$\text{bin-rest } (- \text{ numeral } (\text{Num.BitM } w)) = - \text{ numeral } w$
 $\langle \text{proof} \rangle$

lemma *bin-last-neg-numeral-BitM* [simp]:

$\text{bin-last } (- \text{ numeral } (\text{Num.BitM } w))$
 $\langle \text{proof} \rangle$

FIXME: The rule sets below are very large (24 rules for each operator). Is there a simpler way to do this?

lemma *int-and-numerals* [simp]:

$\text{numeral } (\text{Num.Bit0 } x) \text{ AND } \text{numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ AND } \text{numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ AND } \text{numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ AND } \text{numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ AND } \text{numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ AND } \text{numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ AND } \text{numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ AND } \text{numeral } y) \text{ BIT True}$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ AND } - \text{ numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ AND } - \text{ numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ AND } - \text{ numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ AND } - \text{ numeral } (y + \text{Num.One})) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ AND } - \text{ numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ AND } - \text{ numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ AND } - \text{ numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ AND } - \text{ numeral } (y + \text{Num.One})) \text{ BIT True}$
 $- \text{ numeral } (\text{Num.Bit0 } x) \text{ AND } \text{numeral } (\text{Num.Bit0 } y) = (- \text{ numeral } x \text{ AND } \text{numeral } y) \text{ BIT False}$
 $- \text{ numeral } (\text{Num.Bit0 } x) \text{ AND } \text{numeral } (\text{Num.Bit1 } y) = (- \text{ numeral } x \text{ AND } \text{numeral } y) \text{ BIT False}$
 $- \text{ numeral } (\text{Num.Bit1 } x) \text{ AND } \text{numeral } (\text{Num.Bit0 } y) = (- \text{ numeral } (x + \text{Num.One}) \text{ AND } \text{numeral } y) \text{ BIT False}$
 $- \text{ numeral } (\text{Num.Bit1 } x) \text{ AND } \text{numeral } (\text{Num.Bit1 } y) = (- \text{ numeral } (x + \text{Num.One}) \text{ AND } \text{numeral } y) \text{ BIT True}$

$- \text{numeral } (\text{Num.Bit0 } x) \text{ AND } - \text{numeral } (\text{Num.Bit0 } y) = (- \text{numeral } x \text{ AND } - \text{numeral } y) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ AND } - \text{numeral } (\text{Num.Bit1 } y) = (- \text{numeral } x \text{ AND } - \text{numeral } (y + \text{Num.One})) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ AND } - \text{numeral } (\text{Num.Bit0 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ AND } - \text{numeral } y) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ AND } - \text{numeral } (\text{Num.Bit1 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ AND } - \text{numeral } (y + \text{Num.One})) \text{ BIT True}$
 $(1::\text{int}) \text{ AND numeral } (\text{Num.Bit0 } y) = 0$
 $(1::\text{int}) \text{ AND numeral } (\text{Num.Bit1 } y) = 1$
 $(1::\text{int}) \text{ AND } - \text{numeral } (\text{Num.Bit0 } y) = 0$
 $(1::\text{int}) \text{ AND } - \text{numeral } (\text{Num.Bit1 } y) = 1$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ AND } (1::\text{int}) = 0$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ AND } (1::\text{int}) = 1$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ AND } (1::\text{int}) = 0$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ AND } (1::\text{int}) = 1$
 {proof}

lemma *int-or-numerals* [simp]:

$\text{numeral } (\text{Num.Bit0 } x) \text{ OR numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ OR numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ OR numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ OR numeral } y) \text{ BIT True}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ OR numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ OR numeral } y) \text{ BIT True}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ OR numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ OR numeral } y) \text{ BIT True}$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ OR } - \text{numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ OR } - \text{numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ OR } - \text{numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ OR } - \text{numeral } (y + \text{Num.One})) \text{ BIT True}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ OR } - \text{numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ OR } - \text{numeral } y) \text{ BIT True}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ OR } - \text{numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ OR } - \text{numeral } (y + \text{Num.One})) \text{ BIT True}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ OR numeral } (\text{Num.Bit0 } y) = (- \text{numeral } x \text{ OR numeral } y) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ OR numeral } (\text{Num.Bit1 } y) = (- \text{numeral } x \text{ OR numeral } y) \text{ BIT True}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ OR numeral } (\text{Num.Bit0 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ OR numeral } y) \text{ BIT True}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ OR numeral } (\text{Num.Bit1 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ OR numeral } y) \text{ BIT True}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ OR } - \text{numeral } (\text{Num.Bit0 } y) = (- \text{numeral } x \text{ OR } - \text{numeral } y) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ OR } - \text{numeral } (\text{Num.Bit1 } y) = (- \text{numeral } x \text{ OR } - \text{numeral } (y + \text{Num.One})) \text{ BIT True}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ OR } - \text{numeral } (\text{Num.Bit0 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ OR } - \text{numeral } y) \text{ BIT True}$

$- \text{numeral } (\text{Num.Bit1 } x) \text{ OR } - \text{numeral } (\text{Num.Bit1 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ OR } - \text{numeral } (y + \text{Num.One})) \text{ BIT True}$
 $(1::\text{int}) \text{ OR numeral } (\text{Num.Bit0 } y) = \text{numeral } (\text{Num.Bit1 } y)$
 $(1::\text{int}) \text{ OR numeral } (\text{Num.Bit1 } y) = \text{numeral } (\text{Num.Bit1 } y)$
 $(1::\text{int}) \text{ OR } - \text{numeral } (\text{Num.Bit0 } y) = - \text{numeral } (\text{Num.BitM } y)$
 $(1::\text{int}) \text{ OR } - \text{numeral } (\text{Num.Bit1 } y) = - \text{numeral } (\text{Num.Bit1 } y)$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ OR } (1::\text{int}) = \text{numeral } (\text{Num.Bit1 } x)$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ OR } (1::\text{int}) = \text{numeral } (\text{Num.Bit1 } x)$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ OR } (1::\text{int}) = - \text{numeral } (\text{Num.BitM } x)$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ OR } (1::\text{int}) = - \text{numeral } (\text{Num.Bit1 } x)$
 $\langle \text{proof} \rangle$

lemma *int-xor-numerals [simp]*:

$\text{numeral } (\text{Num.Bit0 } x) \text{ XOR numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ XOR numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ XOR numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ XOR numeral } y) \text{ BIT True}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ XOR numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ XOR numeral } y) \text{ BIT True}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ XOR numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ XOR numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ XOR } - \text{numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ XOR } - \text{numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ XOR } - \text{numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ XOR } - \text{numeral } (y + \text{Num.One})) \text{ BIT True}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ XOR } - \text{numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ XOR } - \text{numeral } y) \text{ BIT True}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ XOR } - \text{numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ XOR } - \text{numeral } (y + \text{Num.One})) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ XOR numeral } (\text{Num.Bit0 } y) = (- \text{numeral } x \text{ XOR numeral } y) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ XOR numeral } (\text{Num.Bit1 } y) = (- \text{numeral } x \text{ XOR numeral } y) \text{ BIT True}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ XOR numeral } (\text{Num.Bit0 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ XOR numeral } y) \text{ BIT True}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ XOR numeral } (\text{Num.Bit1 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ XOR numeral } y) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ XOR } - \text{numeral } (\text{Num.Bit0 } y) = (- \text{numeral } x \text{ XOR } - \text{numeral } y) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ XOR } - \text{numeral } (\text{Num.Bit1 } y) = (- \text{numeral } x \text{ XOR } - \text{numeral } (y + \text{Num.One})) \text{ BIT True}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ XOR } - \text{numeral } (\text{Num.Bit0 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ XOR } - \text{numeral } y) \text{ BIT True}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ XOR } - \text{numeral } (\text{Num.Bit1 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ XOR } - \text{numeral } (y + \text{Num.One})) \text{ BIT False}$
 $(1::\text{int}) \text{ XOR numeral } (\text{Num.Bit0 } y) = \text{numeral } (\text{Num.Bit1 } y)$
 $(1::\text{int}) \text{ XOR numeral } (\text{Num.Bit1 } y) = \text{numeral } (\text{Num.Bit0 } y)$
 $(1::\text{int}) \text{ XOR } - \text{numeral } (\text{Num.Bit0 } y) = - \text{numeral } (\text{Num.BitM } y)$
 $(1::\text{int}) \text{ XOR } - \text{numeral } (\text{Num.Bit1 } y) = - \text{numeral } (\text{Num.Bit0 } (y + \text{Num.One}))$

$\text{numeral } (\text{Num.Bit0 } x) \text{ XOR } (1::\text{int}) = \text{numeral } (\text{Num.Bit1 } x)$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ XOR } (1::\text{int}) = \text{numeral } (\text{Num.Bit0 } x)$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ XOR } (1::\text{int}) = - \text{numeral } (\text{Num.BitM } x)$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ XOR } (1::\text{int}) = - \text{numeral } (\text{Num.Bit0 } (x + \text{Num.One}))$
 $\langle \text{proof} \rangle$

5.1.5 Interactions with arithmetic

lemma *plus-and-or* [rule-format]:

$\text{ALL } y::\text{int}. (x \text{ AND } y) + (x \text{ OR } y) = x + y$
 $\langle \text{proof} \rangle$

lemma *le-int-or*:

$\text{bin-sign } (y::\text{int}) = 0 ==> x <= x \text{ OR } y$
 $\langle \text{proof} \rangle$

lemmas *int-and-le* =

$\text{xtrans}(3) [\text{OF } \text{bbw-ao-absorbs } (2) [\text{THEN } \text{conjunct2}, \text{symmetric}] \text{ le-int-or}]$

lemma *bin-add-not*: $x + \text{NOT } x = (-1::\text{int})$

$\langle \text{proof} \rangle$

5.1.6 Truncating results of bit-wise operations

lemma *bin-trunc-ao*:

$!!x \ y. (\text{bintrunc } n \ x) \text{ AND } (\text{bintrunc } n \ y) = \text{bintrunc } n \ (x \text{ AND } y)$
 $!!x \ y. (\text{bintrunc } n \ x) \text{ OR } (\text{bintrunc } n \ y) = \text{bintrunc } n \ (x \text{ OR } y)$
 $\langle \text{proof} \rangle$

lemma *bin-trunc-xor*:

$!!x \ y. \text{bintrunc } n \ (\text{bintrunc } n \ x \text{ XOR } \text{bintrunc } n \ y) =$
 $\text{bintrunc } n \ (x \text{ XOR } y)$
 $\langle \text{proof} \rangle$

lemma *bin-trunc-not*:

$!!x. \text{bintrunc } n \ (\text{NOT } (\text{bintrunc } n \ x)) = \text{bintrunc } n \ (\text{NOT } x)$
 $\langle \text{proof} \rangle$

lemma *bintr-bintr-i*:

$x = \text{bintrunc } n \ y ==> \text{bintrunc } n \ x = \text{bintrunc } n \ y$
 $\langle \text{proof} \rangle$

lemmas *bin-trunc-and* = *bin-trunc-ao*(1) [THEN *bintr-bintr-i*]

lemmas *bin-trunc-or* = *bin-trunc-ao*(2) [THEN *bintr-bintr-i*]

5.2 Setting and clearing bits

primrec

$bin-sc :: nat \Rightarrow bool \Rightarrow int \Rightarrow int$

where

$Z: bin-sc\ 0\ b\ w = bin-rest\ w\ BIT\ b$

$| Suc: bin-sc\ (Suc\ n)\ b\ w = bin-sc\ n\ b\ (bin-rest\ w)\ BIT\ bin-last\ w$

lemma *bin-nth-sc [simp]*:

$bin-nth\ (bin-sc\ n\ b\ w)\ n \longleftrightarrow b$

<proof>

lemma *bin-sc-sc-same [simp]*:

$bin-sc\ n\ c\ (bin-sc\ n\ b\ w) = bin-sc\ n\ c\ w$

<proof>

lemma *bin-sc-sc-diff*:

$m \sim = n \implies$

$bin-sc\ m\ c\ (bin-sc\ n\ b\ w) = bin-sc\ n\ b\ (bin-sc\ m\ c\ w)$

<proof>

lemma *bin-nth-sc-gen*:

$bin-nth\ (bin-sc\ n\ b\ w)\ m = (if\ m = n\ then\ b\ else\ bin-nth\ w\ m)$

<proof>

lemma *bin-sc-nth [simp]*:

$(bin-sc\ n\ (bin-nth\ w\ n)\ w) = w$

<proof>

lemma *bin-sign-sc [simp]*:

$bin-sign\ (bin-sc\ n\ b\ w) = bin-sign\ w$

<proof>

lemma *bin-sc-bintr [simp]*:

$bintrunc\ m\ (bin-sc\ n\ x\ (bintrunc\ m\ (w))) = bintrunc\ m\ (bin-sc\ n\ x\ w)$

<proof>

lemma *bin-clr-le*:

$bin-sc\ n\ False\ w \leq w$

<proof>

lemma *bin-set-ge*:

$bin-sc\ n\ True\ w \geq w$

<proof>

lemma *bintr-bin-clr-le*:

$bintrunc\ n\ (bin-sc\ m\ False\ w) \leq bintrunc\ n\ w$

<proof>

lemma *bintr-bin-set-ge*:

$\text{bintrunc } n \ (\text{bin-sc } m \ \text{True } w) \geq \text{bintrunc } n \ w$
 $\langle \text{proof} \rangle$

lemma *bin-sc-FP* [simp]: $\text{bin-sc } n \ \text{False } 0 = 0$
 $\langle \text{proof} \rangle$

lemma *bin-sc-TM* [simp]: $\text{bin-sc } n \ \text{True } (-\ 1) = -\ 1$
 $\langle \text{proof} \rangle$

lemmas *bin-sc-simps* = *bin-sc.Z bin-sc.Suc bin-sc-TM bin-sc-FP*

lemma *bin-sc-minus*:
 $0 < n \implies \text{bin-sc } (\text{Suc } (n - 1)) \ b \ w = \text{bin-sc } n \ b \ w$
 $\langle \text{proof} \rangle$

lemmas *bin-sc-Suc-minus* =
trans [OF bin-sc-minus [symmetric] bin-sc.Suc]

lemma *bin-sc-numeral* [simp]:
 $\text{bin-sc } (\text{numeral } k) \ b \ w =$
 $\text{bin-sc } (\text{pred-numeral } k) \ b \ (\text{bin-rest } w) \ \text{BIT } \text{bin-last } w$
 $\langle \text{proof} \rangle$

5.3 Splitting and concatenation

definition *bin-rcat* :: $\text{nat} \Rightarrow \text{int list} \Rightarrow \text{int}$
where

$\text{bin-rcat } n = \text{foldl } (\lambda u \ v. \text{bin-cat } u \ n \ v) \ 0$

fun *bin-rsplit-aux* :: $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{int} \Rightarrow \text{int list} \Rightarrow \text{int list}$
where

$\text{bin-rsplit-aux } n \ m \ c \ bs =$
 $(\text{if } m = 0 \mid n = 0 \text{ then } bs \text{ else}$
 $\text{let } (a, b) = \text{bin-split } n \ c$
 $\text{in } \text{bin-rsplit-aux } n \ (m - n) \ a \ (b \# bs))$

definition *bin-rsplit* :: $\text{nat} \Rightarrow \text{nat} \times \text{int} \Rightarrow \text{int list}$
where

$\text{bin-rsplit } n \ w = \text{bin-rsplit-aux } n \ (\text{fst } w) \ (\text{snd } w) \ []$

fun *bin-rsplittl-aux* :: $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{int} \Rightarrow \text{int list} \Rightarrow \text{int list}$
where

$\text{bin-rsplittl-aux } n \ m \ c \ bs =$
 $(\text{if } m = 0 \mid n = 0 \text{ then } bs \text{ else}$
 $\text{let } (a, b) = \text{bin-split } (\min m \ n) \ c$
 $\text{in } \text{bin-rsplittl-aux } n \ (m - n) \ a \ (b \# bs))$

definition *bin-rsplittl* :: $\text{nat} \Rightarrow \text{nat} \times \text{int} \Rightarrow \text{int list}$
where

$$\text{bin-rsplittl } n \ w = \text{bin-rsplittl-aux } n \ (\text{fst } w) \ (\text{snd } w) \ []$$

declare *bin-rsplit-aux.simps* [simp del]

declare *bin-rsplittl-aux.simps* [simp del]

lemma *bin-sign-cat*:

$$\text{bin-sign } (\text{bin-cat } x \ n \ y) = \text{bin-sign } x$$

⟨proof⟩

lemma *bin-cat-Suc-Bit*:

$$\text{bin-cat } w \ (\text{Suc } n) \ (v \ \text{BIT } b) = \text{bin-cat } w \ n \ v \ \text{BIT } b$$

⟨proof⟩

lemma *bin-nth-cat*:

$$\text{bin-nth } (\text{bin-cat } x \ k \ y) \ n =$$

$$(\text{if } n < k \text{ then } \text{bin-nth } y \ n \text{ else } \text{bin-nth } x \ (n - k))$$

⟨proof⟩

lemma *bin-nth-split*:

$$\text{bin-split } n \ c = (a, b) ==>$$

$$(\text{ALL } k. \text{bin-nth } a \ k = \text{bin-nth } c \ (n + k)) \ \&$$

$$(\text{ALL } k. \text{bin-nth } b \ k = (k < n \ \& \ \text{bin-nth } c \ k))$$

⟨proof⟩

lemma *bin-cat-assoc*:

$$\text{bin-cat } (\text{bin-cat } x \ m \ y) \ n \ z = \text{bin-cat } x \ (m + n) \ (\text{bin-cat } y \ n \ z)$$

⟨proof⟩

lemma *bin-cat-assoc-sym*:

$$\text{bin-cat } x \ m \ (\text{bin-cat } y \ n \ z) = \text{bin-cat } (\text{bin-cat } x \ (m - n) \ y) \ (\text{min } m \ n) \ z$$

⟨proof⟩

lemma *bin-cat-zero* [simp]: $\text{bin-cat } 0 \ n \ w = \text{bintrunc } n \ w$

⟨proof⟩

lemma *bintr-cat1*:

$$\text{bintrunc } (k + n) \ (\text{bin-cat } a \ n \ b) = \text{bin-cat } (\text{bintrunc } k \ a) \ n \ b$$

⟨proof⟩

lemma *bintr-cat*: $\text{bintrunc } m \ (\text{bin-cat } a \ n \ b) =$

$$\text{bin-cat } (\text{bintrunc } (m - n) \ a) \ n \ (\text{bintrunc } (\text{min } m \ n) \ b)$$

⟨proof⟩

lemma *bintr-cat-same* [simp]:

$$\text{bintrunc } n \ (\text{bin-cat } a \ n \ b) = \text{bintrunc } n \ b$$

⟨proof⟩

lemma *cat-bintr* [simp]:

$$\text{bin-cat } a \ n \ (\text{bintrunc } n \ b) = \text{bin-cat } a \ n \ b$$

$\langle \text{proof} \rangle$

lemma *split-bintrunc*:

$\text{bin-split } n \ c = (a, b) \implies b = \text{bintrunc } n \ c$

$\langle \text{proof} \rangle$

lemma *bin-cat-split*:

$\text{bin-split } n \ w = (u, v) \implies w = \text{bin-cat } u \ n \ v$

$\langle \text{proof} \rangle$

lemma *bin-split-cat*:

$\text{bin-split } n \ (\text{bin-cat } v \ n \ w) = (v, \text{bintrunc } n \ w)$

$\langle \text{proof} \rangle$

lemma *bin-split-zero* [simp]: $\text{bin-split } n \ 0 = (0, 0)$

$\langle \text{proof} \rangle$

lemma *bin-split-minus1* [simp]:

$\text{bin-split } n \ (-1) = (-1, \text{bintrunc } n \ (-1))$

$\langle \text{proof} \rangle$

lemma *bin-split-trunc*:

$\text{bin-split } (\text{min } m \ n) \ c = (a, b) \implies$

$\text{bin-split } n \ (\text{bintrunc } m \ c) = (\text{bintrunc } (m - n) \ a, b)$

$\langle \text{proof} \rangle$

lemma *bin-split-trunc1*:

$\text{bin-split } n \ c = (a, b) \implies$

$\text{bin-split } n \ (\text{bintrunc } m \ c) = (\text{bintrunc } (m - n) \ a, \text{bintrunc } m \ b)$

$\langle \text{proof} \rangle$

lemma *bin-cat-num*:

$\text{bin-cat } a \ n \ b = a * 2^n + \text{bintrunc } n \ b$

$\langle \text{proof} \rangle$

lemma *bin-split-num*:

$\text{bin-split } n \ b = (b \text{ div } 2^n, b \text{ mod } 2^n)$

$\langle \text{proof} \rangle$

5.4 Miscellaneous lemmas

lemma *nth-2p-bin*:

$\text{bin-nth } (2^n) \ m = (m = n)$

$\langle \text{proof} \rangle$

lemma *ex-eq-or*:

$(\text{EX } m. n = \text{Suc } m \ \& \ (m = k \mid P \ m)) = (n = \text{Suc } k \mid (\text{EX } m. n = \text{Suc } m \ \& \ P$

$m))$
 $\langle proof \rangle$

lemma *power-BIT*: $2^{\wedge} (Suc\ n) - 1 = (2^{\wedge} n - 1)\ BIT\ True$
 $\langle proof \rangle$

lemma *mod-BIT*:
 $bin\ BIT\ bit\ mod\ 2^{\wedge} Suc\ n = (bin\ mod\ 2^{\wedge} n)\ BIT\ bit$
 $\langle proof \rangle$

lemma *AND-mod*:
fixes $x :: int$
shows $x\ AND\ 2^{\wedge} n - 1 = x\ mod\ 2^{\wedge} n$
 $\langle proof \rangle$

end

6 Bool lists and integers

theory *Bool-List-Representation*
imports *Bits-Int*
begin

definition *map2* :: $('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow 'a\ list \Rightarrow 'b\ list \Rightarrow 'c\ list$
where $map2\ f\ as\ bs = map\ (case\ prod\ f)\ (zip\ as\ bs)$

lemma *map2-Nil* [*simp*, *code*]: $map2\ f\ []\ ys = []$
 $\langle proof \rangle$

lemma *map2-Nil2* [*simp*, *code*]: $map2\ f\ xs\ [] = []$
 $\langle proof \rangle$

lemma *map2-Cons* [*simp*, *code*]: $map2\ f\ (x\ \#\ xs)\ (y\ \#\ ys) = f\ x\ y\ \#\ map2\ f\ xs\ ys$
 $\langle proof \rangle$

6.1 Operations on lists of booleans

primrec *bl-to-bin-aux* :: $bool\ list \Rightarrow int \Rightarrow int$
where
 Nil : $bl\text{-to-bin-aux}\ []\ w = w$
 $| Cons$: $bl\text{-to-bin-aux}\ (b\ \#\ bs)\ w = bl\text{-to-bin-aux}\ bs\ (w\ BIT\ b)$

definition *bl-to-bin* :: $bool\ list \Rightarrow int$
where $bl\text{-to-bin}\ bs = bl\text{-to-bin-aux}\ bs\ 0$

primrec *bin-to-bl-aux* :: $nat \Rightarrow int \Rightarrow bool\ list \Rightarrow bool\ list$
where
 Z : $bin\text{-to-bl-aux}\ 0\ w\ bl = bl$

| *Suc*: *bin-to-bl-aux* (*Suc* *n*) *w* *bl* = *bin-to-bl-aux* *n* (*bin-rest* *w*) ((*bin-last* *w*) # *bl*)

definition *bin-to-bl* :: *nat* $\Rightarrow$ *int* $\Rightarrow$ *bool list*
where *bin-to-bl* *n* *w* = *bin-to-bl-aux* *n* *w* []

primrec *bl-of-nth* :: *nat* $\Rightarrow$ (*nat* $\Rightarrow$ *bool*) $\Rightarrow$ *bool list*
where
Suc: *bl-of-nth* (*Suc* *n*) *f* = *f* *n* # *bl-of-nth* *n* *f*
| *Z*: *bl-of-nth* 0 *f* = []

primrec *takefill* :: '*a* $\Rightarrow$ *nat* $\Rightarrow$ '*a* *list* $\Rightarrow$ '*a* *list*
where
Z: *takefill* *fill* 0 *xs* = []
| *Suc*: *takefill* *fill* (*Suc* *n*) *xs* =
(*case* *xs* *of*
[] $\Rightarrow$ *fill* # *takefill* *fill* *n* *xs*
| *y* # *ys* $\Rightarrow$ *y* # *takefill* *fill* *n* *ys*)

6.2 Arithmetic in terms of bool lists

Arithmetic operations in terms of the reversed bool list, assuming input list(s) the same length, and don't extend them.

primrec *rbl-succ* :: *bool list* $\Rightarrow$ *bool list*
where
Nil: *rbl-succ* *Nil* = *Nil*
| *Cons*: *rbl-succ* (*x* # *xs*) = (*if* *x* *then* *False* # *rbl-succ* *xs* *else* *True* # *xs*)

primrec *rbl-pred* :: *bool list* $\Rightarrow$ *bool list*
where
Nil: *rbl-pred* *Nil* = *Nil*
| *Cons*: *rbl-pred* (*x* # *xs*) = (*if* *x* *then* *False* # *xs* *else* *True* # *rbl-pred* *xs*)

primrec *rbl-add* :: *bool list* $\Rightarrow$ *bool list* $\Rightarrow$ *bool list*
where — result is length of first arg, second arg may be longer
Nil: *rbl-add* *Nil* *x* = *Nil*
| *Cons*: *rbl-add* (*y* # *ys*) *x* =
(*let* *ws* = *rbl-add* *ys* (*tl* *x*)
in (*y* $\neq$ *hd* *x*) # (*if* *hd* *x* $\wedge$ *y* *then* *rbl-succ* *ws* *else* *ws*))

primrec *rbl-mult* :: *bool list* $\Rightarrow$ *bool list* $\Rightarrow$ *bool list*
where — result is length of first arg, second arg may be longer
Nil: *rbl-mult* *Nil* *x* = *Nil*
| *Cons*: *rbl-mult* (*y* # *ys*) *x* =
(*let* *ws* = *False* # *rbl-mult* *ys* *x*
in *if* *y* *then* *rbl-add* *ws* *x* *else* *ws*)

lemma *butlast-power*: (*butlast* $\hat{\hat{}}$ *n*) *bl* = *take* (*length* *bl* - *n*) *bl*
<*proof*>

lemma *bin-to-bl-aux-zero-minus-simp* [simp]:

$$0 < n \implies \text{bin-to-bl-aux } n \ 0 \ bl = \text{bin-to-bl-aux } (n - 1) \ 0 \ (\text{False} \ \# \ bl)$$

<proof>

lemma *bin-to-bl-aux-minus1-minus-simp* [simp]:

$$0 < n \implies \text{bin-to-bl-aux } n \ (-1) \ bl = \text{bin-to-bl-aux } (n - 1) \ (-1) \ (\text{True} \ \# \ bl)$$

<proof>

lemma *bin-to-bl-aux-one-minus-simp* [simp]:

$$0 < n \implies \text{bin-to-bl-aux } n \ 1 \ bl = \text{bin-to-bl-aux } (n - 1) \ 0 \ (\text{True} \ \# \ bl)$$

<proof>

lemma *bin-to-bl-aux-Bit-minus-simp* [simp]:

$$0 < n \implies \text{bin-to-bl-aux } n \ (w \ \text{BIT} \ b) \ bl = \text{bin-to-bl-aux } (n - 1) \ w \ (b \ \# \ bl)$$

<proof>

lemma *bin-to-bl-aux-Bit0-minus-simp* [simp]:

$$0 < n \implies \text{bin-to-bl-aux } n \ (\text{numeral } (\text{Num.Bit0 } w)) \ bl = \text{bin-to-bl-aux } (n - 1) \ (\text{numeral } w) \ (\text{False} \ \# \ bl)$$

<proof>

lemma *bin-to-bl-aux-Bit1-minus-simp* [simp]:

$$0 < n \implies \text{bin-to-bl-aux } n \ (\text{numeral } (\text{Num.Bit1 } w)) \ bl = \text{bin-to-bl-aux } (n - 1) \ (\text{numeral } w) \ (\text{True} \ \# \ bl)$$

<proof>

Link between bin and bool list.

lemma *bl-to-bin-aux-append*: $\text{bl-to-bin-aux } (bs \ @ \ cs) \ w = \text{bl-to-bin-aux } cs \ (\text{bl-to-bin-aux } bs \ w)$
<proof>

lemma *bin-to-bl-aux-append*: $\text{bin-to-bl-aux } n \ w \ bs \ @ \ cs = \text{bin-to-bl-aux } n \ w \ (bs \ @ \ cs)$
<proof>

lemma *bl-to-bin-append*: $\text{bl-to-bin } (bs \ @ \ cs) = \text{bl-to-bin-aux } cs \ (\text{bl-to-bin } bs)$
<proof>

lemma *bin-to-bl-aux-alt*: $\text{bin-to-bl-aux } n \ w \ bs = \text{bin-to-bl } n \ w \ @ \ bs$
<proof>

lemma *bin-to-bl-0* [simp]: $\text{bin-to-bl } 0 \ bs = []$
<proof>

lemma *size-bin-to-bl-aux*: $\text{size } (\text{bin-to-bl-aux } n \ w \ bs) = n + \text{length } bs$
<proof>

lemma *size-bin-to-bl [simp]: size (bin-to-bl n w) = n*
 ⟨proof⟩

lemma *bin-bl-bin': bl-to-bin (bin-to-bl-aux n w bs) = bl-to-bin-aux bs (bintrunc n w)*
 ⟨proof⟩

lemma *bin-bl-bin [simp]: bl-to-bin (bin-to-bl n w) = bintrunc n w*
 ⟨proof⟩

lemma *bl-bin-bl': bin-to-bl (n + length bs) (bl-to-bin-aux bs w) = bin-to-bl-aux n w bs*
 ⟨proof⟩

lemma *bl-bin-bl [simp]: bin-to-bl (length bs) (bl-to-bin bs) = bs*
 ⟨proof⟩

lemma *bl-to-bin-inj: bl-to-bin bs = bl-to-bin cs $\implies$ length bs = length cs $\implies$ bs = cs*
 ⟨proof⟩

lemma *bl-to-bin-False [simp]: bl-to-bin (False # bl) = bl-to-bin bl*
 ⟨proof⟩

lemma *bl-to-bin-Nil [simp]: bl-to-bin [] = 0*
 ⟨proof⟩

lemma *bin-to-bl-zero-aux: bin-to-bl-aux n 0 bl = replicate n False @ bl*
 ⟨proof⟩

lemma *bin-to-bl-zero: bin-to-bl n 0 = replicate n False*
 ⟨proof⟩

lemma *bin-to-bl-minus1-aux: bin-to-bl-aux n (- 1) bl = replicate n True @ bl*
 ⟨proof⟩

lemma *bin-to-bl-minus1: bin-to-bl n (- 1) = replicate n True*
 ⟨proof⟩

lemma *bl-to-bin-rep-F: bl-to-bin (replicate n False @ bl) = bl-to-bin bl*
 ⟨proof⟩

lemma *bin-to-bl-trunc [simp]: n $\leq$ m $\implies$ bin-to-bl n (bintrunc m w) = bin-to-bl n w*
 ⟨proof⟩

lemma *bin-to-bl-aux-bintr:*
bin-to-bl-aux n (bintrunc m bin) bl =

$\text{replicate } (n - m) \text{ False } @ \text{ bin-to-bl-aux } (\min n m) \text{ bin bl}$
 $\langle \text{proof} \rangle$

lemma *bin-to-bl-bintr*:

$\text{bin-to-bl } n \text{ (bintrunc } m \text{ bin)} = \text{replicate } (n - m) \text{ False } @ \text{ bin-to-bl } (\min n m) \text{ bin}$
 $\langle \text{proof} \rangle$

lemma *bl-to-bin-rep-False*: $\text{bl-to-bin } (\text{replicate } n \text{ False}) = 0$

$\langle \text{proof} \rangle$

lemma *len-bin-to-bl-aux*: $\text{length } (\text{bin-to-bl-aux } n \text{ w bs}) = n + \text{length bs}$

$\langle \text{proof} \rangle$

lemma *len-bin-to-bl*: $\text{length } (\text{bin-to-bl } n \text{ w}) = n$

$\langle \text{proof} \rangle$

lemma *sign-bl-bin'*: $\text{bin-sign } (\text{bl-to-bin-aux bs w}) = \text{bin-sign w}$

$\langle \text{proof} \rangle$

lemma *sign-bl-bin*: $\text{bin-sign } (\text{bl-to-bin bs}) = 0$

$\langle \text{proof} \rangle$

lemma *bl-sbin-sign-aux*: $\text{hd } (\text{bin-to-bl-aux } (\text{Suc } n) \text{ w bs}) = (\text{bin-sign } (\text{sbintrunc } n \text{ w}) = -1)$

$\langle \text{proof} \rangle$

lemma *bl-sbin-sign*: $\text{hd } (\text{bin-to-bl } (\text{Suc } n) \text{ w}) = (\text{bin-sign } (\text{sbintrunc } n \text{ w}) = -1)$

$\langle \text{proof} \rangle$

lemma *bin-nth-of-bl-aux*:

$\text{bin-nth } (\text{bl-to-bin-aux bl w}) n =$
 $(n < \text{size bl} \wedge \text{rev bl } ! n \mid n \geq \text{length bl} \wedge \text{bin-nth w } (n - \text{size bl}))$

$\langle \text{proof} \rangle$

lemma *bin-nth-of-bl*: $\text{bin-nth } (\text{bl-to-bin bl}) n = (n < \text{length bl} \wedge \text{rev bl } ! n)$

$\langle \text{proof} \rangle$

lemma *bin-nth-bl*: $n < m \implies \text{bin-nth w } n = \text{nth } (\text{rev } (\text{bin-to-bl m w})) n$

$\langle \text{proof} \rangle$

lemma *nth-rev*: $n < \text{length xs} \implies \text{rev xs } ! n = \text{xs } ! (\text{length xs} - 1 - n)$

$\langle \text{proof} \rangle$

lemma *nth-rev-alt*: $n < \text{length ys} \implies \text{ys } ! n = \text{rev ys } ! (\text{length ys} - \text{Suc } n)$

$\langle \text{proof} \rangle$

lemma *nth-bin-to-bl-aux*:

$n < m + \text{length bl} \implies (\text{bin-to-bl-aux m w bl}) ! n =$

(if $n < m$ then $\text{bin-nth } w \ (m - 1 - n)$ else $\text{bl } ! \ (n - m)$)
 ⟨proof⟩

lemma *nth-bin-to-bl*: $n < m \implies (\text{bin-to-bl } m \ w) \ ! \ n = \text{bin-nth } w \ (m - \text{Suc } n)$
 ⟨proof⟩

lemma *bl-to-bin-lt2p-aux*: $\text{bl-to-bin-aux } bs \ w < (w + 1) * (2 \wedge \text{length } bs)$
 ⟨proof⟩

lemma *bl-to-bin-lt2p-drop*: $\text{bl-to-bin } bs < 2 \wedge \text{length } (\text{dropWhile Not } bs)$
 ⟨proof⟩

lemma *bl-to-bin-lt2p*: $\text{bl-to-bin } bs < 2 \wedge \text{length } bs$
 ⟨proof⟩

lemma *bl-to-bin-ge2p-aux*: $\text{bl-to-bin-aux } bs \ w \geq w * (2 \wedge \text{length } bs)$
 ⟨proof⟩

lemma *bl-to-bin-ge0*: $\text{bl-to-bin } bs \geq 0$
 ⟨proof⟩

lemma *butlast-rest-bin*: $\text{butlast } (\text{bin-to-bl } n \ w) = \text{bin-to-bl } (n - 1) \ (\text{bin-rest } w)$
 ⟨proof⟩

lemma *butlast-bin-rest*: $\text{butlast } bl = \text{bin-to-bl } (\text{length } bl - \text{Suc } 0) \ (\text{bin-rest } (\text{bl-to-bin } bl))$
 ⟨proof⟩

lemma *butlast-rest-bl2bin-aux*:
 $bl \neq [] \implies \text{bl-to-bin-aux } (\text{butlast } bl) \ w = \text{bin-rest } (\text{bl-to-bin-aux } bl \ w)$
 ⟨proof⟩

lemma *butlast-rest-bl2bin*: $\text{bl-to-bin } (\text{butlast } bl) = \text{bin-rest } (\text{bl-to-bin } bl)$
 ⟨proof⟩

lemma *trunc-bl2bin-aux*:
 $\text{bintrunc } m \ (\text{bl-to-bin-aux } bl \ w) =$
 $\text{bl-to-bin-aux } (\text{drop } (\text{length } bl - m) \ bl) \ (\text{bintrunc } (m - \text{length } bl) \ w)$
 ⟨proof⟩

lemma *trunc-bl2bin*: $\text{bintrunc } m \ (\text{bl-to-bin } bl) = \text{bl-to-bin } (\text{drop } (\text{length } bl - m) \ bl)$
 ⟨proof⟩

lemma *trunc-bl2bin-len [simp]*: $\text{bintrunc } (\text{length } bl) \ (\text{bl-to-bin } bl) = \text{bl-to-bin } bl$
 ⟨proof⟩

lemma *bl2bin-drop*: $\text{bl-to-bin } (\text{drop } k \ bl) = \text{bintrunc } (\text{length } bl - k) \ (\text{bl-to-bin } bl)$
 ⟨proof⟩

lemma *nth-rest-power-bin*: $\text{bin-nth } ((\text{bin-rest } \wedge^k) w) n = \text{bin-nth } w (n + k)$
 ⟨proof⟩

lemma *take-rest-power-bin*: $m \leq n \implies \text{take } m (\text{bin-to-bl } n w) = \text{bin-to-bl } m ((\text{bin-rest } \wedge^{n-m}) w)$
 ⟨proof⟩

lemma *hd-butlast*: $\text{size } xs > 1 \implies \text{hd } (\text{butlast } xs) = \text{hd } xs$
 ⟨proof⟩

lemma *last-bin-last'*: $\text{size } xs > 0 \implies \text{last } xs \longleftrightarrow \text{bin-last } (\text{bl-to-bin-aux } xs w)$
 ⟨proof⟩

lemma *last-bin-last*: $\text{size } xs > 0 \implies \text{last } xs \longleftrightarrow \text{bin-last } (\text{bl-to-bin } xs)$
 ⟨proof⟩

lemma *bin-last-last*: $\text{bin-last } w \longleftrightarrow \text{last } (\text{bin-to-bl } (\text{Suc } n) w)$
 ⟨proof⟩

6.3 Links between bit-wise operations and operations on bool lists

lemma *bl-xor-aux-bin*:
 $\text{map2 } (\lambda x y. x \neq y) (\text{bin-to-bl-aux } n v bs) (\text{bin-to-bl-aux } n w cs) =$
 $\text{bin-to-bl-aux } n (v \text{ XOR } w) (\text{map2 } (\lambda x y. x \neq y) bs cs)$
 ⟨proof⟩

lemma *bl-or-aux-bin*:
 $\text{map2 } (op \vee) (\text{bin-to-bl-aux } n v bs) (\text{bin-to-bl-aux } n w cs) =$
 $\text{bin-to-bl-aux } n (v \text{ OR } w) (\text{map2 } (op \vee) bs cs)$
 ⟨proof⟩

lemma *bl-and-aux-bin*:
 $\text{map2 } (op \wedge) (\text{bin-to-bl-aux } n v bs) (\text{bin-to-bl-aux } n w cs) =$
 $\text{bin-to-bl-aux } n (v \text{ AND } w) (\text{map2 } (op \wedge) bs cs)$
 ⟨proof⟩

lemma *bl-not-aux-bin*: $\text{map Not } (\text{bin-to-bl-aux } n w cs) = \text{bin-to-bl-aux } n (\text{NOT } w)$
 (map Not cs)
 ⟨proof⟩

lemma *bl-not-bin*: $\text{map Not } (\text{bin-to-bl } n w) = \text{bin-to-bl } n (\text{NOT } w)$
 ⟨proof⟩

lemma *bl-and-bin*: $\text{map2 } (op \wedge) (\text{bin-to-bl } n v) (\text{bin-to-bl } n w) = \text{bin-to-bl } n (v \text{ AND } w)$
 ⟨proof⟩

lemma *bl-or-bin*: $\text{map2 } (op \vee) (bin\text{-}to\text{-}bl\ n\ v) (bin\text{-}to\text{-}bl\ n\ w) = bin\text{-}to\text{-}bl\ n\ (v\ OR\ w)$

<proof>

lemma *bl-xor-bin*: $\text{map2 } (\lambda x\ y. x \neq y) (bin\text{-}to\text{-}bl\ n\ v) (bin\text{-}to\text{-}bl\ n\ w) = bin\text{-}to\text{-}bl\ n\ (v\ XOR\ w)$

<proof>

lemma *drop-bin2bl-aux*:

$\text{drop } m\ (bin\text{-}to\text{-}bl\text{-}aux\ n\ bin\ bs) =$
 $bin\text{-}to\text{-}bl\text{-}aux\ (n - m)\ bin\ (\text{drop } (m - n)\ bs)$

<proof>

lemma *drop-bin2bl*: $\text{drop } m\ (bin\text{-}to\text{-}bl\ n\ bin) = bin\text{-}to\text{-}bl\ (n - m)\ bin$

<proof>

lemma *take-bin2bl-lem1*: $\text{take } m\ (bin\text{-}to\text{-}bl\text{-}aux\ m\ w\ bs) = bin\text{-}to\text{-}bl\ m\ w$

<proof>

lemma *take-bin2bl-lem*: $\text{take } m\ (bin\text{-}to\text{-}bl\text{-}aux\ (m + n)\ w\ bs) = \text{take } m\ (bin\text{-}to\text{-}bl\ (m + n)\ w)$

<proof>

lemma *bin-split-take*: $bin\text{-}split\ n\ c = (a, b) \implies bin\text{-}to\text{-}bl\ m\ a = \text{take } m\ (bin\text{-}to\text{-}bl\ (m + n)\ c)$

<proof>

lemma *bin-split-take1*:

$k = m + n \implies bin\text{-}split\ n\ c = (a, b) \implies bin\text{-}to\text{-}bl\ m\ a = \text{take } m\ (bin\text{-}to\text{-}bl\ k\ c)$

<proof>

lemma *nth-takefill*: $m < n \implies \text{takefill } fill\ n\ l\ !\ m = (\text{if } m < \text{length } l \text{ then } l\ !\ m \text{ else } fill)$

<proof>

lemma *takefill-alt*: $\text{takefill } fill\ n\ l = \text{take } n\ l\ @\ replicate\ (n - \text{length } l)\ fill$

<proof>

lemma *takefill-replicate* [simp]: $\text{takefill } fill\ n\ (replicate\ m\ fill) = replicate\ n\ fill$

<proof>

lemma *takefill-le'*: $n = m + k \implies \text{takefill } x\ m\ (\text{takefill } x\ n\ l) = \text{takefill } x\ m\ l$

<proof>

lemma *length-takefill* [simp]: $\text{length } (\text{takefill } fill\ n\ l) = n$

<proof>

lemma *take-takefill'*: $\bigwedge w\ n. n = k + m \implies \text{take } k\ (\text{takefill } fill\ n\ w) = \text{takefill } fill\ k\ w$

$\langle \text{proof} \rangle$

lemma *drop-takefill*: $\bigwedge w. \text{drop } k (\text{takefill fill } (m + k) w) = \text{takefill fill } m (\text{drop } k w)$
 $\langle \text{proof} \rangle$

lemma *takefill-le* [simp]: $m \leq n \implies \text{takefill } x m (\text{takefill } x n l) = \text{takefill } x m l$
 $\langle \text{proof} \rangle$

lemma *take-takefill* [simp]: $m \leq n \implies \text{take } m (\text{takefill fill } n w) = \text{takefill fill } m w$
 $\langle \text{proof} \rangle$

lemma *takefill-append*: $\text{takefill fill } (m + \text{length } xs) (xs @ w) = xs @ (\text{takefill fill } m w)$
 $\langle \text{proof} \rangle$

lemma *takefill-same'*: $l = \text{length } xs \implies \text{takefill fill } l xs = xs$
 $\langle \text{proof} \rangle$

lemmas *takefill-same* [simp] = *takefill-same'* [OF refl]

lemma *takefill-bintrunc*: $\text{takefill False } n bl = \text{rev } (\text{bin-to-bl } n (\text{bl-to-bin } (\text{rev } bl)))$
 $\langle \text{proof} \rangle$

lemma *bl-bin-bl-rtf*: $\text{bin-to-bl } n (\text{bl-to-bin } bl) = \text{rev } (\text{takefill False } n (\text{rev } bl))$
 $\langle \text{proof} \rangle$

lemma *bl-bin-bl-rep-drop*:
 $\text{bin-to-bl } n (\text{bl-to-bin } bl) =$
 $\text{replicate } (n - \text{length } bl) \text{ False } @ \text{drop } (\text{length } bl - n) bl$
 $\langle \text{proof} \rangle$

lemma *tf-rev*:
 $n + k = m + \text{length } bl \implies \text{takefill } x m (\text{rev } (\text{takefill } y n bl)) =$
 $\text{rev } (\text{takefill } y m (\text{rev } (\text{takefill } x k (\text{rev } bl))))$
 $\langle \text{proof} \rangle$

lemma *takefill-minus*: $0 < n \implies \text{takefill fill } (\text{Suc } (n - 1)) w = \text{takefill fill } n w$
 $\langle \text{proof} \rangle$

lemmas *takefill-Suc-cases* =
 $\text{list.cases } [\text{THEN } \text{takefill.Suc } [\text{THEN } \text{trans}]]$

lemmas *takefill-Suc-Nil* = *takefill-Suc-cases* (1)

lemmas *takefill-Suc-Cons* = *takefill-Suc-cases* (2)

lemmas *takefill-minus-simps* = *takefill-Suc-cases* [THEN [2]
takefill-minus [symmetric, THEN trans]]

lemma *takefill-numeral-Nil* [simp]:

$$\text{takefill fill (numeral } k) [] = \text{fill } \# \text{ takefill fill (pred-numeral } k) []$$

⟨proof⟩

lemma *takefill-numeral-Cons* [simp]:

$$\text{takefill fill (numeral } k) (x \# xs) = x \# \text{ takefill fill (pred-numeral } k) xs$$

⟨proof⟩

6.4 Links with function *bl-to-bin*

lemma *bl-to-bin-aux-cat*:

$$\bigwedge nv v. \text{bl-to-bin-aux } bs (\text{bin-cat } w \text{ } nv \text{ } v) = \text{bin-cat } w (nv + \text{length } bs) (\text{bl-to-bin-aux } bs \text{ } v)$$

⟨proof⟩

lemma *bin-to-bl-aux-cat*:

$$\bigwedge w bs. \text{bin-to-bl-aux } (nv + nw) (\text{bin-cat } v \text{ } nw \text{ } w) \text{ } bs = \text{bin-to-bl-aux } nv \text{ } v (\text{bin-to-bl-aux } nw \text{ } w \text{ } bs)$$

⟨proof⟩

lemma *bl-to-bin-aux-alt*: $\text{bl-to-bin-aux } bs \text{ } w = \text{bin-cat } w (\text{length } bs) (\text{bl-to-bin } bs)$

⟨proof⟩

lemma *bin-to-bl-cat*:

$$\text{bin-to-bl } (nv + nw) (\text{bin-cat } v \text{ } nw \text{ } w) = \text{bin-to-bl-aux } nv \text{ } v (\text{bin-to-bl } nw \text{ } w)$$

⟨proof⟩

lemmas *bl-to-bin-aux-app-cat* =

$$\text{trans } [OF \text{ bl-to-bin-aux-append bl-to-bin-aux-alt}]$$

lemmas *bin-to-bl-aux-cat-app* =

$$\text{trans } [OF \text{ bin-to-bl-aux-cat bin-to-bl-aux-alt}]$$

lemma *bl-to-bin-app-cat*:

$$\text{bl-to-bin } (bsa @ bs) = \text{bin-cat } (\text{bl-to-bin } bsa) (\text{length } bs) (\text{bl-to-bin } bs)$$

⟨proof⟩

lemma *bin-to-bl-cat-app*:

$$\text{bin-to-bl } (n + nw) (\text{bin-cat } w \text{ } nw \text{ } wa) = \text{bin-to-bl } n \text{ } w @ \text{bin-to-bl } nw \text{ } wa$$

⟨proof⟩

bl-to-bin-app-cat-alt and *bl-to-bin-app-cat* are easily interderivable.

lemma *bl-to-bin-app-cat-alt*: $\text{bin-cat } (\text{bl-to-bin } cs) \text{ } n \text{ } w = \text{bl-to-bin } (cs @ \text{bin-to-bl } n \text{ } w)$

⟨proof⟩

lemma *mask-lem*: $(\text{bl-to-bin } (\text{True } \# \text{ replicate } n \text{ False})) = \text{bl-to-bin } (\text{replicate } n \text{ True}) + 1$

$\langle proof \rangle$

6.5 Function *bl-of-nth*

lemma *length-bl-of-nth [simp]*: $length\ (bl\text{-}of\text{-}nth\ n\ f) = n$
 $\langle proof \rangle$

lemma *nth-bl-of-nth [simp]*: $m < n \implies rev\ (bl\text{-}of\text{-}nth\ n\ f)\ !\ m = f\ m$
 $\langle proof \rangle$

lemma *bl-of-nth-inj*: $(\bigwedge k. k < n \implies f\ k = g\ k) \implies bl\text{-}of\text{-}nth\ n\ f = bl\text{-}of\text{-}nth\ n\ g$
 $\langle proof \rangle$

lemma *bl-of-nth-nth-le*: $n \leq length\ xs \implies bl\text{-}of\text{-}nth\ n\ (nth\ (rev\ xs)) = drop\ (length\ xs - n)\ xs$
 $\langle proof \rangle$

lemma *bl-of-nth-nth [simp]*: $bl\text{-}of\text{-}nth\ (length\ xs)\ (op\ !\ (rev\ xs)) = xs$
 $\langle proof \rangle$

lemma *size-rbl-pred*: $length\ (rbl\text{-}pred\ bl) = length\ bl$
 $\langle proof \rangle$

lemma *size-rbl-succ*: $length\ (rbl\text{-}succ\ bl) = length\ bl$
 $\langle proof \rangle$

lemma *size-rbl-add*: $length\ (rbl\text{-}add\ bl\ cl) = length\ bl$
 $\langle proof \rangle$

lemma *size-rbl-mult*: $length\ (rbl\text{-}mult\ bl\ cl) = length\ bl$
 $\langle proof \rangle$

lemmas *rbl-sizes [simp]* =
 $size\text{-}rbl\text{-}pred\ size\text{-}rbl\text{-}succ\ size\text{-}rbl\text{-}add\ size\text{-}rbl\text{-}mult$

lemmas *rbl-Nils* =
 $rbl\text{-}pred.Nil\ rbl\text{-}succ.Nil\ rbl\text{-}add.Nil\ rbl\text{-}mult.Nil$

lemma *rbl-pred*: $rbl\text{-}pred\ (rev\ (bin\text{-}to\text{-}bl\ n\ bin)) = rev\ (bin\text{-}to\text{-}bl\ n\ (bin - 1))$
 $\langle proof \rangle$

lemma *rbl-succ*: $rbl\text{-}succ\ (rev\ (bin\text{-}to\text{-}bl\ n\ bin)) = rev\ (bin\text{-}to\text{-}bl\ n\ (bin + 1))$
 $\langle proof \rangle$

lemma *rbl-add*:
 $\bigwedge bina\ binb. rbl\text{-}add\ (rev\ (bin\text{-}to\text{-}bl\ n\ bina))\ (rev\ (bin\text{-}to\text{-}bl\ n\ binb)) =$
 $rev\ (bin\text{-}to\text{-}bl\ n\ (bina + binb))$
 $\langle proof \rangle$

lemma *rbl-add-app2*: $\text{length } blb \geq \text{length } bla \implies \text{rbl-add } bla \ (blb \ @ \ blc) = \text{rbl-add } bla \ blb$

<proof>

lemma *rbl-add-take2*:

$\text{length } blb \geq \text{length } bla \implies \text{rbl-add } bla \ (\text{take } (\text{length } bla) \ blb) = \text{rbl-add } bla \ blb$

<proof>

lemma *rbl-add-long*:

$m \geq n \implies \text{rbl-add } (\text{rev } (\text{bin-to-bl } n \ bina)) \ (\text{rev } (\text{bin-to-bl } m \ binb)) = \text{rev } (\text{bin-to-bl } n \ (bina + binb))$

<proof>

lemma *rbl-mult-app2*: $\text{length } blb \geq \text{length } bla \implies \text{rbl-mult } bla \ (blb \ @ \ blc) = \text{rbl-mult } bla \ blb$

<proof>

lemma *rbl-mult-take2*:

$\text{length } blb \geq \text{length } bla \implies \text{rbl-mult } bla \ (\text{take } (\text{length } bla) \ blb) = \text{rbl-mult } bla \ blb$

<proof>

lemma *rbl-mult-gt1*:

$m \geq \text{length } bl \implies \text{rbl-mult } bl \ (\text{rev } (\text{bin-to-bl } m \ binb)) = \text{rbl-mult } bl \ (\text{rev } (\text{bin-to-bl } (\text{length } bl) \ binb))$

<proof>

lemma *rbl-mult-gt*:

$m > n \implies \text{rbl-mult } (\text{rev } (\text{bin-to-bl } n \ bina)) \ (\text{rev } (\text{bin-to-bl } m \ binb)) = \text{rbl-mult } (\text{rev } (\text{bin-to-bl } n \ bina)) \ (\text{rev } (\text{bin-to-bl } n \ binb))$

<proof>

lemmas *rbl-mult-Suc* = *lessI* [THEN *rbl-mult-gt*]

lemma *rbbl-Cons*: $b \# \text{rev } (\text{bin-to-bl } n \ x) = \text{rev } (\text{bin-to-bl } (\text{Suc } n) \ (x \ \text{BIT } b))$

<proof>

lemma *rbl-mult*:

$\text{rbl-mult } (\text{rev } (\text{bin-to-bl } n \ bina)) \ (\text{rev } (\text{bin-to-bl } n \ binb)) = \text{rev } (\text{bin-to-bl } n \ (bina * binb))$

<proof>

lemma *rbl-add-split*:

$P \ (\text{rbl-add } (y \# \ ys) \ (x \# \ xs)) = (\forall \ ws. \ \text{length } ws = \text{length } ys \longrightarrow ws = \text{rbl-add } ys \ xs \longrightarrow (y \longrightarrow ((x \longrightarrow P \ (\text{False} \# \ \text{rbl-succ } ws)) \wedge (\neg x \longrightarrow P \ (\text{True} \# \ ws)))) \wedge (\neg y \longrightarrow P \ (x \# \ ws)))$

$\langle \text{proof} \rangle$

lemma *rbl-mult-split*:

$P \text{ (rbl-mult } (y \# ys) \text{ xs)} =$
 $(\forall ws. \text{length } ws = \text{Suc } (\text{length } ys) \longrightarrow ws = \text{False} \# \text{rbl-mult } ys \text{ xs} \longrightarrow$
 $(y \longrightarrow P \text{ (rbl-add } ws \text{ xs)}) \wedge (\neg y \longrightarrow P \text{ ws}))$
 $\langle \text{proof} \rangle$

6.6 Repeated splitting or concatenation

lemma *sclm*: $\text{size } (\text{concat } (\text{map } (\text{bin-to-bl } n) \text{ xs})) = \text{length } xs * n$
 $\langle \text{proof} \rangle$

lemma *bin-cat-foldl-lem*:

$\text{foldl } (\lambda u. \text{bin-cat } u \text{ n}) \text{ x xs} =$
 $\text{bin-cat } x \text{ (size xs * n) (foldl } (\lambda u. \text{bin-cat } u \text{ n}) \text{ y xs)}$
 $\langle \text{proof} \rangle$

lemma *bin-rcat-bl*: $\text{bin-rcat } n \text{ wl} = \text{bl-to-bin } (\text{concat } (\text{map } (\text{bin-to-bl } n) \text{ wl}))$
 $\langle \text{proof} \rangle$

lemmas *bin-rsplit-aux-simps* = *bin-rsplit-aux.simps bin-rsplittl-aux.simps*

lemmas *rsplit-aux-simps* = *bin-rsplit-aux-simps*

lemmas *th-if-simp1* = *if-split* [where $P = op = l$, THEN *iffD1*, THEN *conjunct1*, THEN *mp*] for l

lemmas *th-if-simp2* = *if-split* [where $P = op = l$, THEN *iffD1*, THEN *conjunct2*, THEN *mp*] for l

lemmas *rsplit-aux-simp1s* = *rsplit-aux-simps* [THEN *th-if-simp1*]

lemmas *rsplit-aux-simp2ls* = *rsplit-aux-simps* [THEN *th-if-simp2*]

lemmas *bin-rsplit-aux-simp2s* [simp] = *rsplit-aux-simp2ls* [unfolded *Let-def*]

lemmas *rbscl* = *bin-rsplit-aux-simp2s* (2)

lemmas *rsplit-aux-0-simps* [simp] =

rsplit-aux-simp1s [OF *disjI1*] *rsplit-aux-simp1s* [OF *disjI2*]

lemma *bin-rsplit-aux-append*: $\text{bin-rsplit-aux } n \text{ m c } (bs @ cs) = \text{bin-rsplit-aux } n \text{ m c } bs @ cs$
 $\langle \text{proof} \rangle$

lemma *bin-rsplittl-aux-append*: $\text{bin-rsplittl-aux } n \text{ m c } (bs @ cs) = \text{bin-rsplittl-aux } n \text{ m c } bs @ cs$
 $\langle \text{proof} \rangle$

lemmas *rsplit-aux-apps* [where $bs = []$] =
bin-rsplit-aux-append bin-rsplittl-aux-append

lemmas *rsplit-def-auxs* = *bin-rsplit-def bin-rsplitl-def*

lemmas *rsplit-aux-alts* = *rsplit-aux-apps*
 [unfolded append-Nil *rsplit-def-auxs* [symmetric]]

lemma *bin-split-minus*: $0 < n \implies \text{bin-split } (\text{Suc } (n - 1)) \ w = \text{bin-split } n \ w$
 ⟨proof⟩

lemmas *bin-split-minus-simp* =
bin-split.Suc [THEN [2] *bin-split-minus* [symmetric, THEN trans]]

lemma *bin-split-pred-simp* [simp]:
 $(0::\text{nat}) < \text{numeral } \text{bin} \implies$
 $\text{bin-split } (\text{numeral } \text{bin}) \ w =$
 $(\text{let } (w1, w2) = \text{bin-split } (\text{numeral } \text{bin} - 1) \ (\text{bin-rest } w)$
 $\text{in } (w1, w2 \text{ BIT } \text{bin-last } w))$
 ⟨proof⟩

lemma *bin-rsplit-aux-simp-alt*:
 $\text{bin-rsplit-aux } n \ m \ c \ bs =$
 $(\text{if } m = 0 \vee n = 0 \text{ then } bs$
 $\text{else let } (a, b) = \text{bin-split } n \ c \text{ in } \text{bin-rsplit } n \ (m - n, a) \ @ \ b \ \# \ bs)$
 ⟨proof⟩

lemmas *bin-rsplit-simp-alt* =
 trans [OF *bin-rsplit-def bin-rsplit-aux-simp-alt*]

lemmas *bthrs* = *bin-rsplit-simp-alt* [THEN [2] trans]

lemma *bin-rsplit-size-sign'* [rule-format]:
 $n > 0 \implies \text{rev } sw = \text{bin-rsplit } n \ (nw, w) \implies \forall v \in \text{set } sw. \text{bintrunc } n \ v = v$
 ⟨proof⟩

lemmas *bin-rsplit-size-sign* = *bin-rsplit-size-sign'* [OF *asm-rl*
rev-rev-ident [THEN trans] *set-rev* [THEN *equalityD2* [THEN *subsetD*]]]

lemma *bin-nth-rsplit* [rule-format] :
 $n > 0 \implies m < n \implies$
 $\forall w \ k \ nw.$
 $\text{rev } sw = \text{bin-rsplit } n \ (nw, w) \longrightarrow$
 $k < \text{size } sw \longrightarrow \text{bin-nth } (sw ! k) \ m = \text{bin-nth } w \ (k * n + m)$
 ⟨proof⟩

lemma *bin-rsplit-all*: $0 < nw \implies nw \leq n \implies \text{bin-rsplit } n \ (nw, w) = [\text{bintrunc } n \ w]$
 ⟨proof⟩

lemma *bin-rsplit-l* [rule-format]:

$\forall \text{ bin. } \text{bin-rsplitl } n \ (m, \text{bin}) = \text{bin-rsplit } n \ (m, \text{bintrunc } m \ \text{bin})$
 $\langle \text{proof} \rangle$

lemma *bin-rsplit-rcat* [rule-format]:

$n > 0 \longrightarrow \text{bin-rsplit } n \ (n * \text{size } ws, \text{bin-rcat } n \ ws) = \text{map } (\text{bintrunc } n) \ ws$
 $\langle \text{proof} \rangle$

lemma *bin-rsplit-aux-len-le* [rule-format] :

$\forall \text{ ws } m. \ n \neq 0 \longrightarrow \text{ws} = \text{bin-rsplit-aux } n \ nw \ w \ bs \longrightarrow$
 $\text{length } ws \leq m \longleftrightarrow nw + \text{length } bs * n \leq m * n$
 $\langle \text{proof} \rangle$

lemma *bin-rsplit-len-le*: $n \neq 0 \longrightarrow \text{ws} = \text{bin-rsplit } n \ (nw, w) \longrightarrow \text{length } ws \leq m$
 $\longleftrightarrow nw \leq m * n$
 $\langle \text{proof} \rangle$

lemma *bin-rsplit-aux-len*:

$n \neq 0 \implies \text{length } (\text{bin-rsplit-aux } n \ nw \ w \ cs) = (nw + n - 1) \text{ div } n + \text{length } cs$
 $\langle \text{proof} \rangle$

lemma *bin-rsplit-len*: $n \neq 0 \implies \text{length } (\text{bin-rsplit } n \ (nw, w)) = (nw + n - 1) \text{ div } n$
 $\langle \text{proof} \rangle$

lemma *bin-rsplit-aux-len-indep*:

$n \neq 0 \implies \text{length } bs = \text{length } cs \implies$
 $\text{length } (\text{bin-rsplit-aux } n \ nw \ v \ bs) =$
 $\text{length } (\text{bin-rsplit-aux } n \ nw \ w \ cs)$
 $\langle \text{proof} \rangle$

lemma *bin-rsplit-len-indep*:

$n \neq 0 \implies \text{length } (\text{bin-rsplit } n \ (nw, v)) = \text{length } (\text{bin-rsplit } n \ (nw, w))$
 $\langle \text{proof} \rangle$

Even more bit operations

instantiation *int* :: *bitss*
begin

definition [*iff*]: $i \ !! \ n \longleftrightarrow \text{bin-nth } i \ n$

definition *lsb* $i = i \ !! \ 0$ **for** $i :: \text{int}$

definition *set-bit* $i \ n \ b = \text{bin-sc } n \ b \ i$

definition

set-bits $f =$
 $(\text{if } \exists n. \forall n' \geq n. \neg f \ n' \text{ then}$
 $\text{let } n = \text{LEAST } n. \forall n' \geq n. \neg f \ n'$
 $\text{in } \text{bl-to-bin } (\text{rev } (\text{map } f \ [0..<n])))$

```

else if  $\exists n. \forall n' \geq n. f\ n'$  then
  let  $n = \text{LEAST } n. \forall n' \geq n. f\ n'$ 
  in  $\text{sbintrunc } n\ (\text{bl-to-bin } (\text{True} \# \text{rev } (\text{map } f\ [0..<n])))$ 
else 0 :: int

```

definition $\text{shiffl } x\ n = x * 2^{\wedge} n$ **for** $x :: \text{int}$

definition $\text{shiftr } x\ n = x \text{ div } 2^{\wedge} n$ **for** $x :: \text{int}$

definition $\text{msb } x \longleftrightarrow x < 0$ **for** $x :: \text{int}$

instance $\langle \text{proof} \rangle$

end

end

7 Type Definition Theorems

theory *Misc-Typedef*

imports *Main*

begin

8 More lemmas about normal type definitions

lemma

```

tdD1: type-definition Rep Abs A  $\implies \forall x. \text{Rep } x \in A$  and
tdD2: type-definition Rep Abs A  $\implies \forall x. \text{Abs } (\text{Rep } x) = x$  and
tdD3: type-definition Rep Abs A  $\implies \forall y. y \in A \longrightarrow \text{Rep } (\text{Abs } y) = y$ 
 $\langle \text{proof} \rangle$ 

```

lemma *td-nat-int*:

```

type-definition int nat (Collect (op <= 0))
 $\langle \text{proof} \rangle$ 

```

context *type-definition*

begin

declare *Rep* [*iff*] *Rep-inverse* [*simp*] *Rep-inject* [*simp*]

lemma *Abs-eqD*: $\text{Abs } x = \text{Abs } y \implies x \in A \implies y \in A \implies x = y$
 $\langle \text{proof} \rangle$

lemma *Abs-inverse'*:

```

 $r : A \implies \text{Abs } r = a \implies \text{Rep } a = r$ 
 $\langle \text{proof} \rangle$ 

```

lemma *Rep-comp-inverse*:

$Rep \circ f = g \implies Abs \circ g = f$
 $\langle proof \rangle$

lemma *Rep-eqD* [elim!]: $Rep\ x = Rep\ y \implies x = y$
 $\langle proof \rangle$

lemma *Rep-inverse'*: $Rep\ a = r \implies Abs\ r = a$
 $\langle proof \rangle$

lemma *comp-Abs-inverse*:
 $f \circ Abs = g \implies g \circ Rep = f$
 $\langle proof \rangle$

lemma *set-Rep*:
 $A = range\ Rep$
 $\langle proof \rangle$

lemma *set-Rep-Abs*: $A = range\ (Rep \circ Abs)$
 $\langle proof \rangle$

lemma *Abs-inj-on*: $inj\text{-}on\ Abs\ A$
 $\langle proof \rangle$

lemma *image*: $Abs\ ` A = UNIV$
 $\langle proof \rangle$

lemmas *td-thm* = *type-definition-axioms*

lemma *fns1*:
 $Rep \circ fa = fr \circ Rep \mid fa \circ Abs = Abs \circ fr \implies Abs \circ fr \circ Rep = fa$
 $\langle proof \rangle$

lemmas *fns1a* = *disjI1* [THEN *fns1*]

lemmas *fns1b* = *disjI2* [THEN *fns1*]

lemma *fns4*:
 $Rep \circ fa \circ Abs = fr \implies$
 $Rep \circ fa = fr \circ Rep \ \& \ fa \circ Abs = Abs \circ fr$
 $\langle proof \rangle$

end

interpretation *nat-int*: *type-definition int nat Collect* (*op* <= 0)
 $\langle proof \rangle$

declare
nat-int.Rep-cases [cases del]
nat-int.Abs-cases [cases del]
nat-int.Rep-induct [induct del]

nat-int.Abs-induct [*induct del*]

8.1 Extended form of type definition predicate

lemma *td-conds*:

$norm \circ norm = norm \implies (fr \circ norm = norm \circ fr) =$
 $(norm \circ fr \circ norm = fr \circ norm \ \& \ norm \circ fr \circ norm = norm \circ fr)$
 $\langle proof \rangle$

lemma *fn-comm-power*:

$fa \circ tr = tr \circ fr \implies fa \wedge n \circ tr = tr \circ fr \wedge n$
 $\langle proof \rangle$

lemmas *fn-comm-power'* =

ext [*THEN fn-comm-power, THEN fun-cong, unfolded o-def*]

locale *td-ext* = *type-definition* +

fixes *norm*

assumes *eq-norm*: $\bigwedge x. Rep \ (Abs \ x) = norm \ x$

begin

lemma *Abs-norm* [*simp*]:

$Abs \ (norm \ x) = Abs \ x$
 $\langle proof \rangle$

lemma *td-th*:

$g \circ Abs = f \implies f \ (Rep \ x) = g \ x$
 $\langle proof \rangle$

lemma *eq-norm'*: $Rep \circ Abs = norm$

$\langle proof \rangle$

lemma *norm-Rep* [*simp*]: $norm \ (Rep \ x) = Rep \ x$

$\langle proof \rangle$

lemmas *td* = *td-thm*

lemma *set-iff-norm*: $w : A \longleftrightarrow w = norm \ w$

$\langle proof \rangle$

lemma *inverse-norm*:

$(Abs \ n = w) = (Rep \ w = norm \ n)$
 $\langle proof \rangle$

lemma *norm-eq-iff*:

$(norm \ x = norm \ y) = (Abs \ x = Abs \ y)$
 $\langle proof \rangle$

lemma *norm-comps*:

$$Abs \circ norm = Abs$$

$$norm \circ Rep = Rep$$

$$norm \circ norm = norm$$

$\langle proof \rangle$

lemmas *norm-norm* [*simp*] = *norm-comps*

lemma *fns5*:

$$Rep \circ fa \circ Abs = fr ==>$$

$$fr \circ norm = fr \ \& \ norm \circ fr = fr$$

$\langle proof \rangle$

lemma *fns2*:

$$Abs \circ fr \circ Rep = fa ==>$$

$$(norm \circ fr \circ norm = fr \circ norm) = (Rep \circ fa = fr \circ Rep)$$

$\langle proof \rangle$

lemma *fns3*:

$$Abs \circ fr \circ Rep = fa ==>$$

$$(norm \circ fr \circ norm = norm \circ fr) = (fa \circ Abs = Abs \circ fr)$$

$\langle proof \rangle$

lemma *fns*:

$$fr \circ norm = norm \circ fr ==>$$

$$(fa \circ Abs = Abs \circ fr) = (Rep \circ fa = fr \circ Rep)$$

$\langle proof \rangle$

lemma *range-norm*:

$$range \ (Rep \circ Abs) = A$$

$\langle proof \rangle$

end

lemmas *td-ext-def'* =

$$td-ext-def \ [unfolded \ type-definition-def \ td-ext-axioms-def]$$

end

9 Miscellaneous lemmas, of at least doubtful value

theory *Word-Miscellaneous*

imports *HOL-Library.Bit Misc-Numeric*

begin

lemma *power-minus-simp*: $0 < n \implies a \wedge n = a * a \wedge (n - 1)$

$\langle proof \rangle$

lemma *funpow-minus-simp*: $0 < n \implies f \hat{\ } n = f \circ f \hat{\ } (n - 1)$
 ⟨proof⟩

lemma *power-numeral*: $a \hat{\ } \text{numeral } k = a * a \hat{\ } (\text{pred-numeral } k)$
 ⟨proof⟩

lemma *funpow-numeral [simp]*: $f \hat{\ } \text{numeral } k = f \circ f \hat{\ } (\text{pred-numeral } k)$
 ⟨proof⟩

lemma *replicate-numeral [simp]*: $\text{replicate } (\text{numeral } k) \ x = x \# \text{replicate } (\text{pred-numeral } k) \ x$
 ⟨proof⟩

lemma *rco-alt*: $(f \circ g) \hat{\ } n \circ f = f \circ (g \circ f) \hat{\ } n$
 ⟨proof⟩

lemma *list-exhaust-size-gt0*:
 assumes $y: \bigwedge a \text{ list. } y = a \# \text{list} \implies P$
 shows $0 < \text{length } y \implies P$
 ⟨proof⟩

lemma *list-exhaust-size-eq0*:
 assumes $y: y = [] \implies P$
 shows $\text{length } y = 0 \implies P$
 ⟨proof⟩

lemma *size-Cons-lem-eq*: $y = xa \# \text{list} \implies \text{size } y = \text{Suc } k \implies \text{size list} = k$
 ⟨proof⟩

lemmas *ls-splits* = *prod.split prod.split-asm if-split-asm*

lemma *not-B1-is-B0*: $y \neq 1 \implies y = 0$
 for $y :: \text{bit}$
 ⟨proof⟩

lemma *B1-ass-B0*:
 fixes $y :: \text{bit}$
 assumes $y: y = 0 \implies y = 1$
 shows $y = 1$
 ⟨proof⟩

lemmas *n2s-ths* [*THEN eq-reflection*] = *add-2-eq-Suc add-2-eq-Suc'*

lemmas *s2n-ths* = *n2s-ths* [*symmetric*]

lemma *and-len*: $xs = ys \implies xs = ys \wedge \text{length } xs = \text{length } ys$
 ⟨proof⟩

lemma *size-if*: $\text{size } (\text{if } p \text{ then } xs \text{ else } ys) = (\text{if } p \text{ then } \text{size } xs \text{ else } \text{size } ys)$
 ⟨proof⟩

lemma *tl-if*: $tl\ (if\ p\ then\ xs\ else\ ys) = (if\ p\ then\ tl\ xs\ else\ tl\ ys)$
 $\langle proof \rangle$

lemma *hd-if*: $hd\ (if\ p\ then\ xs\ else\ ys) = (if\ p\ then\ hd\ xs\ else\ hd\ ys)$
 $\langle proof \rangle$

lemma *if-Not-x*: $(if\ p\ then\ \neg x\ else\ x) = (p = (\neg x))$
 $\langle proof \rangle$

lemma *if-x-Not*: $(if\ p\ then\ x\ else\ \neg x) = (p = x)$
 $\langle proof \rangle$

lemma *if-same-and*: $(If\ p\ x\ y \wedge If\ p\ u\ v) = (if\ p\ then\ x \wedge u\ else\ y \wedge v)$
 $\langle proof \rangle$

lemma *if-same-eq*: $(If\ p\ x\ y = (If\ p\ u\ v)) = (if\ p\ then\ x = u\ else\ y = v)$
 $\langle proof \rangle$

lemma *if-same-eq-not*: $(If\ p\ x\ y = (\neg If\ p\ u\ v)) = (if\ p\ then\ x = (\neg u)\ else\ y = (\neg v))$
 $\langle proof \rangle$

lemma *if-Cons*: $(if\ p\ then\ x \# xs\ else\ y \# ys) = If\ p\ x\ y \# If\ p\ xs\ ys$
 $\langle proof \rangle$

lemma *if-single*: $(if\ xc\ then\ [xab]\ else\ [an]) = [if\ xc\ then\ xab\ else\ an]$
 $\langle proof \rangle$

lemma *if-bool-simps*:
 $If\ p\ True\ y = (p \vee y) \wedge If\ p\ False\ y = (\neg p \wedge y) \wedge$
 $If\ p\ y\ True = (p \longrightarrow y) \wedge If\ p\ y\ False = (p \wedge y)$
 $\langle proof \rangle$

lemmas *if-simps* =
if-x-Not if-Not-x if-cancel if-True if-False if-bool-simps

lemmas *seqr = eq-reflection* [where $x = size\ w$] for w

lemma *the-elemI*: $y = \{x\} \Longrightarrow the-elem\ y = x$
 $\langle proof \rangle$

lemma *nonemptyE*: $S \neq \{\} \Longrightarrow (\bigwedge x. x \in S \Longrightarrow R) \Longrightarrow R$
 $\langle proof \rangle$

lemma *gt-or-eq-0*: $0 < y \vee 0 = y$
for $y :: nat$
 $\langle proof \rangle$

```

lemmas xtr1 = xtrans(1)
lemmas xtr2 = xtrans(2)
lemmas xtr3 = xtrans(3)
lemmas xtr4 = xtrans(4)
lemmas xtr5 = xtrans(5)
lemmas xtr6 = xtrans(6)
lemmas xtr7 = xtrans(7)
lemmas xtr8 = xtrans(8)

```

```

lemmas nat-simps = diff-add-inverse2 diff-add-inverse
lemmas nat-iffs = le-add1 le-add2

```

```

lemma sum-imp-diff:  $j = k + i \implies j - i = k$ 
  for  $k :: \text{nat}$ 
   $\langle \text{proof} \rangle$ 

```

```

lemmas pos-mod-sign2 = zless2 [THEN pos-mod-sign [where  $b = 2::\text{int}$ ]]
lemmas pos-mod-bound2 = zless2 [THEN pos-mod-bound [where  $b = 2::\text{int}$ ]]

```

```

lemma nmod2:  $n \bmod 2 = 0 \vee n \bmod 2 = 1$ 
  for  $n :: \text{int}$ 
   $\langle \text{proof} \rangle$ 

```

```

lemmas eme1p = emep1 [simplified add commute]

```

```

lemma le-diff-eq':  $a \leq c - b \longleftrightarrow b + a \leq c$ 
  for  $a \ b \ c :: \text{int}$ 
   $\langle \text{proof} \rangle$ 

```

```

lemma less-diff-eq':  $a < c - b \longleftrightarrow b + a < c$ 
  for  $a \ b \ c :: \text{int}$ 
   $\langle \text{proof} \rangle$ 

```

```

lemma diff-less-eq':  $a - b < c \longleftrightarrow a < b + c$ 
  for  $a \ b \ c :: \text{int}$ 
   $\langle \text{proof} \rangle$ 

```

```

lemmas m1mod22k = mult-pos-pos [OF zless2 zless2p, THEN zmod-minus1]

```

```

lemma z1pdiv2:  $(2 * b + 1) \text{ div } 2 = b$ 
  for  $b :: \text{int}$ 
   $\langle \text{proof} \rangle$ 

```

```

lemmas zdiv-le-dividend = xtr3 [OF div-by-1 [symmetric] zdiv-mono2,
  simplified int-one-le-iff-zero-less, simplified]

```

```

lemma axbyy:  $a + m + m = b + n + n \implies a = 0 \vee a = 1 \implies b = 0 \vee b = 1 \implies a = b \wedge m = n$ 

```

```

for  $a\ b\ m\ n :: \text{int}$ 
   $\langle \text{proof} \rangle$ 

lemma  $axxmod2$ :  $(1 + x + x) \bmod 2 = 1 \wedge (0 + x + x) \bmod 2 = 0$ 
  for  $x :: \text{int}$ 
   $\langle \text{proof} \rangle$ 

lemma  $axxdiv2$ :  $(1 + x + x) \text{div } 2 = x \wedge (0 + x + x) \text{div } 2 = x$ 
  for  $x :: \text{int}$ 
   $\langle \text{proof} \rangle$ 

lemmas  $iszero-minus =$ 
   $\text{trans } [THEN \text{trans}, OF \text{iszero-def neg-equal-0-iff-equal iszero-def } [symmetric]]$ 

lemmas  $zadd-diff-inverse =$ 
   $\text{trans } [OF \text{diff-add-cancel } [symmetric] \text{add.commute}]$ 

lemmas  $add-diff-cancel2 =$ 
   $\text{add.commute } [THEN \text{diff-eq-eq } [THEN \text{iffD2}]]$ 

lemmas  $rdmods [symmetric] = \text{mod-minus-eq}$ 
   $\text{mod-diff-left-eq mod-diff-right-eq mod-add-left-eq}$ 
   $\text{mod-add-right-eq mod-mult-right-eq mod-mult-left-eq}$ 

lemma  $\text{mod-plus-right}$ :  $(a + x) \bmod m = (b + x) \bmod m \longleftrightarrow a \bmod m = b \bmod m$ 
   $m$ 
  for  $a\ b\ m\ x :: \text{nat}$ 
   $\langle \text{proof} \rangle$ 

lemma  $\text{nat-minus-mod}$ :  $(n - n \bmod m) \bmod m = 0$ 
  for  $m\ n :: \text{nat}$ 
   $\langle \text{proof} \rangle$ 

lemmas  $\text{nat-minus-mod-plus-right} =$ 
   $\text{trans } [OF \text{nat-minus-mod mod-0 } [symmetric],$ 
   $\text{THEN mod-plus-right } [THEN \text{iffD2}], \text{ simplified}]$ 

lemmas  $\text{push-mods}' = \text{mod-add-eq}$ 
   $\text{mod-mult-eq mod-diff-eq}$ 
   $\text{mod-minus-eq}$ 

lemmas  $\text{push-mods} = \text{push-mods}' [THEN \text{eq-reflection}]$ 
lemmas  $\text{pull-mods} = \text{push-mods} [symmetric] \text{rdmods } [THEN \text{eq-reflection}]$ 

lemma  $\text{nat-mod-eq}$ :  $b < n \implies a \bmod n = b \bmod n \implies a \bmod n = b$ 
  for  $a\ b\ n :: \text{nat}$ 
   $\langle \text{proof} \rangle$ 

lemmas  $\text{nat-mod-eq}' = \text{refl } [THEN [2] \text{nat-mod-eq}]$ 

```

lemma *nat-mod-lem*: $0 < n \implies b < n \iff b \bmod n = b$
for $b\ n :: \text{nat}$
<proof>

lemma *mod-nat-add*: $x < z \implies y < z \implies (x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$
for $x\ y\ z :: \text{nat}$
<proof>

lemma *mod-nat-sub*: $x < z \implies (x - y) \bmod z = x - y$
for $x\ y :: \text{nat}$
<proof>

lemma *int-mod-eq*: $0 \leq b \implies b < n \implies a \bmod n = b \bmod n \implies a \bmod n = b$
for $a\ b\ n :: \text{int}$
<proof>

lemmas *int-mod-eq' = mod-pos-pos-trivial*

lemma *int-mod-le*: $0 \leq a \implies a \bmod n \leq a$
for $a :: \text{int}$
<proof>

lemma *mod-add-if-z*:
 $x < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$
for $x\ y\ z :: \text{int}$
<proof>

lemma *mod-sub-if-z*:
 $x < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x - y) \bmod z = (\text{if } y \leq x \text{ then } x - y \text{ else } x - y + z)$
for $x\ y\ z :: \text{int}$
<proof>

lemmas *zmde = mult-div-mod-eq [symmetric, THEN diff-eq-eq [THEN iffD2], symmetric]*

lemmas *mcl = mult-cancel-left [THEN iffD1, THEN make-pos-rule]*

lemma *zdiv-mult-self*: $m \neq 0 \implies (a + m * n) \text{ div } m = a \text{ div } m + n$
for $a\ m\ n :: \text{int}$
<proof>

lemma *mod-power-lem*: $a > 1 \implies a ^ n \bmod a ^ m = (\text{if } m \leq n \text{ then } 0 \text{ else } a ^ n)$
for $a :: \text{int}$
<proof>

lemma *pl-pl-rels*: $a + b = c + d \implies a \geq c \wedge b \leq d \vee a \leq c \wedge b \geq d$
for $a\ b\ c\ d :: \text{nat}$
 $\langle \text{proof} \rangle$

lemmas *pl-pl-rels'* = *add commute* [THEN [2] *trans*, THEN *pl-pl-rels*]

lemma *minus-eq*: $m - k = m \longleftrightarrow k = 0 \vee m = 0$
for $k\ m :: \text{nat}$
 $\langle \text{proof} \rangle$

lemma *pl-pl-mm*: $a + b = c + d \implies a - c = d - b$
for $a\ b\ c\ d :: \text{nat}$
 $\langle \text{proof} \rangle$

lemmas *pl-pl-mm'* = *add commute* [THEN [2] *trans*, THEN *pl-pl-mm*]

lemmas *dme* = *div-mult-mod-eq*
lemmas *dtle* = *xtr3* [OF *dme* [symmetric] *le-add1*]
lemmas *th2* = *order-trans* [OF *order-reft* [THEN [2] *mult-le-mono*] *dtle*]

lemma *td-gal*: $0 < c \implies a \geq b * c \longleftrightarrow a \text{ div } c \geq b$
for $a\ b\ c :: \text{nat}$
 $\langle \text{proof} \rangle$

lemmas *td-gal-lt* = *td-gal* [*simplified not-less* [symmetric], *simplified*]

lemma *div-mult-le*: $a \text{ div } b * b \leq a$
for $a\ b :: \text{nat}$
 $\langle \text{proof} \rangle$

lemmas *sdl* = *split-div-lemma* [THEN *iffD1*, *symmetric*]

lemma *given-quot*: $f > 0 \implies (f * l + (f - 1)) \text{ div } f = l$
for $f\ l :: \text{nat}$
 $\langle \text{proof} \rangle$

lemma *given-quot-alt*: $f > 0 \implies (l * f + f - \text{Suc } 0) \text{ div } f = l$
for $f\ l :: \text{nat}$
 $\langle \text{proof} \rangle$

lemma *diff-mod-le*: $a < d \implies b \text{ dvd } d \implies a - a \text{ mod } b \leq d - b$
for $a\ b\ d :: \text{nat}$
 $\langle \text{proof} \rangle$

lemma *less-le-mult'*: $w * c < b * c \implies 0 \leq c \implies (w + 1) * c \leq b * c$
for $b\ c\ w :: \text{int}$
 $\langle \text{proof} \rangle$

lemma *less-le-mult*: $w * c < b * c \implies 0 \leq c \implies w * c + c \leq b * c$
for $b \ c \ w :: \text{int}$
 $\langle \text{proof} \rangle$

lemmas *less-le-mult-minus* = *iffD2* [*OF le-diff-eq less-le-mult*,
simplified left-diff-distrib]

lemma *gen-minus*: $0 < n \implies f \ n = f \ (\text{Suc } (n - 1))$
 $\langle \text{proof} \rangle$

lemma *mpl-lem*: $j \leq i \implies k < j \implies i - j + k < i$
for $i \ j \ k :: \text{nat}$
 $\langle \text{proof} \rangle$

lemma *nonneg-mod-div*: $0 \leq a \implies 0 \leq b \implies 0 \leq (a \bmod b) \wedge 0 \leq a \text{ div } b$
for $a \ b :: \text{int}$
 $\langle \text{proof} \rangle$

declare *iszero-0* [*intro*]

lemma *min-pm* [*simp*]: $\min \ a \ b + (a - b) = a$
for $a \ b :: \text{nat}$
 $\langle \text{proof} \rangle$

lemma *min-pm1* [*simp*]: $a - b + \min \ a \ b = a$
for $a \ b :: \text{nat}$
 $\langle \text{proof} \rangle$

lemma *rev-min-pm* [*simp*]: $\min \ b \ a + (a - b) = a$
for $a \ b :: \text{nat}$
 $\langle \text{proof} \rangle$

lemma *rev-min-pm1* [*simp*]: $a - b + \min \ b \ a = a$
for $a \ b :: \text{nat}$
 $\langle \text{proof} \rangle$

lemma *min-minus* [*simp*]: $\min \ m \ (m - k) = m - k$
for $m \ k :: \text{nat}$
 $\langle \text{proof} \rangle$

lemma *min-minus'* [*simp*]: $\min \ (m - k) \ m = m - k$
for $m \ k :: \text{nat}$
 $\langle \text{proof} \rangle$

end

10 A type of finite bit strings

theory *Word*

```

imports
  HOL-Library.Type-Length
  HOL-Library.Boolean-Algebra
  Bits-Bit
  Bool-List-Representation
  Misc-Typedef
  Word-Miscellaneous
begin

```

See `Examples/WordExamples.thy` for examples.

10.1 Type definition

```

typedef (overloaded) 'a word = {(0::int) ..< 2 ^ len-of TYPE('a::len0)}
morphisms uint Abs-word <proof>

```

```

lemma uint-nonnegative: 0 ≤ uint w
  <proof>

```

```

lemma uint-bounded: uint w < 2 ^ len-of TYPE('a)
  for w :: 'a::len0 word
  <proof>

```

```

lemma uint-idem: uint w mod 2 ^ len-of TYPE('a) = uint w
  for w :: 'a::len0 word
  <proof>

```

```

lemma word-uint-eq-iff: a = b ⟷ uint a = uint b
  <proof>

```

```

lemma word-uint-eqI: uint a = uint b ⟹ a = b
  <proof>

```

```

definition word-of-int :: int ⇒ 'a::len0 word
  — representation of words using unsigned or signed bins, only difference in these
  is the type class
  where word-of-int k = Abs-word (k mod 2 ^ len-of TYPE('a))

```

```

lemma uint-word-of-int: uint (word-of-int k :: 'a::len0 word) = k mod 2 ^ len-of
  TYPE('a)
  <proof>

```

```

lemma word-of-int-uint: word-of-int (uint w) = w
  <proof>

```

```

lemma split-word-all: (⋀x::'a::len0 word. PROP P x) ≡ (⋀x. PROP P (word-of-int
  x))
  <proof>

```

10.2 Type conversions and casting

definition *sint* :: 'a::len word $\Rightarrow$ int

— treats the most-significant-bit as a sign bit

where *sint-uint*: *sint* *w* = *sbintrunc* (*len-of* *TYPE*('a) - 1) (*uint* *w*)

definition *unat* :: 'a::len0 word $\Rightarrow$ nat

where *unat* *w* = *nat* (*uint* *w*)

definition *uints* :: nat $\Rightarrow$ int set

— the sets of integers representing the words

where *uints* *n* = *range* (*bintrunc* *n*)

definition *sints* :: nat $\Rightarrow$ int set

where *sints* *n* = *range* (*sbintrunc* (*n* - 1))

lemma *uints-num*: *uints* *n* = {*i*. $0 \leq i \wedge i < 2^n$ }

<proof>

lemma *sints-num*: *sints* *n* = {*i*. $-(2^{n-1}) \leq i \wedge i < 2^n$ }

<proof>

definition *unats* :: nat $\Rightarrow$ nat set

where *unats* *n* = {*i*. $i < 2^n$ }

definition *norm-sint* :: nat $\Rightarrow$ int $\Rightarrow$ int

where *norm-sint* *n* *w* = (*w* + 2^{n-1}) mod $2^n - 2^{n-1}$

definition *scast* :: 'a::len word $\Rightarrow$ 'b::len word

— cast a word to a different length

where *scast* *w* = *word-of-int* (*sint* *w*)

definition *ucast* :: 'a::len0 word $\Rightarrow$ 'b::len0 word

where *ucast* *w* = *word-of-int* (*uint* *w*)

instantiation *word* :: (len0) size

begin

definition *word-size*: size (*w* :: 'a word) = *len-of* *TYPE*('a)

instance *<proof>*

end

lemma *word-size-gt-0* [iff]: $0 < \text{size } w$

for *w* :: 'a::len word

<proof>

lemmas *lens-gt-0* = *word-size-gt-0* *len-gt-0*

lemma *lens-not-0* [iff]:
 fixes $w :: 'a::len\ word$
 shows $size\ w \neq 0$
 and $len-of\ TYPE('a) \neq 0$
 ⟨proof⟩

definition *source-size* :: $('a::len0\ word \Rightarrow 'b) \Rightarrow nat$
 — whether a cast (or other) function is to a longer or shorter length
 where [code del]: $source-size\ c = (let\ arb = undefined; x = c\ arb\ in\ size\ arb)$

definition *target-size* :: $('a \Rightarrow 'b::len0\ word) \Rightarrow nat$
 where [code del]: $target-size\ c = size\ (c\ undefined)$

definition *is-up* :: $('a::len0\ word \Rightarrow 'b::len0\ word) \Rightarrow bool$
 where $is-up\ c \longleftrightarrow source-size\ c \leq target-size\ c$

definition *is-down* :: $('a::len0\ word \Rightarrow 'b::len0\ word) \Rightarrow bool$
 where $is-down\ c \longleftrightarrow target-size\ c \leq source-size\ c$

definition *of-bl* :: $bool\ list \Rightarrow 'a::len0\ word$
 where $of-bl\ bl = word-of-int\ (bl-to-bin\ bl)$

definition *to-bl* :: $'a::len0\ word \Rightarrow bool\ list$
 where $to-bl\ w = bin-to-bl\ (len-of\ TYPE('a))\ (uint\ w)$

definition *word-reverse* :: $'a::len0\ word \Rightarrow 'a\ word$
 where $word-reverse\ w = of-bl\ (rev\ (to-bl\ w))$

definition *word-int-case* :: $(int \Rightarrow 'b) \Rightarrow 'a::len0\ word \Rightarrow 'b$
 where $word-int-case\ f\ w = f\ (uint\ w)$

translations

$case\ x\ of\ XCONST\ of-int\ y \Rightarrow b \Rightarrow CONST\ word-int-case\ (\lambda y. b)\ x$
 $case\ x\ of\ (XCONST\ of-int :: 'a)\ y \Rightarrow b \rightarrow CONST\ word-int-case\ (\lambda y. b)\ x$

10.3 Correspondence relation for theorem transfer

definition *cr-word* :: $int \Rightarrow 'a::len0\ word \Rightarrow bool$
 where $cr-word = (\lambda x\ y. word-of-int\ x = y)$

lemma *Quotient-word*:
 $Quotient\ (\lambda x\ y. bintrunc\ (len-of\ TYPE('a))\ x = bintrunc\ (len-of\ TYPE('a))\ y)$
 $word-of-int\ uint\ (cr-word :: - \Rightarrow 'a::len0\ word \Rightarrow bool)$
 ⟨proof⟩

lemma *reflp-word*:
 $reflp\ (\lambda x\ y. bintrunc\ (len-of\ TYPE('a::len0))\ x = bintrunc\ (len-of\ TYPE('a))\ y)$
 ⟨proof⟩

setup-lifting *Quotient-word reflp-word*

TODO: The next lemma could be generated automatically.

lemma *uint-transfer* [*transfer-rule*]:
 (*rel-fun pcr-word op =*) (*bintrunc (len-of TYPE('a'))*) (*uint :: 'a::len0 word ⇒ int*)
<proof>

10.4 Basic code generation setup

definition *Word* :: *int ⇒ 'a::len0 word*
where [*code-post*]: *Word = word-of-int*

lemma [*code abstype*]: *Word (uint w) = w*
<proof>

declare *uint-word-of-int* [*code abstract*]

instantiation *word* :: (*len0*) *equal*
begin

definition *equal-word* :: '*a word ⇒ 'a word ⇒ bool*
where *equal-word k l ⇔ HOL.equal (uint k) (uint l)*

instance
<proof>

end

notation *fcomp* (**infixl** *o>* 60)
notation *scomp* (**infixl** *o→* 60)

instantiation *word* :: ({*len0*, *typerep*}) *random*
begin

definition
random-word i = Random.range i o→ (λk. Pair (
let j = word-of-int (int-of-integer (integer-of-natural k)) :: 'a word
in (j, λ::unit. Code-Evaluation.term-of j)))

instance *<proof>*

end

no-notation *fcomp* (**infixl** *o>* 60)
no-notation *scomp* (**infixl** *o→* 60)

10.5 Type-definition locale instantiations

lemmas *uint-0* = *uint-nonnegative*

lemmas *uint-lt* = *uint-bounded*

lemmas *uint-mod-same* = *uint-idem*

lemma *td-ext-uint*:

td-ext (*uint* :: 'a word $\Rightarrow$ int) *word-of-int* (*uints* (*len-of* TYPE('a::len0)))
 ($\lambda w::\text{int}. w \bmod 2 \wedge \text{len-of TYPE('a)}$)
<proof>

interpretation *word-uint*:

td-ext
uint::'a::len0 word $\Rightarrow$ int
word-of-int
uints (*len-of* TYPE('a::len0))
 $\lambda w. w \bmod 2 \wedge \text{len-of TYPE('a::len0)}$
<proof>

lemmas *td-uint* = *word-uint.td-thm*

lemmas *int-word-uint* = *word-uint.eq-norm*

lemma *td-ext-ubin*:

td-ext (*uint* :: 'a word $\Rightarrow$ int) *word-of-int* (*uints* (*len-of* TYPE('a::len0)))
 (*bintrunc* (*len-of* TYPE('a)))
<proof>

interpretation *word-ubin*:

td-ext
uint::'a::len0 word $\Rightarrow$ int
word-of-int
uints (*len-of* TYPE('a::len0))
bintrunc (*len-of* TYPE('a::len0))
<proof>

10.6 Arithmetic operations

lift-definition *word-succ* :: 'a::len0 word $\Rightarrow$ 'a word **is** $\lambda x. x + 1$

<proof>

lift-definition *word-pred* :: 'a::len0 word $\Rightarrow$ 'a word **is** $\lambda x. x - 1$

<proof>

instantiation *word* :: (len0) {*neg-numeral*, *modulo*, *comm-monoid-mult*, *comm-ring*}
begin

lift-definition *zero-word* :: 'a word **is** 0 *<proof>*

lift-definition *one-word* :: 'a word **is** 1 *<proof>*

lift-definition *plus-word* :: 'a word $\Rightarrow$ 'a word $\Rightarrow$ 'a word **is** op +
 ⟨proof⟩

lift-definition *minus-word* :: 'a word $\Rightarrow$ 'a word $\Rightarrow$ 'a word **is** op −
 ⟨proof⟩

lift-definition *uminus-word* :: 'a word $\Rightarrow$ 'a word **is** uminus
 ⟨proof⟩

lift-definition *times-word* :: 'a word $\Rightarrow$ 'a word $\Rightarrow$ 'a word **is** op *
 ⟨proof⟩

definition *word-div-def*: $a \text{ div } b = \text{word-of-int } (\text{uint } a \text{ div uint } b)$

definition *word-mod-def*: $a \text{ mod } b = \text{word-of-int } (\text{uint } a \text{ mod uint } b)$

instance
 ⟨proof⟩

end

Legacy theorems:

lemma *word-arith-wis* [code]:
shows *word-add-def*: $a + b = \text{word-of-int } (\text{uint } a + \text{uint } b)$
and *word-sub-wi*: $a - b = \text{word-of-int } (\text{uint } a - \text{uint } b)$
and *word-mult-def*: $a * b = \text{word-of-int } (\text{uint } a * \text{uint } b)$
and *word-minus-def*: $- a = \text{word-of-int } (- \text{uint } a)$
and *word-succ-alt*: $\text{word-succ } a = \text{word-of-int } (\text{uint } a + 1)$
and *word-pred-alt*: $\text{word-pred } a = \text{word-of-int } (\text{uint } a - 1)$
and *word-0-wi*: $0 = \text{word-of-int } 0$
and *word-1-wi*: $1 = \text{word-of-int } 1$
 ⟨proof⟩

lemma *wi-homs*:
shows *wi-hom-add*: $\text{word-of-int } a + \text{word-of-int } b = \text{word-of-int } (a + b)$
and *wi-hom-sub*: $\text{word-of-int } a - \text{word-of-int } b = \text{word-of-int } (a - b)$
and *wi-hom-mult*: $\text{word-of-int } a * \text{word-of-int } b = \text{word-of-int } (a * b)$
and *wi-hom-neg*: $- \text{word-of-int } a = \text{word-of-int } (- a)$
and *wi-hom-succ*: $\text{word-succ } (\text{word-of-int } a) = \text{word-of-int } (a + 1)$
and *wi-hom-pred*: $\text{word-pred } (\text{word-of-int } a) = \text{word-of-int } (a - 1)$
 ⟨proof⟩

lemmas *wi-hom-syms* = *wi-homs* [symmetric]

lemmas *word-of-int-homs* = *wi-homs* *word-0-wi* *word-1-wi*

lemmas *word-of-int-hom-syms* = *word-of-int-homs* [symmetric]

instance *word* :: (len) comm-ring-1

<proof>

lemma *word-of-nat*: *of-nat* *n* = *word-of-int* (*int* *n*)
<proof>

lemma *word-of-int*: *of-int* = *word-of-int*
<proof>

definition *udvd* :: '*a*::*len* *word* ⇒ '*a*::*len* *word* ⇒ *bool* (**infixl** *udvd* 50)
where *a* *udvd* *b* = (∃ *n* ≥ 0. *uint* *b* = *n* * *uint* *a*)

10.7 Ordering

instantiation *word* :: (*len0*) *linorder*
begin

definition *word-le-def*: *a* ≤ *b* ⇔ *uint* *a* ≤ *uint* *b*

definition *word-less-def*: *a* < *b* ⇔ *uint* *a* < *uint* *b*

instance
<proof>

end

definition *word-sle* :: '*a*::*len* *word* ⇒ '*a* *word* ⇒ *bool* ((-/ <=s -) [50, 51] 50)
where *a* <=s *b* ⇔ *sint* *a* ≤ *sint* *b*

definition *word-sless* :: '*a*::*len* *word* ⇒ '*a* *word* ⇒ *bool* ((-/ <s -) [50, 51] 50)
where *x* <s *y* ⇔ *x* <=s *y* ∧ *x* ≠ *y*

10.8 Bit-wise operations

instantiation *word* :: (*len0*) *bits*
begin

lift-definition *bitNOT-word* :: '*a* *word* ⇒ '*a* *word* **is** *bitNOT*
<proof>

lift-definition *bitAND-word* :: '*a* *word* ⇒ '*a* *word* ⇒ '*a* *word* **is** *bitAND*
<proof>

lift-definition *bitOR-word* :: '*a* *word* ⇒ '*a* *word* ⇒ '*a* *word* **is** *bitOR*
<proof>

lift-definition *bitXOR-word* :: '*a* *word* ⇒ '*a* *word* ⇒ '*a* *word* **is** *bitXOR*
<proof>

definition *word-test-bit-def*: *test-bit* *a* = *bin-nth* (*uint* *a*)

definition *word-set-bit-def*: $\text{set-bit } a \ n \ x = \text{word-of-int } (\text{bin-sc } n \ x \ (\text{uint } a))$

definition *word-set-bits-def*: $(\text{BITS } n. \ f \ n) = \text{of-bl } (\text{bl-of-nth } (\text{len-of TYPE('a)}) \ f)$

definition *word-lsb-def*: $\text{lsb } a \longleftrightarrow \text{bin-last } (\text{uint } a)$

definition *shiftrl1* :: $'a \ \text{word} \Rightarrow 'a \ \text{word}$
where *shiftrl1* $w = \text{word-of-int } (\text{uint } w \ \text{BIT } \text{False})$

definition *shiftr1* :: $'a \ \text{word} \Rightarrow 'a \ \text{word}$
— shift right as unsigned or as signed, ie logical or arithmetic
where *shiftr1* $w = \text{word-of-int } (\text{bin-rest } (\text{uint } w))$

definition *shiftrl-def*: $w \ll n = (\text{shiftrl1 } ^{\wedge} n) \ w$

definition *shiftr-def*: $w \gg n = (\text{shiftr1 } ^{\wedge} n) \ w$

instance $\langle \text{proof} \rangle$

end

lemma *[code]*:

shows *word-not-def*: $\text{NOT } (a::'a::\text{len0 } \text{word}) = \text{word-of-int } (\text{NOT } (\text{uint } a))$
and *word-and-def*: $(a::'a \ \text{word}) \ \text{AND } b = \text{word-of-int } (\text{uint } a \ \text{AND } \text{uint } b)$
and *word-or-def*: $(a::'a \ \text{word}) \ \text{OR } b = \text{word-of-int } (\text{uint } a \ \text{OR } \text{uint } b)$
and *word-xor-def*: $(a::'a \ \text{word}) \ \text{XOR } b = \text{word-of-int } (\text{uint } a \ \text{XOR } \text{uint } b)$
 $\langle \text{proof} \rangle$

instantiation *word* :: $(\text{len}) \ \text{bitss}$
begin

definition *word-msb-def*: $\text{msb } a \longleftrightarrow \text{bin-sign } (\text{sint } a) = -1$

instance $\langle \text{proof} \rangle$

end

definition *setBit* :: $'a::\text{len0 } \text{word} \Rightarrow \text{nat} \Rightarrow 'a \ \text{word}$
where *setBit* $w \ n = \text{set-bit } w \ n \ \text{True}$

definition *clearBit* :: $'a::\text{len0 } \text{word} \Rightarrow \text{nat} \Rightarrow 'a \ \text{word}$
where *clearBit* $w \ n = \text{set-bit } w \ n \ \text{False}$

10.9 Shift operations

definition *sshiftr1* :: $'a::\text{len } \text{word} \Rightarrow 'a \ \text{word}$
where *sshiftr1* $w = \text{word-of-int } (\text{bin-rest } (\text{sint } w))$

definition *bshiftr1* :: *bool* $\Rightarrow$ '*a*::*len* *word* $\Rightarrow$ '*a* *word*
where *bshiftr1* *b w* = *of-bl* (*b* # *butlast* (*to-bl w*))

definition *sshiftr* :: '*a*::*len* *word* $\Rightarrow$ *nat* $\Rightarrow$ '*a* *word* (**infixl** >>> 55)
where *w* >>> *n* = (*sshiftr1* ^^ *n*) *w*

definition *mask* :: *nat* $\Rightarrow$ '*a*::*len* *word*
where *mask* *n* = (*1* << *n*) - 1

definition *revcast* :: '*a*::*len0* *word* $\Rightarrow$ '*b*::*len0* *word*
where *revcast w* = *of-bl* (*takefill* *False* (*len-of TYPE*('b')) (*to-bl w*))

definition *slice1* :: *nat* $\Rightarrow$ '*a*::*len0* *word* $\Rightarrow$ '*b*::*len0* *word*
where *slice1 n w* = *of-bl* (*takefill* *False* *n* (*to-bl w*))

definition *slice* :: *nat* $\Rightarrow$ '*a*::*len0* *word* $\Rightarrow$ '*b*::*len0* *word*
where *slice n w* = *slice1* (*size w* - *n*) *w*

10.10 Rotation

definition *rotater1* :: '*a* *list* $\Rightarrow$ '*a* *list*
where *rotater1* *ys* =
 (*case* *ys* *of* [] $\Rightarrow$ [] | *x* # *xs* $\Rightarrow$ *last* *ys* # *butlast* *ys*)

definition *rotater* :: *nat* $\Rightarrow$ '*a* *list* $\Rightarrow$ '*a* *list*
where *rotater n* = *rotater1* ^^ *n*

definition *word-rotr* :: *nat* $\Rightarrow$ '*a*::*len0* *word* $\Rightarrow$ '*a*::*len0* *word*
where *word-rotr n w* = *of-bl* (*rotater n* (*to-bl w*))

definition *word-rotl* :: *nat* $\Rightarrow$ '*a*::*len0* *word* $\Rightarrow$ '*a*::*len0* *word*
where *word-rotl n w* = *of-bl* (*rotate n* (*to-bl w*))

definition *word-roti* :: *int* $\Rightarrow$ '*a*::*len0* *word* $\Rightarrow$ '*a*::*len0* *word*
where *word-roti i w* =
 (*if* *i* $\geq$ 0 *then* *word-rotr* (*nat i*) *w* *else* *word-rotl* (*nat* (- *i*)) *w*)

10.11 Split and cat operations

definition *word-cat* :: '*a*::*len0* *word* $\Rightarrow$ '*b*::*len0* *word* $\Rightarrow$ '*c*::*len0* *word*
where *word-cat a b* = *word-of-int* (*bin-cat* (*uint a*) (*len-of TYPE*('b')) (*uint b*))

definition *word-split* :: '*a*::*len0* *word* $\Rightarrow$ '*b*::*len0* *word* $\times$ '*c*::*len0* *word*
where *word-split a* =
 (*case* *bin-split* (*len-of TYPE*('c')) (*uint a*) *of*
 (*u, v*) $\Rightarrow$ (*word-of-int u, word-of-int v*))

definition *word-rcat* :: '*a*::*len0* *word* *list* $\Rightarrow$ '*b*::*len0* *word*
where *word-rcat ws* = *word-of-int* (*bin-rcat* (*len-of TYPE*('a')) (*map uint ws*))

definition *word-rsplit* :: 'a::len0 word $\Rightarrow$ 'b::len word list
 — where *word-rsplit* *w* = map *word-of-int* (*bin-rsplit* (*len-of* *TYPE*('b)) (*len-of* *TYPE*('a), *uint* *w*))

definition *max-word* :: 'a::len word
 — Largest representable machine integer.
 — where *max-word* = *word-of-int* ($2^{\text{len-of } \text{TYPE}('a)} - 1$)

lemmas *of-nth-def* = *word-set-bits-def*

10.12 Theorems about typedefs

lemma *sint-sbintrunc'*: *sint* (*word-of-int* *bin* :: 'a word) = *sbintrunc* (*len-of* *TYPE*('a::len) - 1) *bin*
 <proof>

lemma *uint-sint*: *uint* *w* = *bintrunc* (*len-of* *TYPE*('a)) (*sint* *w*)
 for *w* :: 'a::len word
 <proof>

lemma *bintr-uint*: *len-of* *TYPE*('a) $\leq n \implies$ *bintrunc* *n* (*uint* *w*) = *uint* *w*
 for *w* :: 'a::len0 word
 <proof>

lemma *wi-bintr*:
len-of *TYPE*('a::len0) $\leq n \implies$
word-of-int (*bintrunc* *n* *w*) = (*word-of-int* *w* :: 'a word)
 <proof>

lemma *td-ext-sbin*:
td-ext (*sint* :: 'a word $\Rightarrow$ int) *word-of-int* (*sints* (*len-of* *TYPE*('a::len)))
 (*sbintrunc* (*len-of* *TYPE*('a) - 1))
 <proof>

lemma *td-ext-sint*:
td-ext (*sint* :: 'a word $\Rightarrow$ int) *word-of-int* (*sints* (*len-of* *TYPE*('a::len)))
 ($\lambda w. (w + 2^{\text{len-of } \text{TYPE}('a)} - 1) \bmod 2^{\text{len-of } \text{TYPE}('a)} - 2^{\text{len-of } \text{TYPE}('a)} - 1$)
 <proof>

interpretation *word-sint*:
td-ext
sint :: 'a::len word $\Rightarrow$ int
word-of-int
sints (*len-of* *TYPE*('a::len))
 $\lambda w. (w + 2^{\text{len-of } \text{TYPE}('a)} - 1) \bmod 2^{\text{len-of } \text{TYPE}('a)} - 2^{\text{len-of } \text{TYPE}('a)} - 1$
 <proof>

interpretation *word-sbin*:

td-ext
 $\text{sint} :: 'a :: \text{len } \text{word} \Rightarrow \text{int}$
word-of-int
 $\text{sints } (\text{len-of } \text{TYPE}('a :: \text{len}))$
 $\text{sbintrunc } (\text{len-of } \text{TYPE}('a :: \text{len}) - 1)$
 $\langle \text{proof} \rangle$

lemmas *int-word-sint* = *td-ext-sint* [THEN *td-ext.eq-norm*]

lemmas *td-sint* = *word-sint.td*

lemma *to-bl-def'*: $(\text{to-bl} :: 'a :: \text{len0 } \text{word} \Rightarrow \text{bool list}) = \text{bin-to-bl } (\text{len-of } \text{TYPE}('a))$
 $\circ \text{uint}$
 $\langle \text{proof} \rangle$

lemmas *word-reverse-no-def* [simp] =
word-reverse-def [of numeral *w*] **for** *w*

lemma *uints-mod*: $\text{uints } n = \text{range } (\lambda w. w \bmod 2 ^ n)$
 $\langle \text{proof} \rangle$

lemma *word-numeral-alt*: $\text{numeral } b = \text{word-of-int } (\text{numeral } b)$
 $\langle \text{proof} \rangle$

declare *word-numeral-alt* [symmetric, code-abbrev]

lemma *word-neg-numeral-alt*: $-\text{numeral } b = \text{word-of-int } (-\text{numeral } b)$
 $\langle \text{proof} \rangle$

declare *word-neg-numeral-alt* [symmetric, code-abbrev]

lemma *word-numeral-transfer* [transfer-rule]:
 $(\text{rel-fun } \text{op} = \text{pcr-word}) \text{ numeral numeral}$
 $(\text{rel-fun } \text{op} = \text{pcr-word}) (-\text{numeral}) (-\text{numeral})$
 $\langle \text{proof} \rangle$

lemma *uint-bintrunc* [simp]:
 $\text{uint } (\text{numeral } \text{bin} :: 'a \text{ word}) =$
 $\text{bintrunc } (\text{len-of } \text{TYPE}('a :: \text{len0})) (\text{numeral } \text{bin})$
 $\langle \text{proof} \rangle$

lemma *uint-bintrunc-neg* [simp]:
 $\text{uint } (-\text{numeral } \text{bin} :: 'a \text{ word}) = \text{bintrunc } (\text{len-of } \text{TYPE}('a :: \text{len0})) (-\text{numeral } \text{bin})$
 $\langle \text{proof} \rangle$

lemma *sint-sbintrunc* [simp]:

$\text{sint } (\text{numeral bin} :: 'a \text{ word}) = \text{sbintrunc } (\text{len-of TYPE}('a::\text{len}) - 1) (\text{numeral bin})$
 ⟨proof⟩

lemma *sint-sbintrunc-neg* [simp]:
 $\text{sint } (- \text{numeral bin} :: 'a \text{ word}) = \text{sbintrunc } (\text{len-of TYPE}('a::\text{len}) - 1) (- \text{numeral bin})$
 ⟨proof⟩

lemma *unat-bintrunc* [simp]:
 $\text{unat } (\text{numeral bin} :: 'a::\text{len0 word}) = \text{nat } (\text{bintrunc } (\text{len-of TYPE}('a)) (\text{numeral bin}))$
 ⟨proof⟩

lemma *unat-bintrunc-neg* [simp]:
 $\text{unat } (- \text{numeral bin} :: 'a::\text{len0 word}) = \text{nat } (\text{bintrunc } (\text{len-of TYPE}('a)) (- \text{numeral bin}))$
 ⟨proof⟩

lemma *size-0-eq*: $\text{size } w = 0 \implies v = w$
for $v \ w :: 'a::\text{len0 word}$
 ⟨proof⟩

lemma *uint-ge-0* [iff]: $0 \leq \text{uint } x$
for $x :: 'a::\text{len0 word}$
 ⟨proof⟩

lemma *uint-lt2p* [iff]: $\text{uint } x < 2 ^ \text{len-of TYPE}('a)$
for $x :: 'a::\text{len0 word}$
 ⟨proof⟩

lemma *sint-ge*: $- (2 ^ (\text{len-of TYPE}('a) - 1)) \leq \text{sint } x$
for $x :: 'a::\text{len word}$
 ⟨proof⟩

lemma *sint-lt*: $\text{sint } x < 2 ^ (\text{len-of TYPE}('a) - 1)$
for $x :: 'a::\text{len word}$
 ⟨proof⟩

lemma *sign-uint-Pls* [simp]: $\text{bin-sign } (\text{uint } x) = 0$
 ⟨proof⟩

lemma *uint-m2p-neg*: $\text{uint } x - 2 ^ \text{len-of TYPE}('a) < 0$
for $x :: 'a::\text{len0 word}$
 ⟨proof⟩

lemma *uint-m2p-not-non-neg*: $\neg 0 \leq \text{uint } x - 2 ^ \text{len-of TYPE}('a)$
for $x :: 'a::\text{len0 word}$
 ⟨proof⟩

lemma *lt2p-lem*: $\text{len-of TYPE}('a) \leq n \implies \text{uint } w < 2^{\wedge n}$
for $w :: 'a::\text{len0 word}$
 ⟨proof⟩

lemma *uint-le-0-iff* [simp]: $\text{uint } x \leq 0 \longleftrightarrow \text{uint } x = 0$
 ⟨proof⟩

lemma *uint-nat*: $\text{uint } w = \text{int } (\text{unat } w)$
 ⟨proof⟩

lemma *uint-numeral*: $\text{uint } (\text{numeral } b :: 'a::\text{len0 word}) = \text{numeral } b \bmod 2^{\wedge \text{len-of TYPE}('a)}$
 ⟨proof⟩

lemma *uint-neg-numeral*: $\text{uint } (- \text{numeral } b :: 'a::\text{len0 word}) = - \text{numeral } b \bmod 2^{\wedge \text{len-of TYPE}('a)}$
 ⟨proof⟩

lemma *unat-numeral*: $\text{unat } (\text{numeral } b :: 'a::\text{len0 word}) = \text{numeral } b \bmod 2^{\wedge \text{len-of TYPE}('a)}$
 ⟨proof⟩

lemma *sint-numeral*:
 $\text{sint } (\text{numeral } b :: 'a::\text{len word}) =$
 $(\text{numeral } b +$
 $2^{\wedge (\text{len-of TYPE}('a) - 1)}) \bmod 2^{\wedge \text{len-of TYPE}('a) -}$
 $2^{\wedge (\text{len-of TYPE}('a) - 1)})$
 ⟨proof⟩

lemma *word-of-int-0* [simp, code-post]: $\text{word-of-int } 0 = 0$
 ⟨proof⟩

lemma *word-of-int-1* [simp, code-post]: $\text{word-of-int } 1 = 1$
 ⟨proof⟩

lemma *word-of-int-neg-1* [simp]: $\text{word-of-int } (- 1) = - 1$
 ⟨proof⟩

lemma *word-of-int-numeral* [simp] : $(\text{word-of-int } (\text{numeral } bin) :: 'a::\text{len0 word})$
 $= \text{numeral } bin$
 ⟨proof⟩

lemma *word-of-int-neg-numeral* [simp]:
 $(\text{word-of-int } (- \text{numeral } bin) :: 'a::\text{len0 word}) = - \text{numeral } bin$
 ⟨proof⟩

lemma *word-int-case-wi*:
 $\text{word-int-case } f (\text{word-of-int } i :: 'b \text{ word}) = f (i \bmod 2^{\wedge \text{len-of TYPE}('b::\text{len0})})$

$\langle \text{proof} \rangle$

lemma *word-int-split*:

$P \text{ (word-int-case } f \text{ } x) =$
 $(\forall i. x = (\text{word-of-int } i :: 'b::\text{len0 word}) \wedge 0 \leq i \wedge i < 2^{\text{len-of TYPE('b)}} \longrightarrow P (f i))$
 $\langle \text{proof} \rangle$

lemma *word-int-split-asm*:

$P \text{ (word-int-case } f \text{ } x) =$
 $(\nexists n. x = (\text{word-of-int } n :: 'b::\text{len0 word}) \wedge 0 \leq n \wedge n < 2^{\text{len-of TYPE('b::len0)}} \wedge \neg P (f n))$
 $\langle \text{proof} \rangle$

lemmas *uint-range'* = *word-uint.Rep* [unfolded uints-num mem-Collect-eq]

lemmas *sint-range'* = *word-sint.Rep* [unfolded One-nat-def sints-num mem-Collect-eq]

lemma *uint-range-size*: $0 \leq \text{uint } w \wedge \text{uint } w < 2^{\text{size } w}$
 $\langle \text{proof} \rangle$

lemma *sint-range-size*: $-(2^{\text{size } w - \text{Suc } 0}) \leq \text{sint } w \wedge \text{sint } w < 2^{\text{size } w - \text{Suc } 0}$
 $\langle \text{proof} \rangle$

lemma *sint-above-size*: $2^{\text{size } w - 1} \leq x \implies \text{sint } w < x$
for $w :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemma *sint-below-size*: $x \leq -(2^{\text{size } w - 1}) \implies x \leq \text{sint } w$
for $w :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

10.13 Testing bits

lemma *test-bit-eq-iff*: $\text{test-bit } u = \text{test-bit } v \longleftrightarrow u = v$
for $u \ v :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *test-bit-size* [rule-format]: $w !! n \longrightarrow n < \text{size } w$
for $w :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *word-eq-iff*: $x = y \longleftrightarrow (\forall n < \text{len-of TYPE('a)}. x !! n = y !! n)$
for $x \ y :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *word-eqI*: $(\bigwedge n. n < \text{size } u \longrightarrow u !! n = v !! n) \implies u = v$
for $u :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *word-eqD*: $u = v \implies u !! x = v !! x$
for $u\ v :: 'a::len0\ word$
 $\langle proof \rangle$

lemma *test-bit-bin'*: $w !! n \longleftrightarrow n < size\ w \wedge bin_nth\ (uint\ w)\ n$
 $\langle proof \rangle$

lemmas *test-bit-bin* = *test-bit-bin'* [unfolded word-size]

lemma *bin-nth-uint-imp*: $bin_nth\ (uint\ w)\ n \implies n < len_of\ TYPE('a)$
for $w :: 'a::len0\ word$
 $\langle proof \rangle$

lemma *bin-nth-sint*:
 $len_of\ TYPE('a) \leq n \implies$
 $bin_nth\ (sint\ w)\ n = bin_nth\ (sint\ w)\ (len_of\ TYPE('a) - 1)$
for $w :: 'a::len\ word$
 $\langle proof \rangle$

lemma *td-bl*:
type-definition
 $(to_bl :: 'a::len0\ word \Rightarrow bool\ list)$
of-bl
 $\{bl.\ length\ bl = len_of\ TYPE('a)\}$
 $\langle proof \rangle$

interpretation *word-bl*:
type-definition
 $to_bl :: 'a::len0\ word \Rightarrow bool\ list$
of-bl
 $\{bl.\ length\ bl = len_of\ TYPE('a::len0)\}$
 $\langle proof \rangle$

lemmas *word-bl-Rep'* = *word-bl.Rep* [unfolded mem-Collect-eq, iff]

lemma *word-size-bl*: $size\ w = size\ (to_bl\ w)$
 $\langle proof \rangle$

lemma *to-bl-use-of-bl*: $to_bl\ w = bl \longleftrightarrow w = of_bl\ bl \wedge length\ bl = length\ (to_bl\ w)$
 $\langle proof \rangle$

lemma *to-bl-word-rev*: $to_bl\ (word_reverse\ w) = rev\ (to_bl\ w)$
 $\langle proof \rangle$

lemma *word-rev-rev* [simp]: $word_reverse\ (word_reverse\ w) = w$
 $\langle proof \rangle$

lemma *word-rev-gal*: $\text{word-reverse } w = u \implies \text{word-reverse } u = w$
 ⟨proof⟩

lemma *word-rev-gal'*: $u = \text{word-reverse } w \implies w = \text{word-reverse } u$
 ⟨proof⟩

lemma *length-bl-gt-0* [iff]: $0 < \text{length } (\text{to-bl } x)$
 for $x :: 'a::\text{len } \text{word}$
 ⟨proof⟩

lemma *bl-not-Nil* [iff]: $\text{to-bl } x \neq []$
 for $x :: 'a::\text{len } \text{word}$
 ⟨proof⟩

lemma *length-bl-neq-0* [iff]: $\text{length } (\text{to-bl } x) \neq 0$
 for $x :: 'a::\text{len } \text{word}$
 ⟨proof⟩

lemma *hd-bl-sign-sint*: $\text{hd } (\text{to-bl } w) = (\text{bin-sign } (\text{sint } w) = -1)$
 ⟨proof⟩

lemma *of-bl-drop'*:
 $\text{lend} = \text{length } \text{bl} - \text{len-of TYPE}('a::\text{len}0) \implies$
 $\text{of-bl } (\text{drop } \text{lend } \text{bl}) = (\text{of-bl } \text{bl} :: 'a \text{ word})$
 ⟨proof⟩

lemma *test-bit-of-bl*:
 $(\text{of-bl } \text{bl} :: 'a::\text{len}0 \text{ word}) !! n = (\text{rev } \text{bl} ! n \wedge n < \text{len-of TYPE}('a) \wedge n < \text{length } \text{bl})$
 ⟨proof⟩

lemma *no-of-bl*: $(\text{numeral } \text{bin} :: 'a::\text{len}0 \text{ word}) = \text{of-bl } (\text{bin-to-bl } (\text{len-of TYPE}('a))$
 $(\text{numeral } \text{bin}))$
 ⟨proof⟩

lemma *uint-bl*: $\text{to-bl } w = \text{bin-to-bl } (\text{size } w) (\text{uint } w)$
 ⟨proof⟩

lemma *to-bl-bin*: $\text{bl-to-bin } (\text{to-bl } w) = \text{uint } w$
 ⟨proof⟩

lemma *to-bl-of-bin*: $\text{to-bl } (\text{word-of-int } \text{bin} :: 'a::\text{len}0 \text{ word}) = \text{bin-to-bl } (\text{len-of TYPE}('a))$
 bin
 ⟨proof⟩

lemma *to-bl-numeral* [simp]:
 $\text{to-bl } (\text{numeral } \text{bin} :: 'a::\text{len}0 \text{ word}) =$
 $\text{bin-to-bl } (\text{len-of TYPE}('a)) (\text{numeral } \text{bin})$

$\langle \text{proof} \rangle$

lemma *to-bl-neg-numeral* [simp]:
 $\text{to-bl } (- \text{ numeral bin} :: 'a :: \text{len0 word}) =$
 $\text{bin-to-bl } (\text{len-of TYPE}('a)) (- \text{ numeral bin})$
 $\langle \text{proof} \rangle$

lemma *to-bl-to-bin* [simp] : $\text{bl-to-bin } (\text{to-bl } w) = \text{uint } w$
 $\langle \text{proof} \rangle$

lemma *uint-bl-bin*: $\text{bl-to-bin } (\text{bin-to-bl } (\text{len-of TYPE}('a)) (\text{uint } x)) = \text{uint } x$
for $x :: 'a :: \text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *uints-unats*: $\text{uints } n = \text{int} \text{ ` } \text{unats } n$
 $\langle \text{proof} \rangle$

lemma *unats-uints*: $\text{unats } n = \text{nat} \text{ ` } \text{uints } n$
 $\langle \text{proof} \rangle$

lemmas *bintr-num* =
 $\text{word-ubin.norm-eq-iff} [\text{of numeral } a \text{ numeral } b, \text{symmetric, folded word-numeral-alt}]$
for $a \ b$
lemmas *sbintr-num* =
 $\text{word-sbin.norm-eq-iff} [\text{of numeral } a \text{ numeral } b, \text{symmetric, folded word-numeral-alt}]$
for $a \ b$

lemma *num-of-bintr'*:
 $\text{bintrunc } (\text{len-of TYPE}('a :: \text{len0})) (\text{numeral } a) = (\text{numeral } b) \implies$
 $\text{numeral } a = (\text{numeral } b :: 'a \text{ word})$
 $\langle \text{proof} \rangle$

lemma *num-of-sbintr'*:
 $\text{sbintrunc } (\text{len-of TYPE}('a :: \text{len}) - 1) (\text{numeral } a) = (\text{numeral } b) \implies$
 $\text{numeral } a = (\text{numeral } b :: 'a \text{ word})$
 $\langle \text{proof} \rangle$

lemma *num-abs-bintr*:
 $(\text{numeral } x :: 'a \text{ word}) =$
 $\text{word-of-int } (\text{bintrunc } (\text{len-of TYPE}('a :: \text{len0})) (\text{numeral } x))$
 $\langle \text{proof} \rangle$

lemma *num-abs-sbintr*:
 $(\text{numeral } x :: 'a \text{ word}) =$
 $\text{word-of-int } (\text{sbintrunc } (\text{len-of TYPE}('a :: \text{len}) - 1) (\text{numeral } x))$
 $\langle \text{proof} \rangle$

lemma *ucast-id*: $ucast\ w = w$
 $\langle proof \rangle$

lemma *scast-id*: $scast\ w = w$
 $\langle proof \rangle$

lemma *ucast-bl*: $ucast\ w = of_bl\ (to_bl\ w)$
 $\langle proof \rangle$

lemma *nth-ucast*: $(ucast\ w :: 'a::len0\ word) !!\ n = (w !!\ n \wedge n < len_of\ TYPE('a))$
 $\langle proof \rangle$

lemma *ucast-bintr* [simp]:
 $ucast\ (numeral\ w :: 'a::len0\ word) =$
 $word_of_int\ (bintrunc\ (len_of\ TYPE('a))\ (numeral\ w))$
 $\langle proof \rangle$

lemma *scast-sbintr* [simp]:
 $scast\ (numeral\ w :: 'a::len\ word) =$
 $word_of_int\ (sbintrunc\ (len_of\ TYPE('a) - Suc\ 0)\ (numeral\ w))$
 $\langle proof \rangle$

lemma *source-size*: $source_size\ (c :: 'a::len0\ word \Rightarrow -) = len_of\ TYPE('a)$
 $\langle proof \rangle$

lemma *target-size*: $target_size\ (c :: - \Rightarrow 'b::len0\ word) = len_of\ TYPE('b)$
 $\langle proof \rangle$

lemma *is-down*: $is_down\ c \longleftrightarrow len_of\ TYPE('b) \leq len_of\ TYPE('a)$
for $c :: 'a::len0\ word \Rightarrow 'b::len0\ word$
 $\langle proof \rangle$

lemma *is-up*: $is_up\ c \longleftrightarrow len_of\ TYPE('a) \leq len_of\ TYPE('b)$
for $c :: 'a::len0\ word \Rightarrow 'b::len0\ word$
 $\langle proof \rangle$

lemmas *is-up-down* = *trans* [OF *is-up is-down* [symmetric]]

lemma *down-cast-same* [OF *refl*]: $uc = ucast \implies is_down\ uc \implies uc = scast$
 $\langle proof \rangle$

lemma *word-rev-tf*:
 $to_bl\ (of_bl\ bl :: 'a::len0\ word) =$
 $rev\ (takefill\ False\ (len_of\ TYPE('a))\ (rev\ bl))$

$\langle \text{proof} \rangle$

lemma *word-rep-drop*:

$\text{to-bl } (\text{of-bl } \text{bl} :: 'a :: \text{len0 } \text{word}) =$
 $\text{replicate } (\text{len-of TYPE('a)} - \text{length bl}) \text{ False } @$
 $\text{drop } (\text{length bl} - \text{len-of TYPE('a)}) \text{ bl}$
 $\langle \text{proof} \rangle$

lemma *to-bl-ucast*:

$\text{to-bl } (\text{ucast } (w :: 'b :: \text{len0 } \text{word}) :: 'a :: \text{len0 } \text{word}) =$
 $\text{replicate } (\text{len-of TYPE('a)} - \text{len-of TYPE('b)}) \text{ False } @$
 $\text{drop } (\text{len-of TYPE('b)} - \text{len-of TYPE('a)}) (\text{to-bl } w)$
 $\langle \text{proof} \rangle$

lemma *ucast-up-app* [*OF refl*]:

$uc = \text{ucast} \implies \text{source-size } uc + n = \text{target-size } uc \implies$
 $\text{to-bl } (uc \ w) = \text{replicate } n \text{ False } @ (\text{to-bl } w)$
 $\langle \text{proof} \rangle$

lemma *ucast-down-drop* [*OF refl*]:

$uc = \text{ucast} \implies \text{source-size } uc = \text{target-size } uc + n \implies$
 $\text{to-bl } (uc \ w) = \text{drop } n (\text{to-bl } w)$
 $\langle \text{proof} \rangle$

lemma *scast-down-drop* [*OF refl*]:

$sc = \text{scast} \implies \text{source-size } sc = \text{target-size } sc + n \implies$
 $\text{to-bl } (sc \ w) = \text{drop } n (\text{to-bl } w)$
 $\langle \text{proof} \rangle$

lemma *sint-up-scast* [*OF refl*]: $sc = \text{scast} \implies \text{is-up } sc \implies \text{sint } (sc \ w) = \text{sint } w$
 $\langle \text{proof} \rangle$

lemma *uint-up-ucast* [*OF refl*]: $uc = \text{ucast} \implies \text{is-up } uc \implies \text{uint } (uc \ w) = \text{uint } w$
 $\langle \text{proof} \rangle$

lemma *ucast-up-ucast* [*OF refl*]: $uc = \text{ucast} \implies \text{is-up } uc \implies \text{ucast } (uc \ w) = \text{ucast } w$
 $\langle \text{proof} \rangle$

lemma *scast-up-scast* [*OF refl*]: $sc = \text{scast} \implies \text{is-up } sc \implies \text{scast } (sc \ w) = \text{scast } w$
 $\langle \text{proof} \rangle$

lemma *ucast-of-bl-up* [*OF refl*]: $w = \text{of-bl } \text{bl} \implies \text{size } \text{bl} \leq \text{size } w \implies \text{ucast } w = \text{of-bl } \text{bl}$
 $\langle \text{proof} \rangle$

lemmas *ucast-up-ucast-id* = *trans* [*OF ucast-up-ucast ucast-id*]

lemmas *scast-up-scast-id* = *trans* [*OF scast-up-scast scast-id*]

lemmas *isduu* = *is-up-down* [**where** *c* = *ucast*, *THEN iffD2*]

lemmas *isdus* = *is-up-down* [**where** *c* = *scast*, *THEN iffD2*]

lemmas *ucast-down-ucast-id* = *isduu* [*THEN ucast-up-ucast-id*]

lemmas *scast-down-scast-id* = *isdus* [*THEN ucast-up-ucast-id*]

lemma *up-ucast-surj*:

is-up (*ucast* :: 'b::len0 word $\Rightarrow$ 'a::len0 word) $\Longrightarrow$

surj (*ucast* :: 'a word $\Rightarrow$ 'b word)

<proof>

lemma *up-scast-surj*:

is-up (*scast* :: 'b::len word $\Rightarrow$ 'a::len word) $\Longrightarrow$

surj (*scast* :: 'a word $\Rightarrow$ 'b word)

<proof>

lemma *down-scast-inj*:

is-down (*scast* :: 'b::len word $\Rightarrow$ 'a::len word) $\Longrightarrow$

inj-on (*ucast* :: 'a word $\Rightarrow$ 'b word) *A*

<proof>

lemma *down-ucast-inj*:

is-down (*ucast* :: 'b::len0 word $\Rightarrow$ 'a::len0 word) $\Longrightarrow$

inj-on (*ucast* :: 'a word $\Rightarrow$ 'b word) *A*

<proof>

lemma *of-bl-append-same*: *of-bl* (*X* @ *to-bl w*) = *w*

<proof>

lemma *ucast-down-wi* [*OF refl*]: *uc* = *ucast* $\Longrightarrow$ *is-down uc* $\Longrightarrow$ *uc* (*word-of-int x*) = *word-of-int x*

<proof>

lemma *ucast-down-no* [*OF refl*]: *uc* = *ucast* $\Longrightarrow$ *is-down uc* $\Longrightarrow$ *uc* (*numeral bin*) = *numeral bin*

<proof>

lemma *ucast-down-bl* [*OF refl*]: *uc* = *ucast* $\Longrightarrow$ *is-down uc* $\Longrightarrow$ *uc* (*of-bl bl*) = *of-bl bl*

<proof>

lemmas *slice-def'* = *slice-def* [*unfolded word-size*]

lemmas *test-bit-def'* = *word-test-bit-def* [*THEN fun-cong*]

lemmas *word-log-defs* = *word-and-def word-or-def word-xor-def word-not-def*

10.14 Word Arithmetic

lemma *word-less-alt*: $a < b \longleftrightarrow \text{uint } a < \text{uint } b$
 ⟨proof⟩

lemma *signed-linorder*: *class.linorder word-sle word-sless*
 ⟨proof⟩

interpretation *signed*: *linorder word-sle word-sless*
 ⟨proof⟩

lemma *udvdI*: $0 \leq n \implies \text{uint } b = n * \text{uint } a \implies a \text{ udvd } b$
 ⟨proof⟩

lemmas *word-div-no* [simp] = *word-div-def* [of numeral a numeral b] **for** a b
lemmas *word-mod-no* [simp] = *word-mod-def* [of numeral a numeral b] **for** a b
lemmas *word-less-no* [simp] = *word-less-def* [of numeral a numeral b] **for** a b
lemmas *word-le-no* [simp] = *word-le-def* [of numeral a numeral b] **for** a b
lemmas *word-sless-no* [simp] = *word-sless-def* [of numeral a numeral b] **for** a b
lemmas *word-sle-no* [simp] = *word-sle-def* [of numeral a numeral b] **for** a b

lemma *word-m1-wi*: $-1 = \text{word-of-int } (-1)$
 ⟨proof⟩

lemma *word-0-bl* [simp]: *of-bl* [] = 0
 ⟨proof⟩

lemma *word-1-bl*: *of-bl* [True] = 1
 ⟨proof⟩

lemma *uint-eq-0* [simp]: *uint* 0 = 0
 ⟨proof⟩

lemma *of-bl-0* [simp]: *of-bl* (replicate n False) = 0
 ⟨proof⟩

lemma *to-bl-0* [simp]: *to-bl* (0::'a::len0 word) = replicate (len-of TYPE('a)) False
 ⟨proof⟩

lemma *uint-0-iff*: $\text{uint } x = 0 \longleftrightarrow x = 0$
 ⟨proof⟩

lemma *unat-0-iff*: $\text{unat } x = 0 \longleftrightarrow x = 0$
 ⟨proof⟩

lemma *unat-0* [simp]: *unat* 0 = 0
 ⟨proof⟩

lemma *size-0-same'*: $\text{size } w = 0 \implies w = v$
for v w :: 'a::len0 word

$\langle \text{proof} \rangle$

lemmas *size-0-same* = *size-0-same'* [*unfolded word-size*]

lemmas *unat-eq-0* = *unat-0-iff*

lemmas *unat-eq-zero* = *unat-0-iff*

lemma *unat-gt-0*: $0 < \text{unat } x \longleftrightarrow x \neq 0$

$\langle \text{proof} \rangle$

lemma *ucast-0* [*simp*]: *ucast 0* = 0

$\langle \text{proof} \rangle$

lemma *sint-0* [*simp*]: *sint 0* = 0

$\langle \text{proof} \rangle$

lemma *scast-0* [*simp*]: *scast 0* = 0

$\langle \text{proof} \rangle$

lemma *sint-n1* [*simp*] : *sint* (− 1) = − 1

$\langle \text{proof} \rangle$

lemma *scast-n1* [*simp*]: *scast* (− 1) = − 1

$\langle \text{proof} \rangle$

lemma *uint-1* [*simp*]: *uint* (1::'a::len word) = 1

$\langle \text{proof} \rangle$

lemma *unat-1* [*simp*]: *unat* (1::'a::len word) = 1

$\langle \text{proof} \rangle$

lemma *ucast-1* [*simp*]: *ucast* (1::'a::len word) = 1

$\langle \text{proof} \rangle$

10.15 Transferring goals from words to ints

lemma *word-ths*:

shows *word-succ-p1*: *word-succ a* = *a* + 1

and *word-pred-m1*: *word-pred a* = *a* − 1

and *word-pred-succ*: *word-pred* (*word-succ a*) = *a*

and *word-succ-pred*: *word-succ* (*word-pred a*) = *a*

and *word-mult-succ*: *word-succ a* * *b* = *b* + *a* * *b*

$\langle \text{proof} \rangle$

lemma *uint-cong*: $x = y \implies \text{uint } x = \text{uint } y$

$\langle \text{proof} \rangle$

lemma *uint-word-ariths*:

fixes *a b* :: 'a::len0 word

shows $\text{uint } (a + b) = (\text{uint } a + \text{uint } b) \bmod 2^{\text{len-of TYPE('a::len0)}}$
and $\text{uint } (a - b) = (\text{uint } a - \text{uint } b) \bmod 2^{\text{len-of TYPE('a)}}$
and $\text{uint } (a * b) = \text{uint } a * \text{uint } b \bmod 2^{\text{len-of TYPE('a)}}$
and $\text{uint } (-a) = -\text{uint } a \bmod 2^{\text{len-of TYPE('a)}}$
and $\text{uint } (\text{word-succ } a) = (\text{uint } a + 1) \bmod 2^{\text{len-of TYPE('a)}}$
and $\text{uint } (\text{word-pred } a) = (\text{uint } a - 1) \bmod 2^{\text{len-of TYPE('a)}}$
and $\text{uint } (0 :: 'a \text{ word}) = 0 \bmod 2^{\text{len-of TYPE('a)}}$
and $\text{uint } (1 :: 'a \text{ word}) = 1 \bmod 2^{\text{len-of TYPE('a)}}$
 <proof>

lemma *uint-word-arith-bintrs*:

fixes $a \ b :: 'a::\text{len0 word}$
shows $\text{uint } (a + b) = \text{bintrunc } (\text{len-of TYPE('a)}) (\text{uint } a + \text{uint } b)$
and $\text{uint } (a - b) = \text{bintrunc } (\text{len-of TYPE('a)}) (\text{uint } a - \text{uint } b)$
and $\text{uint } (a * b) = \text{bintrunc } (\text{len-of TYPE('a)}) (\text{uint } a * \text{uint } b)$
and $\text{uint } (-a) = \text{bintrunc } (\text{len-of TYPE('a)}) (-\text{uint } a)$
and $\text{uint } (\text{word-succ } a) = \text{bintrunc } (\text{len-of TYPE('a)}) (\text{uint } a + 1)$
and $\text{uint } (\text{word-pred } a) = \text{bintrunc } (\text{len-of TYPE('a)}) (\text{uint } a - 1)$
and $\text{uint } (0 :: 'a \text{ word}) = \text{bintrunc } (\text{len-of TYPE('a)}) 0$
and $\text{uint } (1 :: 'a \text{ word}) = \text{bintrunc } (\text{len-of TYPE('a)}) 1$
 <proof>

lemma *sint-word-ariths*:

fixes $a \ b :: 'a::\text{len word}$
shows $\text{sint } (a + b) = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) (\text{sint } a + \text{sint } b)$
and $\text{sint } (a - b) = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) (\text{sint } a - \text{sint } b)$
and $\text{sint } (a * b) = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) (\text{sint } a * \text{sint } b)$
and $\text{sint } (-a) = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) (-\text{sint } a)$
and $\text{sint } (\text{word-succ } a) = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) (\text{sint } a + 1)$
and $\text{sint } (\text{word-pred } a) = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) (\text{sint } a - 1)$
and $\text{sint } (0 :: 'a \text{ word}) = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) 0$
and $\text{sint } (1 :: 'a \text{ word}) = \text{sbintrunc } (\text{len-of TYPE('a)} - 1) 1$
 <proof>

lemmas *uint-div-alt* = *word-div-def* [THEN trans [OF uint-cong int-word-uint]]

lemmas *uint-mod-alt* = *word-mod-def* [THEN trans [OF uint-cong int-word-uint]]

lemma *word-pred-0-n1*: $\text{word-pred } 0 = \text{word-of-int } (-1)$

<proof>

lemma *succ-pred-no [simp]*:

$\text{word-succ } (\text{numeral } w) = \text{numeral } w + 1$
 $\text{word-pred } (\text{numeral } w) = \text{numeral } w - 1$
 $\text{word-succ } (-\text{numeral } w) = -\text{numeral } w + 1$
 $\text{word-pred } (-\text{numeral } w) = -\text{numeral } w - 1$

<proof>

lemma *word-sp-01 [simp]*:

$\text{word-succ } (-1) = 0 \wedge \text{word-succ } 0 = 1 \wedge \text{word-pred } 0 = -1 \wedge \text{word-pred } 1 =$

0
 ⟨proof⟩

lemma *word-of-int-Ex*: $\exists y. x = \text{word-of-int } y$
 ⟨proof⟩

10.16 Order on fixed-length words

lemma *word-zero-le* [simp]: $0 \leq y$
 for $y :: 'a::\text{len0 word}$
 ⟨proof⟩

lemma *word-m1-ge* [simp]: $\text{word-pred } 0 \geq y$
 ⟨proof⟩

lemma *word-n1-ge* [simp]: $y \leq -1$
 for $y :: 'a::\text{len0 word}$
 ⟨proof⟩

lemmas *word-not-simps* [simp] =
word-zero-le [THEN *leD*] *word-m1-ge* [THEN *leD*] *word-n1-ge* [THEN *leD*]

lemma *word-gt-0*: $0 < y \longleftrightarrow 0 \neq y$
 for $y :: 'a::\text{len0 word}$
 ⟨proof⟩

lemmas *word-gt-0-no* [simp] = *word-gt-0* [of numeral y] for y

lemma *word-sless-alt*: $a <_s b \longleftrightarrow \text{sint } a < \text{sint } b$
 ⟨proof⟩

lemma *word-le-nat-alt*: $a \leq b \longleftrightarrow \text{unat } a \leq \text{unat } b$
 ⟨proof⟩

lemma *word-less-nat-alt*: $a < b \longleftrightarrow \text{unat } a < \text{unat } b$
 ⟨proof⟩

lemma *wi-less*:
 ($\text{word-of-int } n < (\text{word-of-int } m :: 'a::\text{len0 word})$) =
 ($n \bmod 2^\text{len-of TYPE('a)} < m \bmod 2^\text{len-of TYPE('a)}$)
 ⟨proof⟩

lemma *wi-le*:
 ($\text{word-of-int } n \leq (\text{word-of-int } m :: 'a::\text{len0 word})$) =
 ($n \bmod 2^\text{len-of TYPE('a)} \leq m \bmod 2^\text{len-of TYPE('a)}$)
 ⟨proof⟩

lemma *udvd-nat-alt*: $a \text{ udvd } b \longleftrightarrow (\exists n \geq 0. \text{unat } b = n * \text{unat } a)$

$\langle \text{proof} \rangle$

lemma *udvd-iff-dvd*: $x \text{ udvd } y \longleftrightarrow \text{unat } x \text{ dvd unat } y$
 $\langle \text{proof} \rangle$

lemmas *unat-mono* = *word-less-nat-alt* [THEN *iffD1*]

lemma *unat-minus-one*:
assumes $w \neq 0$
shows $\text{unat } (w - 1) = \text{unat } w - 1$
 $\langle \text{proof} \rangle$

lemma *measure-unat*: $p \neq 0 \implies \text{unat } (p - 1) < \text{unat } p$
 $\langle \text{proof} \rangle$

lemmas *uint-add-ge0* [simp] = *add-nonneg-nonneg* [OF *uint-ge-0 uint-ge-0*]
lemmas *uint-mult-ge0* [simp] = *mult-nonneg-nonneg* [OF *uint-ge-0 uint-ge-0*]

lemma *uint-sub-lt2p* [simp]: $\text{uint } x - \text{uint } y < 2 \wedge \text{len-of TYPE('a)}$
for $x :: 'a::\text{len0 word}$ **and** $y :: 'b::\text{len0 word}$
 $\langle \text{proof} \rangle$

10.17 Conditions for the addition (etc) of two words to overflow

lemma *uint-add-lem*:
 $(\text{uint } x + \text{uint } y < 2 \wedge \text{len-of TYPE('a)}) =$
 $(\text{uint } (x + y) = \text{uint } x + \text{uint } y)$
for $x \ y :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *uint-mult-lem*:
 $(\text{uint } x * \text{uint } y < 2 \wedge \text{len-of TYPE('a)}) =$
 $(\text{uint } (x * y) = \text{uint } x * \text{uint } y)$
for $x \ y :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *uint-sub-lem*: $\text{uint } x \geq \text{uint } y \longleftrightarrow \text{uint } (x - y) = \text{uint } x - \text{uint } y$
 $\langle \text{proof} \rangle$

lemma *uint-add-le*: $\text{uint } (x + y) \leq \text{uint } x + \text{uint } y$
 $\langle \text{proof} \rangle$

lemma *uint-sub-ge*: $\text{uint } (x - y) \geq \text{uint } x - \text{uint } y$
 $\langle \text{proof} \rangle$

lemma *mod-add-if-z*:
 $x < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$

for $x\ y\ z :: \text{int}$
 $\langle \text{proof} \rangle$

lemma *uint-plus-if'*:
 $\text{uint } (a + b) =$
 $(\text{if } \text{uint } a + \text{uint } b < 2 \wedge \text{len-of TYPE('a)} \text{ then } \text{uint } a + \text{uint } b$
 $\text{else } \text{uint } a + \text{uint } b - 2 \wedge \text{len-of TYPE('a)})$
for $a\ b :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *mod-sub-if-z*:
 $x < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x - y) \bmod z = (\text{if } y \leq x \text{ then } x - y \text{ else } x - y + z)$
for $x\ y\ z :: \text{int}$
 $\langle \text{proof} \rangle$

lemma *uint-sub-if'*:
 $\text{uint } (a - b) =$
 $(\text{if } \text{uint } b \leq \text{uint } a \text{ then } \text{uint } a - \text{uint } b$
 $\text{else } \text{uint } a - \text{uint } b + 2 \wedge \text{len-of TYPE('a)})$
for $a\ b :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

10.18 Definition of *uint-arith*

lemma *word-of-int-inverse*:
 $\text{word-of-int } r = a \implies 0 \leq r \implies r < 2 \wedge \text{len-of TYPE('a)} \implies \text{uint } a = r$
for $a :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *uint-split*:
 $P (\text{uint } x) = (\forall i. \text{word-of-int } i = x \wedge 0 \leq i \wedge i < 2 \wedge \text{len-of TYPE('a)} \longrightarrow P\ i)$
for $x :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *uint-split-asm*:
 $P (\text{uint } x) = (\nexists i. \text{word-of-int } i = x \wedge 0 \leq i \wedge i < 2 \wedge \text{len-of TYPE('a)} \wedge \neg P\ i)$
for $x :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemmas *uint-splits* = *uint-split uint-split-asm*

lemmas *uint-arith-simps* =
 $\text{word-le-def word-less-alt}$
 $\text{word-uint.Rep-inject [symmetric]}$
 $\text{uint-sub-if' uint-plus-if'}$

lemma *power-False-cong*: $\text{False} \implies a \wedge b = c \wedge d$

$\langle \text{proof} \rangle$

$\langle ML \rangle$

10.19 More on overflows and monotonicity

lemma *no-plus-overflow-uint-size*: $x \leq x + y \longleftrightarrow \text{uint } x + \text{uint } y < 2^{\text{size } x}$
for $x \ y :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemmas *no-olen-add* = *no-plus-overflow-uint-size* [*unfolded word-size*]

lemma *no-ulen-sub*: $x \geq x - y \longleftrightarrow \text{uint } y \leq \text{uint } x$
for $x \ y :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *no-olen-add'*: $x \leq y + x \longleftrightarrow \text{uint } y + \text{uint } x < 2^{\text{len-of TYPE('a)}}$
for $x \ y :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemmas *olen-add-equiv* = *trans* [*OF no-olen-add no-olen-add'* [*symmetric*]]

lemmas *uint-plus-simple-iff* = *trans* [*OF no-olen-add uint-add-lem*]
lemmas *uint-plus-simple* = *uint-plus-simple-iff* [*THEN iffD1*]
lemmas *uint-minus-simple-iff* = *trans* [*OF no-ulen-sub uint-sub-lem*]
lemmas *uint-minus-simple-alt* = *uint-sub-lem* [*folded word-le-def*]
lemmas *word-sub-le-iff* = *no-ulen-sub* [*folded word-le-def*]
lemmas *word-sub-le* = *word-sub-le-iff* [*THEN iffD2*]

lemma *word-less-sub1*: $x \neq 0 \implies 1 < x \longleftrightarrow 0 < x - 1$
for $x :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemma *word-le-sub1*: $x \neq 0 \implies 1 \leq x \longleftrightarrow 0 \leq x - 1$
for $x :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemma *sub-wrap-lt*: $x < x - z \longleftrightarrow x < z$
for $x \ z :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *sub-wrap*: $x \leq x - z \longleftrightarrow z = 0 \vee x < z$
for $x \ z :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *plus-minus-not-NULL-ab*: $x \leq ab - c \implies c \leq ab \implies c \neq 0 \implies x + c \neq 0$
for $x \ ab \ c :: 'a::\text{len0 word}$

$\langle \text{proof} \rangle$

lemma *plus-minus-no-overflow-ab*: $x \leq ab - c \implies c \leq ab \implies x \leq x + c$
for $x \ ab \ c :: 'a::len0 \ \text{word}$
 $\langle \text{proof} \rangle$

lemma *le-minus'*: $a + c \leq b \implies a \leq a + c \implies c \leq b - a$
for $a \ b \ c :: 'a::len0 \ \text{word}$
 $\langle \text{proof} \rangle$

lemma *le-plus'*: $a \leq b \implies c \leq b - a \implies a + c \leq b$
for $a \ b \ c :: 'a::len0 \ \text{word}$
 $\langle \text{proof} \rangle$

lemmas *le-plus = le-plus'* [rotated]

lemmas *le-minus = leD* [THEN *thin-rl*, THEN *le-minus'*]

lemma *word-plus-mono-right*: $y \leq z \implies x \leq x + z \implies x + y \leq x + z$
for $x \ y \ z :: 'a::len0 \ \text{word}$
 $\langle \text{proof} \rangle$

lemma *word-less-minus-cancel*: $y - x < z - x \implies x \leq z \implies y < z$
for $x \ y \ z :: 'a::len0 \ \text{word}$
 $\langle \text{proof} \rangle$

lemma *word-less-minus-mono-left*: $y < z \implies x \leq y \implies y - x < z - x$
for $x \ y \ z :: 'a::len0 \ \text{word}$
 $\langle \text{proof} \rangle$

lemma *word-less-minus-mono*: $a < c \implies d < b \implies a - b < a \implies c - d < c$
 $\implies a - b < c - d$
for $a \ b \ c \ d :: 'a::len \ \text{word}$
 $\langle \text{proof} \rangle$

lemma *word-le-minus-cancel*: $y - x \leq z - x \implies x \leq z \implies y \leq z$
for $x \ y \ z :: 'a::len0 \ \text{word}$
 $\langle \text{proof} \rangle$

lemma *word-le-minus-mono-left*: $y \leq z \implies x \leq y \implies y - x \leq z - x$
for $x \ y \ z :: 'a::len0 \ \text{word}$
 $\langle \text{proof} \rangle$

lemma *word-le-minus-mono*:
 $a \leq c \implies d \leq b \implies a - b \leq a \implies c - d \leq c \implies a - b \leq c - d$
for $a \ b \ c \ d :: 'a::len \ \text{word}$
 $\langle \text{proof} \rangle$

lemma *plus-le-left-cancel-wrap*: $x + y' < x \implies x + y < x \implies x + y' < x + y$

$\longleftrightarrow y' < y$
for $x\ y\ y' :: 'a::len0\ word$
 $\langle proof \rangle$

lemma *plus-le-left-cancel-nowrap*: $x \leq x + y' \implies x \leq x + y \implies x + y' < x + y \longleftrightarrow y' < y$
for $x\ y\ y' :: 'a::len0\ word$
 $\langle proof \rangle$

lemma *word-plus-mono-right2*: $a \leq a + b \implies c \leq b \implies a \leq a + c$
for $a\ b\ c :: 'a::len0\ word$
 $\langle proof \rangle$

lemma *word-less-add-right*: $x < y - z \implies z \leq y \implies x + z < y$
for $x\ y\ z :: 'a::len0\ word$
 $\langle proof \rangle$

lemma *word-less-sub-right*: $x < y + z \implies y \leq x \implies x - y < z$
for $x\ y\ z :: 'a::len0\ word$
 $\langle proof \rangle$

lemma *word-le-plus-either*: $x \leq y \vee x \leq z \implies y \leq y + z \implies x \leq y + z$
for $x\ y\ z :: 'a::len0\ word$
 $\langle proof \rangle$

lemma *word-less-nowrapI*: $x < z - k \implies k \leq z \implies 0 < k \implies x < x + k$
for $x\ z\ k :: 'a::len0\ word$
 $\langle proof \rangle$

lemma *inc-le*: $i < m \implies i + 1 \leq m$
for $i\ m :: 'a::len\ word$
 $\langle proof \rangle$

lemma *inc-i*: $1 \leq i \implies i < m \implies 1 \leq i + 1 \wedge i + 1 \leq m$
for $i\ m :: 'a::len\ word$
 $\langle proof \rangle$

lemma *udvd-incr-lem*:
 $up < uq \implies up = ua + n * uint\ K \implies$
 $uq = ua + n' * uint\ K \implies up + uint\ K \leq uq$
 $\langle proof \rangle$

lemma *udvd-incr'*:
 $p < q \implies uint\ p = ua + n * uint\ K \implies$
 $uint\ q = ua + n' * uint\ K \implies p + K \leq q$
 $\langle proof \rangle$

lemma *udvd-decr'*:
 $p < q \implies uint\ p = ua + n * uint\ K \implies$

$uint\ q = ua + n' * uint\ K \implies p \leq q - K$
 $\langle proof \rangle$

lemmas $udvd-incr-lem0 = udvd-incr-lem$ [where $ua=0$, unfolded $add-0-left$]

lemmas $udvd-incr0 = udvd-incr'$ [where $ua=0$, unfolded $add-0-left$]

lemmas $udvd-decr0 = udvd-decr'$ [where $ua=0$, unfolded $add-0-left$]

lemma $udvd-minus-le'$: $xy < k \implies z\ udvd\ xy \implies z\ udvd\ k \implies xy \leq k - z$
 $\langle proof \rangle$

lemma $udvd-incr2-K$:

$p < a + s \implies a \leq a + s \implies K\ udvd\ s \implies K\ udvd\ p - a \implies a \leq p \implies$
 $0 < K \implies p \leq p + K \wedge p + K \leq a + s$
 $\langle proof \rangle$

lemma $word-succ-rbl$: $to-bl\ w = bl \implies to-bl\ (word-succ\ w) = rev\ (rbl-succ\ (rev\ bl))$
 $\langle proof \rangle$

lemma $word-pred-rbl$: $to-bl\ w = bl \implies to-bl\ (word-pred\ w) = rev\ (rbl-pred\ (rev\ bl))$
 $\langle proof \rangle$

lemma $word-add-rbl$:

$to-bl\ v = vbl \implies to-bl\ w = wbl \implies$
 $to-bl\ (v + w) = rev\ (rbl-add\ (rev\ vbl)\ (rev\ wbl))$
 $\langle proof \rangle$

lemma $word-mult-rbl$:

$to-bl\ v = vbl \implies to-bl\ w = wbl \implies$
 $to-bl\ (v * w) = rev\ (rbl-mult\ (rev\ vbl)\ (rev\ wbl))$
 $\langle proof \rangle$

lemma $rtb-rbl-ariths$:

$rev\ (to-bl\ w) = ys \implies rev\ (to-bl\ (word-succ\ w)) = rbl-succ\ ys$
 $rev\ (to-bl\ w) = ys \implies rev\ (to-bl\ (word-pred\ w)) = rbl-pred\ ys$
 $rev\ (to-bl\ v) = ys \implies rev\ (to-bl\ w) = xs \implies rev\ (to-bl\ (v * w)) = rbl-mult\ ys\ xs$
 $rev\ (to-bl\ v) = ys \implies rev\ (to-bl\ w) = xs \implies rev\ (to-bl\ (v + w)) = rbl-add\ ys\ xs$
 $\langle proof \rangle$

10.20 Arithmetic type class instantiations

lemmas $word-le-0-iff$ [simp] =

$word-zero-le$ [THEN leD , THEN $linorder-antisym-conv1$]

lemma $word-of-int-nat$: $0 \leq x \implies word-of-int\ x = of-nat\ (nat\ x)$
 $\langle proof \rangle$

lemma *iszero-word-no* [*simp*]:

$$\text{iszero } (\text{numeral } \text{bin} :: 'a::\text{len word}) =$$

$$\text{iszero } (\text{bintrunc } (\text{len-of TYPE('a)}) (\text{numeral } \text{bin}))$$

$$\langle \text{proof} \rangle$$

Use *iszero* to simplify equalities between word numerals.

lemmas *word-eq-numeral-iff-iszero* [*simp*] =
eq-numeral-iff-iszero [**where** $'a = 'a::\text{len word}$]

10.21 Word and nat

lemma *td-ext-unat* [*OF refl*]:

$$n = \text{len-of TYPE('a::len)} \implies$$

$$\text{td-ext } (\text{unat} :: 'a \text{ word} \Rightarrow \text{nat}) \text{ of-nat } (\text{unats } n) (\lambda i. i \bmod 2^n)$$

$$\langle \text{proof} \rangle$$

lemmas *unat-of-nat* = *td-ext-unat* [*THEN td-ext.eq-norm*]

interpretation *word-unat*:

$$\text{td-ext}$$

$$\text{unat} :: 'a::\text{len word} \Rightarrow \text{nat}$$

$$\text{of-nat}$$

$$\text{unats } (\text{len-of TYPE('a::len)})$$

$$\lambda i. i \bmod 2^{\text{len-of TYPE('a::len)}}$$

$$\langle \text{proof} \rangle$$

lemmas *td-unat* = *word-unat.td-thm*

lemmas *unat-lt2p* [*iff*] = *word-unat.Rep* [*unfolded unats-def mem-Collect-eq*]

lemma *unat-le*: $y \leq \text{unat } z \implies y \in \text{unats } (\text{len-of TYPE('a)})$
for $z :: 'a::\text{len word}$

$$\langle \text{proof} \rangle$$

lemma *word-nchotomy*: $\forall w :: 'a::\text{len word}. \exists n. w = \text{of-nat } n \wedge n < 2^{\text{len-of TYPE('a)}}$

$$\langle \text{proof} \rangle$$

lemma *of-nat-eq*: $\text{of-nat } n = w \iff (\exists q. n = \text{unat } w + q * 2^{\text{len-of TYPE('a)}})$
for $w :: 'a::\text{len word}$

$$\langle \text{proof} \rangle$$

lemma *of-nat-eq-size*: $\text{of-nat } n = w \iff (\exists q. n = \text{unat } w + q * 2^{\text{size } w})$

$$\langle \text{proof} \rangle$$

lemma *of-nat-0*: $\text{of-nat } m = (0 :: 'a::\text{len word}) \iff (\exists q. m = q * 2^{\text{len-of TYPE('a)}})$

$\langle \text{proof} \rangle$

lemma *of-nat-2p [simp]*: $\text{of-nat } (2 \wedge \text{len-of TYPE('a)}) = (0 :: 'a :: \text{len word})$
 $\langle \text{proof} \rangle$

lemma *of-nat-gt-0*: $\text{of-nat } k \neq 0 \implies 0 < k$
 $\langle \text{proof} \rangle$

lemma *of-nat-neq-0*: $0 < k \implies k < 2 \wedge \text{len-of TYPE('a :: \text{len})} \implies \text{of-nat } k \neq (0 :: 'a \text{ word})$
 $\langle \text{proof} \rangle$

lemma *Abs-fnat-hom-add*: $\text{of-nat } a + \text{of-nat } b = \text{of-nat } (a + b)$
 $\langle \text{proof} \rangle$

lemma *Abs-fnat-hom-mult*: $\text{of-nat } a * \text{of-nat } b = (\text{of-nat } (a * b) :: 'a :: \text{len word})$
 $\langle \text{proof} \rangle$

lemma *Abs-fnat-hom-Suc*: $\text{word-succ } (\text{of-nat } a) = \text{of-nat } (\text{Suc } a)$
 $\langle \text{proof} \rangle$

lemma *Abs-fnat-hom-0*: $(0 :: 'a :: \text{len word}) = \text{of-nat } 0$
 $\langle \text{proof} \rangle$

lemma *Abs-fnat-hom-1*: $(1 :: 'a :: \text{len word}) = \text{of-nat } (\text{Suc } 0)$
 $\langle \text{proof} \rangle$

lemmas *Abs-fnat-homs* =
Abs-fnat-hom-add Abs-fnat-hom-mult Abs-fnat-hom-Suc
Abs-fnat-hom-0 Abs-fnat-hom-1

lemma *word-arith-nat-add*: $a + b = \text{of-nat } (\text{unat } a + \text{unat } b)$
 $\langle \text{proof} \rangle$

lemma *word-arith-nat-mult*: $a * b = \text{of-nat } (\text{unat } a * \text{unat } b)$
 $\langle \text{proof} \rangle$

lemma *word-arith-nat-Suc*: $\text{word-succ } a = \text{of-nat } (\text{Suc } (\text{unat } a))$
 $\langle \text{proof} \rangle$

lemma *word-arith-nat-div*: $a \text{ div } b = \text{of-nat } (\text{unat } a \text{ div } \text{unat } b)$
 $\langle \text{proof} \rangle$

lemma *word-arith-nat-mod*: $a \text{ mod } b = \text{of-nat } (\text{unat } a \text{ mod } \text{unat } b)$
 $\langle \text{proof} \rangle$

lemmas *word-arith-nat-defs* =
word-arith-nat-add word-arith-nat-mult
word-arith-nat-Suc Abs-fnat-hom-0

Abs-fnat-hom-1 word-arith-nat-div
word-arith-nat-mod

lemma *unat-cong*: $x = y \implies \text{unat } x = \text{unat } y$
 ⟨*proof*⟩

lemmas *unat-word-ariths* = *word-arith-nat-defs*
 [THEN *trans* [OF *unat-cong unat-of-nat*]]

lemmas *word-sub-less-iff* = *word-sub-le-iff*
 [unfolded *linorder-not-less* [symmetric] *Not-eq-iff*]

lemma *unat-add-lem*:
 $\text{unat } x + \text{unat } y < 2^{\text{len-of TYPE('a)}} \longleftrightarrow \text{unat } (x + y) = \text{unat } x + \text{unat } y$
for $x \ y :: 'a::\text{len word}$
 ⟨*proof*⟩

lemma *unat-mult-lem*:
 $\text{unat } x * \text{unat } y < 2^{\text{len-of TYPE('a)}} \longleftrightarrow \text{unat } (x * y) = \text{unat } x * \text{unat } y$
for $x \ y :: 'a::\text{len word}$
 ⟨*proof*⟩

lemmas *unat-plus-if'* =
trans [OF *unat-word-ariths(1) mod-nat-add, simplified*]

lemma *le-no-overflow*: $x \leq b \implies a \leq a + b \implies x \leq a + b$
for $a \ b \ x :: 'a::\text{len0 word}$
 ⟨*proof*⟩

lemmas *un-ui-le* =
trans [OF *word-le-nat-alt* [symmetric] *word-le-def*]

lemma *unat-sub-if-size*:
 $\text{unat } (x - y) =$
 (if $\text{unat } y \leq \text{unat } x$
 then $\text{unat } x - \text{unat } y$
 else $\text{unat } x + 2^{\text{size } x} - \text{unat } y$)
 ⟨*proof*⟩

lemmas *unat-sub-if'* = *unat-sub-if-size* [unfolded *word-size*]

lemma *unat-div*: $\text{unat } (x \text{ div } y) = \text{unat } x \text{ div } \text{unat } y$
for $x \ y :: 'a::\text{len word}$
 ⟨*proof*⟩

lemma *unat-mod*: $\text{unat } (x \text{ mod } y) = \text{unat } x \text{ mod } \text{unat } y$
for $x \ y :: 'a::\text{len word}$
 ⟨*proof*⟩

lemma *uint-div*: $\text{uint } (x \text{ div } y) = \text{uint } x \text{ div uint } y$
for $x \ y :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemma *uint-mod*: $\text{uint } (x \text{ mod } y) = \text{uint } x \text{ mod uint } y$
for $x \ y :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

10.22 Definition of *unat-arith* tactic

lemma *unat-split*: $P (\text{unat } x) \longleftrightarrow (\forall n. \text{of-nat } n = x \wedge n < 2^{\text{len-of TYPE('a)}} \longrightarrow P \ n)$
for $x :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemma *unat-split-asm*: $P (\text{unat } x) \longleftrightarrow (\nexists n. \text{of-nat } n = x \wedge n < 2^{\text{len-of TYPE('a)}} \wedge \neg P \ n)$
for $x :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemmas *of-nat-inverse* =
word-unat.Abs-inverse' [*rotated, unfolded unats-def, simplified*]

lemmas *unat-splits* = *unat-split unat-split-asm*

lemmas *unat-arith-simps* =
word-le-nat-alt word-less-nat-alt
word-unat.Rep-inject [*symmetric*]
unat-sub-if' unat-plus-if' unat-div unat-mod

$\langle \text{ML} \rangle$

lemma *no-plus-overflow-unat-size*: $x \leq x + y \longleftrightarrow \text{unat } x + \text{unat } y < 2^{\text{size } x}$
for $x \ y :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemmas *no-olen-add-nat* =
no-plus-overflow-unat-size [*unfolded word-size*]

lemmas *unat-plus-simple* =
trans [*OF no-olen-add-nat unat-add-lem*]

lemma *word-div-mult*: $0 < y \implies \text{unat } x * \text{unat } y < 2^{\text{len-of TYPE('a)}} \implies x * y \text{ div } y = x$
for $x \ y :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemma *div-lt'*: $i \leq k \text{ div } x \implies \text{unat } i * \text{unat } x < 2^{\text{len-of TYPE('a)}}$

for $i\ k\ x :: 'a::len\ word$
 $\langle proof \rangle$

lemmas $div\text{-}lt'' = order\text{-}less\text{-}imp\text{-}le\ [THEN\ div\text{-}lt']$

lemma $div\text{-}lt\text{-}mult: i < k\ div\ x \implies 0 < x \implies i * x < k$
for $i\ k\ x :: 'a::len\ word$
 $\langle proof \rangle$

lemma $div\text{-}le\text{-}mult: i \leq k\ div\ x \implies 0 < x \implies i * x \leq k$
for $i\ k\ x :: 'a::len\ word$
 $\langle proof \rangle$

lemma $div\text{-}lt\text{-}uint': i \leq k\ div\ x \implies uint\ i * uint\ x < 2^{\wedge len\text{-}of\ TYPE('a)}$
for $i\ k\ x :: 'a::len\ word$
 $\langle proof \rangle$

lemmas $div\text{-}lt\text{-}uint'' = order\text{-}less\text{-}imp\text{-}le\ [THEN\ div\text{-}lt\text{-}uint']$

lemma $word\text{-}le\text{-}exists': x \leq y \implies \exists z. y = x + z \wedge uint\ x + uint\ z < 2^{\wedge len\text{-}of\ TYPE('a)}$
for $x\ y\ z :: 'a::len0\ word$
 $\langle proof \rangle$

lemmas $plus\text{-}minus\text{-}not\text{-}NULL = order\text{-}less\text{-}imp\text{-}le\ [THEN\ plus\text{-}minus\text{-}not\text{-}NULL\text{-}ab]$

lemmas $plus\text{-}minus\text{-}no\text{-}overflow =$
 $order\text{-}less\text{-}imp\text{-}le\ [THEN\ plus\text{-}minus\text{-}no\text{-}overflow\text{-}ab]$

lemmas $mcs = word\text{-}less\text{-}minus\text{-}cancel\ word\text{-}less\text{-}minus\text{-}mono\text{-}left$
 $word\text{-}le\text{-}minus\text{-}cancel\ word\text{-}le\text{-}minus\text{-}mono\text{-}left$

lemmas $word\text{-}l\text{-}diffs = mcs\ [\mathbf{where}\ y = w + x, unfolded\ add\text{-}diff\text{-}cancel]\ \mathbf{for}\ w\ x$
lemmas $word\text{-}diff\text{-}ls = mcs\ [\mathbf{where}\ z = w + x, unfolded\ add\text{-}diff\text{-}cancel]\ \mathbf{for}\ w\ x$
lemmas $word\text{-}plus\text{-}mcs = word\text{-}diff\text{-}ls\ [\mathbf{where}\ y = v + x, unfolded\ add\text{-}diff\text{-}cancel]$
for $v\ x$

lemmas $le\text{-}unat\text{-}uoi = unat\text{-}le\ [THEN\ word\text{-}unat.Abs\text{-}inverse]$

lemmas $thd = refl\ [THEN\ [2]\ split\text{-}div\text{-}lemma\ [THEN\ iffD2], THEN\ conjunct1]$

lemmas $uno\text{-}simps\ [THEN\ le\text{-}unat\text{-}uoi] = mod\text{-}le\text{-}divisor\ div\text{-}le\text{-}dividend\ dtle$

lemma $word\text{-}mod\text{-}div\text{-}equality: (n\ div\ b) * b + (n\ mod\ b) = n$
for $n\ b :: 'a::len\ word$
 $\langle proof \rangle$

lemma $word\text{-}div\text{-}mult\text{-}le: a\ div\ b * b \leq a$
for $a\ b :: 'a::len\ word$

<proof>

lemma *word-mod-less-divisor*: $0 < n \implies m \bmod n < n$
for $m\ n :: 'a::len\ word$
<proof>

lemma *word-of-int-power-hom*: $word\text{-of-int}\ a \wedge n = (word\text{-of-int}\ (a \wedge n) :: 'a::len\ word)$
<proof>

lemma *word-arith-power-alt*: $a \wedge n = (word\text{-of-int}\ (uint\ a \wedge n) :: 'a::len\ word)$
<proof>

lemma *of-bl-length-less*:
 $length\ x = k \implies k < len\text{-of}\ TYPE('a) \implies (of\text{-bl}\ x :: 'a::len\ word) < 2 \wedge k$
<proof>

10.23 Cardinality, finiteness of set of words

instance *word* :: $(len0)\ finite$
<proof>

lemma *card-word*: $CARD('a::len0\ word) = 2 \wedge len\text{-of}\ TYPE('a)$
<proof>

lemma *card-word-size*: $card\ (UNIV :: 'a\ word\ set) = (2 \wedge size\ x)$
for $x :: 'a::len0\ word$
<proof>

10.24 Bitwise Operations on Words

lemmas *bin-log-bintrs* = *bin-trunc-not bin-trunc-xor bin-trunc-and bin-trunc-or*

lemmas *wils1* = *bin-log-bintrs* [THEN *word-ubin.norm-eq-iff* [THEN *iffD1*],
folded word-ubin.eq-norm, THEN eq-reflection]

lemmas *word-log-binary-defs* =
word-and-def word-or-def word-xor-def

lemma *word-wi-log-defs*:
 $NOT\ word\text{-of-int}\ a = word\text{-of-int}\ (NOT\ a)$
 $word\text{-of-int}\ a\ AND\ word\text{-of-int}\ b = word\text{-of-int}\ (a\ AND\ b)$
 $word\text{-of-int}\ a\ OR\ word\text{-of-int}\ b = word\text{-of-int}\ (a\ OR\ b)$
 $word\text{-of-int}\ a\ XOR\ word\text{-of-int}\ b = word\text{-of-int}\ (a\ XOR\ b)$
<proof>

lemma *word-no-log-defs* [simp]:

$NOT\ (numeral\ a) = word-of-int\ (NOT\ (numeral\ a))$
 $NOT\ (-\ numeral\ a) = word-of-int\ (NOT\ (-\ numeral\ a))$
 $numeral\ a\ AND\ numeral\ b = word-of-int\ (numeral\ a\ AND\ numeral\ b)$
 $numeral\ a\ AND\ -\ numeral\ b = word-of-int\ (numeral\ a\ AND\ -\ numeral\ b)$
 $-\ numeral\ a\ AND\ numeral\ b = word-of-int\ (-\ numeral\ a\ AND\ numeral\ b)$
 $-\ numeral\ a\ AND\ -\ numeral\ b = word-of-int\ (-\ numeral\ a\ AND\ -\ numeral\ b)$
 $numeral\ a\ OR\ numeral\ b = word-of-int\ (numeral\ a\ OR\ numeral\ b)$
 $numeral\ a\ OR\ -\ numeral\ b = word-of-int\ (numeral\ a\ OR\ -\ numeral\ b)$
 $-\ numeral\ a\ OR\ numeral\ b = word-of-int\ (-\ numeral\ a\ OR\ numeral\ b)$
 $-\ numeral\ a\ OR\ -\ numeral\ b = word-of-int\ (-\ numeral\ a\ OR\ -\ numeral\ b)$
 $numeral\ a\ XOR\ numeral\ b = word-of-int\ (numeral\ a\ XOR\ numeral\ b)$
 $numeral\ a\ XOR\ -\ numeral\ b = word-of-int\ (numeral\ a\ XOR\ -\ numeral\ b)$
 $-\ numeral\ a\ XOR\ numeral\ b = word-of-int\ (-\ numeral\ a\ XOR\ numeral\ b)$
 $-\ numeral\ a\ XOR\ -\ numeral\ b = word-of-int\ (-\ numeral\ a\ XOR\ -\ numeral\ b)$
 <proof>

Special cases for when one of the arguments equals 1.

lemma *word-bitwise-1-simps* [simp]:

$NOT\ (1::'a::len0\ word) = -2$
 $1\ AND\ numeral\ b = word-of-int\ (1\ AND\ numeral\ b)$
 $1\ AND\ -\ numeral\ b = word-of-int\ (1\ AND\ -\ numeral\ b)$
 $numeral\ a\ AND\ 1 = word-of-int\ (numeral\ a\ AND\ 1)$
 $-\ numeral\ a\ AND\ 1 = word-of-int\ (-\ numeral\ a\ AND\ 1)$
 $1\ OR\ numeral\ b = word-of-int\ (1\ OR\ numeral\ b)$
 $1\ OR\ -\ numeral\ b = word-of-int\ (1\ OR\ -\ numeral\ b)$
 $numeral\ a\ OR\ 1 = word-of-int\ (numeral\ a\ OR\ 1)$
 $-\ numeral\ a\ OR\ 1 = word-of-int\ (-\ numeral\ a\ OR\ 1)$
 $1\ XOR\ numeral\ b = word-of-int\ (1\ XOR\ numeral\ b)$
 $1\ XOR\ -\ numeral\ b = word-of-int\ (1\ XOR\ -\ numeral\ b)$
 $numeral\ a\ XOR\ 1 = word-of-int\ (numeral\ a\ XOR\ 1)$
 $-\ numeral\ a\ XOR\ 1 = word-of-int\ (-\ numeral\ a\ XOR\ 1)$
 <proof>

Special cases for when one of the arguments equals -1.

lemma *word-bitwise-m1-simps* [simp]:

$NOT\ (-1::'a::len0\ word) = 0$
 $(-1::'a::len0\ word)\ AND\ x = x$
 $x\ AND\ (-1::'a::len0\ word) = x$
 $(-1::'a::len0\ word)\ OR\ x = -1$
 $x\ OR\ (-1::'a::len0\ word) = -1$
 $(-1::'a::len0\ word)\ XOR\ x = NOT\ x$
 $x\ XOR\ (-1::'a::len0\ word) = NOT\ x$
 <proof>

lemma *uint-or*: $uint\ (x\ OR\ y) = uint\ x\ OR\ uint\ y$

<proof>

lemma *uint-and*: $\text{uint } (x \text{ AND } y) = \text{uint } x \text{ AND } \text{uint } y$
 $\langle \text{proof} \rangle$

lemma *test-bit-wi* [simp]:
 $(\text{word-of-int } x :: 'a::\text{len0 word}) !! n \longleftrightarrow n < \text{len-of TYPE}('a) \wedge \text{bin-nth } x \ n$
 $\langle \text{proof} \rangle$

lemma *word-test-bit-transfer* [transfer-rule]:
 $(\text{rel-fun pcr-word } (\text{rel-fun op} = \text{op} =))$
 $(\lambda x \ n. \ n < \text{len-of TYPE}('a) \wedge \text{bin-nth } x \ n) (\text{test-bit} :: 'a::\text{len0 word} \Rightarrow -)$
 $\langle \text{proof} \rangle$

lemma *word-ops-nth-size*:
 $n < \text{size } x \implies$
 $(x \text{ OR } y) !! n = (x !! n \mid y !! n) \wedge$
 $(x \text{ AND } y) !! n = (x !! n \wedge y !! n) \wedge$
 $(x \text{ XOR } y) !! n = (x !! n \neq y !! n) \wedge$
 $(\text{NOT } x) !! n = (\neg x !! n)$
for $x :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *word-ao-nth*:
 $(x \text{ OR } y) !! n = (x !! n \mid y !! n) \wedge$
 $(x \text{ AND } y) !! n = (x !! n \wedge y !! n)$
for $x :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *test-bit-numeral* [simp]:
 $(\text{numeral } w :: 'a::\text{len0 word}) !! n \longleftrightarrow$
 $n < \text{len-of TYPE}('a) \wedge \text{bin-nth } (\text{numeral } w) \ n$
 $\langle \text{proof} \rangle$

lemma *test-bit-neg-numeral* [simp]:
 $(- \text{ numeral } w :: 'a::\text{len0 word}) !! n \longleftrightarrow$
 $n < \text{len-of TYPE}('a) \wedge \text{bin-nth } (- \text{ numeral } w) \ n$
 $\langle \text{proof} \rangle$

lemma *test-bit-1* [simp]: $(1 :: 'a::\text{len word}) !! n \longleftrightarrow n = 0$
 $\langle \text{proof} \rangle$

lemma *nth-0* [simp]: $\neg (0 :: 'a::\text{len0 word}) !! n$
 $\langle \text{proof} \rangle$

lemma *nth-minus1* [simp]: $(-1 :: 'a::\text{len0 word}) !! n \longleftrightarrow n < \text{len-of TYPE}('a)$
 $\langle \text{proof} \rangle$

lemmas *bwsimps* =

wi-hom-add
word-wi-log-defs

lemma *word-bw-assocs*:

$(x \text{ AND } y) \text{ AND } z = x \text{ AND } y \text{ AND } z$
 $(x \text{ OR } y) \text{ OR } z = x \text{ OR } y \text{ OR } z$
 $(x \text{ XOR } y) \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$
for $x :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *word-bw-comms*:

$x \text{ AND } y = y \text{ AND } x$
 $x \text{ OR } y = y \text{ OR } x$
 $x \text{ XOR } y = y \text{ XOR } x$
for $x :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *word-bw-lcs*:

$y \text{ AND } x \text{ AND } z = x \text{ AND } y \text{ AND } z$
 $y \text{ OR } x \text{ OR } z = x \text{ OR } y \text{ OR } z$
 $y \text{ XOR } x \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$
for $x :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *word-log-esimps* [simp]:

$x \text{ AND } 0 = 0$
 $x \text{ AND } -1 = x$
 $x \text{ OR } 0 = x$
 $x \text{ OR } -1 = -1$
 $x \text{ XOR } 0 = x$
 $x \text{ XOR } -1 = \text{NOT } x$
 $0 \text{ AND } x = 0$
 $-1 \text{ AND } x = x$
 $0 \text{ OR } x = x$
 $-1 \text{ OR } x = -1$
 $0 \text{ XOR } x = x$
 $-1 \text{ XOR } x = \text{NOT } x$
for $x :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *word-not-dist*:

$\text{NOT } (x \text{ OR } y) = \text{NOT } x \text{ AND } \text{NOT } y$
 $\text{NOT } (x \text{ AND } y) = \text{NOT } x \text{ OR } \text{NOT } y$
for $x :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *word-bw-same*:

$x \text{ AND } x = x$
 $x \text{ OR } x = x$

$x \text{ XOR } x = 0$
for $x :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *word-ao-absorbs* [simp]:

$x \text{ AND } (y \text{ OR } x) = x$
 $x \text{ OR } y \text{ AND } x = x$
 $x \text{ AND } (x \text{ OR } y) = x$
 $y \text{ AND } x \text{ OR } x = x$
 $(y \text{ OR } x) \text{ AND } x = x$
 $x \text{ OR } x \text{ AND } y = x$
 $(x \text{ OR } y) \text{ AND } x = x$
 $x \text{ AND } y \text{ OR } x = x$
for $x :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *word-not-not* [simp]: $\text{NOT NOT } x = x$

for $x :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *word-ao-dist*: $(x \text{ OR } y) \text{ AND } z = x \text{ AND } z \text{ OR } y \text{ AND } z$

for $x :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *word-oa-dist*: $x \text{ AND } y \text{ OR } z = (x \text{ OR } z) \text{ AND } (y \text{ OR } z)$

for $x :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *word-add-not* [simp]: $x + \text{NOT } x = -1$

for $x :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *word-plus-and-or* [simp]: $(x \text{ AND } y) + (x \text{ OR } y) = x + y$

for $x :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *leoa*: $w = x \text{ OR } y \implies y = w \text{ AND } y$

for $x :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *leao*: $w' = x' \text{ AND } y' \implies x' = x' \text{ OR } w'$

for $x' :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *word-ao-equiv*: $w = w \text{ OR } w' \longleftrightarrow w' = w \text{ AND } w'$

for $w w' :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *le-word-or2*: $x \leq x \text{ OR } y$

for $x\ y :: 'a::len0\ word$
 $\langle proof \rangle$

lemmas $le-word-or1 = xtr3\ [OF\ word-bw-comms\ (2)\ le-word-or2]$
lemmas $word-and-le1 = xtr3\ [OF\ word-ao-absorbs\ (4)\ [symmetric]\ le-word-or2]$
lemmas $word-and-le2 = xtr3\ [OF\ word-ao-absorbs\ (8)\ [symmetric]\ le-word-or2]$

lemma $bl-word-not: to-bl\ (NOT\ w) = map\ Not\ (to-bl\ w)$
 $\langle proof \rangle$

lemma $bl-word-xor: to-bl\ (v\ XOR\ w) = map2\ op\ \neq\ (to-bl\ v)\ (to-bl\ w)$
 $\langle proof \rangle$

lemma $bl-word-or: to-bl\ (v\ OR\ w) = map2\ op\ \vee\ (to-bl\ v)\ (to-bl\ w)$
 $\langle proof \rangle$

lemma $bl-word-and: to-bl\ (v\ AND\ w) = map2\ op\ \wedge\ (to-bl\ v)\ (to-bl\ w)$
 $\langle proof \rangle$

lemma $word-lsb-alt: lsb\ w = test-bit\ w\ 0$
for $w :: 'a::len0\ word$
 $\langle proof \rangle$

lemma $word-lsb-1-0\ [simp]: lsb\ (1::'a::len\ word) \wedge \neg\ lsb\ (0::'b::len0\ word)$
 $\langle proof \rangle$

lemma $word-lsb-last: lsb\ w = last\ (to-bl\ w)$
for $w :: 'a::len\ word$
 $\langle proof \rangle$

lemma $word-lsb-int: lsb\ w \longleftrightarrow uint\ w\ mod\ 2 = 1$
 $\langle proof \rangle$

lemma $word-msb-sint: msb\ w \longleftrightarrow sint\ w < 0$
 $\langle proof \rangle$

lemma $msb-word-of-int: msb\ (word-of-int\ x::'a::len\ word) = bin-nth\ x\ (len-of\ TYPE('a) - 1)$
 $\langle proof \rangle$

lemma $word-msb-numeral\ [simp]:$
 $msb\ (numeral\ w::'a::len\ word) = bin-nth\ (numeral\ w)\ (len-of\ TYPE('a) - 1)$
 $\langle proof \rangle$

lemma $word-msb-neg-numeral\ [simp]:$
 $msb\ (-\ numeral\ w::'a::len\ word) = bin-nth\ (-\ numeral\ w)\ (len-of\ TYPE('a) - 1)$
 $\langle proof \rangle$

lemma *word-msb-0* [*simp*]: $\neg \text{msb } (0 :: 'a :: \text{len word})$
 ⟨*proof*⟩

lemma *word-msb-1* [*simp*]: $\text{msb } (1 :: 'a :: \text{len word}) \longleftrightarrow \text{len-of TYPE('a)} = 1$
 ⟨*proof*⟩

lemma *word-msb-nth*: $\text{msb } w = \text{bin-nth } (\text{uint } w) (\text{len-of TYPE('a)} - 1)$
for $w :: 'a :: \text{len word}$
 ⟨*proof*⟩

lemma *word-msb-alt*: $\text{msb } w = \text{hd } (\text{to-bl } w)$
for $w :: 'a :: \text{len word}$
 ⟨*proof*⟩

lemma *word-set-nth* [*simp*]: $\text{set-bit } w \ n (\text{test-bit } w \ n) = w$
for $w :: 'a :: \text{len0 word}$
 ⟨*proof*⟩

lemma *bin-nth-uint'*: $\text{bin-nth } (\text{uint } w) \ n \longleftrightarrow \text{rev } (\text{bin-to-bl } (\text{size } w) (\text{uint } w)) \ ! \ n$
 $\wedge \ n < \text{size } w$
 ⟨*proof*⟩

lemmas *bin-nth-uint* = *bin-nth-uint'* [*unfolded word-size*]

lemma *test-bit-bl*: $w \ ! \ n \longleftrightarrow \text{rev } (\text{to-bl } w) \ ! \ n \wedge \ n < \text{size } w$
 ⟨*proof*⟩

lemma *to-bl-nth*: $n < \text{size } w \implies \text{to-bl } w \ ! \ n = w \ ! \ (\text{size } w - \text{Suc } n)$
 ⟨*proof*⟩

lemma *test-bit-set*: $(\text{set-bit } w \ n \ x) \ ! \ n \longleftrightarrow n < \text{size } w \wedge x$
for $w :: 'a :: \text{len0 word}$
 ⟨*proof*⟩

lemma *test-bit-set-gen*:
 $\text{test-bit } (\text{set-bit } w \ n \ x) \ m = (\text{if } m = n \text{ then } n < \text{size } w \wedge x \text{ else } \text{test-bit } w \ m)$
for $w :: 'a :: \text{len0 word}$
 ⟨*proof*⟩

lemma *of-bl-rep-False*: $\text{of-bl } (\text{replicate } n \ \text{False} \ @ \ bs) = \text{of-bl } bs$
 ⟨*proof*⟩

lemma *msb-nth*: $\text{msb } w = w \ ! \ (\text{len-of TYPE('a)} - 1)$
for $w :: 'a :: \text{len word}$
 ⟨*proof*⟩

lemmas *msb0* = *len-gt-0* [*THEN diff-Suc-less*, *THEN word-ops-nth-size* [*unfolded word-size*]]

lemmas *msb1* = *msb0* [**where** $i = 0$]

lemmas *word-ops-msb* = *msb1* [*unfolded msb-nth* [*symmetric*, *unfolded One-nat-def*]]

lemmas *lsb0* = *len-gt-0* [*THEN word-ops-nth-size* [*unfolded word-size*]]

lemmas *word-ops-lsb* = *lsb0* [*unfolded word-lsb-alt*]

lemma *td-ext-nth* [*OF refl refl refl*, *unfolded word-size*]:

$n = \text{size } w \implies \text{ofn} = \text{set-bits} \implies [w, \text{ofn } g] = l \implies$
 $\text{td-ext test-bit ofn } \{f. \forall i. f \ i \longrightarrow i < n\} (\lambda h \ i. h \ i \wedge i < n)$
for $w :: 'a::\text{len0}$ *word*
 ⟨*proof*⟩

interpretation *test-bit*:

td-ext
 $\text{op } !! :: 'a::\text{len0}$ *word* $\Rightarrow$ *nat* $\Rightarrow$ *bool*
set-bits
 $\{f. \forall i. f \ i \longrightarrow i < \text{len-of TYPE}('a::\text{len0})\}$
 $(\lambda h \ i. h \ i \wedge i < \text{len-of TYPE}('a::\text{len0}))$
 ⟨*proof*⟩

lemmas *td-nth* = *test-bit.td-thm*

lemma *word-set-set-same* [*simp*]: *set-bit* (*set-bit* $w \ n \ x$) $n \ y$ = *set-bit* $w \ n \ y$

for $w :: 'a::\text{len0}$ *word*
 ⟨*proof*⟩

lemma *word-set-set-diff*:

fixes $w :: 'a::\text{len0}$ *word*
assumes $m \neq n$
shows *set-bit* (*set-bit* $w \ m \ x$) $n \ y$ = *set-bit* (*set-bit* $w \ n \ y$) $m \ x$
 ⟨*proof*⟩

lemma *nth-sint*:

fixes $w :: 'a::\text{len}$ *word*
defines $l \equiv \text{len-of TYPE}('a)$
shows *bin-nth* (*sint* w) n = (if $n < l - 1$ then $w \ !! \ n$ else $w \ !! \ (l - 1)$)
 ⟨*proof*⟩

lemma *word-lsb-numeral* [*simp*]:

lsb (*numeral* $\text{bin} :: 'a::\text{len}$ *word*) $\longleftrightarrow$ *bin-last* (*numeral* bin)
 ⟨*proof*⟩

lemma *word-lsb-neg-numeral* [*simp*]:

lsb ($-$ *numeral* $\text{bin} :: 'a::\text{len}$ *word*) $\longleftrightarrow$ *bin-last* ($-$ *numeral* bin)
 ⟨*proof*⟩

lemma *set-bit-word-of-int*: *set-bit* (*word-of-int* x) $n \ b$ = *word-of-int* (*bin-sc* $n \ b \ x$)

⟨*proof*⟩

lemma *word-set-numeral* [*simp*]:

set-bit (numeral bin::<'a::len0 word) n b =
word-of-int (bin-sc n b (numeral bin))
 ⟨proof⟩

lemma *word-set-neg-numeral* [simp]:
set-bit (− numeral bin::<'a::len0 word) n b =
word-of-int (bin-sc n b (− numeral bin))
 ⟨proof⟩

lemma *word-set-bit-0* [simp]: *set-bit* 0 n b = *word-of-int* (bin-sc n b 0)
 ⟨proof⟩

lemma *word-set-bit-1* [simp]: *set-bit* 1 n b = *word-of-int* (bin-sc n b 1)
 ⟨proof⟩

lemma *setBit-no* [simp]: *setBit* (numeral bin) n = *word-of-int* (bin-sc n True (numeral bin))
 ⟨proof⟩

lemma *clearBit-no* [simp]:
clearBit (numeral bin) n = *word-of-int* (bin-sc n False (numeral bin))
 ⟨proof⟩

lemma *to-bl-n1*: *to-bl* (−1::<'a::len0 word) = *replicate* (len-of TYPE('a)) True
 ⟨proof⟩

lemma *word-msb-n1* [simp]: *msb* (−1::<'a::len word)
 ⟨proof⟩

lemma *word-set-nth-iff*: *set-bit* w n b = w $\longleftrightarrow$ w !! n = b $\vee$ n $\geq$ size w
 for w :: 'a::len0 word
 ⟨proof⟩

lemma *test-bit-2p*: (word-of-int (2 ^ n)::'a::len word) !! m $\longleftrightarrow$ m = n $\wedge$ m < len-of TYPE('a)
 ⟨proof⟩

lemma *nth-w2p*: ((2::<'a::len word) ^ n) !! m $\longleftrightarrow$ m = n $\wedge$ m < len-of TYPE('a::len)
 ⟨proof⟩

lemma *uint-2p*: (0::<'a::len word) < 2 ^ n $\implies$ uint (2 ^ n::<'a::len word) = 2 ^ n
 ⟨proof⟩

lemma *word-of-int-2p*: (word-of-int (2 ^ n) :: 'a::len word) = 2 ^ n
 ⟨proof⟩

lemma *bang-is-le*: x !! m $\implies$ 2 ^ m $\leq$ x
 for x :: 'a::len word
 ⟨proof⟩

lemma *word-clr-le*: $w \geq \text{set-bit } w \ n \ \text{False}$
for $w :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *word-set-ge*: $w \leq \text{set-bit } w \ n \ \text{True}$
for $w :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

10.25 Shifting, Rotating, and Splitting Words

lemma *shiftrl-wi [simp]*: $\text{shiftrl } (\text{word-of-int } w) = \text{word-of-int } (w \ \text{BIT } \text{False})$
 $\langle \text{proof} \rangle$

lemma *shiftrl-numeral [simp]*: $\text{shiftrl } (\text{numeral } w) = \text{numeral } (\text{Num.Bit0 } w)$
 $\langle \text{proof} \rangle$

lemma *shiftrl-neg-numeral [simp]*: $\text{shiftrl } (- \text{numeral } w) = - \text{numeral } (\text{Num.Bit0 } w)$
 $\langle \text{proof} \rangle$

lemma *shiftrl-0 [simp]*: $\text{shiftrl } 0 = 0$
 $\langle \text{proof} \rangle$

lemma *shiftrl-def-u*: $\text{shiftrl } w = \text{word-of-int } (\text{uint } w \ \text{BIT } \text{False})$
 $\langle \text{proof} \rangle$

lemma *shiftrl-def-s*: $\text{shiftrl } w = \text{word-of-int } (\text{sint } w \ \text{BIT } \text{False})$
 $\langle \text{proof} \rangle$

lemma *shiftr1-0 [simp]*: $\text{shiftr1 } 0 = 0$
 $\langle \text{proof} \rangle$

lemma *sshiftr1-0 [simp]*: $\text{sshiftr1 } 0 = 0$
 $\langle \text{proof} \rangle$

lemma *sshiftr1-n1 [simp]*: $\text{sshiftr1 } (-1) = -1$
 $\langle \text{proof} \rangle$

lemma *shiftrl-0 [simp]*: $(0 :: 'a::\text{len0 word}) << n = 0$
 $\langle \text{proof} \rangle$

lemma *shiftr-0 [simp]*: $(0 :: 'a::\text{len0 word}) >> n = 0$
 $\langle \text{proof} \rangle$

lemma *sshiftr-0 [simp]*: $0 >>> n = 0$
 $\langle \text{proof} \rangle$

lemma *sshiftr-n1 [simp]*: $-1 >>> n = -1$

<proof>

lemma *nth-shiftl1*: *shiftl1 w !! n* $\longleftrightarrow$ *n* < *size w* $\wedge$ *n* > 0 $\wedge$ *w !! (n - 1)*
<proof>

lemma *nth-shiftl'*: (*w* << *m*) *!! n* $\longleftrightarrow$ *n* < *size w* $\wedge$ *n* >= *m* $\wedge$ *w !! (n - m)*
for *w* :: 'a::len0 word
<proof>

lemmas *nth-shiftl* = *nth-shiftl'* [*unfolded word-size*]

lemma *nth-shiftr1*: *shiftr1 w !! n* = *w !! Suc n*
<proof>

lemma *nth-shiftr*: (*w* >> *m*) *!! n* = *w !! (n + m)*
for *w* :: 'a::len0 word
<proof>

lemma *uint-shiftr1*: *uint (shiftr1 w)* = *bin-rest (uint w)*
<proof>

lemma *nth-sshiftr1*: *sshiftr1 w !! n* = (*if n* = *size w* - 1 *then w !! n* *else w !! Suc n*)
<proof>

lemma *nth-sshiftr* [*rule-format*] :
 $\forall n. \text{sshiftr } w \text{ } m \text{ } !! n =$
 $(n < \text{size } w \wedge (\text{if } n + m \geq \text{size } w \text{ then } w !! (\text{size } w - 1) \text{ else } w !! (n + m)))$
<proof>

lemma *shiftr1-div-2*: *uint (shiftr1 w)* = *uint w div 2*
<proof>

lemma *sshiftr1-div-2*: *sint (sshiftr1 w)* = *sint w div 2*
<proof>

lemma *shiftr-div-2n*: *uint (shiftr w n)* = *uint w div 2 ^ n*
<proof>

lemma *sshiftr-div-2n*: *sint (sshiftr w n)* = *sint w div 2 ^ n*
<proof>

10.25.1 shift functions in terms of lists of bools

lemmas *bshiftr1-numeral* [*simp*] =
bshiftr1-def [**where** *w=numeral w*, *unfolded to-bl-numeral*] **for** *w*

lemma *bshiftr1-bl*: $to\text{-}bl\ (bshiftr1\ b\ w) = b\ \# \text{butlast}\ (to\text{-}bl\ w)$
 $\langle proof \rangle$

lemma *shiftrl1-of-bl*: $shiftrl1\ (of\text{-}bl\ bl) = of\text{-}bl\ (bl\ @\ [False])$
 $\langle proof \rangle$

lemma *shiftrl1-bl*: $shiftrl1\ w = of\text{-}bl\ (to\text{-}bl\ w\ @\ [False])$
for $w :: 'a::len0\ word$
 $\langle proof \rangle$

lemma *bl-shiftrl1*: $to\text{-}bl\ (shiftrl1\ w) = tl\ (to\text{-}bl\ w)\ @\ [False]$
for $w :: 'a::len\ word$
 $\langle proof \rangle$

lemma *bl-shiftrl1'*: $to\text{-}bl\ (shiftrl1\ w) = tl\ (to\text{-}bl\ w\ @\ [False])$
 $\langle proof \rangle$

lemma *shiftr1-bl*: $shiftr1\ w = of\text{-}bl\ (\text{butlast}\ (to\text{-}bl\ w))$
 $\langle proof \rangle$

lemma *bl-shiftr1*: $to\text{-}bl\ (shiftr1\ w) = False\ \# \text{butlast}\ (to\text{-}bl\ w)$
for $w :: 'a::len\ word$
 $\langle proof \rangle$

lemma *bl-shiftr1'*: $to\text{-}bl\ (shiftr1\ w) = \text{butlast}\ (False\ \# \text{to}\text{-}bl\ w)$
 $\langle proof \rangle$

lemma *shiftrl1-rev*: $shiftrl1\ w = \text{word}\text{-}\text{reverse}\ (shiftr1\ (\text{word}\text{-}\text{reverse}\ w))$
 $\langle proof \rangle$

lemma *shiftrl-rev*: $shiftrl\ w\ n = \text{word}\text{-}\text{reverse}\ (shiftr\ (\text{word}\text{-}\text{reverse}\ w)\ n)$
 $\langle proof \rangle$

lemma *rev-shiftrl*: $\text{word}\text{-}\text{reverse}\ w\ <<\ n = \text{word}\text{-}\text{reverse}\ (w\ >>\ n)$
 $\langle proof \rangle$

lemma *shiftr-rev*: $w\ >>\ n = \text{word}\text{-}\text{reverse}\ (\text{word}\text{-}\text{reverse}\ w\ <<\ n)$
 $\langle proof \rangle$

lemma *rev-shiftr*: $\text{word}\text{-}\text{reverse}\ w\ >>\ n = \text{word}\text{-}\text{reverse}\ (w\ <<\ n)$
 $\langle proof \rangle$

lemma *bl-sshiftr1*: $to\text{-}bl\ (sshiftr1\ w) = hd\ (to\text{-}bl\ w)\ \# \text{butlast}\ (to\text{-}bl\ w)$
for $w :: 'a::len\ word$
 $\langle proof \rangle$

lemma *drop-shiftr*: $\text{drop}\ n\ (to\text{-}bl\ (w\ >>\ n)) = \text{take}\ (size\ w - n)\ (to\text{-}bl\ w)$

for $w :: 'a::len\ word$
 $\langle proof \rangle$

lemma *drop-sshiftr*: $drop\ n\ (to_bl\ (w\ >>>\ n)) = take\ (size\ w - n)\ (to_bl\ w)$
for $w :: 'a::len\ word$
 $\langle proof \rangle$

lemma *take-shiftr*: $n \leq size\ w \implies take\ n\ (to_bl\ (w\ >>\ n)) = replicate\ n\ False$
 $\langle proof \rangle$

lemma *take-sshiftr'* [*rule-format*] :
 $n \leq size\ w \longrightarrow hd\ (to_bl\ (w\ >>>\ n)) = hd\ (to_bl\ w) \wedge$
 $take\ n\ (to_bl\ (w\ >>>\ n)) = replicate\ n\ (hd\ (to_bl\ w))$
for $w :: 'a::len\ word$
 $\langle proof \rangle$

lemmas $hd_sshiftr = take_sshiftr'\ [THEN\ conjunct1]$
lemmas $take_sshiftr = take_sshiftr'\ [THEN\ conjunct2]$

lemma *atd-lem*: $take\ n\ xs = t \implies drop\ n\ xs = d \implies xs = t\ @\ d$
 $\langle proof \rangle$

lemmas $bl_shiftr = atd_lem\ [OF\ take_shiftr\ drop_shiftr]$
lemmas $bl_sshiftr = atd_lem\ [OF\ take_sshiftr\ drop_sshiftr]$

lemma *shiftl-of-bl*: $of_bl\ bl\ <<\ n = of_bl\ (bl\ @\ replicate\ n\ False)$
 $\langle proof \rangle$

lemma *shiftl-bl*: $w\ <<\ n = of_bl\ (to_bl\ w\ @\ replicate\ n\ False)$
for $w :: 'a::len0\ word$
 $\langle proof \rangle$

lemmas *shiftl-numeral* [*simp*] = *shiftl-def* [**where** $w=numeral\ w$] **for** w

lemma *bl-shiftl*: $to_bl\ (w\ <<\ n) = drop\ n\ (to_bl\ w)\ @\ replicate\ (min\ (size\ w)\ n)\ False$
 $\langle proof \rangle$

lemma *shiftl-zero-size*: $size\ x \leq n \implies x\ <<\ n = 0$
for $x :: 'a::len0\ word$
 $\langle proof \rangle$

lemma *shiftl1-2t*: $shiftl1\ w = 2 * w$
for $w :: 'a::len\ word$
 $\langle proof \rangle$

lemma *shiftl1-p*: $shiftl1\ w = w + w$

for $w :: 'a::len\ word$
 $\langle proof \rangle$

lemma *shiffl-t2n*: $shiffl\ w\ n = 2^{\wedge} n * w$
for $w :: 'a::len\ word$
 $\langle proof \rangle$

lemma *shiftr1-bintr* [simp]:
 $(shiftr1\ (numeral\ w) :: 'a::len0\ word) =$
 $word-of-int\ (bin-rest\ (bintrunc\ (len-of\ TYPE('a))\ (numeral\ w)))$
 $\langle proof \rangle$

lemma *sshiftr1-sbintr* [simp]:
 $(sshiftr1\ (numeral\ w) :: 'a::len\ word) =$
 $word-of-int\ (bin-rest\ (sbintrunc\ (len-of\ TYPE('a) - 1)\ (numeral\ w)))$
 $\langle proof \rangle$

lemma *shiftr-no* [simp]:
 $(numeral\ w :: 'a::len0\ word) >> n = word-of-int$
 $((bin-rest\ \wedge^{\wedge} n)\ (bintrunc\ (len-of\ TYPE('a))\ (numeral\ w)))$
 $\langle proof \rangle$

lemma *sshiftr-no* [simp]:
 $(numeral\ w :: 'a::len\ word) >>> n = word-of-int$
 $((bin-rest\ \wedge^{\wedge} n)\ (sbintrunc\ (len-of\ TYPE('a) - 1)\ (numeral\ w)))$
 $\langle proof \rangle$

lemma *shiftr1-bl-of*:
 $length\ bl \leq len-of\ TYPE('a) \implies$
 $shiftr1\ (of-bl\ bl :: 'a::len0\ word) = of-bl\ (butlast\ bl)$
 $\langle proof \rangle$

lemma *shiftr-bl-of*:
 $length\ bl \leq len-of\ TYPE('a) \implies$
 $(of-bl\ bl :: 'a::len0\ word) >> n = of-bl\ (take\ (length\ bl - n)\ bl)$
 $\langle proof \rangle$

lemma *shiftr-bl*: $x >> n \equiv of-bl\ (take\ (len-of\ TYPE('a) - n)\ (to-bl\ x))$
for $x :: 'a::len0\ word$
 $\langle proof \rangle$

lemma *msb-shift*: $msb\ w \longleftrightarrow w >> (len-of\ TYPE('a) - 1) \neq 0$
for $w :: 'a::len\ word$
 $\langle proof \rangle$

lemma *zip-replicate*: $n \geq length\ ys \implies zip\ (replicate\ n\ x)\ ys = map\ (\lambda y. (x, y))\ ys$

$\langle \text{proof} \rangle$

lemma *align-lem-or* [rule-format] :

$\forall x m. \text{length } x = n + m \longrightarrow \text{length } y = n + m \longrightarrow$
 $\text{drop } m \ x = \text{replicate } n \ \text{False} \longrightarrow \text{take } m \ y = \text{replicate } m \ \text{False} \longrightarrow$
 $\text{map2 } op \mid x \ y = \text{take } m \ x \ @ \ \text{drop } m \ y$
 $\langle \text{proof} \rangle$

lemma *align-lem-and* [rule-format] :

$\forall x m. \text{length } x = n + m \longrightarrow \text{length } y = n + m \longrightarrow$
 $\text{drop } m \ x = \text{replicate } n \ \text{False} \longrightarrow \text{take } m \ y = \text{replicate } m \ \text{False} \longrightarrow$
 $\text{map2 } op \wedge x \ y = \text{replicate } (n + m) \ \text{False}$
 $\langle \text{proof} \rangle$

lemma *aligned-bl-add-size* [OF refl]:

$\text{size } x - n = m \implies n \leq \text{size } x \implies \text{drop } m \ (\text{to-bl } x) = \text{replicate } n \ \text{False} \implies$
 $\text{take } m \ (\text{to-bl } y) = \text{replicate } m \ \text{False} \implies$
 $\text{to-bl } (x + y) = \text{take } m \ (\text{to-bl } x) \ @ \ \text{drop } m \ (\text{to-bl } y)$
 $\langle \text{proof} \rangle$

10.25.2 Mask

lemma *nth-mask* [OF refl, simp]: $m = \text{mask } n \implies \text{test-bit } m \ i \longleftrightarrow i < n \wedge i < \text{size } m$
 $\langle \text{proof} \rangle$

lemma *mask-bl*: $\text{mask } n = \text{of-bl } (\text{replicate } n \ \text{True})$
 $\langle \text{proof} \rangle$

lemma *mask-bin*: $\text{mask } n = \text{word-of-int } (\text{bintrunc } n \ (-1))$
 $\langle \text{proof} \rangle$

lemma *and-mask-bintr*: $w \ \text{AND} \ \text{mask } n = \text{word-of-int } (\text{bintrunc } n \ (\text{uint } w))$
 $\langle \text{proof} \rangle$

lemma *and-mask-wi*: $\text{word-of-int } i \ \text{AND} \ \text{mask } n = \text{word-of-int } (\text{bintrunc } n \ i)$
 $\langle \text{proof} \rangle$

lemma *and-mask-wi'*:

$\text{word-of-int } i \ \text{AND} \ \text{mask } n = (\text{word-of-int } (\text{bintrunc } (\min \ \text{LENGTH}('a) \ n) \ i) :: 'a::\text{len word})$
 $\langle \text{proof} \rangle$

lemma *and-mask-no*: $\text{numeral } i \ \text{AND} \ \text{mask } n = \text{word-of-int } (\text{bintrunc } n \ (\text{numeral } i))$
 $\langle \text{proof} \rangle$

lemma *bl-and-mask'*:

$\text{to-bl } (w \ \text{AND} \ \text{mask } n :: 'a::\text{len word}) =$

$\text{replicate } (\text{len-of } \text{TYPE}('a) - n) \text{ False } @$
 $\text{drop } (\text{len-of } \text{TYPE}('a) - n) (\text{to-bl } w)$
 $\langle \text{proof} \rangle$

lemma *and-mask-mod-2p*: $w \text{ AND mask } n = \text{word-of-int } (\text{uint } w \text{ mod } 2^n)$
 $\langle \text{proof} \rangle$

lemma *and-mask-lt-2p*: $\text{uint } (w \text{ AND mask } n) < 2^n$
 $\langle \text{proof} \rangle$

lemma *eq-mod-iff*: $0 < n \implies b = b \text{ mod } n \iff 0 \leq b \wedge b < n$
for $b \text{ n} :: \text{int}$
 $\langle \text{proof} \rangle$

lemma *mask-eq-iff*: $w \text{ AND mask } n = w \iff \text{uint } w < 2^n$
 $\langle \text{proof} \rangle$

lemma *and-mask-dvd*: $2^n \text{ dvd uint } w \iff w \text{ AND mask } n = 0$
 $\langle \text{proof} \rangle$

lemma *and-mask-dvd-nat*: $2^n \text{ dvd unat } w \iff w \text{ AND mask } n = 0$
 $\langle \text{proof} \rangle$

lemma *word-2p-lem*: $n < \text{size } w \implies w < 2^n = (\text{uint } w < 2^n)$
for $w :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemma *less-mask-eq*: $x < 2^n \implies x \text{ AND mask } n = x$
for $x :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemmas *mask-eq-iff-w2p* = *trans* [*OF* *mask-eq-iff* *word-2p-lem* [*symmetric*]]

lemmas *and-mask-less'* = *iffD2* [*OF* *word-2p-lem* *and-mask-lt-2p*, *simplified word-size*]

lemma *and-mask-less-size*: $n < \text{size } x \implies x \text{ AND mask } n < 2^n$
 $\langle \text{proof} \rangle$

lemma *word-mod-2p-is-mask* [*OF* *refl*]: $c = 2^n \implies c > 0 \implies x \text{ mod } c = x$
 $\text{AND mask } n$
for $c \text{ x} :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemma *mask-egs*:
 $(a \text{ AND mask } n) + b \text{ AND mask } n = a + b \text{ AND mask } n$
 $a + (b \text{ AND mask } n) \text{ AND mask } n = a + b \text{ AND mask } n$
 $(a \text{ AND mask } n) - b \text{ AND mask } n = a - b \text{ AND mask } n$
 $a - (b \text{ AND mask } n) \text{ AND mask } n = a - b \text{ AND mask } n$
 $a * (b \text{ AND mask } n) \text{ AND mask } n = a * b \text{ AND mask } n$

$(b \text{ AND } \text{mask } n) * a \text{ AND } \text{mask } n = b * a \text{ AND } \text{mask } n$
 $(a \text{ AND } \text{mask } n) + (b \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = a + b \text{ AND } \text{mask } n$
 $(a \text{ AND } \text{mask } n) - (b \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = a - b \text{ AND } \text{mask } n$
 $(a \text{ AND } \text{mask } n) * (b \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = a * b \text{ AND } \text{mask } n$
 $-(a \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = -a \text{ AND } \text{mask } n$
 $\text{word-succ } (a \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = \text{word-succ } a \text{ AND } \text{mask } n$
 $\text{word-pred } (a \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = \text{word-pred } a \text{ AND } \text{mask } n$
 $\langle \text{proof} \rangle$

lemma *mask-power-eq*: $(x \text{ AND } \text{mask } n) \wedge^k \text{ AND } \text{mask } n = x \wedge^k \text{ AND } \text{mask } n$
 $\langle \text{proof} \rangle$

10.25.3 Revcast

lemmas *revcast-def'* = *revcast-def* [*simplified*]
lemmas *revcast-def''* = *revcast-def'* [*simplified word-size*]
lemmas *revcast-no-def* [*simp*] = *revcast-def'* [**where** $w = \text{numeral } w$, *unfolded word-size*]
for w

lemma *to-bl-revcast*:
 $\text{to-bl } (\text{revcast } w :: 'a::\text{len0 word}) = \text{takefill False } (\text{len-of TYPE('a)}) (\text{to-bl } w)$
 $\langle \text{proof} \rangle$

lemma *revcast-rev-ucast* [*OF refl refl refl*]:
 $cs = [rc, uc] \implies rc = \text{revcast } (\text{word-reverse } w) \implies uc = \text{ucast } w \implies$
 $rc = \text{word-reverse } uc$
 $\langle \text{proof} \rangle$

lemma *revcast-ucast*: $\text{revcast } w = \text{word-reverse } (\text{ucast } (\text{word-reverse } w))$
 $\langle \text{proof} \rangle$

lemma *ucast-revcast*: $\text{ucast } w = \text{word-reverse } (\text{revcast } (\text{word-reverse } w))$
 $\langle \text{proof} \rangle$

lemma *ucast-rev-revcast*: $\text{ucast } (\text{word-reverse } w) = \text{word-reverse } (\text{revcast } w)$
 $\langle \text{proof} \rangle$

linking revcast and cast via shift

lemmas *wsst-TYs* = *source-size target-size word-size*

lemma *revcast-down-uu* [*OF refl*]:
 $rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies rc \text{ } w = \text{ucast } (w >> n)$
for $w :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemma *revcast-down-us* [*OF refl*]:
 $rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies rc \text{ } w = \text{ucast } (w >>> n)$
for $w :: 'a::\text{len word}$

<proof>

lemma *revcast-down-su* [*OF refl*]:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies rc \ w = \text{scast } (w \gg n)$

for $w :: 'a::\text{len word}$

<proof>

lemma *revcast-down-ss* [*OF refl*]:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies rc \ w = \text{scast } (w \ggg n)$

for $w :: 'a::\text{len word}$

<proof>

lemma *cast-down-rev*:

$uc = \text{ucast} \implies \text{source-size } uc = \text{target-size } uc + n \implies uc \ w = \text{revcast } (w \ll n)$

for $w :: 'a::\text{len word}$

<proof>

lemma *revcast-up* [*OF refl*]:

$rc = \text{revcast} \implies \text{source-size } rc + n = \text{target-size } rc \implies$

$rc \ w = (\text{ucast } w :: 'a::\text{len word}) \ll n$

<proof>

lemmas *rc1* = *revcast-up* [*THEN*

revcast-rev-ucast [*symmetric*, *THEN trans*, *THEN word-rev-gal*, *symmetric*]]

lemmas *rc2* = *revcast-down-uu* [*THEN*

revcast-rev-ucast [*symmetric*, *THEN trans*, *THEN word-rev-gal*, *symmetric*]]

lemmas *ucast-up* =

rc1 [*simplified rev-shiftr* [*symmetric*] *revcast-ucast* [*symmetric*]]

lemmas *ucast-down* =

rc2 [*simplified rev-shiftr* *revcast-ucast* [*symmetric*]]

10.25.4 Slices

lemma *slice1-no-bin* [*simp*]:

$\text{slice1 } n \ (\text{numeral } w :: 'b \ \text{word}) = \text{of-bl } (\text{takefill } \text{False } n \ (\text{bin-to-bl } (\text{len-of TYPE('b::len0))} \ (\text{numeral } w)))$

<proof>

lemma *slice-no-bin* [*simp*]:

$\text{slice } n \ (\text{numeral } w :: 'b \ \text{word}) = \text{of-bl } (\text{takefill } \text{False } (\text{len-of TYPE('b::len0)} - n) \ (\text{bin-to-bl } (\text{len-of TYPE('b::len0))} \ (\text{numeral } w)))$

<proof>

lemma *slice1-0* [*simp*] : $\text{slice1 } n \ 0 = 0$

<proof>

lemma *slice-0* [*simp*] : *slice* *n* 0 = 0
<proof>

lemma *slice-take'*: *slice* *n* *w* = *of-bl* (*take* (*size* *w* - *n*) (*to-bl* *w*))
<proof>

lemmas *slice-take* = *slice-take'* [*unfolded word-size*]

— *shiftr* to a word of the same size is just *slice*, *slice* is just *shiftr* then *ucast*

lemmas *shiftr-slice* = *trans* [*OF shiftr-bl* [*THEN meta-eq-to-obj-eq*] *slice-take* [*symmetric*]]

lemma *slice-shiftr*: *slice* *n* *w* = *ucast* (*w* >> *n*)
<proof>

lemma *nth-slice*: (*slice* *n* *w* :: 'a::len0 word) !! *m* = (*w* !! (*m* + *n*) ∧ *m* < *len-of* *TYPE*('a))
<proof>

lemma *slice1-down-alt'*:
sl = *slice1* *n* *w* ⇒ *fs* = *size* *sl* ⇒ *fs* + *k* = *n* ⇒
to-bl *sl* = *takefill* *False* *fs* (*drop* *k* (*to-bl* *w*))
<proof>

lemma *slice1-up-alt'*:
sl = *slice1* *n* *w* ⇒ *fs* = *size* *sl* ⇒ *fs* = *n* + *k* ⇒
to-bl *sl* = *takefill* *False* *fs* (*replicate* *k* *False* @ (*to-bl* *w*))
<proof>

lemmas *sd1* = *slice1-down-alt'* [*OF refl refl, unfolded word-size*]

lemmas *su1* = *slice1-up-alt'* [*OF refl refl, unfolded word-size*]

lemmas *slice1-down-alt* = *le-add-diff-inverse* [*THEN sd1*]

lemmas *slice1-up-alt* =
le-add-diff-inverse [*symmetric, THEN su1*]
le-add-diff-inverse2 [*symmetric, THEN su1*]

lemma *ucast-slice1*: *ucast* *w* = *slice1* (*size* *w*) *w*
<proof>

lemma *ucast-slice*: *ucast* *w* = *slice* 0 *w*
<proof>

lemma *slice-id*: *slice* 0 *t* = *t*
<proof>

lemma *revcast-slice1* [*OF refl*]: *rc* = *revcast* *w* ⇒ *slice1* (*size* *rc*) *w* = *rc*
<proof>

lemma *slice1-tf-tf'*:

to-bl (*slice1* *n* *w* :: 'a::len0 word) =
 rev (*takefill* *False* (*len-of* *TYPE*('a)) (*rev* (*takefill* *False* *n* (*to-bl* *w*))))
 ⟨*proof*⟩

lemmas *slice1-tf-tf* = *slice1-tf-tf'* [*THEN* *word-bl.Rep-inverse'*, *symmetric*]

lemma *rev-slice1*:

n + *k* = *len-of* *TYPE*('a) + *len-of* *TYPE*('b) $\implies$
 slice1 *n* (*word-reverse* *w* :: 'b::len0 word) =
 word-reverse (*slice1* *k* *w* :: 'a::len0 word)
 ⟨*proof*⟩

lemma *rev-slice*:

n + *k* + *len-of* *TYPE*('a::len0) = *len-of* *TYPE*('b::len0) $\implies$
 slice *n* (*word-reverse* (*w*::'b word)) = *word-reverse* (*slice* *k* *w* :: 'a word)
 ⟨*proof*⟩

lemmas *sym-notr* =

not-iff [*THEN* *iffD2*, *THEN* *not-sym*, *THEN* *not-iff* [*THEN* *iffD1*]]

— problem posed by TPHOLs referee: criterion for overflow of addition of signed integers

lemma *soft-test*:

(*sint* *x* + *sint* *y* = *sint* (*x* + *y*)) =
 (((*x* + *y*) *XOR* *x*) *AND* ((*x* + *y*) *XOR* *y*)) >> (*size* *x* - 1) = 0
 for *x* *y* :: 'a::len word
 ⟨*proof*⟩

10.26 Split and cat

lemmas *word-split-bin'* = *word-split-def*

lemmas *word-cat-bin'* = *word-cat-def*

lemma *word-rsplit-no*:

(*word-rsplit* (*numeral* *bin* :: 'b::len0 word) :: 'a word list) =
 map *word-of-int* (*bin-rsplit* (*len-of* *TYPE*('a::len))
 (*len-of* *TYPE*('b), *bintrunc* (*len-of* *TYPE*('b)) (*numeral* *bin*)))
 ⟨*proof*⟩

lemmas *word-rsplit-no-cl* [*simp*] = *word-rsplit-no*
 [*unfolded* *bin-rsplitl-def* *bin-rsplit-l* [*symmetric*]]

lemma *test-bit-cat*:

wc = *word-cat* *a* *b* $\implies$ *wc* !! *n* = (*n* < *size* *wc* $\wedge$
 (*if* *n* < *size* *b* then *b* !! *n* else *a* !! (*n* - *size* *b*)))
 ⟨*proof*⟩

lemma *word-cat-bl*: $\text{word-cat } a \ b = \text{of-bl } (\text{to-bl } a \ @ \ \text{to-bl } b)$
 $\langle \text{proof} \rangle$

lemma *of-bl-append*:
 $(\text{of-bl } (xs \ @ \ ys) :: 'a::\text{len word}) = \text{of-bl } xs * 2^{(\text{length } ys)} + \text{of-bl } ys$
 $\langle \text{proof} \rangle$

lemma *of-bl-False [simp]*: $\text{of-bl } (\text{False} \# xs) = \text{of-bl } xs$
 $\langle \text{proof} \rangle$

lemma *of-bl-True [simp]*: $(\text{of-bl } (\text{True} \# xs) :: 'a::\text{len word}) = 2^{\text{length } xs} + \text{of-bl } xs$
 $\langle \text{proof} \rangle$

lemma *of-bl-Cons*: $\text{of-bl } (x \# xs) = \text{of-bool } x * 2^{\text{length } xs} + \text{of-bl } xs$
 $\langle \text{proof} \rangle$

lemma *split-uint-lem*: $\text{bin-split } n \ (\text{uint } w) = (a, b) \implies$
 $a = \text{bintrunc } (\text{len-of TYPE('a)} - n) \ a \wedge b = \text{bintrunc } (\text{len-of TYPE('a)}) \ b$
for $w :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *word-split-bl'*:
 $\text{std} = \text{size } c - \text{size } b \implies (\text{word-split } c = (a, b)) \implies$
 $(a = \text{of-bl } (\text{take } \text{std } (\text{to-bl } c)) \wedge b = \text{of-bl } (\text{drop } \text{std } (\text{to-bl } c)))$
 $\langle \text{proof} \rangle$

lemma *word-split-bl*: $\text{std} = \text{size } c - \text{size } b \implies$
 $(a = \text{of-bl } (\text{take } \text{std } (\text{to-bl } c)) \wedge b = \text{of-bl } (\text{drop } \text{std } (\text{to-bl } c))) \longleftrightarrow$
 $\text{word-split } c = (a, b)$
 $\langle \text{proof} \rangle$

lemma *word-split-bl-eq*:
 $(\text{word-split } c :: ('c::\text{len0 word} \times 'd::\text{len0 word})) =$
 $(\text{of-bl } (\text{take } (\text{len-of TYPE('a)} - \text{len-of TYPE('d)})) (\text{to-bl } c)),$
 $\text{of-bl } (\text{drop } (\text{len-of TYPE('a)} - \text{len-of TYPE('d)}) (\text{to-bl } c)))$
for $c :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemma *test-bit-split'*:
 $\text{word-split } c = (a, b) \longrightarrow$
 $(\forall n \ m.$
 $b !! n = (n < \text{size } b \wedge c !! n) \wedge$
 $a !! m = (m < \text{size } a \wedge c !! (m + \text{size } b)))$
 $\langle \text{proof} \rangle$

lemma *test-bit-split*:
 $\text{word-split } c = (a, b) \implies$
 $(\forall n::\text{nat}. b !! n \longleftrightarrow n < \text{size } b \wedge c !! n) \wedge$
 $(\forall m::\text{nat}. a !! m \longleftrightarrow m < \text{size } a \wedge c !! (m + \text{size } b))$

$\langle \text{proof} \rangle$

lemma *test-bit-split-eq*:

$\text{word-split } c = (a, b) \longleftrightarrow$
 $((\forall n::\text{nat}. b !! n = (n < \text{size } b \wedge c !! n)) \wedge$
 $(\forall m::\text{nat}. a !! m = (m < \text{size } a \wedge c !! (m + \text{size } b))))$
 $\langle \text{proof} \rangle$

lemma *word-cat-id*: $\text{word-cat } a \ b = b$

$\langle \text{proof} \rangle$

lemma *word-cat-hom*:

$\text{len-of TYPE('a::len0)} \leq \text{len-of TYPE('b::len0)} + \text{len-of TYPE('c::len0)} \implies$
 $(\text{word-cat } (\text{word-of-int } w :: 'b \text{ word}) (b :: 'c \text{ word}) :: 'a \text{ word}) =$
 $\text{word-of-int } (\text{bin-cat } w (\text{size } b) (\text{uint } b))$
 $\langle \text{proof} \rangle$

lemma *word-cat-split-alt*: $\text{size } w \leq \text{size } u + \text{size } v \implies \text{word-split } w = (u, v) \implies$

$\text{word-cat } u \ v = w$

$\langle \text{proof} \rangle$

lemmas *word-cat-split-size* = *sym* [THEN [2] *word-cat-split-alt* [symmetric]]

10.26.1 Split and slice

lemma *split-slices*: $\text{word-split } w = (u, v) \implies u = \text{slice } (\text{size } v) \ w \wedge v = \text{slice } 0 \ w$

$\langle \text{proof} \rangle$

lemma *slice-cat1* [OF *refl*]:

$wc = \text{word-cat } a \ b \implies \text{size } wc \geq \text{size } a + \text{size } b \implies \text{slice } (\text{size } b) \ wc = a$
 $\langle \text{proof} \rangle$

lemmas *slice-cat2* = *trans* [OF *slice-id* *word-cat-id*]

lemma *cat-slices*:

$a = \text{slice } n \ c \implies b = \text{slice } 0 \ c \implies n = \text{size } b \implies$
 $\text{size } a + \text{size } b \geq \text{size } c \implies \text{word-cat } a \ b = c$
 $\langle \text{proof} \rangle$

lemma *word-split-cat-alt*:

$w = \text{word-cat } u \ v \implies \text{size } u + \text{size } v \leq \text{size } w \implies \text{word-split } w = (u, v)$
 $\langle \text{proof} \rangle$

lemmas *word-cat-bl-no-bin* [simp] =

word-cat-bl [where *a=numeral a* and *b=numeral b*, unfolded to-bl-numeral]
 for *a b*

lemmas *word-split-bl-no-bin* [simp] =

word-split-bl-eq [where *c=numeral c*, unfolded to-bl-numeral] for *c*

This odd result arises from the fact that the statement of the result implies

that the decoded words are of the same type, and therefore of the same length, as the original word.

lemma *word-rsplit-same*: $\text{word-rsplit } w = [w]$
 $\langle \text{proof} \rangle$

lemma *word-rsplit-empty-iff-size*: $\text{word-rsplit } w = [] \longleftrightarrow \text{size } w = 0$
 $\langle \text{proof} \rangle$

lemma *test-bit-rsplit*:
 $sw = \text{word-rsplit } w \implies m < \text{size } (hd \ sw) \implies$
 $k < \text{length } sw \implies (\text{rev } sw ! k) !! m = w !! (k * \text{size } (hd \ sw) + m)$
for $sw :: 'a::\text{len word list}$
 $\langle \text{proof} \rangle$

lemma *word-rcat-bl*: $\text{word-rcat } wl = \text{of-bl } (\text{concat } (\text{map } \text{to-bl } wl))$
 $\langle \text{proof} \rangle$

lemma *size-rcat-lem'*: $\text{size } (\text{concat } (\text{map } \text{to-bl } wl)) = \text{length } wl * \text{size } (hd \ wl)$
 $\langle \text{proof} \rangle$

lemmas *size-rcat-lem* = *size-rcat-lem'* [unfolded word-size]

lemmas *td-gal-lt-len* = *len-gt-0* [THEN *td-gal-lt*]

lemma *nth-rcat-lem*:
 $n < \text{length } (wl :: 'a \text{ word list}) * \text{len-of TYPE } ('a::\text{len}) \implies$
 $\text{rev } (\text{concat } (\text{map } \text{to-bl } wl)) ! n =$
 $\text{rev } (\text{to-bl } (\text{rev } wl ! (n \text{ div } \text{len-of TYPE } ('a)))) ! (n \text{ mod } \text{len-of TYPE } ('a))$
 $\langle \text{proof} \rangle$

lemma *test-bit-rcat*:
 $sw = \text{size } (hd \ wl) \implies rc = \text{word-rcat } wl \implies rc !! n =$
 $(n < \text{size } rc \wedge n \text{ div } sw < \text{size } wl \wedge (\text{rev } wl) ! (n \text{ div } sw) !! (n \text{ mod } sw))$
for $wl :: 'a::\text{len word list}$
 $\langle \text{proof} \rangle$

lemma *foldl-eq-foldr*: $\text{foldl } op + x \ xs = \text{foldr } op + (x \ \# \ xs) \ 0$
for $x :: 'a::\text{comm-monoid-add}$
 $\langle \text{proof} \rangle$

lemmas *test-bit-cong* = *arg-cong* [where $f = \text{test-bit}$, THEN *fun-cong*]

lemmas *test-bit-rsplit-alt* =
 $\text{trans } [OF \ \text{nth-rev-alt } [THEN \ \text{test-bit-cong}]]$
 $\text{test-bit-rsplit } [OF \ \text{refl asm-rl diff-Suc-less}]$

— lazy way of expressing that u and v , and su and sv , have same types

lemma *word-rsplit-len-indep* [OF *refl refl refl refl*]:
 $[u, v] = p \implies [su, sv] = q \implies \text{word-rsplit } u = su \implies$

$word_rsplit\ v = sv \implies length\ su = length\ sv$
 ⟨proof⟩

lemma *length-word-rsplit-size*:
 $n = len_of\ TYPE('a::len) \implies$
 $length\ (word_rsplit\ w :: 'a\ word\ list) \leq m \longleftrightarrow size\ w \leq m * n$
 ⟨proof⟩

lemmas *length-word-rsplit-lt-size* =
 $length_word_rsplit_size\ [unfolded\ Not_eq_iff\ linorder_not_less\ [symmetric]]$

lemma *length-word-rsplit-exp-size*:
 $n = len_of\ TYPE('a::len) \implies$
 $length\ (word_rsplit\ w :: 'a\ word\ list) = (size\ w + n - 1) \mathit{div}\ n$
 ⟨proof⟩

lemma *length-word-rsplit-even-size*:
 $n = len_of\ TYPE('a::len) \implies size\ w = m * n \implies$
 $length\ (word_rsplit\ w :: 'a\ word\ list) = m$
 ⟨proof⟩

lemmas *length-word-rsplit-exp-size'* = *refl* [THEN *length-word-rsplit-exp-size*]

lemmas *tdle* = *iffD2* [OF *split-div-lemma refl*, THEN *conjunct1*]
lemmas *dtle* = *xtr4* [OF *tdle mult.commute*]

lemma *word-rcat-rsplit*: $word_rcat\ (word_rsplit\ w) = w$
 ⟨proof⟩

lemma *size-word-rsplit-rcat-size*:
 $word_rcat\ ws = frcw \implies size\ frcw = length\ ws * len_of\ TYPE('a)$
 $\implies length\ (word_rsplit\ frcw :: 'a\ word\ list) = length\ ws$
for $ws :: 'a::len\ word\ list$ **and** $frcw :: 'b::len0\ word$
 ⟨proof⟩

lemma *msreus*:
 $0 < n \implies (k * n + m) \mathit{div}\ n = m \mathit{div}\ n + k$
 $(k * n + m) \mathit{mod}\ n = m \mathit{mod}\ n$
for $n :: nat$
 ⟨proof⟩

lemma *word-rsplit-rcat-size* [OF *refl*]:
 $word_rcat\ ws = frcw \implies$
 $size\ frcw = length\ ws * len_of\ TYPE('a) \implies word_rsplit\ frcw = ws$
for $ws :: 'a::len\ word\ list$
 ⟨proof⟩

10.27 Rotation

lemmas *rotater-0'* [simp] = *rotater-def* [where $n = 0$, simplified]

lemmas *word-rot-defs* = *word-roti-def* *word-rotr-def* *word-rotl-def*

lemma *rotate-eq-mod*: $m \bmod \text{length } xs = n \bmod \text{length } xs \implies \text{rotate } m \text{ } xs = \text{rotate } n \text{ } xs$
 ⟨proof⟩

lemmas *rotate-eqs* =
trans [OF *rotate0* [THEN *fun-cong*] *id-apply*]
rotate-rotate [symmetric]
rotate-id
rotate-conv-mod
rotate-eq-mod

10.27.1 Rotation of list to right

lemma *rotate1-rl'*: $\text{rotater1 } (l @ [a]) = a \# l$
 ⟨proof⟩

lemma *rotate1-rl* [simp]: $\text{rotater1 } (\text{rotate1 } l) = l$
 ⟨proof⟩

lemma *rotate1-lr* [simp]: $\text{rotate1 } (\text{rotater1 } l) = l$
 ⟨proof⟩

lemma *rotater1-rev'*: $\text{rotater1 } (\text{rev } xs) = \text{rev } (\text{rotate1 } xs)$
 ⟨proof⟩

lemma *rotater-rev'*: $\text{rotater } n (\text{rev } xs) = \text{rev } (\text{rotate } n \text{ } xs)$
 ⟨proof⟩

lemma *rotater-rev*: $\text{rotater } n \text{ } ys = \text{rev } (\text{rotate } n (\text{rev } ys))$
 ⟨proof⟩

lemma *rotater-drop-take*:
 $\text{rotater } n \text{ } xs =$
 $\text{drop } (\text{length } xs - n \bmod \text{length } xs) \text{ } xs @$
 $\text{take } (\text{length } xs - n \bmod \text{length } xs) \text{ } xs$
 ⟨proof⟩

lemma *rotater-Suc* [simp]: $\text{rotater } (\text{Suc } n) \text{ } xs = \text{rotater1 } (\text{rotater } n \text{ } xs)$
 ⟨proof⟩

lemma *rotate-inv-plus* [rule-format]:
 $\forall k. k = m + n \longrightarrow \text{rotater } k (\text{rotate } n \text{ } xs) = \text{rotater } m \text{ } xs \wedge$
 $\text{rotate } k (\text{rotater } n \text{ } xs) = \text{rotate } m \text{ } xs \wedge$
 $\text{rotater } n (\text{rotate } k \text{ } xs) = \text{rotate } m \text{ } xs \wedge$

$\text{rotate } n \ (\text{rotater } k \ xs) = \text{rotater } m \ xs$
 $\langle \text{proof} \rangle$

lemmas $\text{rotate-inv-rel} = \text{le-add-diff-inverse2} \ [\text{symmetric}, \text{THEN rotate-inv-plus}]$

lemmas $\text{rotate-inv-eq} = \text{order-refl} \ [\text{THEN rotate-inv-rel}, \text{simplified}]$

lemmas $\text{rotate-lr} \ [\text{simp}] = \text{rotate-inv-eq} \ [\text{THEN conjunct1}]$

lemmas $\text{rotate-rl} \ [\text{simp}] = \text{rotate-inv-eq} \ [\text{THEN conjunct2}, \text{THEN conjunct1}]$

lemma $\text{rotate-gal}: \text{rotater } n \ xs = ys \longleftrightarrow \text{rotate } n \ ys = xs$
 $\langle \text{proof} \rangle$

lemma $\text{rotate-gal}': \text{ys} = \text{rotater } n \ xs \longleftrightarrow \text{rotate } n \ ys = xs$
 $\langle \text{proof} \rangle$

lemma $\text{length-rotater} \ [\text{simp}]: \text{length} \ (\text{rotater } n \ xs) = \text{length } xs$
 $\langle \text{proof} \rangle$

lemma $\text{restrict-to-left}: x = y \implies x = z \longleftrightarrow y = z$
 $\langle \text{proof} \rangle$

lemmas $\text{rrs0} = \text{rotate-eqs} \ [\text{THEN restrict-to-left},$
 $\text{simplified rotate-gal} \ [\text{symmetric}] \text{ rotate-gal}' \ [\text{symmetric}]]$

lemmas $\text{rrs1} = \text{rrs0} \ [\text{THEN refl} \ [\text{THEN rev-iffD1}]]$

lemmas $\text{rotater-eqs} = \text{rrs1} \ [\text{simplified length-rotater}]$

lemmas $\text{rotater-0} = \text{rotater-eqs} \ (1)$

lemmas $\text{rotater-add} = \text{rotater-eqs} \ (2)$

10.27.2 map, map2, commuting with rotate(r)

lemma $\text{butlast-map}: xs \neq [] \implies \text{butlast} \ (\text{map } f \ xs) = \text{map } f \ (\text{butlast } xs)$
 $\langle \text{proof} \rangle$

lemma $\text{rotater1-map}: \text{rotater1} \ (\text{map } f \ xs) = \text{map } f \ (\text{rotater1 } xs)$
 $\langle \text{proof} \rangle$

lemma $\text{rotater-map}: \text{rotater } n \ (\text{map } f \ xs) = \text{map } f \ (\text{rotater } n \ xs)$
 $\langle \text{proof} \rangle$

lemma $\text{but-last-zip} \ [\text{rule-format}] :$
 $\forall \text{ys}. \text{length } xs = \text{length } ys \longrightarrow xs \neq [] \longrightarrow$
 $\text{last} \ (\text{zip } xs \ ys) = (\text{last } xs, \text{last } ys) \wedge$
 $\text{butlast} \ (\text{zip } xs \ ys) = \text{zip} \ (\text{butlast } xs) \ (\text{butlast } ys)$
 $\langle \text{proof} \rangle$

lemma $\text{but-last-map2} \ [\text{rule-format}] :$
 $\forall \text{ys}. \text{length } xs = \text{length } ys \longrightarrow xs \neq [] \longrightarrow$
 $\text{last} \ (\text{map2 } f \ xs \ ys) = f \ (\text{last } xs) \ (\text{last } ys) \wedge$

$\text{butlast } (\text{map2 } f \text{ } xs \text{ } ys) = \text{map2 } f \text{ } (\text{butlast } xs) \text{ } (\text{butlast } ys)$
 $\langle \text{proof} \rangle$

lemma *rotater1-zip*:

$\text{length } xs = \text{length } ys \implies$
 $\text{rotater1 } (\text{zip } xs \text{ } ys) = \text{zip } (\text{rotater1 } xs) \text{ } (\text{rotater1 } ys)$
 $\langle \text{proof} \rangle$

lemma *rotater1-map2*:

$\text{length } xs = \text{length } ys \implies$
 $\text{rotater1 } (\text{map2 } f \text{ } xs \text{ } ys) = \text{map2 } f \text{ } (\text{rotater1 } xs) \text{ } (\text{rotater1 } ys)$
 $\langle \text{proof} \rangle$

lemmas *lrth* =

$\text{box-equals } [\text{OF } \text{asm-rl } \text{length-rotater } [\text{symmetric}]$
 $\text{length-rotater } [\text{symmetric}],$
 $\text{THEN } \text{rotater1-map2}]$

lemma *rotater-map2*:

$\text{length } xs = \text{length } ys \implies$
 $\text{rotater } n \text{ } (\text{map2 } f \text{ } xs \text{ } ys) = \text{map2 } f \text{ } (\text{rotater } n \text{ } xs) \text{ } (\text{rotater } n \text{ } ys)$
 $\langle \text{proof} \rangle$

lemma *rotate1-map2*:

$\text{length } xs = \text{length } ys \implies$
 $\text{rotate1 } (\text{map2 } f \text{ } xs \text{ } ys) = \text{map2 } f \text{ } (\text{rotate1 } xs) \text{ } (\text{rotate1 } ys)$
 $\langle \text{proof} \rangle$

lemmas *lth* = $\text{box-equals } [\text{OF } \text{asm-rl } \text{length-rotate } [\text{symmetric}]$
 $\text{length-rotate } [\text{symmetric}], \text{ THEN } \text{rotate1-map2}]$

lemma *rotate-map2*:

$\text{length } xs = \text{length } ys \implies$
 $\text{rotate } n \text{ } (\text{map2 } f \text{ } xs \text{ } ys) = \text{map2 } f \text{ } (\text{rotate } n \text{ } xs) \text{ } (\text{rotate } n \text{ } ys)$
 $\langle \text{proof} \rangle$

lemma *to-bl-rotl*: $\text{to-bl } (\text{word-rotl } n \text{ } w) = \text{rotate } n \text{ } (\text{to-bl } w)$
 $\langle \text{proof} \rangle$

lemmas *blrs0* = $\text{rotate-egs } [\text{THEN } \text{to-bl-rotl } [\text{THEN } \text{trans}]]$

lemmas *word-rotl-egs* =

$\text{blrs0 } [\text{simplified } \text{word-bl-Rep}' \text{ word-bl.Rep-inject } \text{to-bl-rotl } [\text{symmetric}]]$

lemma *to-bl-rotr*: $\text{to-bl } (\text{word-rotr } n \text{ } w) = \text{rotater } n \text{ } (\text{to-bl } w)$
 $\langle \text{proof} \rangle$

lemmas *brrs0* = $\text{rotater-egs } [\text{THEN } \text{to-bl-rotr } [\text{THEN } \text{trans}]]$

```

lemmas word-rotr-eqs =
  brrs0 [simplified word-bl-Rep' word-bl.Rep-inject to-bl-rotr [symmetric]]

declare word-rotr-eqs (1) [simp]
declare word-rotl-eqs (1) [simp]

lemma word-rot-rl [simp]: word-rotl k (word-rotr k v) = v
  and word-rot-lr [simp]: word-rotr k (word-rotl k v) = v
  ⟨proof⟩

lemma word-rot-gal: word-rotr n v = w  $\longleftrightarrow$  word-rotl n w = v
  and word-rot-gal': w = word-rotr n v  $\longleftrightarrow$  v = word-rotl n w
  ⟨proof⟩

lemma word-rotr-rev: word-rotr n w = word-reverse (word-rotl n (word-reverse
w))
  ⟨proof⟩

lemma word-roti-0 [simp]: word-roti 0 w = w
  ⟨proof⟩

lemmas abl-cong = arg-cong [where f = of-bl]

lemma word-roti-add: word-roti (m + n) w = word-roti m (word-roti n w)
  ⟨proof⟩

lemma word-roti-conv-mod': word-roti n w = word-roti (n mod int (size w)) w
  ⟨proof⟩

lemmas word-roti-conv-mod = word-roti-conv-mod' [unfolded word-size]

10.27.3 “Word rotation commutes with bit-wise operations

locale word-rotate
begin

lemmas word-rot-defs' = to-bl-rotl to-bl-rotr

lemmas blwl-syms [symmetric] = bl-word-not bl-word-and bl-word-or bl-word-xor

lemmas lbl-lbl = trans [OF word-bl-Rep' word-bl-Rep' [symmetric]]

lemmas ths-map2 [OF lbl-lbl] = rotate-map2 rotater-map2

lemmas ths-map [where xs = to-bl v] = rotate-map rotater-map for v

lemmas th1s [simplified word-rot-defs' [symmetric]] = ths-map2 ths-map

lemma word-rot-logs:

```

```

word-rotl n (NOT v) = NOT word-rotl n v
word-rotr n (NOT v) = NOT word-rotr n v
word-rotl n (x AND y) = word-rotl n x AND word-rotl n y
word-rotr n (x AND y) = word-rotr n x AND word-rotr n y
word-rotl n (x OR y) = word-rotl n x OR word-rotl n y
word-rotr n (x OR y) = word-rotr n x OR word-rotr n y
word-rotl n (x XOR y) = word-rotl n x XOR word-rotl n y
word-rotr n (x XOR y) = word-rotr n x XOR word-rotr n y
⟨proof⟩
end

```

lemmas *word-rot-logs* = *word-rotate.word-rot-logs*

lemmas *bl-word-rotl-dt* = *trans* [*OF to-bl-rotl rotate-drop-take*,
simplified word-bl-Rep']

lemmas *bl-word-rotr-dt* = *trans* [*OF to-bl-rotr rotater-drop-take*,
simplified word-bl-Rep']

lemma *bl-word-roti-dt'*:
 $n = \text{nat } ((- i) \bmod \text{int } (\text{size } (w :: 'a::\text{len word}))) \implies$
 $\text{to-bl } (\text{word-roti } i \ w) = \text{drop } n \ (\text{to-bl } w) @ \text{take } n \ (\text{to-bl } w)$
 ⟨proof⟩

lemmas *bl-word-roti-dt* = *bl-word-roti-dt'* [*unfolded word-size*]

lemmas *word-rotl-dt* = *bl-word-rotl-dt* [*THEN word-bl.Rep-inverse'* [*symmetric*]]
lemmas *word-rotr-dt* = *bl-word-rotr-dt* [*THEN word-bl.Rep-inverse'* [*symmetric*]]
lemmas *word-roti-dt* = *bl-word-roti-dt* [*THEN word-bl.Rep-inverse'* [*symmetric*]]

lemma *word-rotx-0* [*simp*] : *word-rotr i 0 = 0* $\wedge$ *word-rotl i 0 = 0*
 ⟨proof⟩

lemma *word-roti-0'* [*simp*] : *word-roti n 0 = 0*
 ⟨proof⟩

lemmas *word-rotr-dt-no-bin'* [*simp*] =
word-rotr-dt [**where** *w=numeral w*, *unfolded to-bl-numeral*] **for** *w*

lemmas *word-rotl-dt-no-bin'* [*simp*] =
word-rotl-dt [**where** *w=numeral w*, *unfolded to-bl-numeral*] **for** *w*

declare *word-roti-def* [*simp*]

10.28 Maximum machine word

lemma *word-int-cases*:

fixes $x :: 'a::len0 \text{ word}$
obtains n **where** $x = \text{word-of-int } n$ **and** $0 \leq n$ **and** $n < 2^{\text{len-of TYPE('a)}}$
 $\langle \text{proof} \rangle$

lemma *word-nat-cases* [*cases type: word*]:
fixes $x :: 'a::len \text{ word}$
obtains n **where** $x = \text{of-nat } n$ **and** $n < 2^{\text{len-of TYPE('a)}}$
 $\langle \text{proof} \rangle$

lemma *max-word-eq*: $(\text{max-word}::'a::len \text{ word}) = 2^{\text{len-of TYPE('a)}} - 1$
 $\langle \text{proof} \rangle$

lemma *max-word-max* [*simp,intro!*]: $n \leq \text{max-word}$
 $\langle \text{proof} \rangle$

lemma *word-of-int-2p-len*: $\text{word-of-int } (2^{\text{len-of TYPE('a)}}) = (0::'a::len0 \text{ word})$
 $\langle \text{proof} \rangle$

lemma *word-pow-0*: $(2::'a::len \text{ word})^{\text{len-of TYPE('a)}} = 0$
 $\langle \text{proof} \rangle$

lemma *max-word-wrap*: $x + 1 = 0 \implies x = \text{max-word}$
 $\langle \text{proof} \rangle$

lemma *max-word-minus*: $\text{max-word} = (-1::'a::len \text{ word})$
 $\langle \text{proof} \rangle$

lemma *max-word-bl* [*simp*]: $\text{to-bl } (\text{max-word}::'a::len \text{ word}) = \text{replicate } (\text{len-of TYPE('a)})$
True
 $\langle \text{proof} \rangle$

lemma *max-test-bit* [*simp*]: $(\text{max-word}::'a::len \text{ word}) !! n \longleftrightarrow n < \text{len-of TYPE('a)}$
 $\langle \text{proof} \rangle$

lemma *word-and-max* [*simp*]: $x \text{ AND } \text{max-word} = x$
 $\langle \text{proof} \rangle$

lemma *word-or-max* [*simp*]: $x \text{ OR } \text{max-word} = \text{max-word}$
 $\langle \text{proof} \rangle$

lemma *word-ao-dist2*: $x \text{ AND } (y \text{ OR } z) = x \text{ AND } y \text{ OR } x \text{ AND } z$
for $x \ y \ z :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *word-oa-dist2*: $x \text{ OR } y \text{ AND } z = (x \text{ OR } y) \text{ AND } (x \text{ OR } z)$
for $x \ y \ z :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *word-and-not* [*simp*]: $x \text{ AND NOT } x = 0$

for $x :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *word-or-not* [simp]: $x \text{ OR } NOT \ x = \text{max-word}$
 $\langle \text{proof} \rangle$

lemma *word-boolean*: *boolean* (*op AND*) (*op OR*) *bitNOT* 0 *max-word*
 $\langle \text{proof} \rangle$

interpretation *word-bool-alg*: *boolean op AND op OR bitNOT* 0 *max-word*
 $\langle \text{proof} \rangle$

lemma *word-xor-and-or*: $x \text{ XOR } y = x \text{ AND } NOT \ y \text{ OR } NOT \ x \text{ AND } y$
for $x \ y :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

interpretation *word-bool-alg*: *boolean-xor op AND op OR bitNOT* 0 *max-word op XOR*
 $\langle \text{proof} \rangle$

lemma *shiftr-x-0* [iff]: $x \gg 0 = x$
for $x :: 'a::len0 \text{ word}$
 $\langle \text{proof} \rangle$

lemma *shiftr-x-0* [simp]: $x \ll 0 = x$
for $x :: 'a::len \text{ word}$
 $\langle \text{proof} \rangle$

lemma *shiftr-1* [simp]: $(1::'a::len \text{ word}) \ll n = 2^n$
 $\langle \text{proof} \rangle$

lemma *uint-lt-0* [simp]: $\text{uint } x < 0 = \text{False}$
 $\langle \text{proof} \rangle$

lemma *shiftr-1-1* [simp]: $\text{shiftr1 } (1::'a::len \text{ word}) = 0$
 $\langle \text{proof} \rangle$

lemma *shiftr-1* [simp]: $(1::'a::len \text{ word}) \gg n = (\text{if } n = 0 \text{ then } 1 \text{ else } 0)$
 $\langle \text{proof} \rangle$

lemma *word-less-1* [simp]: $x < 1 \longleftrightarrow x = 0$
for $x :: 'a::len \text{ word}$
 $\langle \text{proof} \rangle$

lemma *to-bl-mask*:
 $\text{to-bl } (\text{mask } n :: 'a::len \text{ word}) =$
 $\text{replicate } (\text{len-of TYPE('a)} - n) \text{ False } @$
 $\text{replicate } (\min (\text{len-of TYPE('a)}) \ n) \text{ True}$
 $\langle \text{proof} \rangle$

lemma *map-replicate-True*:

$n = \text{length } xs \implies$
 $\text{map } (\lambda(x,y). x \wedge y) (\text{zip } xs (\text{replicate } n \text{ True})) = xs$
 $\langle \text{proof} \rangle$

lemma *map-replicate-False*:

$n = \text{length } xs \implies \text{map } (\lambda(x,y). x \wedge y)$
 $(\text{zip } xs (\text{replicate } n \text{ False})) = \text{replicate } n \text{ False}$
 $\langle \text{proof} \rangle$

lemma *bl-and-mask*:

fixes $w :: 'a::\text{len word}$
and $n :: \text{nat}$
defines $n' \equiv \text{len-of TYPE('a)} - n$
shows $\text{to-bl } (w \text{ AND mask } n) = \text{replicate } n' \text{ False @ drop } n' (\text{to-bl } w)$
 $\langle \text{proof} \rangle$

lemma *drop-rev-takefill*:

$\text{length } xs \leq n \implies$
 $\text{drop } (n - \text{length } xs) (\text{rev } (\text{takefill } \text{False } n (\text{rev } xs))) = xs$
 $\langle \text{proof} \rangle$

lemma *map-nth-0 [simp]*: $\text{map } (op !! (0 :: 'a::\text{len0 word})) xs = \text{replicate } (\text{length } xs)$
 False
 $\langle \text{proof} \rangle$

lemma *uint-plus-if-size*:

$\text{uint } (x + y) =$
 $(\text{if } \text{uint } x + \text{uint } y < 2^{\text{size } x}$
 $\text{then } \text{uint } x + \text{uint } y$
 $\text{else } \text{uint } x + \text{uint } y - 2^{\text{size } x})$
 $\langle \text{proof} \rangle$

lemma *unat-plus-if-size*:

$\text{unat } (x + y) =$
 $(\text{if } \text{unat } x + \text{unat } y < 2^{\text{size } x}$
 $\text{then } \text{unat } x + \text{unat } y$
 $\text{else } \text{unat } x + \text{unat } y - 2^{\text{size } x})$
for $x y :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemma *word-neq-0-conv*: $w \neq 0 \longleftrightarrow 0 < w$

for $w :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemma *max-lt*: $\text{unat } (\text{max } a b \text{ div } c) = \text{unat } (\text{max } a b) \text{ div } \text{unat } c$

for $c :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemma *uint-sub-if-size*:

uint ($x - y$) =
 (*if* *uint* $y \leq$ *uint* x
 then *uint* $x -$ *uint* y
 else *uint* $x -$ *uint* $y + 2^{\text{size } x}$)
 ⟨*proof*⟩

lemma *unat-sub*: $b \leq a \implies \text{unat } (a - b) = \text{unat } a - \text{unat } b$

⟨*proof*⟩

lemmas *word-less-sub1-numberof* [*simp*] = *word-less-sub1* [of numeral w] **for** w

lemmas *word-le-sub1-numberof* [*simp*] = *word-le-sub1* [of numeral w] **for** w

lemma *word-of-int-minus*: *word-of-int* ($2^{\text{len-of TYPE('a)}}$ - i) = (*word-of-int* ($-i$))::'*a*::len *word*)

⟨*proof*⟩

lemmas *word-of-int-inj* =

word-uint.Abs-inject [*unfolded uints-num, simplified*]

lemma *word-le-less-eq*: $x \leq y \longleftrightarrow x = y \vee x < y$

for $x \ y :: 'a :: \text{len } \text{word}$

⟨*proof*⟩

lemma *mod-plus-cong*:

fixes $b \ b' :: \text{int}$

assumes $1: b = b'$

and $2: x \bmod b' = x' \bmod b'$

and $3: y \bmod b' = y' \bmod b'$

and $4: x' + y' = z'$

shows $(x + y) \bmod b = z' \bmod b'$

⟨*proof*⟩

lemma *mod-minus-cong*:

fixes $b \ b' :: \text{int}$

assumes $b = b'$

and $x \bmod b' = x' \bmod b'$

and $y \bmod b' = y' \bmod b'$

and $x' - y' = z'$

shows $(x - y) \bmod b = z' \bmod b'$

⟨*proof*⟩

lemma *word-induct-less*: $P \ 0 \implies (\bigwedge n. n < m \implies P \ n \implies P \ (1 + n)) \implies P \ m$

for $P :: 'a :: \text{len } \text{word} \Rightarrow \text{bool}$

⟨*proof*⟩

lemma *word-induct*: $P \ 0 \implies (\bigwedge n. P \ n \implies P \ (1 + n)) \implies P \ m$

for $P :: 'a :: \text{len } \text{word} \Rightarrow \text{bool}$

$\langle \text{proof} \rangle$

lemma *word-induct2* [*induct type*]: $P\ 0 \implies (\bigwedge n. 1 + n \neq 0 \implies P\ n \implies P\ (1 + n)) \implies P\ n$
for $P :: 'b::\text{len word} \Rightarrow \text{bool}$
 $\langle \text{proof} \rangle$

10.29 Recursion combinator for words

definition *word-rec* :: $'a \Rightarrow ('b::\text{len word} \Rightarrow 'a \Rightarrow 'a) \Rightarrow 'b\ \text{word} \Rightarrow 'a$
where *word-rec* *forZero forSuc* $n = \text{rec-nat forZero (forSuc} \circ \text{of-nat) (unat } n)$

lemma *word-rec-0*: $\text{word-rec } z\ s\ 0 = z$
 $\langle \text{proof} \rangle$

lemma *word-rec-Suc*: $1 + n \neq 0 \implies \text{word-rec } z\ s\ (1 + n) = s\ n\ (\text{word-rec } z\ s\ n)$
for $n :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

lemma *word-rec-Pred*: $n \neq 0 \implies \text{word-rec } z\ s\ n = s\ (n - 1)\ (\text{word-rec } z\ s\ (n - 1))$
 $\langle \text{proof} \rangle$

lemma *word-rec-in*: $f\ (\text{word-rec } z\ (\lambda-. f)\ n) = \text{word-rec } (f\ z)\ (\lambda-. f)\ n$
 $\langle \text{proof} \rangle$

lemma *word-rec-in2*: $f\ n\ (\text{word-rec } z\ f\ n) = \text{word-rec } (f\ 0\ z)\ (f \circ \text{op} + 1)\ n$
 $\langle \text{proof} \rangle$

lemma *word-rec-twice*:
 $m \leq n \implies \text{word-rec } z\ f\ n = \text{word-rec } (\text{word-rec } z\ f\ (n - m))\ (f \circ \text{op} + (n - m))\ m$
 $\langle \text{proof} \rangle$

lemma *word-rec-id*: $\text{word-rec } z\ (\lambda-. \text{id})\ n = z$
 $\langle \text{proof} \rangle$

lemma *word-rec-id-eq*: $\forall m < n. f\ m = \text{id} \implies \text{word-rec } z\ f\ n = z$
 $\langle \text{proof} \rangle$

lemma *word-rec-max*:
 $\forall m \geq n. m \neq -1 \longrightarrow f\ m = \text{id} \implies \text{word-rec } z\ f\ (-1) = \text{word-rec } z\ f\ n$
 $\langle \text{proof} \rangle$

lemma *unatSuc*: $1 + n \neq 0 \implies \text{unat } (1 + n) = \text{Suc } (\text{unat } n)$
for $n :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

declare *bin-to-bl-def* [*simp*]

$\langle ML \rangle$

hide-const (**open**) *Word*

end

theory *WordBitwise*

imports *Word*

begin

Helper constants used in defining addition

definition *xor3* :: *bool* $\Rightarrow$ *bool* $\Rightarrow$ *bool* $\Rightarrow$ *bool*
where *xor3* *a b c* = (*a* = (*b* = *c*))

definition *carry* :: *bool* $\Rightarrow$ *bool* $\Rightarrow$ *bool* $\Rightarrow$ *bool*
where *carry* *a b c* = ((*a* $\wedge$ (*b* $\vee$ *c*)) $\vee$ (*b* $\wedge$ *c*))

lemma *carry-simps*:

carry *True a b* = (*a* $\vee$ *b*)
carry *a True b* = (*a* $\vee$ *b*)
carry *a b True* = (*a* $\vee$ *b*)
carry *False a b* = (*a* $\wedge$ *b*)
carry *a False b* = (*a* $\wedge$ *b*)
carry *a b False* = (*a* $\wedge$ *b*)
 $\langle proof \rangle$

lemma *xor3-simps*:

xor3 *True a b* = (*a* = *b*)
xor3 *a True b* = (*a* = *b*)
xor3 *a b True* = (*a* = *b*)
xor3 *False a b* = (*a* $\neq$ *b*)
xor3 *a False b* = (*a* $\neq$ *b*)
xor3 *a b False* = (*a* $\neq$ *b*)
 $\langle proof \rangle$

Breaking up word equalities into equalities on their bit lists. Equalities are generated and manipulated in the reverse order to *to-bl*.

lemma *word-eq-rbl-eq*: *x* = *y* $\longleftrightarrow$ *rev* (*to-bl* *x*) = *rev* (*to-bl* *y*)
 $\langle proof \rangle$

lemma *rbl-word-or*: *rev* (*to-bl* (*x OR y*)) = *map2* *op* $\vee$ (*rev* (*to-bl* *x*)) (*rev* (*to-bl* *y*))
 $\langle proof \rangle$

lemma *rbl-word-and*: *rev* (*to-bl* (*x AND y*)) = *map2* *op* $\wedge$ (*rev* (*to-bl* *x*)) (*rev* (*to-bl* *y*))
 $\langle proof \rangle$

lemma *rbl-word-xor*: $\text{rev } (\text{to-bl } (x \text{ XOR } y)) = \text{map2 } \text{op} \neq (\text{rev } (\text{to-bl } x)) (\text{rev } (\text{to-bl } y))$
 ⟨proof⟩

lemma *rbl-word-not*: $\text{rev } (\text{to-bl } (\text{NOT } x)) = \text{map } \text{Not } (\text{rev } (\text{to-bl } x))$
 ⟨proof⟩

lemma *bl-word-sub*: $\text{to-bl } (x - y) = \text{to-bl } (x + (- y))$
 ⟨proof⟩

lemma *rbl-word-1*: $\text{rev } (\text{to-bl } (1 :: 'a::\text{len0 word})) = \text{takefill } \text{False } (\text{len-of TYPE('a)})$
 [True]
 ⟨proof⟩

lemma *rbl-word-if*: $\text{rev } (\text{to-bl } (\text{if } P \text{ then } x \text{ else } y)) = \text{map2 } (\text{If } P) (\text{rev } (\text{to-bl } x))$
 ($\text{rev } (\text{to-bl } y)$)
 ⟨proof⟩

lemma *rbl-add-carry-Cons*:
 ($\text{if } \text{car} \text{ then } \text{rbl-succ} \text{ else } \text{id}$) ($\text{rbl-add } (x \# xs) (y \# ys)$) =
 $\text{xor3 } x \ y \ \text{car} \# (\text{if } \text{carry } x \ y \ \text{car} \text{ then } \text{rbl-succ} \text{ else } \text{id}) (\text{rbl-add } xs \ ys)$
 ⟨proof⟩

lemma *rbl-add-suc-carry-fold*:
 $\text{length } xs = \text{length } ys \implies$
 $\forall \text{car. } (\text{if } \text{car} \text{ then } \text{rbl-succ} \text{ else } \text{id}) (\text{rbl-add } xs \ ys) =$
 $(\text{foldr } (\lambda(x, y) \ \text{res } \text{car. } \text{xor3 } x \ y \ \text{car} \# \text{res } (\text{carry } x \ y \ \text{car})) (\text{zip } xs \ ys) (\lambda-. []))$
 car
 ⟨proof⟩

lemma *to-bl-plus-carry*:
 $\text{to-bl } (x + y) =$
 $\text{rev } (\text{foldr } (\lambda(x, y) \ \text{res } \text{car. } \text{xor3 } x \ y \ \text{car} \# \text{res } (\text{carry } x \ y \ \text{car}))$
 $(\text{rev } (\text{zip } (\text{to-bl } x) (\text{to-bl } y))) (\lambda-. []) \ \text{False})$
 ⟨proof⟩

definition *rbl-plus cin xs ys* =
 $\text{foldr } (\lambda(x, y) \ \text{res } \text{car. } \text{xor3 } x \ y \ \text{car} \# \text{res } (\text{carry } x \ y \ \text{car})) (\text{zip } xs \ ys) (\lambda-. []) \ \text{cin}$

lemma *rbl-plus-simps*:
 $\text{rbl-plus } \text{cin } (x \# xs) (y \# ys) = \text{xor3 } x \ y \ \text{cin} \# \text{rbl-plus } (\text{carry } x \ y \ \text{cin}) \ xs \ ys$
 $\text{rbl-plus } \text{cin } [] \ ys = []$
 $\text{rbl-plus } \text{cin } xs \ [] = []$
 ⟨proof⟩

lemma *rbl-word-plus*: $\text{rev } (\text{to-bl } (x + y)) = \text{rbl-plus } \text{False } (\text{rev } (\text{to-bl } x)) (\text{rev } (\text{to-bl } y))$
 ⟨proof⟩

definition $rbl-succ2\ b\ xs = (if\ b\ then\ rbl-succ\ xs\ else\ xs)$

lemma $rbl-succ2-simps$:

$rbl-succ2\ b\ [] = []$

$rbl-succ2\ b\ (x \# xs) = (b \neq x) \# rbl-succ2\ (x \wedge b)\ xs$

$\langle proof \rangle$

lemma $twos-complement$: $- x = word-succ\ (NOT\ x)$

$\langle proof \rangle$

lemma $rbl-word-neg$: $rev\ (to-bl\ (-\ x)) = rbl-succ2\ True\ (map\ Not\ (rev\ (to-bl\ x)))$

$\langle proof \rangle$

lemma $rbl-word-cat$:

$rev\ (to-bl\ (word-cat\ x\ y :: 'a::len0\ word)) =$

$takefill\ False\ (len-of\ TYPE('a))\ (rev\ (to-bl\ y) @ rev\ (to-bl\ x))$

$\langle proof \rangle$

lemma $rbl-word-slice$:

$rev\ (to-bl\ (slice\ n\ w :: 'a::len0\ word)) =$

$takefill\ False\ (len-of\ TYPE('a))\ (drop\ n\ (rev\ (to-bl\ w)))$

$\langle proof \rangle$

lemma $rbl-word-ucast$:

$rev\ (to-bl\ (ucast\ x :: 'a::len0\ word)) = takefill\ False\ (len-of\ TYPE('a))\ (rev\ (to-bl\ x))$

$\langle proof \rangle$

lemma $rbl-shifftl$:

$rev\ (to-bl\ (w << n)) = takefill\ False\ (size\ w)\ (replicate\ n\ False @ rev\ (to-bl\ w))$

$\langle proof \rangle$

lemma $rbl-shiftr$:

$rev\ (to-bl\ (w >> n)) = takefill\ False\ (size\ w)\ (drop\ n\ (rev\ (to-bl\ w)))$

$\langle proof \rangle$

definition $drop-nonempty\ v\ n\ xs = (if\ n < length\ xs\ then\ drop\ n\ xs\ else\ [last\ (v \# xs)])$

lemma $drop-nonempty-simps$:

$drop-nonempty\ v\ (Suc\ n)\ (x \# xs) = drop-nonempty\ x\ n\ xs$

$drop-nonempty\ v\ 0\ (x \# xs) = (x \# xs)$

$drop-nonempty\ v\ n\ [] = [v]$

$\langle proof \rangle$

definition $takefill-last\ x\ n\ xs = takefill\ (last\ (x \# xs))\ n\ xs$

lemma $takefill-last-simps$:

$takefill-last\ z\ (Suc\ n)\ (x\ \# \ xs) = x\ \# \ takefill-last\ x\ n\ xs$
 $takefill-last\ z\ 0\ xs = []$
 $takefill-last\ z\ n\ [] = replicate\ n\ z$
 <proof>

lemma *rbl-sshiftr*:

$rev\ (to-bl\ (w\ >>> n)) = takefill-last\ False\ (size\ w)\ (drop-nonempty\ False\ n\ (rev\ (to-bl\ w)))$
 <proof>

lemma *nth-word-of-int*:

$(word-of-int\ x :: 'a::len0\ word) !! n = (n < len-of\ TYPE('a) \wedge bin-nth\ x\ n)$
 <proof>

lemma *nth-scst*:

$(scst\ (x :: 'a::len\ word) :: 'b::len\ word) !! n =$
 $(n < len-of\ TYPE('b) \wedge$
 $(if\ n < len-of\ TYPE('a) - 1\ then\ x !! n$
 $else\ x !! (len-of\ TYPE('a) - 1)))$
 <proof>

lemma *rbl-word-scst*:

$rev\ (to-bl\ (scst\ x :: 'a::len\ word)) = takefill-last\ False\ (len-of\ TYPE('a))\ (rev\ (to-bl\ x))$
 <proof>

definition *rbl-mul* :: *bool list* $\Rightarrow$ *bool list* $\Rightarrow$ *bool list*

where $rbl-mul\ xs\ ys = foldr\ (\lambda x\ sm.\ rbl-plus\ False\ (map\ (op\ \wedge\ x)\ ys)\ (False\ \# \ sm))\ xs\ []$

lemma *rbl-mul-simps*:

$rbl-mul\ (x\ \# \ xs)\ ys = rbl-plus\ False\ (map\ (op\ \wedge\ x)\ ys)\ (False\ \# \ rbl-mul\ xs\ ys)$
 $rbl-mul\ []\ ys = []$
 <proof>

lemma *takefill-le2*: $length\ xs \leq n \implies takefill\ x\ m\ (takefill\ x\ n\ xs) = takefill\ x\ m\ xs$

<proof>

lemma *take-rbl-plus*: $\forall n\ b.\ take\ n\ (rbl-plus\ b\ xs\ ys) = rbl-plus\ b\ (take\ n\ xs)\ (take\ n\ ys)$

<proof>

lemma *word-rbl-mul-induct*:

$length\ xs \leq size\ y \implies$
 $rbl-mul\ xs\ (rev\ (to-bl\ y)) = take\ (length\ xs)\ (rev\ (to-bl\ (of-bl\ (rev\ xs) * y)))$
for $y :: 'a::len\ word$
 <proof>

lemma *rbl-word-mul*: $\text{rev } (\text{to-bl } (x * y)) = \text{rbl-mul } (\text{rev } (\text{to-bl } x)) (\text{rev } (\text{to-bl } y))$
for $x :: 'a::\text{len word}$
 $\langle \text{proof} \rangle$

Breaking up inequalities into bitlist properties.

definition

$\text{rev-bl-order } F \text{ } xs \text{ } ys =$
 $(\text{length } xs = \text{length } ys \wedge$
 $((xs = ys \wedge F)$
 $\vee (\exists n < \text{length } xs. \text{drop } (\text{Suc } n) \text{ } xs = \text{drop } (\text{Suc } n) \text{ } ys$
 $\wedge \neg xs ! n \wedge ys ! n)))$

lemma *rev-bl-order-simps*:

$\text{rev-bl-order } F [] [] = F$
 $\text{rev-bl-order } F (x \# xs) (y \# ys) = \text{rev-bl-order } ((y \wedge \neg x) \vee ((y \vee \neg x) \wedge F))$
 $xs \text{ } ys$
 $\langle \text{proof} \rangle$

lemma *rev-bl-order-rev-simp*:

$\text{length } xs = \text{length } ys \implies$
 $\text{rev-bl-order } F (xs @ [x]) (ys @ [y]) = ((y \wedge \neg x) \vee ((y \vee \neg x) \wedge \text{rev-bl-order } F$
 $xs \text{ } ys))$
 $\langle \text{proof} \rangle$

lemma *rev-bl-order-bl-to-bin*:

$\text{length } xs = \text{length } ys \implies$
 $\text{rev-bl-order } \text{True } xs \text{ } ys = (\text{bl-to-bin } (\text{rev } xs) \leq \text{bl-to-bin } (\text{rev } ys)) \wedge$
 $\text{rev-bl-order } \text{False } xs \text{ } ys = (\text{bl-to-bin } (\text{rev } xs) < \text{bl-to-bin } (\text{rev } ys))$
 $\langle \text{proof} \rangle$

lemma *word-le-rbl*: $x \leq y \longleftrightarrow \text{rev-bl-order } \text{True } (\text{rev } (\text{to-bl } x)) (\text{rev } (\text{to-bl } y))$
for $x \text{ } y :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *word-less-rbl*: $x < y \longleftrightarrow \text{rev-bl-order } \text{False } (\text{rev } (\text{to-bl } x)) (\text{rev } (\text{to-bl } y))$
for $x \text{ } y :: 'a::\text{len0 word}$
 $\langle \text{proof} \rangle$

lemma *word-sint-msb-eq*: $\text{sint } x = \text{uint } x - (\text{if } \text{msb } x \text{ then } 2^{\text{size } x} \text{ else } 0)$
 $\langle \text{proof} \rangle$

lemma *word-sle-msb-le*: $x \leq_s y \longleftrightarrow (\text{msb } y \longrightarrow \text{msb } x) \wedge ((\text{msb } x \wedge \neg \text{msb } y)$
 $\vee x \leq y)$
 $\langle \text{proof} \rangle$

lemma *word-sless-msb-less*: $x <_s y \longleftrightarrow (\text{msb } y \longrightarrow \text{msb } x) \wedge ((\text{msb } x \wedge \neg \text{msb } y)$
 $\vee x < y)$
 $\langle \text{proof} \rangle$

definition $\text{map-last } f \text{ } xs = (\text{if } xs = [] \text{ then } [] \text{ else } \text{butlast } xs @ [f \text{ (last } xs)])$

lemma map-last-simps :

$\text{map-last } f \text{ } [] = []$
 $\text{map-last } f \text{ } [x] = [f \text{ } x]$
 $\text{map-last } f \text{ } (x \# y \# zs) = x \# \text{map-last } f \text{ } (y \# zs)$
 $\langle \text{proof} \rangle$

lemma word-sle-rbl :

$x \leq_s y \iff \text{rev-bl-order True (map-last Not (rev (to-bl x))) (map-last Not (rev (to-bl y)))}$
 $\langle \text{proof} \rangle$

lemma word-sless-rbl :

$x <_s y \iff \text{rev-bl-order False (map-last Not (rev (to-bl x))) (map-last Not (rev (to-bl y)))}$
 $\langle \text{proof} \rangle$

Lemmas for unpacking $\text{rev (to-bl } n)$ for numerals n and also for irreducible values and expressions.

lemma $\text{rev-bin-to-bl-simps}$:

$\text{rev (bin-to-bl 0 } x) = []$
 $\text{rev (bin-to-bl (Suc } n) (\text{numeral (num.Bit0 } nm)) = False \# \text{rev (bin-to-bl } n (\text{numeral } nm))$
 $\text{rev (bin-to-bl (Suc } n) (\text{numeral (num.Bit1 } nm)) = True \# \text{rev (bin-to-bl } n (\text{numeral } nm))$
 $\text{rev (bin-to-bl (Suc } n) (\text{numeral (num.One)})) = True \# \text{replicate } n \text{ False}$
 $\text{rev (bin-to-bl (Suc } n) (- \text{numeral (num.Bit0 } nm)) = False \# \text{rev (bin-to-bl } n (- \text{numeral } nm))$
 $\text{rev (bin-to-bl (Suc } n) (- \text{numeral (num.Bit1 } nm)) =$
 $\text{True \# rev (bin-to-bl } n (- \text{numeral (nm + num.One)}))$
 $\text{rev (bin-to-bl (Suc } n) (- \text{numeral (num.One)})) = True \# \text{replicate } n \text{ True}$
 $\text{rev (bin-to-bl (Suc } n) (- \text{numeral (num.Bit0 } nm + \text{num.One}))) =$
 $\text{True \# rev (bin-to-bl } n (- \text{numeral (nm + num.One)}))$
 $\text{rev (bin-to-bl (Suc } n) (- \text{numeral (num.Bit1 } nm + \text{num.One}))) =$
 $\text{False \# rev (bin-to-bl } n (- \text{numeral (nm + num.One)}))$
 $\text{rev (bin-to-bl (Suc } n) (- \text{numeral (num.One + num.One)})) =$
 $\text{False \# rev (bin-to-bl } n (- \text{numeral num.One}))$
 $\langle \text{proof} \rangle$

lemma to-bl-upt : $\text{to-bl } x = \text{rev (map (op !! } x) [0 \text{ ..< size } x])$

$\langle \text{proof} \rangle$

lemma rev-to-bl-upt : $\text{rev (to-bl } x) = \text{map (op !! } x) [0 \text{ ..< size } x]$

$\langle \text{proof} \rangle$

lemma $\text{upt-eq-list-intros}$:

$j \leq i \implies [i \text{ ..< } j] = []$
 $i = x \implies x < j \implies [x + 1 \text{ ..< } j] = xs \implies [i \text{ ..< } j] = (x \# xs)$

$\langle proof \rangle$

10.30 Tactic definition

$\langle ML \rangle$

end

theory *Bit-Comparison*
imports *Bits-Int*
begin

lemma *AND-lower* [*simp*]:
 fixes $x :: int$ **and** $y :: int$
 assumes $0 \leq x$
 shows $0 \leq x \text{ AND } y$
 $\langle proof \rangle$

lemma *OR-lower* [*simp*]:
 fixes $x :: int$ **and** $y :: int$
 assumes $0 \leq x$ $0 \leq y$
 shows $0 \leq x \text{ OR } y$
 $\langle proof \rangle$

lemma *XOR-lower* [*simp*]:
 fixes $x :: int$ **and** $y :: int$
 assumes $0 \leq x$ $0 \leq y$
 shows $0 \leq x \text{ XOR } y$
 $\langle proof \rangle$

lemma *AND-upper1* [*simp*]:
 fixes $x :: int$ **and** $y :: int$
 assumes $0 \leq x$
 shows $x \text{ AND } y \leq x$
 $\langle proof \rangle$

lemmas *AND-upper1'* [*simp*] = *order-trans* [*OF AND-upper1*]
lemmas *AND-upper1''* [*simp*] = *order-le-less-trans* [*OF AND-upper1*]

lemma *AND-upper2* [*simp*]:
 fixes $x :: int$ **and** $y :: int$
 assumes $0 \leq y$
 shows $x \text{ AND } y \leq y$
 $\langle proof \rangle$

lemmas *AND-upper2'* [*simp*] = *order-trans* [*OF AND-upper2*]
lemmas *AND-upper2''* [*simp*] = *order-le-less-trans* [*OF AND-upper2*]

lemma *OR-upper*:
fixes $x :: \text{int}$ **and** $y :: \text{int}$
assumes $0 \leq x$ $x < 2^n$ $y < 2^n$
shows $x \text{ OR } y < 2^n$
 $\langle \text{proof} \rangle$

lemma *XOR-upper*:
fixes $x :: \text{int}$ **and** $y :: \text{int}$
assumes $0 \leq x$ $x < 2^n$ $y < 2^n$
shows $x \text{ XOR } y < 2^n$
 $\langle \text{proof} \rangle$

end

References

- [1] Jeremy Dawson. Isabelle theories for machine words. In Michael Goldsmith and Bill Roscoe, editors, *Seventh International Workshop on Automated Verification of Critical Systems (AVOCS'07)*, Electronic Notes in Theoretical Computer Science, page 15, Oxford, September 2007. Elsevier. to appear.